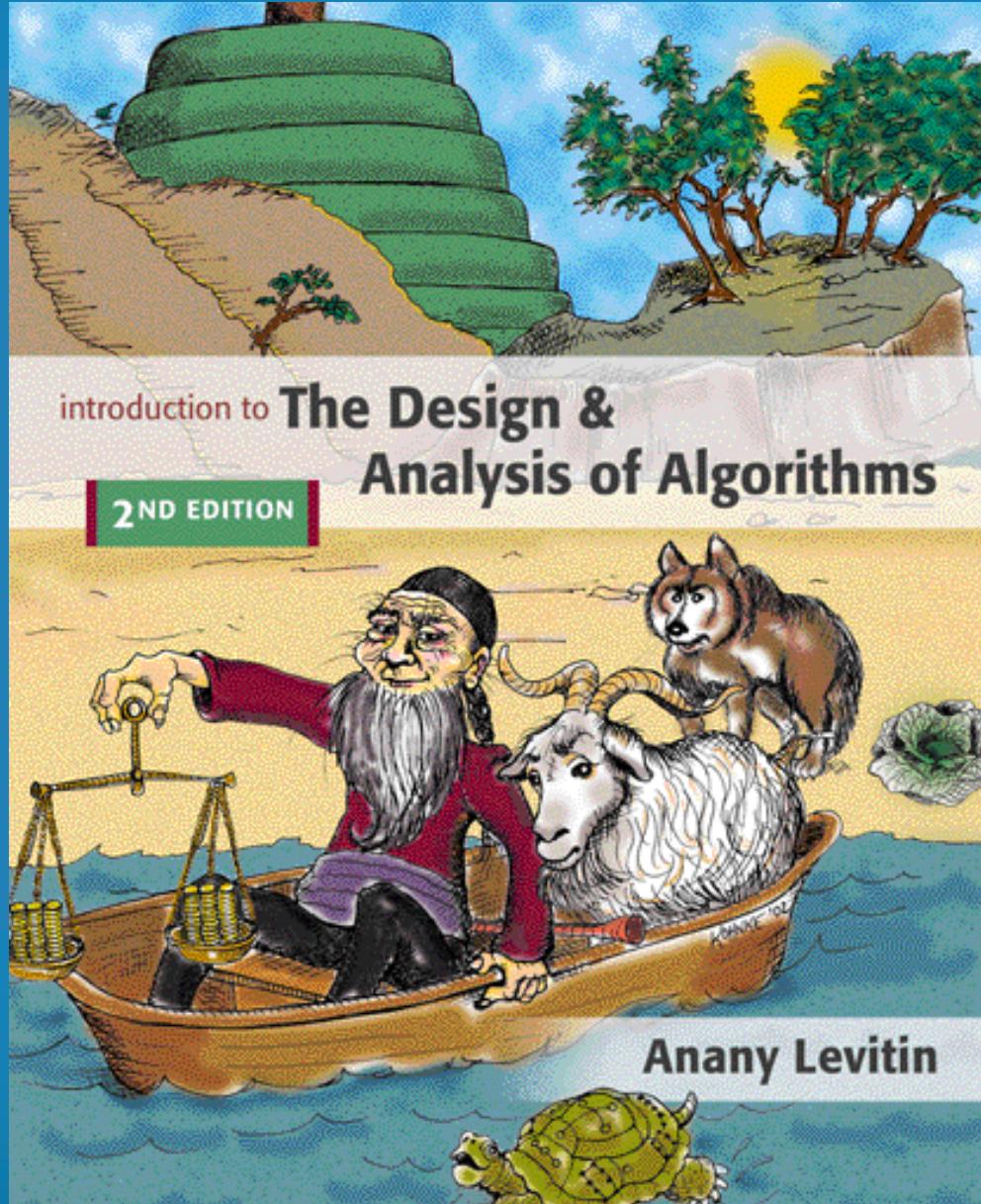


Chapter 1

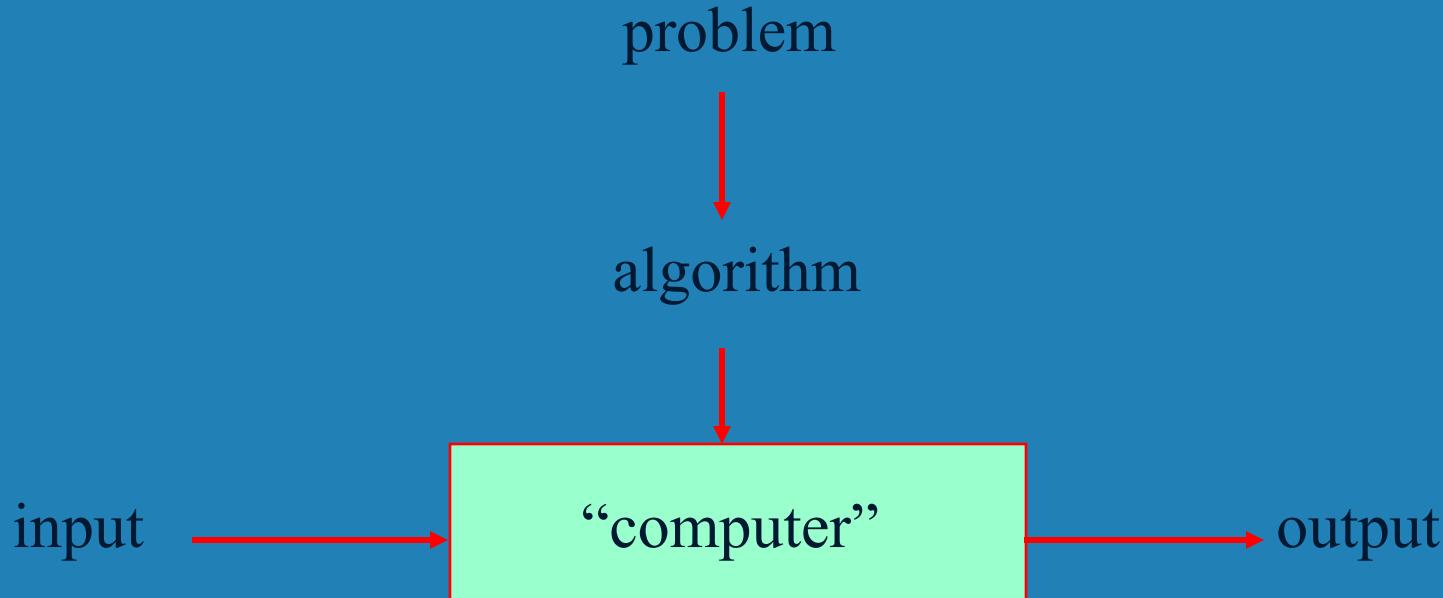
Introduction



What is an algorithm?



An **algorithm** is a sequence of unambiguous instructions for solving a problem, i.e., for obtaining a required output for any legitimate input in a finite amount of time.



Algorithm



- ➊ An algorithm is a sequence of unambiguous instructions for solving a problem, i.e., for obtaining a required output for any legitimate input in a finite amount of time.

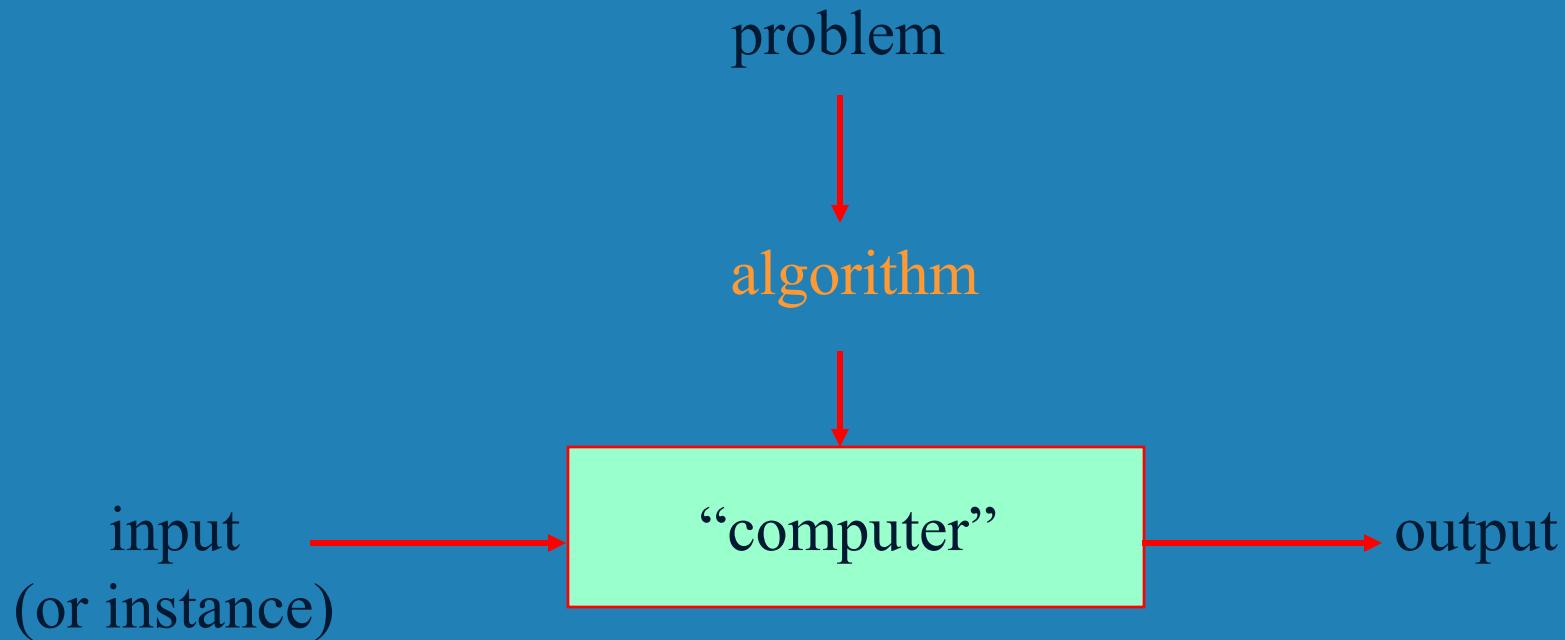
- Can be represented various forms
- Unambiguity/clearness
- Effectiveness
- Finiteness/termination
- Correctness

Historical Perspective



- Euclid's algorithm for finding the greatest common divisor
- Muhammad ibn Musa al-Khwarizmi – 9th century mathematician
www.lib.virginia.edu/science/parshall/khwariz.html

Notion of algorithm and problem



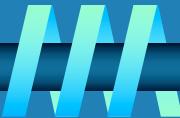
algebraic solution
(different from a conventional solution)

Example of computational problem: sorting



- **Statement of problem:**
 - *Input:* A sequence of n numbers $\langle a_1, a_2, \dots, a_n \rangle$
 - *Output:* A reordering of the input sequence $\langle a'_1, a'_2, \dots, a'_n \rangle$ so that $a'_i \leq a'_j$ whenever $i < j$
- **Instance:** The sequence $\langle 5, 3, 2, 8, 3 \rangle$
- **Algorithms:**
 - Selection sort
 - Insertion sort
 - Merge sort
 - (many others)

Selection Sort



- **Input:** array $a[1], \dots, a[n]$
- **Output:** array a sorted in non-decreasing order
- **Algorithm:**

```
for  $i=1$  to  $n$ 
```

```
    swap  $a[i]$  with smallest of  $a[i], \dots, a[n]$ 
```

- Is this unambiguous? Effective?
- See also pseudocode, section 3.1

Some Well-known Computational Problems



- **Sorting**
- **Searching**
- **Shortest paths in a graph**
- **Minimum spanning tree**
- **Primality testing**
- **Traveling salesman problem**
- **Knapsack problem**
- **Chess**
- **Towers of Hanoi**
- **Program termination**

Some of these problems don't have efficient algorithms, or algorithms at all!

Basic Issues Related to Algorithms



- How to design algorithms
- How to express algorithms
- Proving correctness
- Efficiency (or complexity) analysis
 - Theoretical analysis
 - Empirical analysis
- Optimality

Algorithm design strategies



- **Brute force**
- **Divide and conquer**
- **Decrease and conquer**
- **Transform and conquer**
- **Greedy approach**
- **Dynamic programming**
- **Backtracking and branch-and-bound**
- **Space and time tradeoffs**

Analysis of Algorithms



- ➊ How good is the algorithm?

- Correctness
- Time efficiency
- Space efficiency

- ➋ Does there exist a better algorithm?

- Lower bounds
- Optimality

What is an algorithm?



- **Recipe, process, method, technique, procedure, routine,... with the following requirements:**

1. Finiteness

- terminates after a finite number of steps

2. Definiteness

- rigorously and unambiguously specified

3. Clearly specified input

- valid inputs are clearly specified

4. Clearly specified/expected output

- can be proved to produce the correct output given a valid input

5. Effectiveness

- steps are sufficiently simple and basic

Why study algorithms?



- ➊ **Theoretical importance**

- **the core of computer science**

- ➋ **Practical importance**

- **A practitioner's toolkit of known algorithms**
- **Framework for designing and analyzing algorithms for new problems**

Example: Google's PageRank Technology

Euclid's Algorithm



Problem: Find $\gcd(m,n)$, the greatest common divisor of two nonnegative, not both zero integers m and n

Examples: $\gcd(60,24) = 12$, $\gcd(60,0) = 60$, $\gcd(0,0) = ?$

Euclid's algorithm is based on repeated application of equality

$$\gcd(m,n) = \gcd(n, m \bmod n)$$

until the second number becomes 0, which makes the problem trivial.

Example: $\gcd(60,24) = \gcd(24,12) = \gcd(12,0) = 12$

Two descriptions of Euclid's algorithm



- Step 1** If $n = 0$, return m and stop; otherwise go to Step 2
- Step 2** Divide m by n and assign the value of the remainder to r
- Step 3** Assign the value of n to m and the value of r to n . Go to Step 1.

while $n \neq 0$ **do**

$r \leftarrow m \bmod n$

$m \leftarrow n$

$n \leftarrow r$

return m

Other methods for computing $\gcd(m,n)$



Consecutive integer checking algorithm

Step 1 Assign the value of $\min\{m,n\}$ to t

Step 2 Divide m by t . If the remainder is 0, go to Step 3;
otherwise, go to Step 4

Step 3 Divide n by t . If the remainder is 0, return t and stop;
otherwise, go to Step 4

Step 4 Decrease t by 1 and go to Step 2

Is this slower than Euclid's algorithm?

How much slower?

$O(n)$, if $n \leq m$, vs $O(\log n)$

Other methods for $\text{gcd}(m,n)$ [cont.]



Middle-school procedure

Step 1 Find the prime factorization of m

Step 2 Find the prime factorization of n

Step 3 Find all the common prime factors

**Step 4 Compute the product of all the common prime factors
and return it as $\text{gcd}(m,n)$**

Is this an algorithm?

How efficient is it?

Time complexity: $O(\sqrt{n})$

Sieve of Eratosthenes



Input: Integer $n \geq 2$

Output: List of primes less than or equal to n

for $p \leftarrow 2$ **to** n **do** $A[p] \leftarrow p$

for $p \leftarrow 2$ **to** n **do**

if $A[p] \neq 0$ // p hasn't been previously eliminated from the list

$j \leftarrow p * p$

while $j \leq n$ **do**

$A[j] \leftarrow 0$ //mark element as eliminated

$j \leftarrow j + p$

Example: 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20

Time complexity: $O(n)$

Two main issues related to algorithms



- ➊ How to design algorithms
- ➋ How to analyze algorithm efficiency



Algorithm design techniques/strategies



- **Brute force**
- **Divide and conquer**
- **Decrease and conquer**
- **Transform and conquer**
- **Space and time tradeoffs**
- **Greedy approach**
- **Dynamic programming**
- **Iterative improvement**
- **Backtracking**
- **Branch and bound**

Analysis of algorithms



- How good is the algorithm?
 - time efficiency
 - space efficiency
 - correctness ignored in this course

- Does there exist a better algorithm?
 - lower bounds
 - optimality

Important problem types



- sorting
- searching
- string processing
- graph problems
- combinatorial problems
- geometric problems
- numerical problems



Sorting (I)



- **Rearrange the items of a given list in ascending order.**
 - Input: A sequence of n numbers $\langle a_1, a_2, \dots, a_n \rangle$
 - Output: A reordering $\langle a'_1, a'_2, \dots, a'_n \rangle$ of the input sequence such that $a'_1 \leq a'_2 \leq \dots \leq a'_n$.
- **Why sorting?**
 - Help searching
 - Algorithms often use sorting as a key subroutine.
- **Sorting key**
 - A specially chosen piece of information used to guide sorting. E.g., sort student records by names.

Sorting (II)



- **Examples of sorting algorithms**
 - Selection sort
 - Bubble sort
 - Insertion sort
 - Merge sort
 - Heap sort ...
- **Evaluate sorting algorithm complexity: the number of key comparisons.**
- **Two properties**
 - Stability: A sorting algorithm is called stable if it preserves the relative order of any two equal elements in its input.
 - In place : A sorting algorithm is in place if it does not require extra memory, except, possibly for a few memory units.

Selection Sort



Algorithm *SelectionSort(A[0..n-1])*

//The algorithm sorts a given array by selection sort

//Input: An array A[0..n-1] of orderable elements

//Output: Array A[0..n-1] sorted in ascending order

for i \leftarrow 0 to n – 2 do

 min \leftarrow i

 for j \leftarrow i + 1 to n – 1 do

 if A[j] < A[min]

 min \leftarrow j

 swap A[i] and A[min]

Searching



- Find a given value, called a search key, in a given set.
- Examples of searching algorithms
 - Sequential search
 - Binary search ...

Input: sorted array $a_i < \dots < a_j$ and key x ;

$m \leftarrow (i+j)/2$;

while $i < j$ and $x \neq a_m$ do

if $x < a_m$ then $j \leftarrow m-1$

else $i \leftarrow m+1$;

if $x = a_m$ then output a_m ;

Time: $O(\log n)$

String Processing



- A string is a sequence of characters from an alphabet.
- Text strings: letters, numbers, and special characters.
- String matching: searching for a given word/pattern in a text.

Examples:

- (i) searching for a word or phrase on WWW or in a Word document
- (ii) searching for a short read in the reference genomic sequence

Graph Problems



- **Informal definition**
 - A graph is a collection of points called **vertices**, some of which are connected by line segments called **edges**.
- **Modeling real-life problems**
 - Modeling WWW
 - Communication networks
 - Project scheduling ...
- **Examples of graph algorithms**
 - Graph traversal algorithms
 - Shortest-path algorithms
 - Topological sorting

Fundamental data structures



- **list**
 - **array**
 - **linked list**
 - **string**
- **stack**
- **queue**
- **priority queue/heap**
- **graph**
- **tree and binary tree**
- **set and dictionary**

Linear Data Structures



Arrays

- A sequence of n items of the same data type that are stored contiguously in computer memory and made accessible by specifying a value of the array's index.

Linked List

- A sequence of zero or more nodes each containing two kinds of information: some data and one or more links called pointers to other nodes of the linked list.
- Singly linked list (next pointer)
- Doubly linked list (next + previous pointers)



Arrays

- fixed length (need preliminary reservation of memory)
- contiguous memory locations
- direct access
- Insert/delete

Linked Lists

- dynamic length
- arbitrary memory locations
- access by following links
- Insert/delete

Stacks and Queues



● **Stacks**

- A stack of plates
 - insertion/deletion can be done only at the top.
 - LIFO
- Two operations (push and pop)

● **Queues**

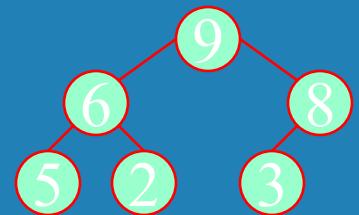
- A queue of customers waiting for services
 - Insertion/enqueue from the rear and deletion/dequeue from the front.
 - FIFO
- Two operations (enqueue and dequeue)

Priority Queue and Heap



Priority queues (implemented using heaps)

- A data structure for maintaining a **set** of elements, each associated with a key/priority, with the following operations
 - Finding the element with the highest priority
 - Deleting the element with the highest priority
 - Inserting a new element
- Scheduling jobs on a shared computer



Graphs



● Formal definition

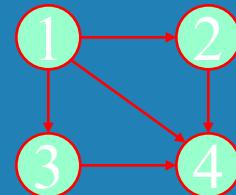
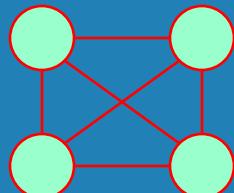
- A graph $G = \langle V, E \rangle$ is defined by a pair of two sets: a finite set V of items called vertices and a set E of vertex pairs called edges.

● Undirected and directed graphs (digraphs).

● What's the maximum number of edges in an undirected graph with $|V|$ vertices?

● Complete, dense, and sparse graphs

- A graph with every pair of its vertices connected by an edge is called complete, $K_{|V|}$



Graph Representation



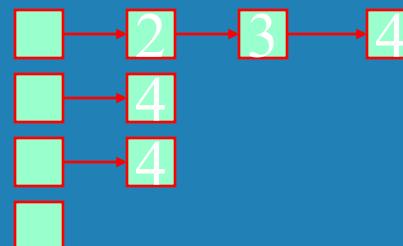
- **Adjacency matrix**

- $n \times n$ boolean matrix if $|V|$ is n .
- The element on the i th row and j th column is 1 if there's an edge from i th vertex to the j th vertex; otherwise 0.
- The adjacency matrix of an undirected graph is symmetric.

- **Adjacency linked lists**

- A collection of linked lists, one for each vertex, that contain all the vertices adjacent to the list's vertex.
- Which data structure would you use if the graph is a 100-node star shape?

0	1	1	1
0	0	0	1
0	0	0	1
0	0	0	0

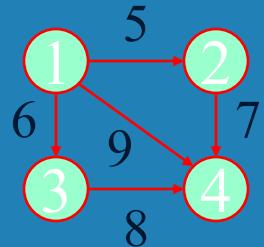


Weighted Graphs



- **Weighted graphs**

- **Graphs or digraphs with numbers assigned to the edges.**



Graph Properties -- Paths and Connectivity



• Paths

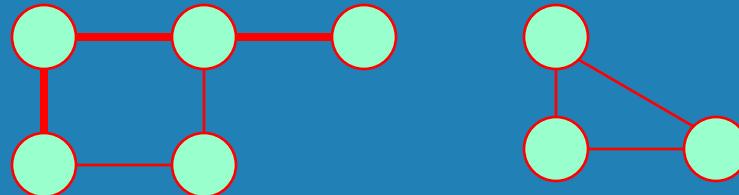
- A path from vertex u to v of a graph G is defined as a sequence of adjacent (connected by an edge) vertices that starts with u and ends with v .
- Simple paths: All edges of a path are distinct.
- Path lengths: the number of edges, or the number of vertices – 1.

• Connected graphs

- A graph is said to be connected if for every pair of its vertices u and v there is a path from u to v .

• Connected component

- The maximum connected subgraph of a given graph.



Graph Properties -- Acyclicity

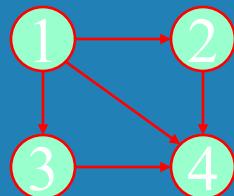


✿ Cycle

- A simple path of a positive length that starts and ends at the same vertex.

✿ Acyclic graph

- A graph without cycles
- DAG (Directed Acyclic Graph)



Trees



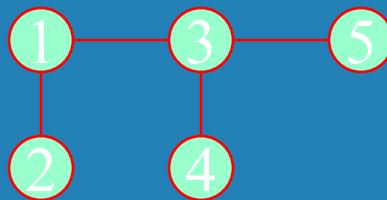
• Trees

- A tree (or free tree) is a connected acyclic graph.
- Forest: a graph that has no cycles but is not necessarily connected.

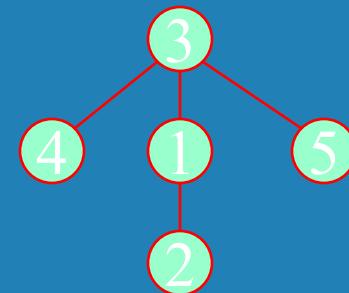
• Properties of trees

- For every two vertices in a tree there always exists exactly one simple path from one of these vertices to the other. Why?
 - Rooted trees: The above property makes it possible to select an arbitrary vertex in a free tree and consider it as the root of the so called rooted tree.
 - Levels in a rooted tree.

■ $|E| = |V| - 1$



rooted



Rooted Trees (I)

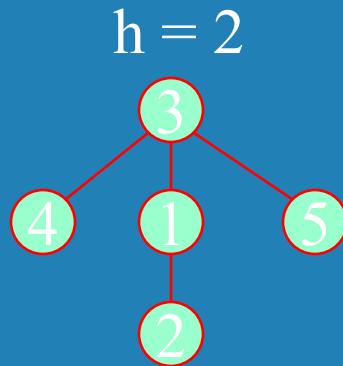


- **Ancestors**
 - For any vertex v in a tree T , all the vertices on the simple path from the root to that vertex are called ancestors.
- **Descendants**
 - All the vertices for which a vertex v is an ancestor are said to be descendants of v .
- **Parent, child and siblings**
 - If (u, v) is the last edge of the simple path from the root to vertex v , u is said to be the parent of v and v is called a child of u .
 - Vertices that have the same parent are called siblings.
- **Leaves**
 - A vertex without children is called a leaf.
- **Subtree**
 - A vertex v with all its descendants is called the subtree of T rooted at v .

Rooted Trees (II)



- **Depth of a vertex**
 - The length of the simple path from the root to the vertex.
- **Height of a tree**
 - The length of the longest simple path from the root to a leaf.



Ordered Trees



- **Ordered trees**
 - An ordered tree is a rooted tree in which all the children of each vertex are ordered.
- **Binary trees**
 - A binary tree is an ordered tree in which every vertex has no more than two children and each child is designated as either a left child or a right child of its parent.
- **Binary search trees**
 - Each vertex is assigned a number.
 - A number assigned to each parental vertex is larger than all the numbers in its left subtree and smaller than all the numbers in its right subtree.
- $\lfloor \log_2 n \rfloor \leq h \leq n - 1$, where h is the height of a binary tree and n the size.

