Compiling for Parallelism & Locality

Last time

- Dynamic compilation
- End of lectures on low level optimizations

Today

- Data dependences and loops
- Parallelism and locality

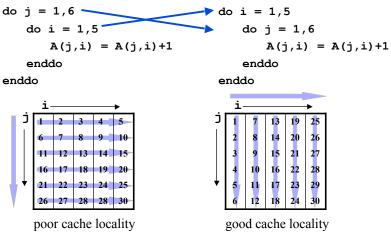
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Example 1: Loop Permutation for Improved Locality

Sample code: Assume Fortran's Column Major Order array layout



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Example 2: Parallelization

Can we parallelize the following loops?

do i = 1,100

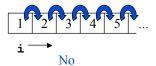
$$A(i) = A(i)+1$$

enddo



do
$$i = 1,100$$

A(i) = A(i-1)+1
enddo



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Data Dependences

Recall

- A data dependence defines ordering relationship two between statements
- In executing statements, data dependences must be respected to preserve correctness

Example

$$s_1$$
 a := 5;
 s_2 b := a + 1;
 s_3 a := 6;
 s_1 a := 5;
 s_3 a := 6;
 s_2 b := a + 1;

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Data Dependences and Loops

How do we identify dependences in loops?

do i = 1,5

$$A(i) = A(i-1)+1$$

enddo

Simple view

- Imagine that all loops are fully unrolled
- Examine data dependences as before

Problems

- Impractical
- -Lose loop structure

A(1) = A(0)+1

$$A(2) = A(1)+1$$

$$A(3) = A(2)+1$$

$$A(4) = A(3) + 1$$

$$A(5) = A(4)+1$$

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Dependence Analysis for Loops

Big picture

- To improve data locality and parallelism we often focus on loops
- To transform loops, we must understand data dependences in loops
- Since we can't represent all iterations of a loop, we need some abstractions
- The basic question: does a transformation preserve all dependences?

Today and Next Time

- Basic abstractions and machinery

Wednesday

- Its application to loop transformations

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Data Dependence Terminology

We say statement s2 depends on s1

True (flow) dependence: s₁ writes memory that s₂ later reads
 Anti-dependence: s₁ reads memory that s₂ later writes
 Output dependences: s₁ writes memory that s₂ later writes
 Input dependences: s₁ reads memory that s₂ later reads

Notation: $s_1 \delta s_2$

- $-s_1$ is called the **source** of the dependence
- s_2 is called the **sink** or **target**
- s₁ must be executed before s₂

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Dependences and Loops

Loop-independent dependences

Loop-carried dependences

do
$$i = 1,100$$

$$A(i) = B(i)+1$$

$$C(i) = A(i-1)*2$$
enddo

Dependences that cross loop iterations

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Iteration Spaces

Idea

- Explicitly represent the iterations of a loop nest

Example

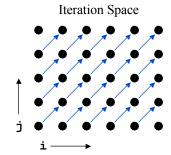
```
do i = 1,6

do j = 1,5

A(i,j) = A(i-1,j-1)+1

enddo

enddo
```



Iteration Space

- A set of tuples that represents the iterations of a loop
- Can visualize the dependences in an iteration space

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Protein String Matching Example

```
q = k_1
r = k_2
score = 0
for i=0 to n1-1
 h[0,-1] = p[0,-1] = 0
  f[0,-1] = -q
  for j=0 to n0-1
    f[i,j] = max(f[i,j-1],h[i,j-1]-q)-r
    EE[i,j] = max(EE[i-1,j],HH[i-1,j],-q)-r
    h[i,j] = p[i,j-1] + pam2[aa1[i],aa0[j]]
    h[i,j] = max( max(0,EE[i,j]), max(f[i,j],h[i,j])
    p[i,j] = HH[i-1,j]
    HH[i,j] = h[i,j]
    score[i,j] = max(score[i,j-1],h[i,j])
  endfor
endfor
return score[n1-1,n0-1]
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```

Distance Vectors

Idea

- Concisely describe dependence relationships between iterations of an iteration space
- For each dimension of an iteration space, the distance is the number of iterations between accesses to the same memory location

Definition

$$- \mathbf{v} = \mathbf{i}^{\mathrm{T}} - \mathbf{i}^{\mathrm{S}}$$

Example

Distance Vector: (2,1) inner loop
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j i ——

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Distance Vectors and Loop Transformations

Idea

- Any transformation we perform on the loop must respect the dependences

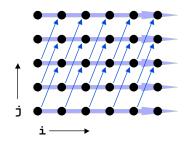
Example

do i = 1,6
do j = 1,5

$$A(i,j) = A(i-1,j-2)+1$$

enddo
enddo

la ama?



Can we permute the i and j loops?

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Distance Vectors and Loop Transformations

Idea

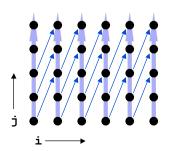
- Any transformation we perform on the loop must respect the dependences

Example

do j = 1,5
do i = 1,6

$$A(i,j) = A(i-1,j-2)+1$$

enddo
enddo



Can we permute the i and j loops?

- Yes

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