

# Lecture 1

Alex Hassett  
Applied Analysis

May 29, 2025

**Theorem 1.1.** Let  $p$ ,  $q$ , and  $r$  be logical statements (meaning they can either be true ( $T$ ) or false ( $F$ )). The logical equivalences are as follows

1. Identity laws

- (a)  $p \wedge T \equiv p$
- (b)  $p \vee F \equiv p$

2. Domination laws

- (a)  $p \vee T \equiv T$
- (b)  $p \wedge F \equiv F$

3. Idempotent laws

- (a)  $p \vee p \equiv p$
- (b)  $p \wedge p \equiv p$

4. Double negation law

$$\neg(\neg p) \equiv p$$

5. Commutative laws

- (a)  $p \vee q \equiv q \vee p$
- (b)  $p \wedge q \equiv q \wedge p$

6. Associative laws

- (a)  $(p \vee q) \vee r \equiv p \vee (q \vee r)$
- (b)  $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$

7. Distributive laws

- (a)  $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$
- (b)  $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

8. De Morgan's laws

- (a)  $\neg(p \wedge q) \equiv \neg p \vee \neg q$
- (b)  $\neg(p \vee q) \equiv \neg p \wedge \neg q$

9. Absorption laws

- (a)  $p \vee (p \wedge q) \equiv p$
- (b)  $p \wedge (p \vee q) \equiv p$

10. Negation laws

- (a)  $p \vee \neg p \equiv T$
- (b)  $p \wedge \neg p \equiv F$

**Theorem 1.2.** Let  $p$  and  $q$  be logical statements. Then  $p \implies q$ , read "p implies q", is the statement "if  $p$ , then  $q$ ". Note that  $p \implies q \equiv \neg p \vee q$ .

**Corollary 1.3.** Let  $p$  and  $q$  be logical statements. Let  $p \equiv r \implies t$  where  $r$  and  $t$  are logical statements. Then the following formal English statements are all logically equivalent to  $p \implies q$ .

1. If  $p$ , then  $q$ .
2. Suppose  $p$ . Then  $q$ .
3.  $p$  implies  $q$ .
4. Suppose  $p$  such that  $q$ .
5. Let  $p$  be such that  $q$ .
6. If  $r$  such that  $t$ , then  $q$ .

Note that the above equivalencies are only a few among the hundreds of equivalent ways to express the statement  $p \implies q$  in formal English.

**Definition 1.4.** Let  $S(p, q)$  be the statement  $p \implies q$ , i.e. let  $S(p, q) \equiv p \implies q$ . To non-trivially prove that  $S(p, q)$  is true,  $S(p, q) \equiv T$ , one must suppose  $p \equiv T$ , then show that  $q \equiv T$ . In formal English, this would be written "suppose  $p$  is true, then show that  $q$  is also true."

**Definition 1.5.** A set is a collection of objects. For a set  $X$ , we say that  $x$  is an element of  $X$  by writing  $x \in X$ , read " $x$  in  $X$ ". We say  $x$  is not an element of  $X$  by writing  $x \notin X$ , read " $x$  not in  $X$ ".

**Definition 1.6.** A set  $Y$  is a subset of a set  $X$ , written  $Y \subset X$ , if for every  $y \in Y$ , we have that  $y \in X$ . We say that  $Y$  is a proper subset of  $X$  if  $Y \subset X$  and  $Y \neq X$ .

**Theorem 1.7.** Let  $X$  be a set. Then  $X \subset X$ .

*Proof.* If  $X = \emptyset$ , then  $X \subset X$  because  $\emptyset \subset \emptyset$ . Now suppose  $X \neq \emptyset$  and let  $x \in X$ . Then  $x \in X$ . Therefore  $X \subset X$ .  $\square$

**Definition 1.8.** Let  $\zeta$ , read "zeta", be a logical statement. Then the negation of  $\zeta$ , written  $\neg\zeta$  is the logical opposite of  $\zeta$ , i.e. its negation. If  $\zeta \equiv T$ , then  $\neg\zeta \equiv F$  and vice versa.

**Definition 1.9.** Let  $x$  and  $y$  be defined in a first-order language. Then  $x$  is said to be defined equivalently to  $y$ , written  $x \approx y$ , if the definition of  $x$  is logically equivalent to the definition of  $y$ . Some examples of this are

1. Let  $A$  be a set. Let  $x \in A$  and let  $y \in A$ . Then  $x \approx y$ .
2. Let  $x$  be the statement "if  $p$ , then  $q$ ". Let  $y$  be the statement "p implies q". Then  $x \approx y$ . Note that we can also say that  $x \equiv y$ , because  $x$  and  $y$  are propositional (logical) statements.

**Corollary 1.10.** Let  $x$  and  $y$  be defined equivalently. Then  $x$  is equivalent to  $y$ , but  $x$  does not necessarily equal  $y$ .