

Let $G = (V, E)$ be an undirected graph. Unless we say otherwise, a graph has no loops or parallel edges.

- (a) Prove that if $|V| \geq 2$ there are two distinct nodes u and v such that degree of u is equal to degree of v . Recall that the degree of a node x is the number of edges incident to x .
- (b) Prove that if G has at least one edge then there is a path between two distinct nodes u and v such that degree of u is equal to degree of v .

Solution: (a) Proof:

Let $|V| = k$, assume that, there's no two distinctive nodes with the same degree, the degree for each node, respectively, should be $0, 1, \dots, k$. Say for $u', v' \in V$, $\deg(u') = 0$ and $\deg(v') = k$, which means that u' is adjacent to no other nodes while v' is adjacent to all other nodes in the graph. It is impossible, so there must be two nodes u, v to make $\deg(u) = \deg(v)$.

As a result, for all $|V| \geq 2$, there are two nodes u, v to make $\deg(u) = \deg(v)$.

(b) Proof:

Since $|E| \geq 1$, the number of connected component is at least one and $|V| \geq 2$. And so we can arbitrarily choose one of the connected component in the graph without losing generality and drop all the other vertices and edges in the other components and still keep the structure of the one.

So we get the new graph $G' \subseteq G$. Recall that we've just proved that for all $|V| \geq 2$, there are two nodes u, v to make $\deg(u) = \deg(v)$. We know such $u, v \in G' \subseteq G$ in the same connected exists. So naturally, there is a path between u, v .

As a result, if G has at least one edge then there is a path between two distinct nodes u and v such that degree of u is equal to degree of v . ■