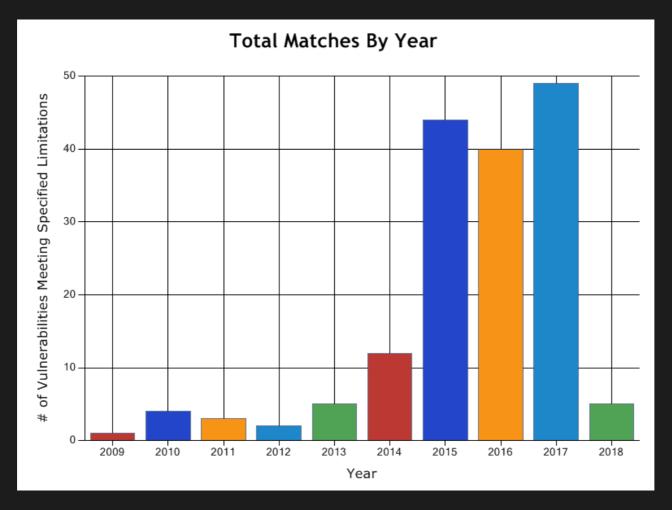


# Type Confusion: Discovery, Abuse, Protection

Mathias Payer, @gannimo http://hexhive.github.io

## Type confusion leads to RCE





## Attack surface is huge

Google Chrome: 76 MLoC

Gnome: 9 MLoC

Xorg: 1 MLoC

glibc: 2 MLoC

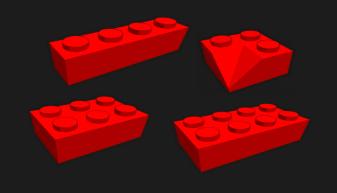
Linux kernel: 17 MLoC

## Control-Flow Hijack Attack (and CFI)

### Problem: broken abstractions?



```
C/C++
void log(int a) {
  printf("Log: %d", a);
}
void (*fun)(int) = &log;
void init() {
  fun(15);
}
```

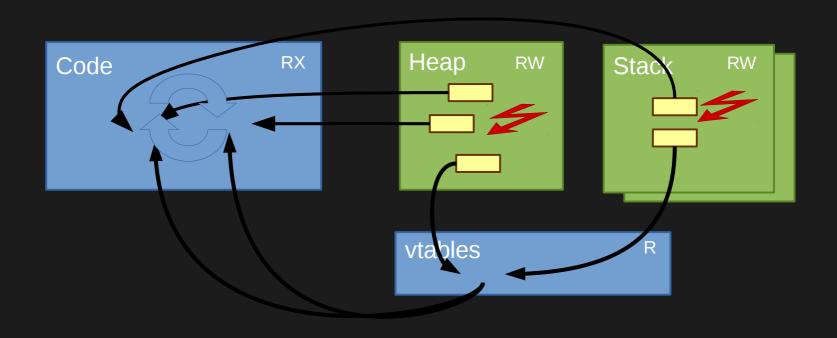




```
log: ASM
...
fun:
   .quad log
init:
   ...
   movl $15, %edi
   movq fun(%rip), %rax
   call *%rax
```



## Attacker model: hijacking control-flow



## Control-Flow Integrity (CFI)\*

Restrict a program's dynamic control-flow to the static CFG

- Requires static analysis
- Dynamic enforcement mechanism

<sup>\*</sup> **Control-Flow Integrity.** Martin Abadi, Mihai Budiu, Ulfar Erlingsson, Jay Ligatti. ACM CCS '05 \* **Control-Flow Integrity: Protection, Security, and Performance.** Nathan Burow, Scott A. Carr, Joseph Nash, Per Larsen, Michael Franz, Stefan Brunthaler, Mathias Payer. ACM CSUR '17

## Class hierarchy depth

Impl. Count	Chromium		Firefox	
[1-10]	13,751	(99.33%)	4,632	(99.90%)
>10	78	(0.57%)	47	(0.10%)
Max	78		107	

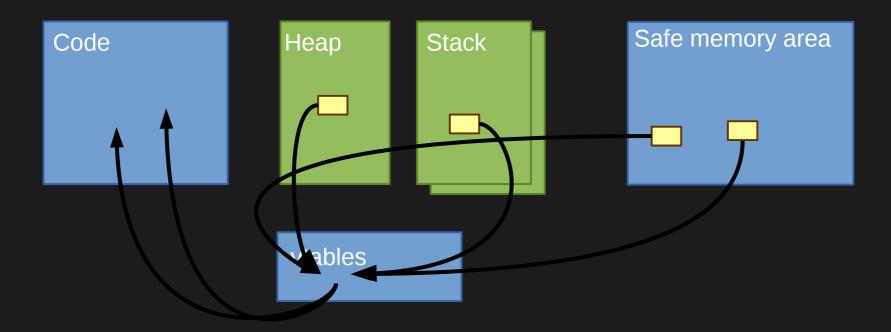
CFI prohibits use of corrupted pointer.

Can we do better?

## **Object Type Integrity**

## Object Type Integrity (OTI)\*

Enforce integrity of vtable pointer, use protected dispatch



\* CFIXX: Object Type Integrity for C++ Virtual Dispatch. Nathan Burow, Derrick McKee, Scott A. Carr, and Mathias Payer. In ISOC NDSS '18

### **CFIXX** instrumentation

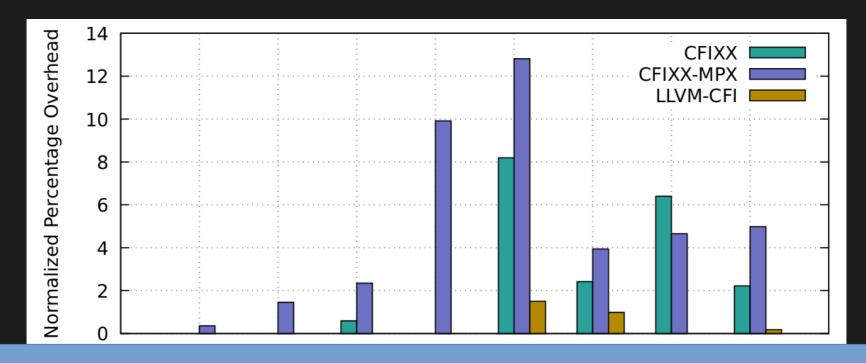
#### C++ dynamic dispatch has single target

- Only constructor allowed to write vtable pointer
- Deallocation invalidates vtable pointer
- Dispatch uses vtable pointer

### Enforcing OTI protects against

- VTable injection even with correct method signature
- Swap vtable even in the same hierarchy
- Fake object creation (COOP)

### CFIXX performance



Chromium: 2.03% (Octane), 1.99% (Kraken)

### CFI and CFIXX summary

#### CFI makes attacks harder

- Effectiveness depends on analysis and complexity
- Deployed in Microsoft Edge, Google Chrome on Linux
- Limitation: large equivalence classes

### Object Type Integrity (CFIXX)

- Protect object instead of dispatch
- Single valid target per object

Future direction: type check / overhead

Source: https://github.com/HexHive/CFIXX



## C++ Casting

## C++ casting operations

### dynamic\_cast<ToClass>(Object)

- Runtime check based on allocated type (vtable pointer)
- Not used in performance critical code

### static\_cast<ToClass>(Object)

- Class hierarchy check at compile time

### (ToClass) (Object)

- C-style cast, no check at all



### Static cast

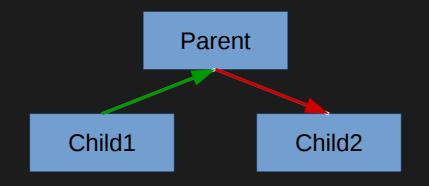
a = static\_cast<Greeter\*>(b);

```
movq -24(%rbp), %rax # Load pointer
# Type "check"
movq %rax, -40(%rbp) # Store pointer
```

### Dynamic cast, optimized

## Type Confusion

### Type confusion arises through illegal downcasts



```
Child1 *c = new Child1();
Parent *p = static_cast<Parent*>(c); \
Child2 *d = static_cast<Child2*>(p); \( \)
```

### Type confusion

```
class P {
                                vtable*?
  int x;
                         Pptr
                                    X
class C: P {
                                   У?
  int y;
                                vtable*
  virtual void print();
};
                                    X
P * Pptr = new P;
C *Cptr = static_cast<C*>Pptr; // Type Conf.
Cptr->y = 0x43; // Memory safety violation!
Cptr->print(); // Control-flow hijacking
```

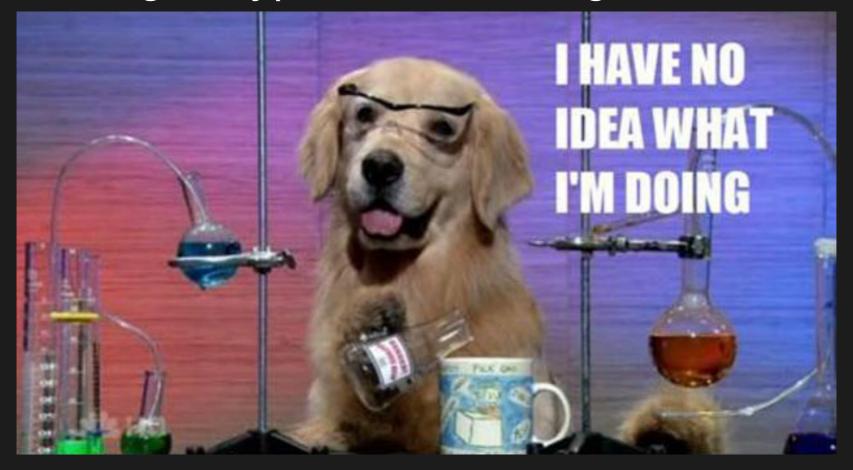
### Exploit primitive

Control two pointers of different types to single memory area Different interpretation of fields leads to "opportunities"



https://googleprojectzero.blogspot.ch/2015/07/one-perfect-bug-exploiting-type\_20.html https://blogs.technet.microsoft.com/mmpc/2015/06/17/understanding-type-confusion-vu lnerabilities-cve-2015-0336/

## Searching for type confusion bugs: SEGFAULT



## Type Sanitization

### Type safety for C++

A static cast is checked only at compile time

Fast but no runtime guarantees

Dynamic casts are checked at runtime

- High overhead, limited to polymorphic classes

#### Our core idea:

- Conceptually check all casts dynamically
- Aggressively optimize design and implementation

\* TypeSanitizer: Practical Type Confusion Detection. Istvan Haller, Yuseok Jeon, Hui Peng, Mathias Payer, Herbert Bos, Cristiano Giuffrida, Erik van der Kouwe. In CCS'16 \* HexType: Efficient Detection of Type Confusion Errors for C++. Yuseok Jeon, Priyam Biswas, Scott A. Carr, Byoungyoung Lee, and Mathias Payer. In CCS'17

## Making type checks explicit

Enforce runtime check at all cast sites

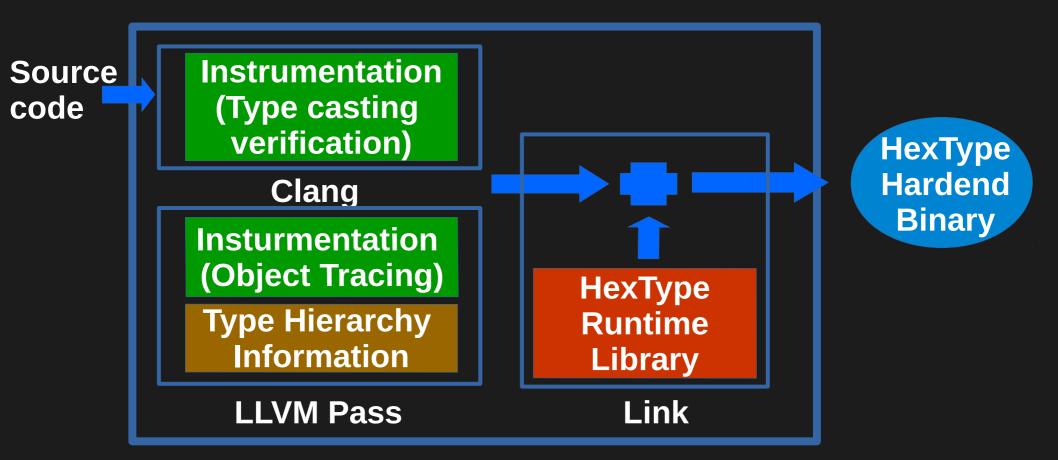
- static\_cast<ToClass>(Object)
- dynamic\_cast<ToClass>(Object)
- reinterpret\_cast<ToClass>(Object)
- (ToClass)(Object)

Build global type hierarchy

Keep track of the allocation type of each object

- Must instrument all forms of allocation
- Requires disjoint metadata

## HexType: design



### HexType: go full coverage!

#### Cover "new" object allocations

Obscure allocation cases for, e.g., arrays, stack

### Support placement\_new

Custom allocators don't call malloc/new

### Support reinterpret\_cast

Repurpose and revive existing objects

### HexType: aggressive optimization

Limit tracing to unsafe types

Remove tracing of types that are never cast

Limit checking to unsafe casts

Remove statically verifiable casts

No more RTTI for dynamic casts

Replace dynamic casts with fast lookup

## Simple exploitation demo

```
class Base { ... };
                                   Greater
class Exec: public Base {
  public:
    virtual void exec(const char *prg) {
      system(prg);
class Greeter: public Base {
  public:
    virtual void sayHi(const char *str) {
      std::cout << str << std::endl;</pre>
```

Base

Exec

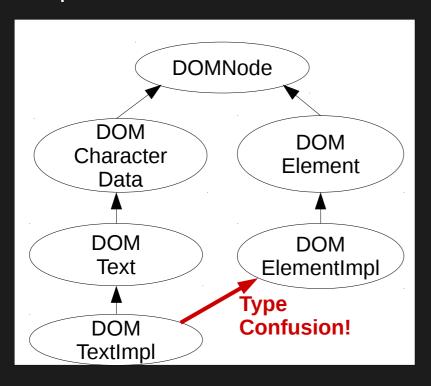
### Simple exploitation demo

```
vtable*
int main() {
  Base *b1 = new Greeter();
                                       vtable*
  Base *b2 = new Exec();
  Greeter *q;
                                             ExecT
  g = static_cast<Greeter*>(b1);
  g->sayHi("Greeter says hi!");
                                  // g[0][0](str);
  g = static_cast<Greeter*>(b2);
  q->sayHi("/usr/bin/xcalc");
                                  // q[0][0](str);
  delete b1;
  delete b2;
  return 0;
```

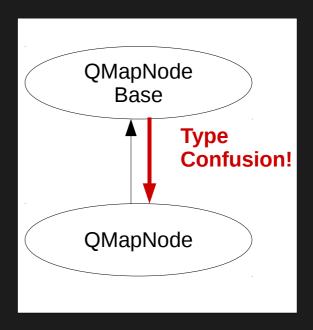
GreeterT

### Low hanging fruits: four new vulnerabilities

Apache Xerces C++



Qt base library



## Fuzz all the Things!

### Combine AFL/libFuzzer with HexType

### AFL/libFuzzer and HexType play surprisingly well together

- Compile software with HexType, trap on type confusion
- Let fuzzing do its magic
- Triage type confusion reports
- \$\$\$





## First two weeks of fuzzing

QTcore: two new type confusion bugs (not exploitable, fixed)

Xerces C++: one new type confusion (reported, fixed)



### One more week of fuzzing with libFuzzer

ChakraCore: two cases of type confusion (reported)

MySQL 5.7: five cases of type confusion (reported, serious)

Other C++ software: evaluation ongoing

- Let us know what we should target next
- Have spare fuzzing capacity? Let's team up!

### But what about Firefox?

FF-Octane: 5,506,850 type confusion reports

FF-Dramaeo-JS: 15,216,798 type confusion reports

FF-Dramaeo-dom: 7,240,272,959 type confusion reports

Large amount of duplicates and false positives

Firefox code is messy, few actual bugs but lots of code smell

## Conclusion

## Ongoing work

### Fuzz all the things!

More software, better test cases, deeper coverage

### Selective fuzzing

- Select which types to test (DOM anyone?)
- Extend type check to dereference



### Conclusion

Type confusion is fundamental in today's exploits Existing solutions are incomplete, partial, slow

HexType: type sanitizer for C++

- Trap upon type confusion, not memory safety violation
- Reasonable slowdown for testing (~50%)
- Integrated with AFL/libFuzzer for broad bug discovery

https://github.com/HexHive/HexType

Twitter: @gannimo