HYPERPILL

Fuzzing for Hypervisor-bugs by leveraging the Hardware Virtualization Interface

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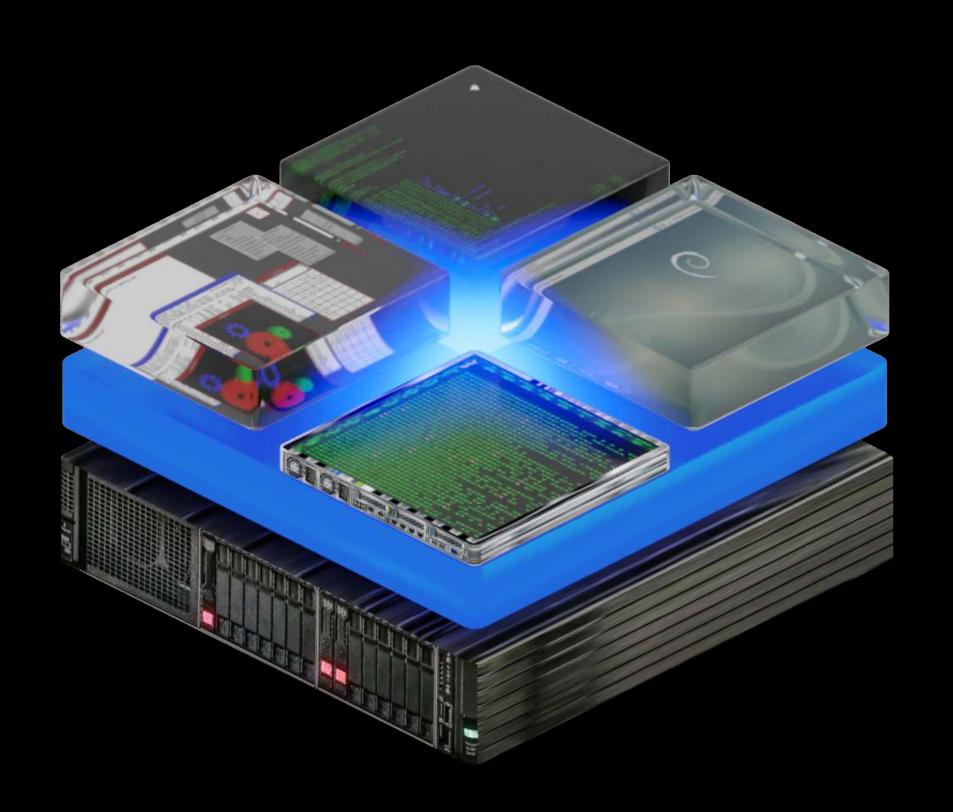
Qiang Liu—Boston University

Manuel Egele Zhejiang University

Mathias Payer — Amazon (all work completed prior to joining Amazon)



Motivation



Applications:

Cloud
Personal Computing
Development
Security Research
Mobile

Automotive

Fuzzing Hypervisors: Challenges

Hypervisor inputs feature complex semantics

There are many hypervisors

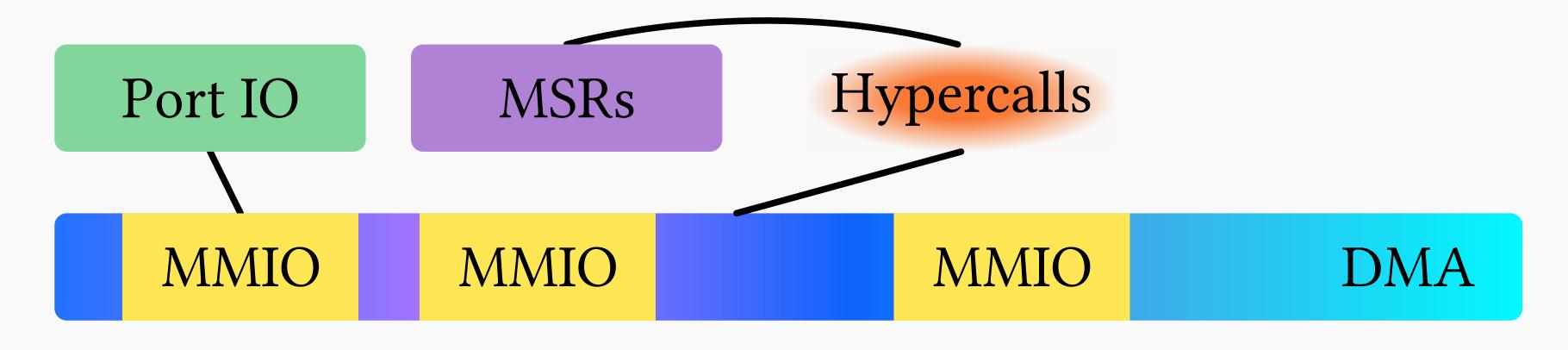
Implementations are diverse

Hypervisors are low-level systems software

All Hypervisors use an identical CPU virtualization interface

The Input Space

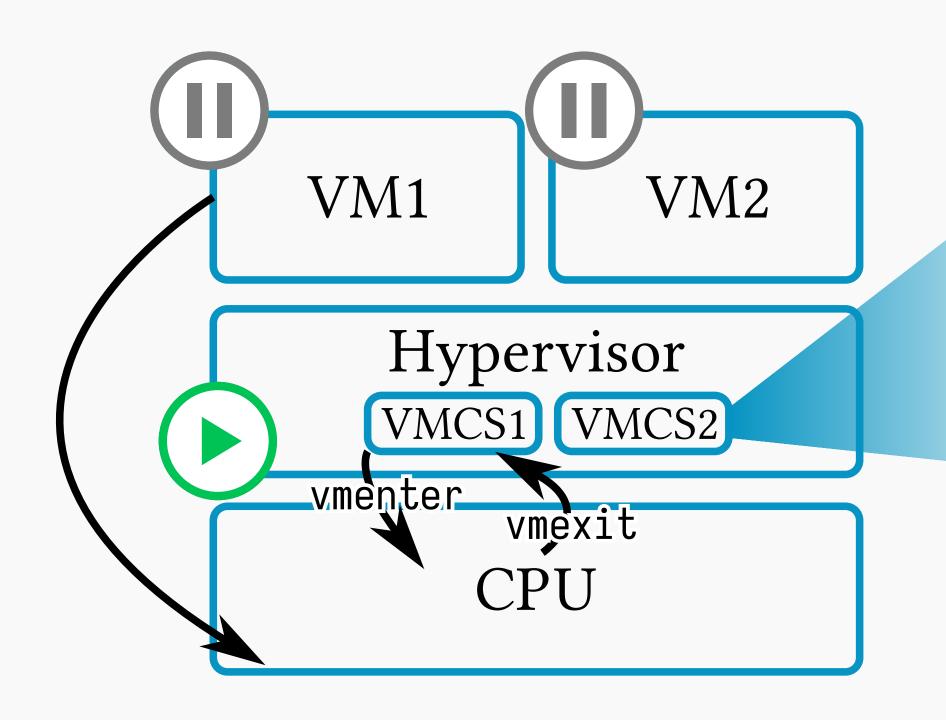
What interfaces do hypervisors expose to VMs?



The combined hypervisor input space is enormous

Hypervisor semantics are complex and variable

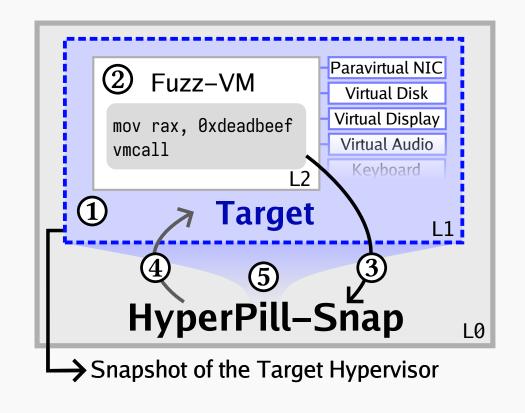
The Hardware Virtualization Interface

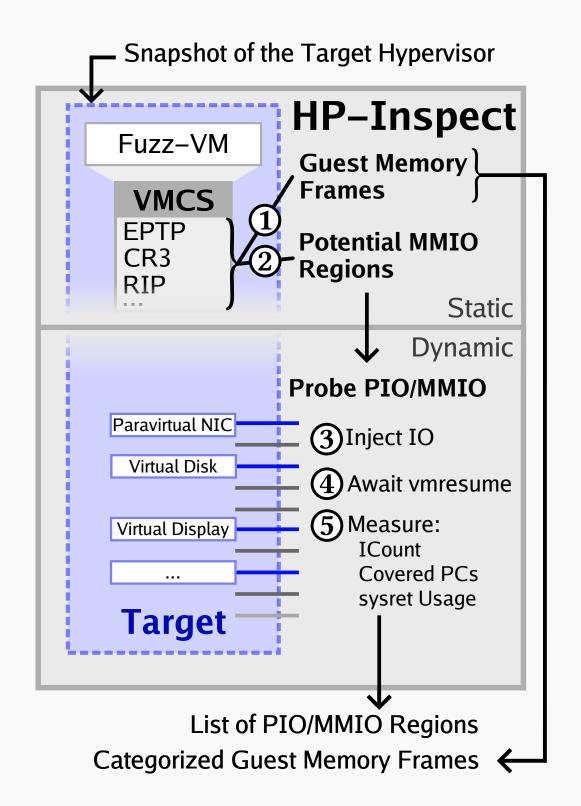


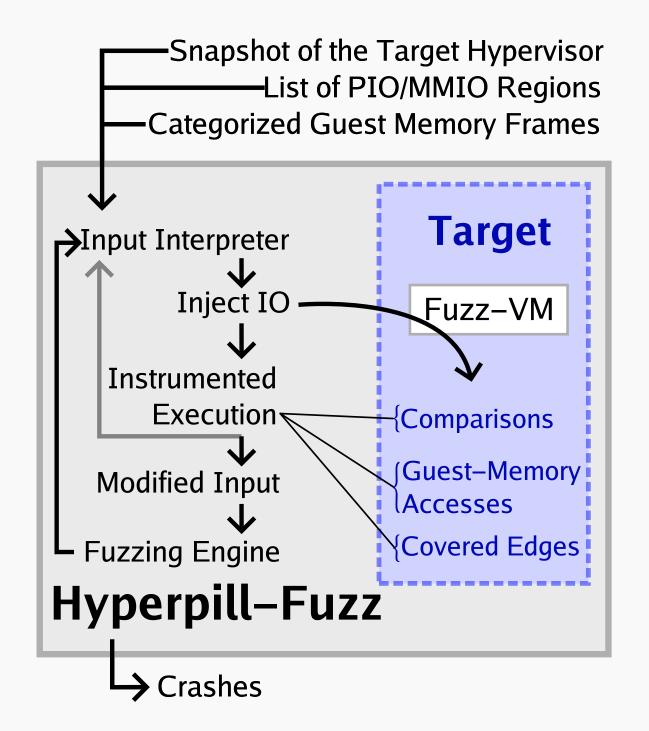
- Guest vCPU State
- Memory allocated to the Guest
- MMIO Regions
- PIO Regions
- Enabled MSRs
- VM Exit Reasons

Every Hypervisor must use an identical CPU virtualization interface. HyperPill leverages this fact to fuzz any hypervisor - generically.

HyperPill







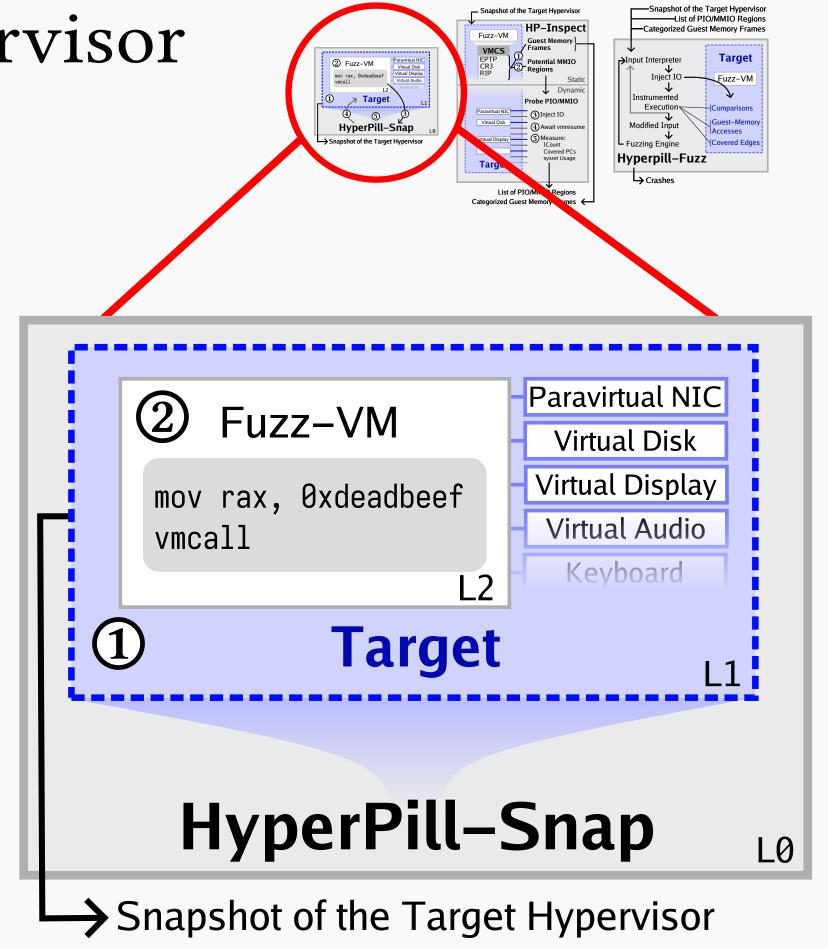
Step 1: Snapshotting the Hypervisor

Problem: Make a snapshot of the hypervisor just as it is about to handle a VM Exit.

Solution:

Run the hypervisor nested in HyperPill-Snap.

Invoke a special hypercall from the VM to trigger a VM Exit and tell HyperPill-Snap to collect a snapshot.



Step 2: Inspecting the Snapshot

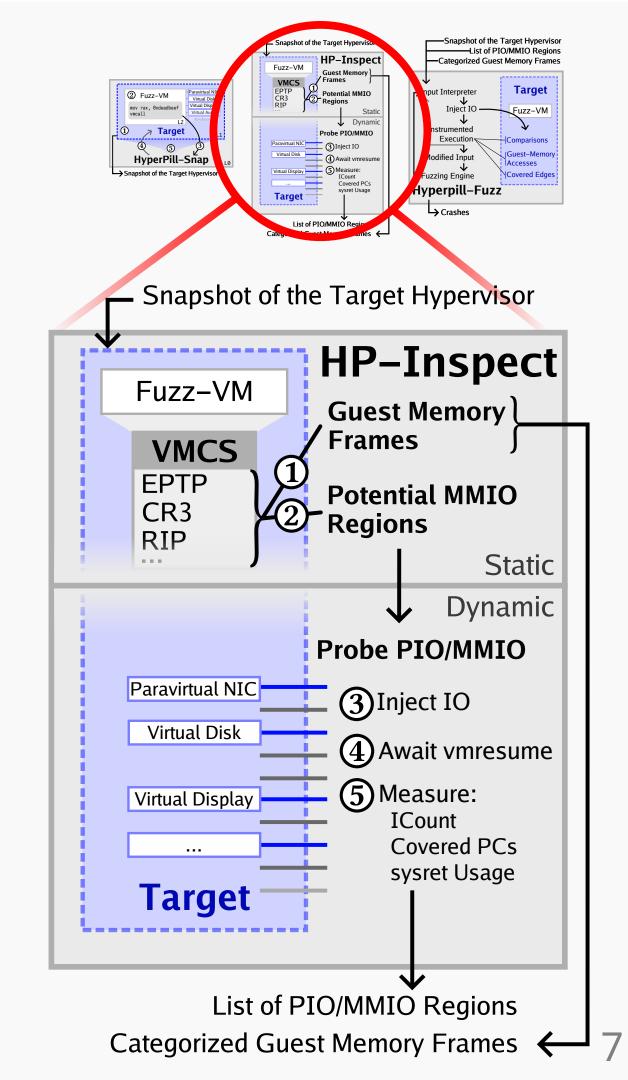
The snapshot contains the entire hypervisor state, including the VMCS the hypervisor sets up for the VM.

VMCS Inspection/Probing to Identify:

MMIO Regions

PIO Regions

Memory allocated to the Guest



Step 3: Fuzzing

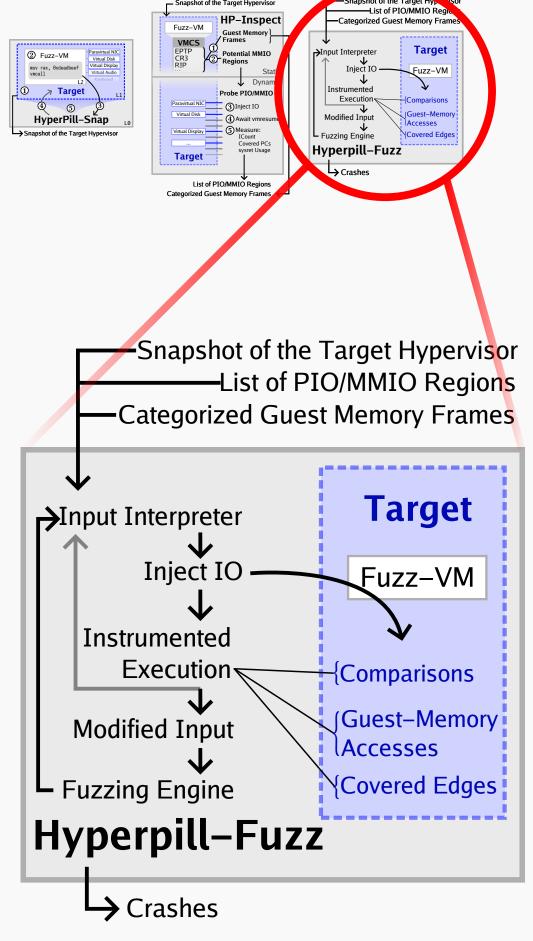
Fuzzer input is a sequence of IO Operations



Modify VMCS to reflect VMExit for the IO Operation. Resume the Hypervisor Snapshot in an emulator.

When the Hypervisor resumes the VM, immediately inject the next VM-Exit

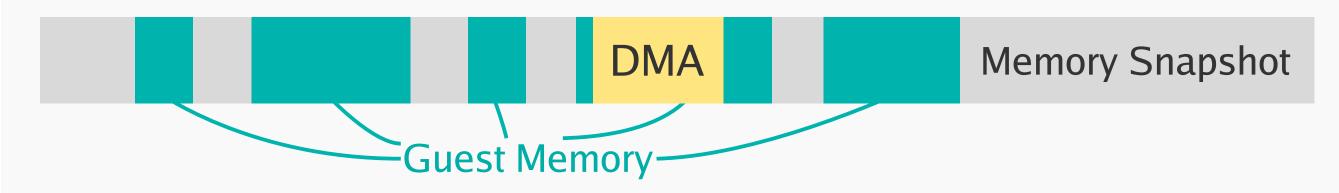




Step 3: Fuzzing DMA

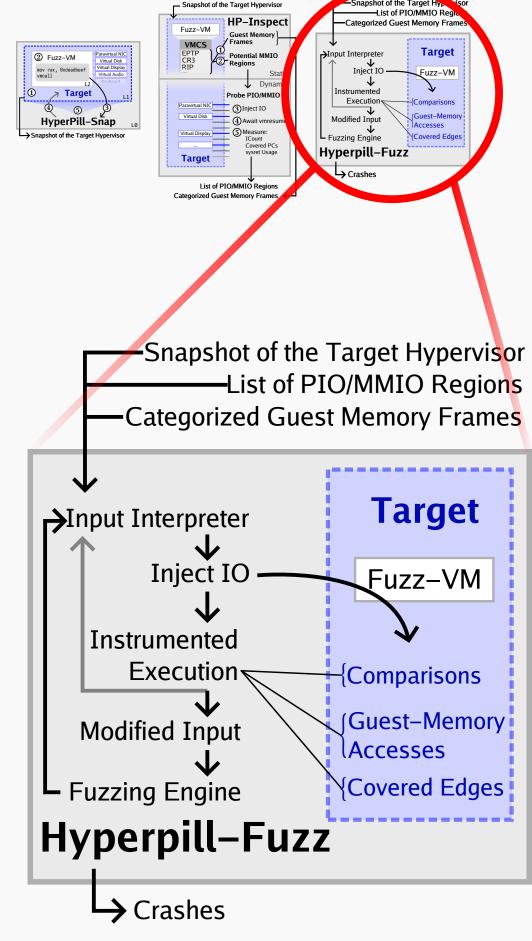
DMA allows hypervisor to read from guest memory at any time while handling VMExits

In step 2, HyperPill identified all Guest Memory pages within the snapshot



During emulation, if any instruction reads from the Guest Memory, HyperPill fills the memory with fuzzer data





Step 3: Fuzzing DMA

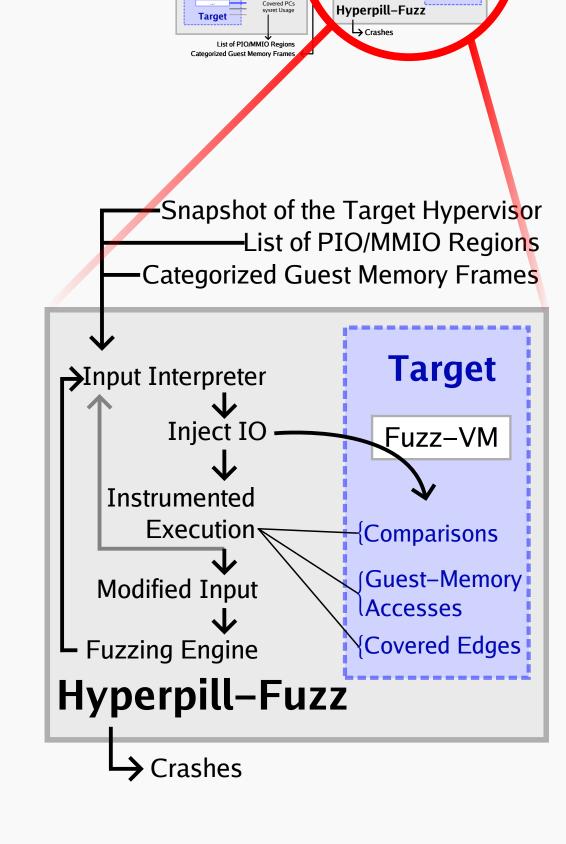
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PIO MMIO MSR's HyperCalls DMA

Guest Memory

During emulation, if any instruction reads from the Guest Memory, HyperPill fills the memory with fuzzer data



Results

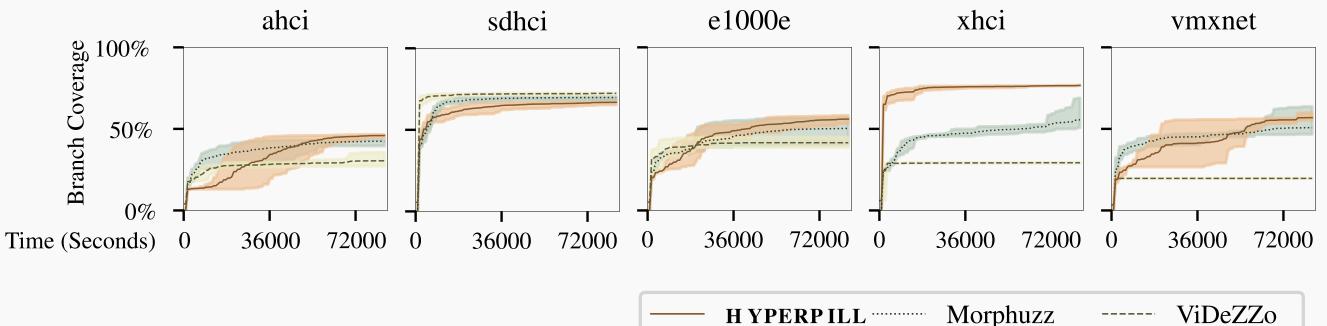
Three Hypervisors

HyperPill achieves higher coverage for 10/12 devices.

Inherent performance penalty due to use of full-system snapshotting and emulation.

Byte-level DMA sensitivity allows better performance for DMA-heavy devices.

		Morphuzz		ViDeZZo		H YPERP ILL	
		12 Cores 24 Hours					
Device	·	Branch Coverage (Executions/Second)					
Block							
ahci	√	42.43%	(25.68)	30.42%	(562.24)	45.90%	(26.18)
nvme		29.12%	(23.82)			36.44%	(14.45)
sdhci	√	69.81%	(22.98)	72.37%	(107.22)	66.85%	(32.34)
virtio-scsi	\checkmark	27.96%	(23.83)	11.73%	(217.28)	48.83%	(51.68)
Display							
cirrus		88.10%	(19.06)	83.42%	(138.78)	88.67%	(32.18)
qxl	√					59.68%	(26.96)
virtio-gpu	√	24.37%	(26.21)	2.77%	(222.42)	45.52%	(36.53)
Networking							
e1000e	√	50.27%	(24.83)	41.52%	(53.04)	55.99%	(42.22)
igb	\checkmark	29.73%	(25.63)			35.93%	(60.85)
vmxnet	\checkmark	50.75%	(27.01)	19.64%	(145.73)	56.89%	(48.14)
USB							
ehci	√	73.76%	(24.58)	74.38%	(177.08)	73.32%	(10.46)
xhci	√	55.54%	(28.83)	29.25%	(1061.36)	76.64%	(69.26)
Geo. Mean		45.20%	(24.65)	28.00%	(203.07)	55.45%	(33.20)



Results

Three Hypervisors

```
/* map PRDT
if (!(prdt = dma_memory_map(ad->hba->as, prdt_addr, &prdt len,
                            DMA DIRECTION TO DEVICE,
                            MEMTXATTRS UNSPECIFIED))){
   trace_ahci_populate s
                                               port no);
   return -1;
/* Get entries in the PRDT, init a qemu sglist accordingly */
if (prdtl > 0) {
   AHCI_SG *tbl = (AHCI_SG *)prdt;
   for (i = 0; i < problem + 1++) {
       tbl_entry_si e = prdt_tbl_entry_size(&tbl[i]);
        if (offset < (sum +thl entry size)) {
            off idx
```

```
Branch Coverage (Executions/Second)

42.43% (25.68) 30.42% (562.24) 45.90% (26.18)
29.12% (23.82) 36.44% (14.45)
69.81% (22.98) 72.37% (107.22) 66.85% (32.34)
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```

```
leaq (%r15,%r12), %r14
incq %r14
movq %r14, %rdi
movq %rbx, %rsi
```

HyperPill DMA Hooking

Source-Based DMA Hooking

QEMU

Arbitrary memory-access in e1000e_start_xmit
Heap-overflow in usb_mouse_poll
Heap-overflow in virtqueue_alloc_element
Heap-overflow in qxl_cookie_new
Heap-overflow in igb_tx_pkt_switch
Out-of-bounds memory access in nvme_process_sq
Out-of-bounds memory access in nvme_io_mgmt_send
DoS via arbitrary-sized allocation in qxl
DoS via arbitrary-sized allocation in virtio_gpu
DoS in process_ncq_command
DoS in icmp_input

Bugs

Hyper-V

Heap-corruption in EthernetCard::HandleTransmitSetupFrame
Abort in EthernetCard::PollForTransmitDataTimer
Abort after IdeChannel::EnlightenedHddCommand
EthernetCard::SetupEthernetCardModeFromRegisters
Out-of-bounds write in GuestStateAccess::SetDeviceInfo
Abort after PitDevice::NotifyIoPortRead
Abort in I8042Device::HandleCommand
Abort after HvCallDetachDevice
Abort after HvCallGetGpaPagesAccessState

macOS Virtualization Framework

Memory-privilege violation in xHCI
Out-of-bounds write in virtio-gpu
Out-of-bounds write in virtio-audio
Out-of-bounds access in virtio-block
Out-of-bounds access in virtio-console
Out-of-bounds access in virtio-net

Reshape hypervisors by modifying the CPU virtualization interface No modification to hypervisors code needed!

Fuzz any hypervisor across its PIO, MMIO, DMA, and Hypercall interfaces

More precise than source-level reshaping

HERPERPIL

https://github.com/HexHive/HyperPill