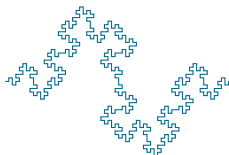


Presentation Title

Cameron Bracken
Humboldt State University



June 18, 2014

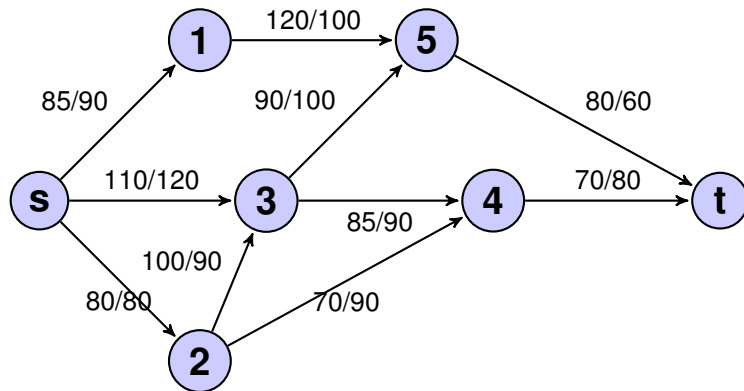
INTRODUCTION

INTRODUCTION

Greedy Heuristic Algorithm

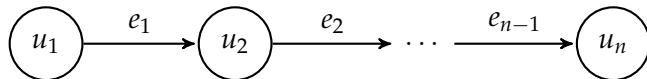
Optimal solution to a path

EXAMPLE



PHYSICAL SYSTEM

Modeling the physical system, of an EV and a path.



► Path

- + Charging stations, with charging rate ($R_{CH}(u_i)$)
- + Road segments, with speed limit ($v_{min}(e_i), v_{max}(e_i)$) and distance ($D(e_i)$)

► EV

- + Driving consumes energy accordingly to the speed of the EV, defined by: ($R_{CO}(e_i)$)
- + Further two constants from the EV are important to model, namely, battery capacity (B_{max}) and initial battery (B_{cur})

OPTIMISATION PROBLEM

Formulating a optimization problem, which when solved will yield a optimal solution.

- ▶ Objective: Move from u_1 to u_n using minimum time .
 - + Time can be used driving or charging.
 - min: $\sum_{i=1}^{n-1} \left(\frac{D(e_i)}{v_{e_i}} + CT_{u_i} \right)$
- ▶ Physical constraints:
 - + Each edge must be driven at a speed within the speed limit:
 - $\forall_{i \in 1 \dots n-1} : v_{min}(e_i) \leq v_{e_i} \leq v_{max}(e_i)$
 - + Time can only be positive.
 - $\forall_{i \in 1 \dots n} : 0 \leq CT_{u_i}$
 - + The energy is the battery must always be between 0 and B_{max}

BATTERY CONSTRAINT

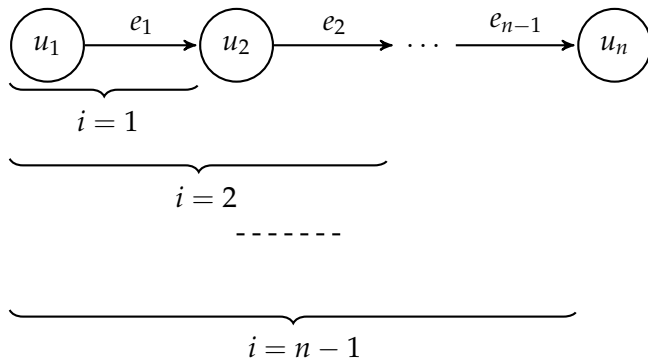
The battery constraint of the optimization problem can be split into two parts

- ▶ No road segment can be passed without having the required energy
- ▶ No overcharging at any charging station.

Energy can be..

- ▶ Spend: $\forall_{i \in 1 \dots n-1} : ES(e_i) = D(e_i) \times R_{CO}(v_{e_i})$
- ▶ Acquired: $\forall_{i \in 1 \dots n} : EA(u_i) = R_{CH}(u_i) \times CT_{u_i}$
- ▶ Already in the battery: B_{cur}

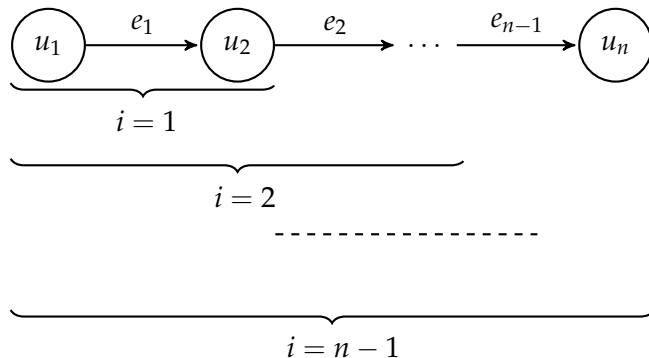
No road segment can be passed without having the required energy



► $\forall_{i \in 1 \dots n-1} : 0 \leq B_{cur} + \sum_{j=1}^i EA(u_j) - \sum_{j=1}^i ES(e_j) \leq B_{max}$

BATTERY CONSTRAINT

No overcharging at any charging station.



$$\blacktriangleright \forall_{i \in 1 \dots n-1} : 0 \leq B_{cur} + \sum_{j=1}^{i+1} EA(u_j) - \sum_{j=1}^i ES(e_j) \leq B_{max}$$

LINEAR PROGRAMMING

NP-complete problem.

Linearization and linear programming for approximate solution.

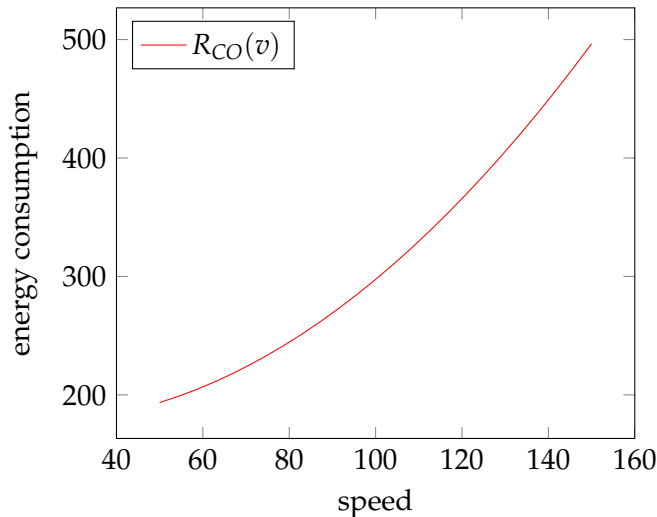
Two functions of the optimization problem are non linear functions.

- ▶ Consumption rate ($R_{CO}(v_{e_i})$)
- ▶ Driving time ($\frac{D(e_i)}{v_{e_i}}$)

LINEARIZATION EXAMPLE

Function for energy consumption before linearization.

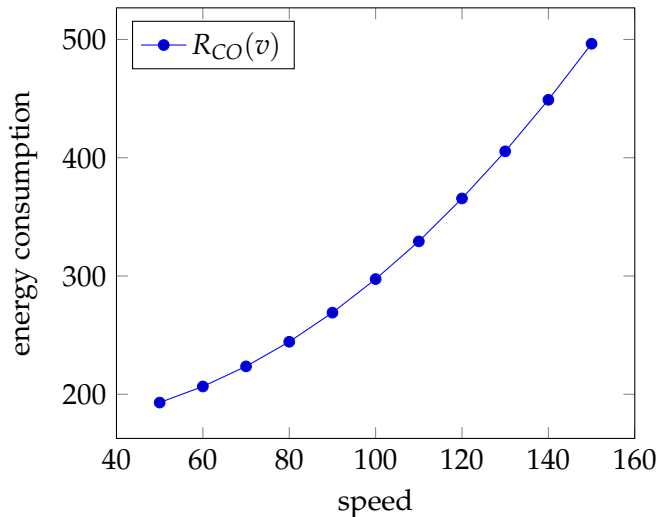
$$R_{CO}(v) = 0.019 * x^2 - 0.770 * x + 184.4$$



LINEARIZATION EXAMPLE

Function for energy consumption after linearization.

$$R_{CO}(v) = 0.019 * x^2 - 0.770 * x + 184.4$$



LINEARIZATION

- ▶ For all linear function their slope and the y-intercept is precomputed.
- ▶ For every edge in the path exactly one line segment needs to be chosen. Thus a binary matrix is introduced of size $n \times m$, where n = edges in the path and m = linear pieces of each line.