

## 1 Compactness

We assume there are countably many sentence symbols.

A finite set  $\{\varphi_1, \dots, \varphi_n\}$  of wff's is satisfiable iff  $\varphi_1 \wedge \dots \wedge \varphi_n$  is not a contradiction.

We say that  $\{\varphi_1, \dots, \varphi_n\}$  is consistent when it's satisfiable.

We say that  $\mathcal{T}$  of wff's is consistent when  $\mathcal{T}$  is finitely satisfiable.

i.e.  $\varphi_1 \wedge \dots \wedge \varphi_n$  is not contradiction for any  $\varphi_i \in \mathcal{T}$ .

Let  $\mathcal{T}$  be a consistent set of wff's. We extend  $\mathcal{T}$  to a maximal consistent set  $\Delta$  of wff's.

**Lemma 1.0.1** *Let  $S$  be a consistent set of wff's, and let  $\varphi$  a wff. Then  $S \cup \{\varphi\}$  or  $S \cup \{\neg\varphi\}$  is consistent.*

**Lemma 1.0.2** *Let  $S$  be a consistent set of wff's, then the following are equivalent:*

1.  $S$  is a maximal consistent set of wff's
2. For any wff  $\varphi$  either  $\varphi \in S$  or  $\neg\varphi \in S$