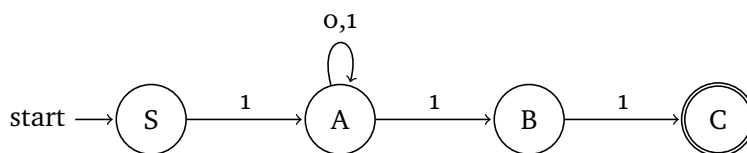


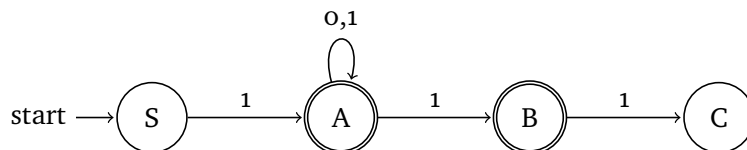
- (a) Describe a concrete example of a machine  $N$  to show that  $L(N_{\text{comp}}) \neq \overline{L(N)}$ .  
You need to explain for your machine  $N$  what  $\overline{L(N)}$  and  $L(N_{\text{comp}})$  are.
- (b) Define an NFA that accepts  $\overline{L(N)} - L(N_{\text{comp}})$ , and explain how it works.
- (c) Define an NFA that accepts  $L(N_{\text{comp}}) - \overline{L(N)}$ , and explain how it works.

**Solution:**

(a)



**Figure 1.**



**Figure 2.**

Let  $L(N) = L(1(0+1)^*11)$ , as in figure 1, then  $\overline{L(N)}$  does not end with 11. By swapping the accept and non-accept states which is shows as figure 2,  $L(N_{\text{comp}})$  accepts 011, however,  $\overline{L(N)}$  does not end with 11, so that  $L(N_{\text{comp}}) \neq \overline{L(N)}$ . ■

**Lemma 1.** Let  $A, B \subseteq Q$  and  $\delta(S, w_1) = \{A, B\}$ ,  $\delta(S, w_2) = \{A\}$ ,  $\delta(S, w_3) = \{B\}$ ,  $\delta(S, w_4) = \emptyset$ . If  $N$  accepts both  $w_1$  and  $w_2$ , then  $N'$  that accepts  $\overline{L(N)}$  accepts both  $w_3$  and  $w_4$ , and  $N_{\text{comp}}$  accepts both  $w_1$  and  $w_3$ . Denote  $L_1 = \{w : \delta(S, w) = \{A, B\}\}$ ,  $L_2 = \{w : \delta(S, w) = \{A\}\}$ ,  $L_3 = \{w : \delta(S, w) = \{B\}\}$ ,  $L_4 = \{w : \delta(S, w) = \emptyset\}$ .

**Proof:** This lemma is true by the definition of NFA. ■

(b) By Lemma 1,

$$\overline{L(N)} - L(N_{\text{comp}}) = L_3 + L_4 - (L_1 + L_3) = L_4 = \emptyset$$

(c) By Lemma 1,

$$L(N_{\text{comp}}) - \overline{L(N)} = L_1 + L_3 - (L_3 + L_4) = L_1 = A + B$$

■