

Step 1) Arrouge the network from the largest to the smallest.

Step 2) Select the largest N/w.

Step 3) Select the subject of largest n/w. dest step beken the last n/w.

LAN-L(58 hus) iv) Link A" (2 host) ii) LAN-2(27 hash) V) LINKIS (men) 11) LAN-1 (12 hook) VI) LANK C (---)

				<u> </u>	1/16	32 /	164	128	256
45-27	Subject	256 12	8 64	32	16	8	4	2	1
	Host	1/24 /2	5 /26/	127	1/28	120	30	/31	/32
	Misk		LANI	LAN2	LANS	Li	* A/B/C		-> Canv

assign if address to gry hist becoz 2 mlJeures olways never not for Networking & Directal Browledge

y Gren Rouge 174,168.10.0/24

Hetwark 19	Subnet Mark	Host	Newspale
174.148.10.0	(255.255.255.192)	64-2	LAN-1
174.168.10.64			unuxo
174.168.10.128	_4		
174.168.10.192			

53)

			1
Network Id	Gubret Mask	Hot	Netwood
174.168.10.0	/26	64-2	LAN-L
174-18.10.64	129	32-2	LAN-2
174.168.10.96	/27	3.2-2	unuxd
174.168.10.128	/26	64-2	المستعرب المستعرب
			· ·
92.168.10.192	/26 -		

S=4)

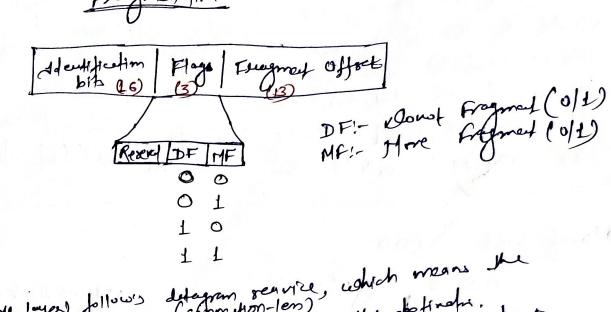
174.168.10:36	128	16-2	LAN-3
174.18.10.112	128	15-2	unered

3=9

179.158.20.212 /30 4-2 Line	A
110 /51	
179.18810.14 /30 4-2 Line	B
174.163.10.120 /30 4-2 Live	(C
174.168.10.124 /30 4-2 Unu	sef]

& & com do any no. of subnettry of would.

Fregnerform



It IP I uper (Network I uper) follows detagram seawice, which wears the definder. (connection-len) the definder. de detagram anough follows any moste to reach the destination. It designed to death the destination is done, each time header to death of each frame.

a) A datagram of length 5000 Bytes with 20 Bytes of headen in it greathed a geortes. The sworter has to forward the sestagram on the link whore MTU is forward the song fragments the rocke has to
Too bytes. How many fragments the rocke has to
do? Wetenime the total leight of each fragmental found,
the M bit and the fragmentation offset. 60/n: 5000 (20+4980), [4980] 2 8 ceiling value

FP2 FP3 FP4 FP5 FP6 FP7 FP8
(680+20) (680+20) (680+20) (680+20) (680+20) (680+20) LLLL FP1 85 170° 255 340 425 510 595 rlbit Fragmath O offret

Mo. of hosb = 500 $n \sim 3$ (2n = 500)Network Mark = 258.258.254.0 (Addres Moste) 171. 43.16.37/23 171.43, 0001000 7 171.43.16,0 171.43.19.255 Dut Brown all. 7 512-2 - 510 The anythole, 500 year, 10 remaining 1/25. > 2 ibis reserved 1.2.3.4/23 1. 211111 1.2.3.258

Classless Addressing CIDR (Classless Intendent Routing) Q. b. G. xd/n prefir leyt G1:-46.82.1.7/8 31.14.87.21/12 0 1010010 CLAR No. of networks = 2 129 4 2 4 2 1 No. of hosts in a network = 232-n Rules: 1) All ip addresses should be contingues there? Demand of the address should be in the form of 2n (pize of the block) First i/p allows in the block should be evenly divisible by the size of the block.

Out can be charged by pooley the mist

n-least significant bits to be of the wis is The cheek this is Grouple: 74.10,7.32 4774.10,7.47 B 74.10.732 = @~ 0- 16 ip addams = 24 74.10.7. 0010000 & Determe the searce of the adjusters of the 171.43.16.37 /29 001,00101 171.43.18.32 70700100 171.43, 11, 63 2111

Representation com be done by classing any &p-adless from the energy 171.43.10.60/27 & Entredy information from om Addens 1) No. of addresses in the block of N = 2 To find the tist oddens, keep he is refined bits of to ob.

and sed the (32-n) rightmust bits all to 0/3. (32-n) nothing bib all to 1).

(32-n) nothing bib all to 1). so Divily a n/w mpo sub-nepusals. y consider class éci ilp allers (oste from studets)
218.13.42.0 related t LLLLLLL 218.13.42.128 => subret id 1 2nd 218.13.42.0 -218.13.42.127 subject it of Double Broaden and 1. 218. 13. 42.200 Direct broader 258.13.41.0