Programming Language—Common Lisp

14. Conses

14.1 Cons Concepts

A cons is a compound data *object* having two components called the car and the cdr.

car	cons	rplacd	
$\operatorname{\mathbf{cdr}}$	rplaca		

Figure 14–1. Some defined names relating to conses.

Depending on context, a group of connected *conses* can be viewed in a variety of different ways. A variety of operations is provided to support each of these various views.

14.1.1 Conses as Trees

A **tree** is a binary recursive data structure made up of *conses* and *atoms*: the *conses* are themselves also *trees* (sometimes called "subtrees" or "branches"), and the *atoms* are terminal nodes (sometimes called *leaves*). Typically, the *leaves* represent data while the branches establish some relationship among that data.

l			• .	
caaaar	caddar	$\operatorname{\mathbf{cdar}}$	nsubst	
caaadr	cadddr	cddaar	${f nsubst-if}$	
caaar	caddr	cddadr	${f nsubst-if-not}$	
caadar	cadr	cddar	${f nthcdr}$	
caaddr	cdaaar	cdddar	${f sublis}$	
caadr	cdaadr	$\operatorname{\mathbf{cddddr}}$	${f subst}$	
caar	cdaar	$\operatorname{\mathbf{cdddr}}$	$\mathbf{subst} ext{-}\mathbf{if}$	
cadaar	cdadar	$\operatorname{\mathbf{cddr}}$	${f subst-if-not}$	
cadadr	cdaddr	$\operatorname{copy-tree}$	tree-equal	
cadar	cdadr	nsublis		

Figure 14-2. Some defined names relating to trees.

14.1.1.1 General Restrictions on Parameters that must be Trees

Except as explicitly stated otherwise, for any *standardized function* that takes a *parameter* that is required to be a *tree*, the consequences are undefined if that *tree* is circular.

14.1.2 Conses as Lists

A *list* is a chain of *conses* in which the *car* of each *cons* is an *element* of the *list*, and the *cdr* of each *cons* is either the next link in the chain or a terminating *atom*.

A **proper list** is a list terminated by the empty list. The empty list is a proper list, but is not a cons.

An improper list is a list that is not a proper list; that is, it is a circular list or a dotted list.

A **dotted list** is a list that has a terminating atom that is not the empty list. A non-nil atom by itself is not considered to be a list of any kind—not even a dotted list.

A $circular\ list$ is a chain of conses that has no termination because some cons in the chain is the cdr of a later cons.

append	last	nbutlast	rest
$\operatorname{butlast}$	ldiff	nconc	revappend
${f copy-alist}$	${f list}$	\mathbf{ninth}	\mathbf{second}
$\operatorname{copy-list}$	list*	nreconc	$\mathbf{seventh}$
${f eighth}$	${f list} ext{-length}$	\mathbf{nth}	${f sixth}$
${f endp}$	${f make-list}$	${f nthcdr}$	${f tailp}$
${f fifth}$	\mathbf{member}	pop	tenth
${f first}$	member-if	push	${f third}$
fourth	${\bf member\text{-}if\text{-}not}$	pushnew	

Figure 14–3. Some defined names relating to lists.

14.1.2.1 Lists as Association Lists

An **association list** is a *list* of *conses* representing an association of *keys* with *values*, where the *car* of each *cons* is the *key* and the *cdr* is the *value* associated with that *key*.

acons	assoc-if	pairlis	rassoc-if	
assoc	${\it assoc-if-not}$	rassoc	${f rassoc ext{-}if ext{-}not}$	

Figure 14-4. Some defined names related to assocation lists.

14.1.2.2 Lists as Sets

Lists are sometimes viewed as sets by considering their elements unordered and by assuming there is no duplication of elements.

adjoin	nset-difference	set-difference	union
intersection	nset-exclusive-or	${f set} ext{-exclusive-or}$	
nintersection	nunion	$\mathbf{subsetp}$	

Figure 14-5. Some defined names related to sets.

14.1.2.3 General Restrictions on Parameters that must be Lists

Except as explicitly specified otherwise, any *standardized function* that takes a *parameter* that is required to be a *list* should be prepared to signal an error of *type* **type-error** if the *value* received is a *dotted list*.

Except as explicitly specified otherwise, for any *standardized function* that takes a *parameter* that is required to be a *list*, the consequences are undefined if that *list* is *circular*.

list System Class

Class Precedence List:

list, sequence, t

Description:

A *list* is a chain of *conses* in which the *car* of each *cons* is an *element* of the *list*, and the *cdr* of each *cons* is either the next link in the chain or a terminating *atom*.

A **proper list** is a chain of *conses* terminated by the **empty list**, (), which is itself a *proper list*. A **dotted list** is a *list* which has a terminating *atom* that is not the *empty list*. A **circular list** is a chain of *conses* that has no termination because some *cons* in the chain is the *cdr* of a later *cons*.

Dotted lists and circular lists are also lists, but usually the unqualified term "list" within this specification means proper list. Nevertheless, the type list unambiguously includes dotted lists and circular lists.

For each element of a list there is a cons. The empty list has no elements and is not a cons.

The types cons and null form an exhaustive partition of the type list.

See Also:

Section 2.4.1 (Left-Parenthesis), Section 22.1.3.5 (Printing Lists and Conses)

null System Class

Class Precedence List:

null, symbol, list, sequence, t

Description:

The only object of type null is nil, which represents the empty list and can also be notated ().

See Also:

Section 2.3.4 (Symbols as Tokens), Section 2.4.1 (Left-Parenthesis), Section 22.1.3.3 (Printing Symbols)

CONS System Class

Class Precedence List:

cons, list, sequence, t

Description:

A cons is a compound object having two components, called the car and cdr. These form a dotted pair. Each component can be any object.

Compound Type Specifier Kind:

Specializing.

Compound Type Specifier Syntax:

(cons [car-typespec [cdr-typespec]])

Compound Type Specifier Arguments:

car-typespec—a type specifier, or the symbol *. The default is the symbol *.

 $\it cdr$ -typespec—a $\it type \ specifier$, or the $\it symbol \ *$. The default is the $\it symbol \ *$.

Compound Type Specifier Description:

This denotes the set of *conses* whose car is constrained to be of type car-typespec and whose cdr is constrained to be of type cdr-typespec. (If either car-typespec or cdr-typespec is *, it is as if the type t had been denoted.)

See Also:

Section 2.4.1 (Left-Parenthesis), Section 22.1.3.5 (Printing Lists and Conses)

atom Type

Supertypes:

atom, t

Description:

It is equivalent to (not cons).

cons Function

Syntax:

```
cons\ object-1\ object-2\ 	o\ cons
```

Arguments and Values:

```
object-1—an object.object-2—an object.cons—a cons.
```

Description:

Creates a fresh cons, the car of which is object-1 and the cdr of which is object-2.

Examples:

```
(cons 1 2) \rightarrow (1 . 2)

(cons 1 nil) \rightarrow (1)

(cons nil 2) \rightarrow (NIL . 2)

(cons nil nil) \rightarrow (NIL)

(cons 1 (cons 2 (cons 3 (cons 4 nil)))) \rightarrow (1 2 3 4)

(cons 'a 'b) \rightarrow (A . B)

(cons 'a (cons 'b (cons 'c '()))) \rightarrow (A B C)

(cons 'a '(b c d)) \rightarrow (A B C D)
```

See Also:

list

Notes:

If object-2 is a list, cons can be thought of as producing a new list which is like it but has object-1 prepended.

consp Function

Syntax:

```
\mathbf{consp}\ \textit{object}\ \rightarrow \textit{generalized-boolean}
```

Arguments and Values:

```
object—an object.
```

generalized-boolean—a generalized boolean.

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Description:

Returns true if object is of type cons; otherwise, returns false.

Examples:

```
(consp nil) \rightarrow false
(consp (cons 1 2)) \rightarrow true
The empty list is not a cons, so
(consp '()) \equiv (consp 'nil) \rightarrow false
```

See Also:

listp

Notes:

```
(consp object) \equiv (typep object 'cons) \equiv (not (typep object 'atom)) \equiv (typep object '(not atom))
```

atom Function

Syntax:

 $\mathbf{atom}\ \mathit{object}\ \rightarrow \mathit{generalized-boolean}$

Arguments and Values:

```
object—an object.
```

generalized-boolean—a generalized boolean.

Description:

Returns true if object is of type atom; otherwise, returns false.

Examples:

```
 \begin{array}{l} (\texttt{atom 'sss}) \to true \\ (\texttt{atom (cons 1 2)}) \to false \\ (\texttt{atom nil}) \to true \\ (\texttt{atom '()}) \to true \\ (\texttt{atom 3}) \to true \\ \end{array}
```

Notes:

```
(atom object) ≡ (typep object 'atom) ≡ (not (consp object))
≡ (not (typep object 'cons)) ≡ (typep object '(not cons))
```

rplaca, rplacd

Function

Syntax:

```
\begin{array}{ll} \mathbf{rplaca} \ \textit{cons object} & \rightarrow \textit{cons} \\ \mathbf{rplacd} \ \textit{cons object} & \rightarrow \textit{cons} \\ \end{array}
```

Pronunciation:

```
\begin{split} \mathbf{rplaca} \colon \left[ \ _{\mathbf{I}}\mathbf{r}\bar{\mathbf{e}}^{\ \mathbf{I}} \ \mathbf{plak}\epsilon \ \right] & \text{or} \left[ \ _{\mathbf{I}}\mathbf{r}\epsilon^{\ \mathbf{I}} \ \mathbf{plak}\epsilon \ \right] \\ \mathbf{rplacd} \colon \left[ \ _{\mathbf{I}}\mathbf{r}\bar{\mathbf{e}}^{\ \mathbf{I}} \ \mathbf{plak}d\epsilon \ \right] & \text{or} \left[ \ _{\mathbf{I}}\mathbf{r}\epsilon^{\ \mathbf{I}} \ \mathbf{plak}d\bar{\mathbf{e}} \ \right] & \text{or} \left[ \ _{\mathbf{I}}\mathbf{r}\epsilon^{\ \mathbf{I}} \ \mathbf{plak}d\bar{\mathbf{e}} \ \right] \end{split}
```

Arguments and Values:

```
cons—a cons.
object—an object.
```

Description:

rplaca replaces the car of the cons with object. rplacd replaces the cdr of the cons with object.

Examples:

```
(defparameter *some-list* (list* 'one 'two 'three 'four)) \rightarrow *some-list* *some-list* \rightarrow (ONE TWO THREE . FOUR) (rplaca *some-list* 'uno) \rightarrow (UNO TWO THREE . FOUR) *some-list* \rightarrow (UNO TWO THREE . FOUR) (rplacd (last *some-list*) (list 'IV)) \rightarrow (THREE IV) *some-list* \rightarrow (UNO TWO THREE IV)
```

Side Effects:

The *cons* is modified.

Should signal an error of *type* **type-error** if *cons* is not a *cons*.

car, cdr, caar, cadr, cdar, cddr, caaar, caadr, cadar, ...

car, cdr, caar, cadr, cdar, cddr, caaar, caadr, cadar, caddr, cdaar, cdadr, cddar, caddr, caaaar, caaadr, caadar, caddar, caddar, caddar, caddar, caddar, cdaaar, cdaadr, cdadar, cddadr, cddar, cdddar, cdddar

Syntax:

```
\rightarrow object
\operatorname{car} X
                                          (setf (car x) new-object)
                \rightarrow object
                                          (\mathbf{setf} (\mathbf{cdr} \ x) \ \mathit{new-object})
\operatorname{\mathbf{cdr}} x
                \rightarrow object
                                          (setf (caar x) new-object)
\operatorname{\mathbf{caar}} X
\operatorname{cadr} X
                \rightarrow object
                                          (setf (cadr x) new-object)
                \rightarrow object
                                          (setf (cdar x) new-object)
cdar x
\operatorname{cddr} x
                \rightarrow object
                                          (setf (cddr x) new-object)
caaar x
                \rightarrow object
                                          (setf (caaar x) new-object)
\operatorname{caadr} X
                \rightarrow object
                                          (setf (caadr x) new-object)
                                          (setf (cadar x) new-object)

ightarrow object
\mathbf{cadar}\ X
\mathbf{caddr}\ X
                                          (setf (caddr x) new-object)
                \rightarrow object
\operatorname{cdaar} x
                                          (setf (cdaar x) new-object)
                \rightarrow object
\operatorname{cdadr} X

ightarrow object
                                          (setf (cdadr x) new-object)
                                          (setf (cddar x) new-object)
\operatorname{cddar} X
                \rightarrow object
\mathbf{cdddr}\ X
                \rightarrow object
                                          (setf (cdddr x) new-object)

ightarrow object
                                          (setf (caaaar x) new-object)
caaaar x

ightarrow object
                                          (setf (caaadr x) new-object)
\operatorname{caaadr} X

ightarrow object
                                          (setf (caadar x) new-object)
caadar x
\mathbf{caaddr} \ x \rightarrow object
                                          (setf (caaddr x) new-object)
cadaar x \rightarrow object
                                          (setf (cadaar x) new-object)
cadadr x \rightarrow object
                                          (setf (cadadr x) new-object)
caddar x \rightarrow object
                                          (setf (caddar x) new-object)
\mathbf{cadddr} \ x \rightarrow object
                                          (setf (cadddr x) new-object)

ightarrow object
                                          (setf (cdaaar x) new-object)
cdaaar x
cdaadr x \rightarrow object
                                          (setf (cdaadr x) new-object)
cdadar x \rightarrow object
                                          (setf (cdadar x) new-object)
\operatorname{cdaddr} x \to \operatorname{object}
                                          (setf (cdaddr x) new-object)
                                          (setf (cddaar x) new-object)
cddaar x \rightarrow object
\mathbf{cddadr}\ x \quad \to \mathit{object}
                                          (setf (cddadr x) new-object)
\mathbf{cdddar} \ x \quad \to \mathit{object}
                                          (setf (cdddar x) new-object)
\operatorname{\mathbf{cddddr}} x \to \operatorname{\mathit{object}}
                                          (setf (cddddr x) new-object)
```

Pronunciation:

cadr: $[ka_1 d\epsilon r]$

car, cdr, caar, cadr, cdar, caaar, caadr, cadar, ...

```
caddr: [ \ ^{1}kad\epsilon_{1}d\epsilon_{1}] or [ \ ^{1}ka_{1}d\dot{u}d\epsilon_{1}]
cdr: [ \ ^{1}k\dot{u}_{1}d\epsilon_{1}]
cddr: [ \ ^{1}k\dot{u}d\epsilon_{1}d\epsilon_{1}] or [ \ ^{1}k\epsilon_{1}d\dot{u}d\epsilon_{1}]
```

Arguments and Values:

```
x—a list.
object—an object.
new-object—an object.
```

Description:

If x is a cons, car returns the car of that cons. If x is nil, car returns nil.

If x is a *cons*, **cdr** returns the *cdr* of that *cons*. If x is **nil**, **cdr** returns **nil**.

Functions are provided which perform compositions of up to four car and cdr operations. Their names consist of a C, followed by two, three, or four occurrences of A or D, and finally an R. The series of A's and D's in each function's name is chosen to identify the series of car and cdr operations that is performed by the function. The order in which the A's and D's appear is the inverse of the order in which the corresponding operations are performed. Figure 14–6 defines the relationships precisely.

car, cdr, caar, cadr, cddr, caaar, caadr, cadar, ...

This place	Is equivalent to this place
(caar x)	(car (car x))
(cadr X)	$(\operatorname{car} (\operatorname{cdr} X))$
(cdar x)	$(\operatorname{cdr} (\operatorname{car} X))$
(cddr x)	(cdr (cdr <i>x</i>))
(caaar X)	(car (car x)))
(caadr X)	(car (cdr x)))
(cadar X)	(car (cdr (car x)))
(caddr X)	(car (cdr (cdr x)))
(cdaar X)	(cdr (car (car x)))
(cdadr X)	$(\operatorname{cdr} (\operatorname{cdr} x))$
(cddar X)	(cdr (cdr (car x)))
(cdddr X)	(cdr (cdr (cdr x)))
(caaaar X)	(car (car (car x))))
(caaadr X)	(car (car (cdr x))))
(caadar X)	(car (cdr (car x))))
(caaddr X)	(car (cdr (cdr x))))
(cadaar X)	(car (cdr (car (car x))))
(cadadr X)	(car (cdr (cdr x))))
(caddar X)	(car (cdr (cdr (car x))))
(cadddr X)	(car (cdr (cdr x))))
(cdaaar x)	(cdr (car (car x))))
(cdaadr X)	(cdr (car (cdr x))))
(cdadar x)	(cdr (car (cdr (car X))))
(cdaddr X)	(cdr (cdr (cdr x))))
(cddaar X)	(cdr (cdr (car x))))
(cddadr X)	(cdr (cdr (cdr x))))
(cdddar X)	(cdr (cdr (car x))))
(cddddr X)	(cdr (cdr (cdr x))))

Figure 14–6. CAR and CDR variants

setf can also be used with any of these functions to change an existing component of x, but setf will not make new components. So, for example, the car of a cons can be assigned with setf of car, but the car of nil cannot be assigned with setf of car. Similarly, the car of the car of a cons whose car is a cons can be assigned with setf of caar, but neither nilnor a cons whose car is nil can be assigned with setf of caar.

The argument x is permitted to be a dotted list or a circular list.

Examples:

```
\begin{array}{l} (\text{car nil}) \, \rightarrow \, \text{NIL} \\ (\text{cdr '(1 . 2)}) \, \rightarrow \, 2 \\ (\text{cdr '(1 2)}) \, \rightarrow \, (2) \end{array}
```

```
\begin{array}{l} (\texttt{cadr '(1 2)}) \rightarrow \texttt{2} \\ (\texttt{car '(a b c)}) \rightarrow \texttt{A} \\ (\texttt{cdr '(a b c)}) \rightarrow (\texttt{B C}) \end{array}
```

Exceptional Situations:

The functions **car** and **cdr** should signal **type-error** if they receive an argument which is not a *list*. The other functions (**caar**, **cadr**, ... **cddddr**) should behave for the purpose of error checking as if defined by appropriate calls to **car** and **cdr**.

See Also:

rplaca, first, rest

Notes:

The car of a cons can also be altered by using **rplaca**, and the cdr of a cons can be altered by using **rplacd**.

```
\begin{array}{lll} (\operatorname{car} x) & \equiv (\operatorname{first} x) \\ (\operatorname{cadr} x) & \equiv (\operatorname{second} x) \equiv (\operatorname{car} (\operatorname{cdr} x)) \\ (\operatorname{caddr} x) & \equiv (\operatorname{third} x) \equiv (\operatorname{car} (\operatorname{cdr} (\operatorname{cdr} x))) \\ (\operatorname{cadddr} x) & \equiv (\operatorname{fourth} x) \equiv (\operatorname{car} (\operatorname{cdr} (\operatorname{cdr} (\operatorname{cdr} x)))) \end{array}
```

copy-tree

Function

Syntax:

```
	ext{copy-tree} tree 	o new-tree
```

Arguments and Values:

```
tree—a tree.
new-tree—a tree.
```

Description:

Creates a copy of a tree of conses.

If *tree* is not a *cons*, it is returned; otherwise, the result is a new *cons* of the results of calling **copy-tree** on the *car* and *cdr* of *tree*. In other words, all *conses* in the *tree* represented by *tree* are copied recursively, stopping only when non-*conses* are encountered.

copy-tree does not preserve circularities and the sharing of substructure.

Examples:

```
(setq object (list (cons 1 "one")
```

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```
(cons 2 (list 'a 'b 'c))))

ightarrow ((1 . "one") (2 A B C))
 (setq object-too object) 
ightarrow ((1 . "one") (2 A B C))
 (setq copy-as-list (copy-list object))
 (setq copy-as-alist (copy-alist object))
 (setq copy-as-tree (copy-tree object))
 (eq object object-too) 
ightarrow true
 (eq copy-as-tree object) 
ightarrow false
 (eql copy-as-tree object) 
ightarrow false
 (equal copy-as-tree object) 
ightarrow true
 (setf (first (cdr (second object))) "a"
        (car (second object)) "two"
        (car object) '(one . 1)) 
ightarrow (ONE . 1)
 \texttt{object} \, \rightarrow \, \texttt{((ONE . 1) ("two" "a" B C))}
 object-too 
ightarrow ((ONE . 1) ("two" "a" B C))
 copy-as-list 
ightarrow ((1 . "one") ("two" "a" B C))
 copy-as-alist 
ightarrow ((1 . "one") (2 "a" B C))
 copy-as-tree 
ightarrow ((1 . "one") (2 A B C))
```

See Also:

tree-equal

sublis, nsublis

Function

Syntax:

```
sublis alist tree &key key test test-not \rightarrow new-tree nsublis alist tree &key key test test-not \rightarrow new-tree
```

Arguments and Values:

```
alist—an association list.
```

tree—a tree.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

new-tree—a tree.

sublis, nsublis

Description:

sublis makes substitutions for *objects* in *tree* (a structure of *conses*). nsublis is like sublis but destructively modifies the relevant parts of the *tree*.

sublis looks at all subtrees and leaves of *tree*; if a subtree or leaf appears as a key in *alist* (that is, the key and the subtree or leaf *satisfy the test*), it is replaced by the *object* with which that key is associated. This operation is non-destructive. In effect, sublis can perform several subst operations simultaneously.

If **sublis** succeeds, a new copy of *tree* is returned in which each occurrence of such a subtree or leaf is replaced by the *object* with which it is associated. If no changes are made, the original tree is returned. The original *tree* is left unchanged, but the result tree may share cells with it.

nsublis is permitted to modify tree but otherwise returns the same values as sublis.

Examples:

```
(sublis '((x . 100) (z . zprime))
          '(plus x (minus g z x p) 4 . x))

ightarrow (PLUS 100 (MINUS G ZPRIME 100 P) 4 . 100)
 (sublis '(((+ x y) . (- x y)) ((- x y) . (+ x y)))
          '(* (/ (+ x y) (+ x p)) (- x y))
          :test #'equal)
\rightarrow (* (/ (- X Y) (+ X P)) (+ X Y))
 (setq tree1 '(1 (1 2) ((1 2 3)) (((1 2 3 4)))))
\rightarrow (1 (1 2) ((1 2 3)) (((1 2 3 4))))
 (sublis '((3 . "three")) tree1)

ightarrow (1 (1 2) ((1 2 "three")) (((1 2 "three" 4))))
 (sublis '((t . "string"))
           (sublis '((1 . "") (4 . 44)) tree1)
           :key #'stringp)
\rightarrow ("string" ("string" 2) (("string" 2 3)) ((("string" 2 3 44))))
tree1 \rightarrow (1 (1 2) ((1 2 3)) (((1 2 3 4))))
 (setq tree2 ', ("one" ("one" "two") (("one" "Two" "three"))))
\rightarrow ("one" ("one" "two") (("one" "Two" "three")))
 (sublis '(("two" . 2)) tree2)

ightarrow ("one" ("one" "two") (("one" "Two" "three")))
tree2 \rightarrow ("one" ("one" "two") (("one" "Two" "three")))
 (sublis '(("two" . 2)) tree2 :test 'equal)

ightarrow ("one" ("one" 2) (("one" "Two" "three")))
 (nsublis '((t . 'temp))
            :key #'(lambda (x) (or (atom x) (< (list-length x) 3))))</pre>

ightarrow ((QUOTE TEMP) (QUOTE TEMP) QUOTE TEMP)
```

Side Effects:

nsublis modifies tree.

See Also:

subst, Section 3.2.1 (Compiler Terminology), Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

Because the side-effecting variants (e.g., nsublis) potentially change the path that is being traversed, their effects in the presence of shared or circular structure structure may vary in surprising ways when compared to their non-side-effecting alternatives. To see this, consider the following side-effect behavior, which might be exhibited by some implementations:

subst, subst-if, subst-if-not, nsubst, nsubst-if, nsubst-if-not Function

Syntax:

```
subst new old tree &key key test test-not \rightarrow new-tree subst-if new predicate tree &key key \rightarrow new-tree subst-if-not new predicate tree &key key \rightarrow new-tree nsubst new old tree &key key test test-not \rightarrow new-tree nsubst-if-not new predicate tree &key key \rightarrow new-tree nsubst-if-not new predicate tree &key key \rightarrow new-tree
```

Arguments and Values:

```
new—an object.
```

subst, subst-if, subst-if-not, nsubst, nsubst-if, ...

predicate—a symbol that names a function, or a function of one argument that returns a generalized boolean value.

```
tree—a tree.
```

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

new-tree—a tree.

Description:

subst, subst-if, and subst-if-not perform substitution operations on *tree*. Each function searches *tree* for occurrences of a particular *old* item of an element or subexpression that *satisfies the test*.

nsubst, nsubst-if, and nsubst-if-not are like subst, subst-if, and subst-if-not respectively, except that the original *tree* is modified.

subst makes a copy of *tree*, substituting *new* for every subtree or leaf of *tree* (whether the subtree or leaf is a *car* or a *cdr* of its parent) such that *old* and the subtree or leaf *satisfy the test*.

nsubst is a destructive version of **subst**. The list structure of *tree* is altered by destructively replacing with *new* each leaf of the *tree* such that *old* and the leaf *satisfy the test*.

For **subst**, **subst-if**, and **subst-if-not**, if the functions succeed, a new copy of the tree is returned in which each occurrence of such an element is replaced by the *new* element or subexpression. If no changes are made, the original *tree* may be returned. The original *tree* is left unchanged, but the result tree may share storage with it.

For **nsubst**, **nsubst-if**, and **nsubst-if-not** the original *tree* is modified and returned as the function result, but the result may not be **eq** to *tree*.

Examples:

```
 \begin{array}{l} (\text{setq tree1 '(1 (1 2) (1 2 3) (1 2 3 4)})) \to (1 (1 2) (1 2 3) (1 2 3 4)) \\ (\text{subst "two" 2 tree1}) \to (1 (1 "two") (1 "two" 3) (1 "two" 3 4)) \\ (\text{subst "five" 5 tree1}) \to (1 (1 2) (1 2 3) (1 2 3 4)) \\ (\text{eq tree1 (subst "five" 5 tree1})) \to implementation-dependent \\ (\text{subst 'tempest 'hurricane} \\ \text{'(shakespeare wrote (the hurricane))}) \\ \to (\text{SHAKESPEARE WROTE (THE TEMPEST)}) \\ (\text{subst 'foo 'nil '(shakespeare wrote (twelfth night))}) \\ \to (\text{SHAKESPEARE WROTE (TWELFTH NIGHT . FOO) . FOO)} \\ (\text{subst '(a . cons) '(old . pair)} \\ \text{'((old . spice) ((old . shoes) old . pair) (old . pair))} \\ \end{array}
```

```
:test #'equal) 

\rightarrow ((OLD . SPICE) ((OLD . SHOES) A . CONS) (A . CONS)) 

(subst-if 5 #'listp tree1) \rightarrow 5 

(subst-if-not '(x) #'consp tree1) 

\rightarrow (1 X) 

tree1 \rightarrow (1 (1 2) (1 2 3) (1 2 3 4)) 

(nsubst 'x 3 tree1 :key #'(lambda (y) (and (listp y) (third y)))) 

\rightarrow (1 (1 2) X X) 

tree1 \rightarrow (1 (1 2) X X)
```

Side Effects:

nsubst, nsubst-if, and nsubst-if-not might alter the tree structure of tree.

See Also:

substitute, nsubstitute, Section 3.2.1 (Compiler Terminology), Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

The functions subst-if-not and nsubst-if-not are deprecated.

One possible definition of **subst**:

tree-equal

tree-equal

Function

Syntax:

tree-equal tree-1 tree-2 &key test test-not \rightarrow generalized-boolean

Arguments and Values:

```
tree-1—a tree.
tree-2—a tree.
```

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

generalized-boolean—a generalized boolean.

Description:

tree-equal tests whether two trees are of the same shape and have the same leaves. **tree-equal** returns *true* if *tree-1* and *tree-2* are both *atoms* and *satisfy the test*, or if they are both *conses* and the *car* of *tree-1* is **tree-equal** to the *car* of *tree-2* and the *cdr* of *tree-1* is **tree-equal** to the *cdr* of *tree-2*. Otherwise, **tree-equal** returns *false*.

tree-equal recursively compares conses but not any other objects that have components.

The first argument to the :test or :test-not function is *tree-1* or a *car* or *cdr* of *tree-1*; the second argument is *tree-2* or a *car* or *cdr* of *tree-2*.

Examples:

```
(setq tree1 '(1 (1 2))  tree2 \text{ '(1 (1 2))} \rightarrow (1 \text{ (1 2)})  (tree-equal tree1 tree2) \rightarrow true (eql tree1 tree2) \rightarrow false (setq tree1 '('a ('b 'c)) \rightarrow ('a ('b 'c)) \rightarrow ((QUOTE A) ((QUOTE B) (QUOTE C))) (tree-equal tree1 tree2 :test 'eq) \rightarrow true
```

Exceptional Situations:

The consequences are undefined if both *tree-1* and *tree-2* are circular.

See Also:

equal, Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

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copy-list Function

Syntax:

```
copy-list list \rightarrow copy
```

Arguments and Values:

list—a proper list or a dotted list.

```
copy—a list.
```

Description:

Returns a copy of list. If list is a dotted list, the resulting list will also be a dotted list.

Only the *list structure* of *list* is copied; the *elements* of the resulting list are the *same* as the corresponding *elements* of the given *list*.

Examples:

Exceptional Situations:

The consequences are undefined if *list* is a *circular list*.

See Also:

copy-alist, copy-seq, copy-tree

Notes:

The copy created is **equal** to *list*, but not **eq**.

list, list*

list, list*

Syntax:

```
\begin{tabular}{ll} list & constraints & c
```

Arguments and Values:

```
object—an object.
list—a list.
result—an object.
```

Description:

list returns a list containing the supplied objects.

list* is like list except that the last argument to list becomes the car of the last cons constructed, while the last argument to list* becomes the cdr of the last cons constructed. Hence, any given call to list* always produces one fewer conses than a call to list with the same number of arguments.

If the last argument to list* is a list, the effect is to construct a new list which is similar, but which has additional elements added to the front corresponding to the preceding arguments of list*.

If list* receives only one object, that object is returned, regardless of whether or not it is a list.

Examples:

```
(list 1) \rightarrow (1)

(list* 1) \rightarrow 1

(setq a 1) \rightarrow 1

(list a 2) \rightarrow (1 2)

'(a 2) \rightarrow (A 2)

(list 'a 2) \rightarrow (A 2)

(list* a 2) \rightarrow (1 . 2)

(list) \rightarrow NIL ; i.e., ()

(setq a '(1 2)) \rightarrow (1 2)

(eq a (list* a)) \rightarrow true

(list 3 4 'a (car '(b . c)) (+ 6 -2)) \rightarrow (3 4 A B 4)

(list* 'a 'b 'c 'd) \equiv (cons 'a (cons 'b (cons 'c 'd))) \rightarrow (A B C . D)

(list* 'a 'b 'c '(d e f)) \rightarrow (A B C D E F)
```

See Also:

cons

Notes:

```
(list* x) \equiv x
```

list-length

Function

Syntax:

```
list-length list \rightarrow length
```

Arguments and Values:

```
list—a proper list or a circular list.
```

length—a non-negative integer, or nil.

Description:

Returns the length of list if list is a proper list. Returns nil if list is a circular list.

Examples:

```
\begin{array}{l} (\mbox{list-length '(a b c d)}) \rightarrow 4 \\ (\mbox{list-length '(a (b c) d)}) \rightarrow 3 \\ (\mbox{list-length '()}) \rightarrow 0 \\ (\mbox{list-length nil}) \rightarrow 0 \\ (\mbox{defun circular-list (\&rest elements)}) \\ (\mbox{let ((cycle (copy-list elements)))} \\ (\mbox{nconc cycle cycle)})) \\ (\mbox{list-length (circular-list 'a 'b)}) \rightarrow \mbox{NIL} \\ (\mbox{list-length (circular-list)}) \rightarrow 0 \\ \end{array}
```

Exceptional Situations:

Should signal an error of type type-error if list is not a proper list or a circular list.

See Also:

length

Notes:

list-length could be implemented as follows:

```
(slow x (cdr slow))) ;Slow pointer: leaps by 1.
  (nil)
;; If fast pointer hits the end, return the count.
(when (endp fast) (return n))
(when (endp (cdr fast)) (return (+ n 1)))
;; If fast pointer eventually equals slow pointer,
;; then we must be stuck in a circular list.
;; (A deeper property is the converse: if we are
;; stuck in a circular list, then eventually the
;; fast pointer will equal the slow pointer.
;; That fact justifies this implementation.)
(when (and (eq fast slow) (> n 0)) (return nil))))
```

listp Function

Syntax:

listp object \rightarrow generalized-boolean

Arguments and Values:

object—an object.

 ${\it generalized-boolean} {--} a {\it generalized boolean}.$

Description:

Returns true if object is of type list; otherwise, returns false.

Examples:

```
\begin{array}{ll} ({\tt listp\ nil}) \,\to\, true \\ ({\tt listp\ (cons\ 1\ 2)}) \,\to\, true \\ ({\tt listp\ (make-array\ 6)}) \,\to\, false \\ ({\tt listp\ t}) \,\to\, false \end{array}
```

See Also:

consp

Notes:

If object is a cons, **listp** does not check whether object is a proper list; it returns true for any kind of list.

```
(listp object) \equiv (typep object 'list) \equiv (typep object '(or cons null))
```

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make-list Function

Syntax:

make-list size &key initial-element \rightarrow list

Arguments and Values:

```
size—a non-negative integer.
initial-element—an object. The default is nil.
list—a list.
```

Description:

Returns a list of length given by size, each of the elements of which is initial-element.

Examples:

```
(make-list 5) \rightarrow (NIL NIL NIL NIL NIL) (make-list 3 :initial-element 'rah) \rightarrow (RAH RAH RAH) (make-list 2 :initial-element '(1 2 3)) \rightarrow ((1 2 3)) (make-list 0) \rightarrow NIL ;i.e., () (make-list 0 :initial-element 'new-element) \rightarrow NIL
```

Exceptional Situations:

Should signal an error of type type-error if size is not a non-negative integer.

See Also:

cons, list

push

Syntax:

```
\mathbf{push} \ \textit{item place} \ \rightarrow \textit{new-place-value}
```

Arguments and Values:

```
item—an object.
place—a place, the value of which may be any object.
new-place-value—a list (the new value of place).
```

Description:

push prepends *item* to the *list* that is stored in *place*, stores the resulting *list* in *place*, and returns the *list*.

For information about the evaluation of subforms of place, see Section 5.1.1.1 (Evaluation of Subforms to Places).

Examples:

```
\begin{array}{l} (\text{setq llst '(nil)}) \rightarrow (\text{NIL}) \\ (\text{push 1 (car llst)}) \rightarrow (1) \\ \text{llst} \rightarrow ((1)) \\ (\text{push 1 (car llst)}) \rightarrow (1 \ 1) \\ \text{llst} \rightarrow ((1 \ 1)) \\ (\text{setq x '(a (b c) d)}) \rightarrow (\text{A (B C) D)} \\ (\text{push 5 (cadr x)}) \rightarrow (\text{5 B C)} \\ \text{x} \rightarrow (\text{A (5 B C) D)} \end{array}
```

Side Effects:

The contents of *place* are modified.

(setf place (cons item place))

See Also:

```
pop, pushnew, Section 5.1 (Generalized Reference)
```

Notes:

```
The effect of (push item\ place) is equivalent to
```

except that the *subforms* of *place* are evaluated only once, and *item* is evaluated before *place*.

pop Macro

Syntax:

```
\mathbf{pop} place \rightarrow element
```

Arguments and Values:

```
place—a place, the value of which is a list (possibly, but necessarily, a dotted list or circular list). element—an object (the car of the contents of place).
```

Description:

pop reads the value of place, remembers the car of the list which was retrieved, writes the cdr of the list back into the place, and finally yields the car of the originally retrieved list.

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For information about the *evaluation* of *subforms* of *place*, see Section 5.1.1.1 (Evaluation of Subforms to Places).

Examples:

```
(setq stack '(a b c)) \rightarrow (A B C)
(pop stack) \rightarrow A
stack \rightarrow (B C)
(setq llst '((1 2 3 4))) \rightarrow ((1 2 3 4))
(pop (car llst)) \rightarrow 1
llst \rightarrow ((2 3 4))
```

Side Effects:

The contents of *place* are modified.

See Also:

push, pushnew, Section 5.1 (Generalized Reference)

Notes:

The effect of $({\tt pop}\ {\it place})$ is roughly equivalent to

```
(prog1 (car place) (setf place (cdr place)))
```

except that the latter would evaluate any subforms of place three times, while pop evaluates them only once.

first, second, third, fourth, fifth, sixth, seventh, eighth, ninth, tenth Accessor

Syntax:

```
first list
                  \rightarrow object
                                         (setf (first list) new-object)
                  \rightarrow object
                                         (setf (second list) new-object)
second list
third list
                 \rightarrow object
                                         (setf (third list) new-object)
fourth list
                  \rightarrow object
                                         (setf (fourth list) new-object)
fifth list
                  \rightarrow object
                                         (setf (fifth list) new-object)
                                         (setf (sixth list) new-object)
sixth list
                  \rightarrow object
seventh list \rightarrow object
                                         (setf (seventh list) new-object)

ightarrow object 
ightarrow object
                                         (setf (eighth list) new-object)
eighth list
ninth list
                                         (setf (ninth list) new-object)
                 \rightarrow \textit{object}
                                         (setf (tenth list) new-object)
tenth list
```

Arguments and Values:

list—a list, which might be a dotted list or a circular list.

first, second, third, fourth, fifth, sixth, seventh, ...

object, new-object—an object.

Description:

The functions first, second, third, fourth, fifth, sixth, seventh, eighth, ninth, and tenth access the first, second, third, fourth, fifth, sixth, seventh, eighth, ninth, and tenth elements of list, respectively. Specifically,

```
(first list)
                \equiv (car list)
(second list)
                \equiv (car (cdr list))
(third list)
               \equiv (car (cddr list))
(fourth list)
               \equiv (car (cdddr list))
                ≡ (car (cddddr list))
(fifth list)
                ≡ (car (cdr (cddddr list)))
(sixth list)
(seventh | list) \equiv (car (cddr (cddddr | list)))
                ≡ (car (cdddr (cddddr list)))
(eighth list)
(ninth list)
                ≡ (car (cddddr (cddddr list)))
(tenth list)
                ≡ (car (cdr (cddddr (cddddr list))))
```

setf can also be used with any of these functions to change an existing component. The same equivalences apply. For example:

```
(setf (fifth list) new-object) ≡ (setf (car (cddddr list)) new-object)
```

Examples:

```
(setq lst '(1 2 3 (4 5 6) ((V)) vi 7 8 9 10)) \rightarrow (1 2 3 (4 5 6) ((V)) VI 7 8 9 10) (first lst) \rightarrow 1 (tenth lst) \rightarrow 10 (fifth lst) \rightarrow ((V)) (second (fourth lst)) \rightarrow 5 (sixth '(1 2 3)) \rightarrow NIL (setf (fourth lst) "four") \rightarrow "four" lst \rightarrow (1 2 3 "four" ((V)) VI 7 8 9 10)
```

See Also:

car, nth

Notes:

first is functionally equivalent to **car**, **second** is functionally equivalent to **cadd**, **third** is functionally equivalent to **caddr**, and **fourth** is functionally equivalent to **cadddr**.

The ordinal numbering used here is one-origin, as opposed to the zero-origin numbering used by **nth**:

```
(fifth x) \equiv (nth 4 x)
```

nth Accessor

Syntax:

```
nth n list \rightarrow object (setf (nth n list) new-object)
```

Arguments and Values:

```
n—a non-negative integer.
list—a list, which might be a dotted list or a circular list.
object—an object.
new-object—an object.
```

Description:

```
nth locates the nth element of list, where the car of the list is the "zeroth" element. Specifically, (nth n list) \equiv (car (nthcdr n list))

nth may be used to specify a place to setf. Specifically, (setf (nth n list) new-object) \equiv (setf (car (nthcdr n list)) new-object)
```

Examples:

```
(nth 0 '(foo bar baz)) \rightarrow F00 (nth 1 '(foo bar baz)) \rightarrow BAR (nth 3 '(foo bar baz)) \rightarrow NIL (setq 0-to-3 (list 0 1 2 3)) \rightarrow (0 1 2 3) (setf (nth 2 0-to-3) "two") \rightarrow "two" 0-to-3 \rightarrow (0 1 "two" 3)
```

See Also:

 $elt,\,first,\,nthcdr$

endp

Syntax:

endp list \rightarrow generalized-boolean

Arguments and Values:

list—a list, which might be a dotted list or a circular list.

generalized-boolean—a generalized boolean.

Description:

Returns true if list is the empty list. Returns false if list is a cons.

Examples:

```
\begin{array}{l} (\texttt{endp nil}) \to true \\ (\texttt{endp '(1 2)}) \to false \\ (\texttt{endp (cddr '(1 2))}) \to true \end{array}
```

Exceptional Situations:

Should signal an error of type type-error if list is not a list.

Notes:

The purpose of **endp** is to test for the end of *proper list*. Since **endp** does not descend into a *cons*, it is well-defined to pass it a *dotted list*. However, if shorter "lists" are iteratively produced by calling **cdr** on such a *dotted list* and those "lists" are tested with **endp**, a situation that has undefined consequences will eventually result when the *non-nil atom* (which is not in fact a *list*) finally becomes the argument to **endp**. Since this is the usual way in which **endp** is used, it is conservative programming style and consistent with the intent of **endp** to treat **endp** as simply a function on *proper lists* which happens not to enforce an argument type of *proper list* except when the argument is *atomic*.

null

Syntax:

 $null\ object\ o boolean$

Arguments and Values:

```
object—an object.
```

boolean—a boolean.

Description:

Returns t if object is the empty list; otherwise, returns nil.

Examples:

```
\begin{array}{l} (\text{null '()}) \rightarrow \text{T} \\ (\text{null nil}) \rightarrow \text{T} \\ (\text{null t}) \rightarrow \text{NIL} \\ (\text{null 1}) \rightarrow \text{NIL} \end{array}
```

See Also:

not

Notes:

null is intended to be used to test for the *empty list* whereas **not** is intended to be used to invert a *boolean* (or *generalized boolean*). Operationally, **null** and **not** compute the same result; which to use is a matter of style.

```
(null object) ≡ (typep object 'null) ≡ (eq object '())
```

nconc Function

Syntax:

 \mathbf{nconc} &rest lists o concatenated-list

Arguments and Values:

list—each but the last must be a *list* (which might be a *dotted list* but must not be a *circular list*); the last *list* may be any *object*.

concatenated-list—a list.

Description:

Returns a *list* that is the concatenation of *lists*. If no *lists* are supplied, (nconc) returns nil. nconc is defined using the following recursive relationship:

```
(nconc) \rightarrow ()

(nconc \ nil \ . \ lists) \equiv (nconc \ . \ lists)

(nconc \ list) \rightarrow list

(nconc \ list-1 \ list-2) \equiv (progn \ (rplacd \ (last \ list-1) \ list-2) \ . \ list-1)

(nconc \ list-1 \ list-2 \ . \ lists) \equiv (nconc \ (nconc \ list-1 \ list-2) \ . \ lists)
```

Examples:

```
\begin{array}{l} (\texttt{nconc}) \, \to \, \texttt{NIL} \\ (\texttt{setq x '(a b c)}) \, \to \, (\texttt{A B C}) \\ (\texttt{setq y '(d e f)}) \, \to \, (\texttt{D E F}) \\ (\texttt{nconc x y}) \, \to \, (\texttt{A B C D E F}) \\ \texttt{x} \, \to \, (\texttt{A B C D E F}) \end{array}
```

Note, in the example, that the value of x is now different, since its last *cons* has been **rplacd**'d to the value of y. If (nconc x y) were evaluated again, it would yield a piece of a *circular list*, whose printed representation would be (A B C D E F D E F D E F ...), repeating forever; if the *print-circle* switch were *non-nil*, it would be printed as (A B C . #1=(D E F . #1#)).

Side Effects:

The *lists* are modified rather than copied.

See Also:

append, concatenate

append

Syntax:

```
{f append \& rest \it lists} 
ightarrow {\it result}
```

Arguments and Values:

list—each must be a proper list except the last, which may be any object.

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result—an object. This will be a list unless the last list was not a list and all preceding lists were null.

Description:

append returns a new *list* that is the concatenation of the copies. *lists* are left unchanged; the *list* structure of each of *lists* except the last is copied. The last argument is not copied; it becomes the cdr of the final dotted pair of the concatenation of the preceding *lists*, or is returned directly if there are no preceding non-empty lists.

Examples:

```
(append '(a b c) '(d e f) '() '(g)) \rightarrow (A B C D E F G) (append '(a b c) 'd) \rightarrow (A B C . D) (setq lst '(a b c)) \rightarrow (A B C D) (append lst '(d)) \rightarrow (A B C D) lst \rightarrow (A B C) (append) \rightarrow NIL (append 'a) \rightarrow A
```

See Also:

nconc, concatenate

revappend, nreconc

Function

Syntax:

```
revappend list tail \rightarrow result-list
nreconc list tail \rightarrow result-list
```

Arguments and Values:

```
list—a proper list.
tail—an object.
result-list—an object.
```

Description:

revappend constructs a $copy_2$ of list, but with the *elements* in reverse order. It then appends (as if by **nconc**) the tail to that reversed list and returns the result.

nreconc reverses the order of *elements* in *list* (as if by **nreverse**). It then appends (as if by **nconc**) the *tail* to that reversed list and returns the result.

The resulting list shares list structure with tail.

revappend, nreconc

Examples:

```
(let ((list-1 (list 1 2 3))
        (list-2 (list 'a 'b 'c)))
   (print (revappend list-1 list-2))
   (print (equal list-1 '(1 2 3)))
   (print (equal list-2 '(a b c))))
\triangleright T
\triangleright T
\rightarrow \ \mathtt{T}
 (revappend '(1 2 3) '()) \rightarrow (3 2 1)
 (revappend '(1 2 3) '(a . b)) 
ightarrow (3 2 1 A . B)
 (revappend '() '(a b c)) 
ightarrow (A B C)
 (revappend '(1 2 3) 'a) 
ightarrow (3 2 1 . A)
 (revappend '() 'a) 
ightarrow A ;degenerate case
 (let ((list-1 '(1 2 3))
        (list-2 '(a b c)))
    (print (nreconc list-1 list-2))
   (print (equal list-1 '(1 2 3)))
   (print (equal list-2 '(a b c))))
▷ NIL
\triangleright T
\rightarrow T
```

Side Effects:

revappend does not modify either of its *arguments*. **nreconc** is permitted to modify *list* but not *tail*.

Although it might be implemented differently, **nreconc** is constrained to have side-effect behavior equivalent to:

```
(nconc (nreverse list) tail)
```

See Also:

reverse, nreverse, nconc

Notes:

The following functional equivalences are true, although good *implementations* will typically use a faster algorithm for achieving the same effect:

```
(revappend list tail) ≡ (nconc (reverse list) tail)
(nreconc list tail) ≡ (nconc (nreverse list) tail)
```

butlast, nbutlast

Function

Syntax:

```
butlast list &optional n 	o result-list
nbutlast list &optional n 	o result-list
```

Arguments and Values:

list—a list, which might be a dotted list but must not be a circular list.

n—a non-negative integer.

result-list—a list.

Description:

butlast returns a copy of *list* from which the last n conses have been omitted. If n is not supplied, its value is 1. If there are fewer than n conses in *list*, nil is returned and, in the case of nbutlast, *list* is not modified.

nbutlast is like **butlast**, but **nbutlast** may modify *list*. It changes the cdr of the cons n+1 from the end of the *list* to nil.

Examples:

```
(\mathsf{setq}\; \mathsf{1st}\; \, \mathsf{'(1\;2\;3\;4\;5\;6\;7\;8\;9)}) \;\rightarrow\; (\mathsf{1\;2\;3\;4\;5\;6\;7\;8\;9})
(butlast lst) \rightarrow (1 2 3 4 5 6 7 8)
(butlast 1st 5) \rightarrow (1 2 3 4)
(butlast 1st (+ 5 5)) \rightarrow NIL
lst \rightarrow (1 2 3 4 5 6 7 8 9)
(nbutlast 1st 3) \rightarrow (1 2 3 4 5 6)
lst \rightarrow (1 2 3 4 5 6)
(nbutlast 1st 99) \rightarrow NIL
lst \rightarrow (1 2 3 4 5 6)
(butlast '(a b c d)) \rightarrow (A B C)
(butlast '((a b) (c d))) \rightarrow ((A B))
(butlast '(a)) 
ightarrow NIL
(butlast nil) 
ightarrow NIL
(setq foo (list 'a 'b 'c 'd)) 
ightarrow (A B C D)
(nbutlast foo) 
ightarrow (A B C)
foo 
ightarrow (A B C)
```

```
\begin{array}{l} \text{(nbutlast (list 'a))} \, \to \, \text{NIL} \\ \text{(nbutlast '())} \, \to \, \text{NIL} \end{array}
```

Exceptional Situations:

Should signal an error of type **type-error** if list is not a proper list or a dotted list. Should signal an error of type **type-error** if n is not a non-negative integer.

Notes:

```
(butlast list n) \equiv (ldiff list (last list n))
```

last

Syntax:

last list &optional n o tail

Arguments and Values:

list—a list, which might be a dotted list but must not be a circular list.

n—a non-negative integer. The default is 1.

tail—an object.

Description:

last returns the last n conses (not the last n elements) of list). If list is (), last returns ().

If n is zero, the atom that terminates list is returned. If n is greater than or equal to the number of cons cells in list, the result is list.

Examples:

```
\begin{array}{l} (\text{last nil}) \, \to \, \text{NIL} \\ (\text{last '(1 2 3)}) \, \to \, (3) \\ (\text{last '(1 2 . 3)}) \, \to \, (2 . 3) \\ (\text{setq x (list 'a 'b 'c 'd)}) \, \to \, (\text{A B C D}) \\ (\text{last x}) \, \to \, (\text{D}) \\ (\text{rplacd (last x) (list 'e 'f)}) \, \, \text{x} \, \to \, (\text{A B C D E F}) \\ (\text{last x}) \, \to \, (\text{F}) \\ (\text{last '(a b c)}) \, \to \, (\text{C}) \\ (\text{last '(a b c) 1)} \, \to \, (\text{C}) \end{array}
```

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Exceptional Situations:

The consequences are undefined if list is a $circular\ list$. Should signal an error of $type\ \mathbf{type\text{-error}}$ if n is not a non-negative integer.

See Also:

butlast, nth

Notes:

The following code could be used to define last.

ldiff, tailp

Function

Syntax:

```
\begin{array}{ll} \textbf{ldiff} \textit{ list object} & \rightarrow \textit{result-list} \\ \textbf{tailp object list} & \rightarrow \textit{generalized-boolean} \end{array}
```

Arguments and Values:

```
list—a list, which might be a dotted list.
object—an object.
result-list—a list.
generalized-boolean—a generalized boolean.
```

ldiff, tailp

Description:

If object is the same as some tail of list, tailp returns true; otherwise, it returns false.

If object is the same as some tail of list, ldiff returns a fresh list of the elements of list that precede object in the list structure of list; otherwise, it returns a $copy_2$ of list.

Examples:

```
(let ((lists '#((a b c) (a b c . d))))
   (dotimes (i (length lists)) ()
     (let ((list (aref lists i)))
       (format t "~2&list=~S ~21T(tailp object list)~
                   ~44T(ldiff list object)~%" list)
         (let ((objects (vector list (cddr list) (copy-list (cddr list))
                                 '(f g h) '() 'd 'x)))
           (dotimes (j (length objects)) ()
             (let ((object (aref objects j)))
               (format t "~& object=~S ~21T~S ~44T~S"
                       object (tailp object list) (ldiff list object)))))))
\triangleright
▷ list=(A B C)
                        (tailp object list)
                                                (ldiff list object)
▷ object=(A B C)
                                                NIL
▷ object=(C)
                       T
                                                (A B)
▷ object=(C)
                       NIL
                                                (A B C)
                       NIL
▷ object=(F G H)
                                                (A B C)
                                                (ABC)
▷ object=NIL
                       Т
\triangleright object=D
                       NIL
                                                (ABC)
▷ object=X
                       NIL
                                                (ABC)
\triangleright
▷ list=(A B C . D)
                        (tailp object list)
                                                (ldiff list object)
▷ object=(A B C . D)
                       Τ
                                                NIL
▷ object=(C . D)
                       Т
                                                (A B)
▷ object=(C . D)
                       NIL
                                                (A B C . D)
▷ object=(F G H)
                                                (A B C . D)
                       NIL
                                                (A B C . D)
▷ object=NIL
                       NIL
                                                (ABC)
▷ object=D
                       Т
▷ object=X
                       NIL
                                                (A B C . D)

ightarrow NIL
```

Side Effects:

Neither ldiff nor tailp modifies either of its arguments.

Exceptional Situations:

Should be prepared to signal an error of type type-error if list is not a proper list or a dotted list.

See Also:

set-difference

Notes:

If the *list* is a *circular list*, **tailp** will reliably *yield* a *value* only if the given *object* is in fact a *tail* of *list*. Otherwise, the consequences are unspecified: a given *implementation* which detects the circularity must return *false*, but since an *implementation* is not obliged to detect such a *situation*, **tailp** might just loop indefinitely without returning in that case.

tailp could be defined as follows:

nthcdr Function

Syntax:

 $\mathbf{nthcdr}\ \mathit{n}\ \mathit{list}\ o \mathit{tail}$

Arguments and Values:

```
n—a non-negative integer.
```

list—a list, which might be a $dotted\ list$ or a $circular\ list$.

tail—an object.

Description:

Returns the tail of list that would be obtained by calling cdr n times in succession.

Examples:

```
(nthcdr 0 '()) → NIL
(nthcdr 3 '()) → NIL
(nthcdr 0 '(a b c)) → (A B C)
(nthcdr 2 '(a b c)) → (C)
(nthcdr 4 '(a b c)) → ()
(nthcdr 1 '(0 . 1)) → 1

(locally (declare (optimize (safety 3)))
  (nthcdr 3 '(0 . 1)))
Error: Attempted to take CDR of 1.
```

Exceptional Situations:

Should signal an error of type type-error if n is not a non-negative integer.

For n being an integer greater than 1, the error checking done by (nthcdr n list) is the same as for (nthcdr $(-n \ 1)$ (cdr list)); see the function cdr.

See Also:

cdr, nth, rest

rest

Syntax:

```
\operatorname{rest} list 	o tail (\operatorname{setf} (rest list) new-tail)
```

Arguments and Values:

```
list—a list, which might be a dotted list or a circular list. tail—an object.
```

Description:

rest performs the same operation as cdr, but mnemonically complements first. Specifically,

```
(rest list) ≡ (cdr list)
(setf (rest list) new-tail) ≡ (setf (cdr list) new-tail)
```

Examples:

```
(rest '(1 2)) 
ightarrow (2)
```

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```
\begin{array}{l} (\text{rest '(1 . 2)}) \rightarrow 2 \\ (\text{rest '(1)}) \rightarrow \text{NIL} \\ (\text{setq *cons* '(1 . 2)}) \rightarrow (\text{1 . 2}) \\ (\text{setf (rest *cons*) "two"}) \rightarrow \text{"two"} \\ *\text{cons*} \rightarrow (\text{1 . "two"}) \end{array}
```

See Also:

cdr, nthcdr

Notes:

 \mathbf{rest} is often preferred stylistically over \mathbf{cdr} when the argument is to being subjectively viewed as a list rather than as a cons.

member, member-if, member-if-not

Function

Syntax:

```
member item list &key key test test-not \rightarrow tail member-if predicate list &key key \rightarrow tail member-if-not predicate list &key key \rightarrow tail
```

Arguments and Values:

```
item—an object.
```

list—a proper list.

predicate—a designator for a function of one argument that returns a generalized boolean.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

tail—a list.

Description:

member, member-if, and member-if-not each search *list* for *item* or for a top-level element that satisfies the test. The argument to the predicate function is an element of list.

If some element *satisfies the test*, the tail of *list* beginning with this element is returned; otherwise nil is returned.

list is searched on the top level only.

Examples:

Exceptional Situations:

Should be prepared to signal an error of type type-error if list is not a proper list.

See Also:

find, position, Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

The function member-if-not is deprecated.

In the following

```
(member 'a '(g (a y) c a d e a f)) \rightarrow (A D E A F)
```

the value returned by **member** is *identical* to the portion of the *list* beginning with a. Thus **rplaca** on the result of **member** can be used to alter the part of the *list* where a was found (assuming a check has been made that **member** did not return **nil**).

mapc, mapcar, mapcan, mapl, maplist, mapcon Function

Syntax:

```
mapc function &rest lists^+ 	o  list-1 mapcar function &rest lists^+ 	o  result-list mapcan function &rest lists^+ 	o  concatenated-results mapl function &rest lists^+ 	o  list-1 maplist function &rest lists^+ 	o  result-list
```

mapc, mapcar, mapcan, mapl, maplist, mapcon

mapcon function &rest lists $^+$ \rightarrow concatenated-results

Arguments and Values:

function—a designator for a function that must take as many arguments as there are lists.

```
list—a proper list.
```

list-1—the first *list* (which must be a *proper list*).

result-list—a list.

concatenated-results—a list.

Description:

The mapping operation involves applying *function* to successive sets of arguments in which one argument is obtained from each *sequence*. Except for **mapc** and **mapl**, the result contains the results returned by *function*. In the cases of **mapc** and **mapl**, the resulting *sequence* is *list*.

function is called first on all the elements with index 0, then on all those with index 1, and so on. result-type specifies the type of the resulting sequence. If function is a symbol, it is coerced to a function as if by symbol-function.

mapcar operates on successive *elements* of the *lists. function* is applied to the first *element* of each *list*, then to the second *element* of each *list*, and so on. The iteration terminates when the shortest *list* runs out, and excess elements in other lists are ignored. The value returned by mapcar is a *list* of the results of successive calls to *function*.

mapc is like mapcar except that the results of applying function are not accumulated. The list argument is returned.

maplist is like mapcar except that function is applied to successive sublists of the lists. function is first applied to the lists themselves, and then to the cdr of each list, and then to the cdr of each list, and so on.

mapl is like maplist except that the results of applying function are not accumulated; list-1 is returned.

mapcan and mapcon are like mapcar and maplist respectively, except that the results of applying function are combined into a list by the use of nconc rather than list. That is,

and similarly for the relationship between mapcan and mapcar.

Examples:

```
(mapcar #'car '((1 a) (2 b) (3 c))) \rightarrow (1 2 3) (mapcar #'abs '(3 -4 2 -5 -6)) \rightarrow (3 4 2 5 6)
```

```
(mapcar #'cons '(a b c) '(1 2 3)) \rightarrow ((A . 1) (B . 2) (C . 3))
 (maplist #'append '(1 2 3 4) '(1 2) '(1 2 3))
\rightarrow ((1 2 3 4 1 2 1 2 3) (2 3 4 2 2 3))
 (maplist #'(lambda (x) (cons 'foo x)) '(a b c d))

ightarrow ((F00 A B C D) (F00 B C D) (F00 C D) (F00 D))
 (maplist #'(lambda (x) (if (member (car x) (cdr x)) 0 1)) '(a b a c d b c))
\rightarrow (0 0 1 0 1 1 1)
; An entry is 1 if the corresponding element of the input
; list was the last instance of that element in the input list.
 (setq dummy nil) 
ightarrow NIL
 (mapc #'(lambda (&rest x) (setq dummy (append dummy x)))
          (1 2 3 4)
         '(a b c d e)
         (x y z)) \rightarrow (1 2 3 4)
 dummy \rightarrow (1 A X 2 B Y 3 C Z)
 (\mathtt{setq}\ \mathtt{dummy}\ \mathtt{nil})\ \to\ \mathtt{NIL}
 (mapl #'(lambda (x) (push x dummy)) '(1 2 3 4)) \rightarrow (1 2 3 4)
 \mathtt{dummy} \,\rightarrow\, (\texttt{(4)} \,\, \texttt{(3 4)} \,\, \texttt{(2 3 4)} \,\, \texttt{(1 2 3 4))}
 (mapcan #'(lambda (x y) (if (null x) nil (list x y)))
            '(nil nil nil d e)
            '(1 2 3 4 5 6)) \rightarrow (D 4 E 5)
 (mapcan #'(lambda (x) (and (numberp x) (list x)))
            '(a 1 b c 3 4 d 5))
\rightarrow (1 3 4 5)
In this case the function serves as a filter; this is a standard Lisp idiom using mapcan.
 (mapcon #'list '(1 2 3 4)) \rightarrow ((1 2 3 4) (2 3 4) (3 4) (4))
```

Exceptional Situations:

Should be prepared to signal an error of type type-error if any list is not a proper list.

See Also:

dolist, map, Section 3.6 (Traversal Rules and Side Effects)

acons

Syntax:

 $\mathbf{acons}\ \mathit{key}\ \mathit{datum}\ \mathit{alist}\ \to \mathit{new-alist}$

Arguments and Values:

```
key—an object.
datum—an object.
alist—an association list.
new-alist—an association list.
```

Description:

Creates a *fresh cons*, the cdr of which is *alist* and the car of which is another $fresh\ cons$, the car of which is key and the cdr of which is datum.

Examples:

```
\label{eq:cons_section} \begin{array}{l} (\text{setq alist '()}) \rightarrow \text{NIL} \\ (\text{acons 1 "one" alist}) \rightarrow ((1 \ . \ "one")) \\ \text{alist} \rightarrow \text{NIL} \\ (\text{setq alist (acons 1 "one" (acons 2 "two" alist))}) \rightarrow ((1 \ . \ "one") \ (2 \ . \ "two")) \\ (\text{assoc 1 alist}) \rightarrow (1 \ . \ "one") \\ (\text{setq alist (acons 1 "uno" alist)}) \rightarrow ((1 \ . \ "uno") \ (1 \ . \ "one") \ (2 \ . \ "two")) \\ (\text{assoc 1 alist}) \rightarrow (1 \ . \ "uno") \\ \end{array}
```

See Also:

assoc, pairlis

Notes:

 $(acons key datum alist) \equiv (cons (cons key datum) alist)$

assoc, assoc-if, assoc-if-not

Function

Syntax:

```
assoc item alist &key key test test-not \rightarrow entry assoc-if predicate alist &key key \rightarrow entry
```

assoc, assoc-if, assoc-if-not

assoc-if-not predicate alist &key key \rightarrow entry

Arguments and Values:

```
item—an object.
```

alist—an association list.

predicate—a designator for a function of one argument that returns a generalized boolean.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

entry—a cons that is an element of alist, or nil.

Description:

assoc, assoc-if, and assoc-if-not return the first cons in alist whose car satisfies the test, or nil if no such cons is found.

For assoc, assoc-if, and assoc-if-not, if nil appears in alist in place of a pair, it is ignored.

Examples:

```
(setq values '((x . 100) (y . 200) (z . 50))) \rightarrow ((X . 100) (Y . 200) (Z . 50))
 (assoc 'y values) 
ightarrow (Y . 200)
 (rplacd (assoc 'y values) 201) 
ightarrow (Y . 201)
 (assoc 'y values) 
ightarrow (Y . 201)
 (setq alist '((1 . "one")(2 . "two")(3 . "three")))

ightarrow ((1 . "one") (2 . "two") (3 . "three"))
 (assoc 2 alist) \rightarrow (2 . "two")
 (assoc-if #'evenp alist) 
ightarrow (2 . "two")
 (assoc-if-not \#'(lambda(x) (< x 3)) alist) \to (3 . "three")
 (\mathtt{setq\ alist\ `(("one"\ .\ 1)("two"\ .\ 2)))}\ \rightarrow\ (("one"\ .\ 1)\ ("two"\ .\ 2))
 (assoc "one" alist) 
ightarrow NIL
 (assoc "one" alist :test #'equalp) 
ightarrow ("one" . 1)
 (assoc "two" alist :key \#'(lambda(x) (char x 2))) \rightarrow NIL
 (assoc #\o alist :key #'(lambda(x) (char x 2))) \rightarrow ("two" . 2)
 (assoc 'r '((a . b) (c . d) (r . x) (s . y) (r . z))) 
ightarrow (R . X)
 (assoc 'goo '((foo . bar) (zoo . goo))) 
ightarrow NIL
 (assoc '2 '((1 a b c) (2 b c d) (-7 x y z))) 
ightarrow (2 B C D)
(setq alist '(("one" . 1) ("2" . 2) ("three" . 3))) \rightarrow (("one" . 1) ("2" . 2) ("three" . 3))
 (assoc-if-not #'alpha-char-p alist
                  :key \#'(lambda (x) (char x 0))) \rightarrow ("2" . 2)
```

Exceptional Situations:

Should be prepared to signal an error of type type-error if alist is not an association list.

See Also:

rassoc, find, member, position, Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

The function assoc-if-not is deprecated.

It is possible to **rplacd** the result of **assoc**, provided that it is not **nil**, in order to "update" alist.

The two expressions

```
(assoc item list :test fn)
and
(find item list :test fn :key #'car)
```

are equivalent in meaning with one exception: if **nil** appears in *alist* in place of a pair, and *item* is **nil**, **find** will compute the *car* of the **nil** in *alist*, find that it is equal to *item*, and return **nil**, whereas **assoc** will ignore the **nil** in *alist* and continue to search for an actual *cons* whose *car* is **nil**.

copy-alist Function

Syntax:

 $\operatorname{copy-alist}$ alist \to new-alist

Arguments and Values:

alist—an association list.

new-alist—an association list.

Description:

copy-alist returns a copy of alist.

The *list structure* of *alist* is copied, and the *elements* of *alist* which are *conses* are also copied (as *conses* only). Any other *objects* which are referred to, whether directly or indirectly, by the *alist* continue to be shared.

Examples:

See Also:

copy-list

pairlis

Syntax:

pairlis keys data &optional alist \rightarrow new-alist

Arguments and Values:

```
keys—a proper list.

data—a proper list.

alist—an association list. The default is the empty list.

new-alist—an association list.
```

Description:

Returns an association list that associates elements of keys to corresponding elements of data. The consequences are undefined if keys and data are not of the same length.

If *alist* is supplied, **pairlis** returns a modified *alist* with the new pairs prepended to it. The new pairs may appear in the resulting *association list* in either forward or backward order. The result of

```
(pairlis '(one two) '(1 2) '((three . 3) (four . 19)))
might be
((one . 1) (two . 2) (three . 3) (four . 19))
```

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```
or ((two . 2) (one . 1) (three . 3) (four . 19))
```

Examples:

Exceptional Situations:

Should be prepared to signal an error of type type-error if keys and data are not proper lists.

See Also:

acons

rassoc, rassoc-if, rassoc-if-not

Function

Syntax:

```
rassoc item alist &key key test test-not \rightarrow entry rassoc-if predicate alist &key key \rightarrow entry rassoc-if-not predicate alist &key key \rightarrow entry
```

Arguments and Values:

```
item—an object.
```

alist—an association list.

predicate—a designator for a function of one argument that returns a generalized boolean.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

entry—a cons that is an element of the alist, or nil.

Description:

rassoc, rassoc-if, and rassoc-if-not return the first cons whose cdr satisfies the test. If no such cons is found, nil is returned.

If nil appears in alist in place of a pair, it is ignored.

Examples:

```
(setq alist '((1 . "one") (2 . "two") (3 . 3))) 
 \rightarrow ((1 . "one") (2 . "two") (3 . 3))
(rassoc 3 alist) \rightarrow (3 . 3)
(rassoc "two" alist) \rightarrow NIL
(rassoc "two" alist :test 'equal) \rightarrow (2 . "two")
(rassoc 1 alist :key #'(lambda (x) (if (numberp x) (/ x 3)))) \rightarrow (3 . 3)
(rassoc 'a '((a . b) (b . c) (c . a) (z . a))) \rightarrow (C . A)
(rassoc-if #'stringp alist) \rightarrow (1 . "one")
(rassoc-if-not #'vectorp alist) \rightarrow (3 . 3)
```

See Also:

assoc, Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

The function rassoc-if-not is deprecated.

It is possible to rplaca the result of rassoc, provided that it is not nil, in order to "update" alist.

The expressions

```
(rassoc item list :test fn)
and
(find item list :test fn :key #'cdr)
```

are equivalent in meaning, except when the item is nil and nil appears in place of a pair in the alist. See the function assoc.

get-properties

Function

Syntax:

get-properties plist indicator-list \rightarrow indicator, value, tail

Arguments and Values:

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```
plist—a property list.

indicator-list—a proper list (of indicators).

indicator—an object that is an element of indicator-list.

value—an object.

tail—a list.
```

Description:

get-properties is used to look up any of several property list entries all at once.

It searches the *plist* for the first entry whose *indicator* is *identical* to one of the *objects* in *indicator-list*. If such an entry is found, the *indicator* and *value* returned are the *property indicator* and its associated *property value*, and the *tail* returned is the *tail* of the *plist* that begins with the found entry (*i.e.*, whose *car* is the *indicator*). If no such entry is found, the *indicator*, *value*, and *tail* are all nil.

Examples:

```
(setq x '()) \rightarrow NIL

(setq *indicator-list* '(prop1 prop2)) \rightarrow (PROP1 PROP2)

(getf x 'prop1) \rightarrow NIL

(setf (getf x 'prop1) 'val1) \rightarrow VAL1

(eq (getf x 'prop1) 'val1) \rightarrow true

(get-properties x *indicator-list*) \rightarrow PROP1, VAL1, (PROP1 VAL1)

x \rightarrow (PROP1 VAL1)
```

See Also:

get, getf

getf

Syntax:

```
\begin{tabular}{ll} {\bf getf} \ \it plist \ \it indicator \ \& optional \ \it default \ \rightarrow value \\ \end{tabular} \begin{tabular}{ll} {\bf getf} \ \it place \ \it indicator \ \& optional \ \it default) \ \it new-value) \end{tabular}
```

Arguments and Values:

```
plist—a property list.
place—a place, the value of which is a property list.
```

getf

```
indicator—an object.
default—an object. The default is nil.
value—an object.
new-value—an object.
```

Description:

getf finds a property on the plist whose property indicator is identical to indicator, and returns its corresponding property value. If there are multiple properties₁ with that property indicator, getf uses the first such property. If there is no property with that property indicator, default is returned.

setf of getf may be used to associate a new object with an existing indicator in the property list held by place, or to create a new association if none exists. If there are multiple properties with that property indicator, setf of getf associates the new-value with the first such property. When a getf form is used as a setf place, any default which is supplied is evaluated according to normal left-to-right evaluation rules, but its value is ignored.

setf of getf is permitted to either write the value of place itself, or modify of any part, car or cdr, of the list structure held by place.

Examples:

```
(setq x '()) 
ightarrow NIL
  (\texttt{getf x 'prop1}) \ \rightarrow \ \texttt{NIL}
  (getf x 'prop1 7) 
ightarrow 7
  (getf x 'prop1) \rightarrow NIL
  (setf (getf x 'prop1) 'val1) 
ightarrow VAL1
  (eq (getf x 'prop1) 'val1) 
ightarrow true
  (getf x 'prop1) 
ightarrow VAL1
  (getf x 'prop1 7) 
ightarrow VAL1
  x \rightarrow (PROP1 VAL1)
;; Examples of implementation variation permitted.
  (setq foo (list 'a 'b 'c 'd 'e 'f)) 
ightarrow (A B C D E F)
  (setq bar (cddr foo)) 
ightarrow (C D E F)
  (remf foo 'c) 
ightarrow true
  foo \rightarrow (A B E F)
 bar
\begin{array}{ccc} \rightarrow & (\texttt{C} \ \texttt{D} \ \texttt{E} \ \texttt{F}) \\ \stackrel{or}{\rightarrow} & (\texttt{C}) \end{array}
\stackrel{\overrightarrow{or}}{\rightarrow} (NIL)
\stackrel{\overrightarrow{or}}{\rightarrow} (C NIL)
\stackrel{-}{\rightarrow} (C D)
```

See Also:

get, get-properties, setf, Section 5.1.2.2 (Function Call Forms as Places)

Notes:

There is no way (using **getf**) to distinguish an absent property from one whose value is **default**; but see **get-properties**.

Note that while supplying a *default* argument to **getf** in a **setf** situation is sometimes not very interesting, it is still important because some macros, such as **push** and **incf**, require a *place* argument which data is both *read* from and *written* to. In such a context, if a *default* argument is to be supplied for the *read* situation, it must be syntactically valid for the *write* situation as well. For example,

```
(let ((plist '()))
  (incf (getf plist 'count 0))
  plist) → (COUNT 1)
```

 ${f remf}$

Syntax:

 \mathbf{remf} place indicator \rightarrow generalized-boolean

Arguments and Values:

```
place—a place.
indicator—an object.
generalized-boolean—a generalized boolean.
```

Description:

remf removes from the property list stored in place a property₁ with a property indicator identical to indicator. If there are multiple properties₁ with the identical key, remf only removes the first such property. remf returns false if no such property was found, or true if a property was found.

The property indicator and the corresponding property value are removed in an undefined order by destructively splicing the property list. **remf** is permitted to either **setf** place or to **setf** any part, **car** or **cdr**, of the *list structure* held by that place.

For information about the *evaluation* of *subforms* of *place*, see Section 5.1.1.1 (Evaluation of Subforms to Places).

Examples:

```
(setq x (cons () ())) \rightarrow (NIL)
(setf (getf (car x) 'prop1) 'val1) \rightarrow VAL1
(remf (car x) 'prop1) \rightarrow true
(remf (car x) 'prop1) \rightarrow false
```

Side Effects:

The property list stored in *place* is modified.

See Also:

remprop, getf

intersection, nintersection

Function

Syntax:

```
intersection list-1 list-2 &key key test test-not \rightarrow result-list nintersection list-1 list-2 &key key test test-not \rightarrow result-list
```

Arguments and Values:

```
list-1—a proper list.
```

list-2—a proper list.

test—a designator for a function of two arguments that returns a generalized boolean.

 $\textbf{\textit{test-not}} {--} a \ \textit{designator} \ \text{for a } \textit{function} \ \text{of two} \ \textit{arguments} \ \text{that returns a} \ \textit{generalized boolean}.$

key—a designator for a function of one argument, or nil.

result-list—a list.

Description:

intersection and **nintersection** return a *list* that contains every element that occurs in both *list-1* and *list-2*.

nintersection is the destructive version of **intersection**. It performs the same operation, but may destroy *list-1* using its cells to construct the result. *list-2* is not destroyed.

The intersection operation is described as follows. For all possible ordered pairs consisting of one element from list-1 and one element from list-2, :test or :test-not are used to determine whether they satisfy the test. The first argument to the :test or :test-not function is an element of list-1; the second argument is an element of list-2. If :test or :test-not is not supplied, eql is used. It is an error if :test and :test-not are supplied in the same function call.

intersection, nintersection

If :key is supplied (and not nil), it is used to extract the part to be tested from the *list* element. The argument to the :key function is an element of either *list-1* or *list-2*; the :key function typically returns part of the supplied element. If :key is not supplied or nil, the *list-1* and *list-2* elements are used.

For every pair that *satisfies the test*, exactly one of the two elements of the pair will be put in the result. No element from either *list* appears in the result that does not *satisfy the test* for an element from the other *list*. If one of the *lists* contains duplicate elements, there may be duplication in the result.

There is no guarantee that the order of elements in the result will reflect the ordering of the arguments in any particular way. The result *list* may share cells with, or be **eq** to, either *list-1* or *list-2* if appropriate.

Examples:

```
(setq list1 (list 1 1 2 3 4 a b c "A" "B" "C" "d")
       list2 (list 1 4 5 b c d "a" "B" "c" "D"))

ightarrow (1 4 5 B C D "a" "B" "c" "D")
 (intersection list1 list2) \rightarrow (C B 4 1 1)
 (intersection list1 list2 :test 'equal) 
ightarrow ("B" C B 4 1 1)
 (intersection list1 list2 :test #'equalp) → ("d" "C" "B" "A" C B 4 1 1)
 (nintersection list1 list2) \rightarrow (1 1 4 B C)
 list1 \rightarrow implementation-dependent; e.g., (1 1 4 B C)
 list2 
ightarrow implementation-dependent ; e.g., (1 4 5 B C D "a" "B" "c" "D")
 (setq list1 (copy-list '((1 . 2) (2 . 3) (3 . 4) (4 . 5))))
\rightarrow ((1 . 2) (2 . 3) (3 . 4) (4 . 5))
 (setq list2 (copy-list '((1 . 3) (2 . 4) (3 . 6) (4 . 8))))
\rightarrow ((1 . 3) (2 . 4) (3 . 6) (4 . 8))
 (nintersection list1 list2 :key #'cdr) \rightarrow ((2 . 3) (3 . 4))
 list1 \rightarrow implementation-dependent ; e.g., ((1 . 2) (2 . 3) (3 . 4))
 list2 \rightarrow implementation-dependent; e.g., ((1 . 3) (2 . 4) (3 . 6) (4 . 8))
```

Side Effects:

nintersection can modify list-1, but not list-2.

Exceptional Situations:

Should be prepared to signal an error of type type-error if list-1 and list-2 are not proper lists.

See Also:

union, Section 3.2.1 (Compiler Terminology), Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

Since the **nintersection** side effect is not required, it should not be used in for-effect-only posi-

tions in portable code.

adjoin Function

Syntax:

adjoin item list &key key test test-not \rightarrow new-list

Arguments and Values:

```
item—an object.
```

list—a proper list.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

new-list—a list.

Description:

Tests whether *item* is the same as an existing element of *list*. If the *item* is not an existing element, **adjoin** adds it to *list* (as if by **cons**) and returns the resulting *list*; otherwise, nothing is added and the original *list* is returned.

The *test*, *test-not*, and *key* affect how it is determined whether *item* is the same as an *element* of *list*. For details, see Section 17.2.1 (Satisfying a Two-Argument Test).

Examples:

```
(setq slist '()) \rightarrow NIL (adjoin 'a slist) \rightarrow (A) slist \rightarrow NIL (setq slist (adjoin '(test-item 1) slist)) \rightarrow ((TEST-ITEM 1)) (adjoin '(test-item 1) slist) \rightarrow ((TEST-ITEM 1) (TEST-ITEM 1)) (adjoin '(test-item 1) slist :test 'equal) \rightarrow ((TEST-ITEM 1)) (adjoin '(new-test-item 1) slist :key #'cadr) \rightarrow ((TEST-ITEM 1)) (adjoin '(new-test-item 1) slist) \rightarrow ((NEW-TEST-ITEM 1) (TEST-ITEM 1))
```

Exceptional Situations:

Should be prepared to signal an error of type type-error if list is not a proper list.

See Also:

pushnew, Section 3.6 (Traversal Rules and Side Effects)

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Notes:

```
The :test-not parameter is deprecated.
```

pushnew

Syntax:

 $\begin{array}{l} \mathbf{pushnew} \ \textit{item place \&key key test test-not} \\ \rightarrow \textit{new-place-value} \end{array}$

Arguments and Values:

item—an object.

place—a place, the value of which is a proper list.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

new-place-value—a list (the new value of place).

Description:

pushnew tests whether *item* is the same as any existing element of the *list* stored in *place*. If *item* is not, it is prepended to the *list*, and the new *list* is stored in *place*.

pushnew returns the new list that is stored in place.

Whether or not *item* is already a member of the *list* that is in *place* is determined by comparisons using :test or :test-not. The first argument to the :test or :test-not function is *item*; the second argument is an element of the *list* in *place* as returned by the :key function (if supplied).

If :key is supplied, it is used to extract the part to be tested from both *item* and the *list* element, as for adjoin.

The argument to the :key function is an element of the *list* stored in *place*. The :key function typically returns part part of the element of the *list*. If :key is not supplied or **nil**, the *list* element is used.

For information about the *evaluation* of *subforms* of *place*, see Section 5.1.1.1 (Evaluation of Subforms to Places).

It is *implementation-dependent* whether or not **pushnew** actually executes the storing form for its *place* in the situation where the *item* is already a member of the *list* held by *place*.

Examples:

```
 \begin{array}{l} (\text{setq x '(a (b c) d)}) \to (\text{A (B C) D}) \\ (\text{pushnew 5 (cadr x)}) \to (\text{5 B C}) \\ \text{x} \to (\text{A (5 B C) D}) \\ (\text{pushnew 'b (cadr x)}) \to (\text{5 B C}) \\ \text{x} \to (\text{A (5 B C) D}) \\ (\text{setq 1st '((1) (1 2) (1 2 3))}) \to ((1) (1 2) (1 2 3)) \\ (\text{pushnew '(2) 1st}) \to ((2) (1) (1 2) (1 2 3)) \\ (\text{pushnew '(1) 1st}) \to ((1) (2) (1) (1 2) (1 2 3)) \\ (\text{pushnew '(1) 1st :test 'equal}) \to ((1) (2) (1) (1 2) (1 2 3)) \\ (\text{pushnew '(1) 1st :key \#'car)} \to ((1) (2) (1) (1 2) (1 2 3)) \\ \end{array}
```

Side Effects:

The contents of *place* may be modified.

See Also:

```
push, adjoin, Section 5.1 (Generalized Reference)
```

Notes:

```
The effect of (pushnew item place :test p)
is roughly equivalent to (setf place (adjoin item place :test p))
except that the subforms of place are evaluated only once, and item is evaluated before place.
```

set-difference, nset-difference

Function

Syntax:

```
set-difference list-1 list-2 & key key test test-not \rightarrow result-list
nset-difference list-1 list-2 & key key test test-not \rightarrow result-list
```

Arguments and Values:

```
list-1—a proper list.

list-2—a proper list.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.
```

set-difference, nset-difference

key—a designator for a function of one argument, or nil. result-list—a list.

Description:

set-difference returns a list of elements of list-1 that do not appear in list-2.

nset-difference is the destructive version of set-difference. It may destroy list-1.

For all possible ordered pairs consisting of one element from *list-1* and one element from *list-2*, the :test or :test-not function is used to determine whether they satisfy the test. The first argument to the :test or :test-not function is the part of an element of *list-1* that is returned by the :key function (if supplied); the second argument is the part of an element of *list-2* that is returned by the :key function (if supplied).

If :key is supplied, its argument is a *list-1* or *list-2* element. The :key function typically returns part of the supplied element. If :key is not supplied, the *list-1* or *list-2* element is used.

An element of *list-1* appears in the result if and only if it does not match any element of *list-2*.

There is no guarantee that the order of elements in the result will reflect the ordering of the arguments in any particular way. The result *list* may share cells with, or be eq to, either of *list-1* or *list-2*, if appropriate.

Examples:

```
(setq lst1 (list "A" "b" "C" "d")
       lst2 (list "a" "B" "C" "d")) 
ightarrow ("a" "B" "C" "d")
 (set-difference lst1 lst2) 
ightarrow ("d" "C" "b" "A")
 (set-difference lst1 lst2 :test 'equal) 
ightarrow ("b" "A")
 (set-difference lst1 lst2 :test #'equalp) 
ightarrow NIL
 (nset-difference lst1 lst2 :test #'string=) \rightarrow ("A" "b")
 (setq lst1 '(("a" . "b") ("c" . "d") ("e" . "f")))
\rightarrow (("a" . "b") ("c" . "d") ("e" . "f"))
(setq lst2 '(("c" . "a") ("e" . "b") ("d" . "a")))
\rightarrow (("c" . "a") ("e" . "b") ("d" . "a"))
(nset-difference lst1 lst2 :test #'string= :key #'cdr)

ightarrow (("c" . "d") ("e" . "f"))
lst1 \rightarrow (("a" . "b") ("c" . "d") ("e" . "f"))
lst2 \rightarrow (("c" . "a") ("e" . "b") ("d" . "a"))
;; Remove all flavor names that contain "c" or "w".
 (set-difference '("strawberry" "chocolate" "banana"
                   "lemon" "pistachio" "rhubarb")
           '(#\c #\w)
           :test #'(lambda (s c) (find c s)))

ightarrow ("banana" "rhubarb" "lemon") ;One possible ordering.
```

Side Effects:

nset-difference may destroy list-1.

Exceptional Situations:

Should be prepared to signal an error of type type-error if list-1 and list-2 are not proper lists.

See Also:

Section 3.2.1 (Compiler Terminology), Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

set-exclusive-or, nset-exclusive-or

Function

Syntax:

set-exclusive-or list-1 list-2 &key key test test-not \rightarrow result-list nset-exclusive-or list-1 list-2 &key key test test-not \rightarrow result-list

Arguments and Values:

list-1—a proper list.

list-2—a proper list.

 $test-a\ designator$ for a function of two arguments that returns a $generalized\ boolean$.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

result-list—a list.

Description:

set-exclusive-or returns a list of elements that appear in exactly one of list-1 and list-2.

nset-exclusive-or is the destructive version of set-exclusive-or.

For all possible ordered pairs consisting of one element from *list-1* and one element from *list-2*, the :test or :test-not function is used to determine whether they satisfy the test.

If :key is supplied, it is used to extract the part to be tested from the list-1 or list-2 element. The first argument to the :test or :test-not function is the part of an element of list-1 extracted by the :key function (if supplied); the second argument is the part of an element of list-2 extracted by the :key function (if supplied). If :key is not supplied or nil, the list-1 or list-2 element is used.

The result contains precisely those elements of *list-1* and *list-2* that appear in no matching pair.

The result list of set-exclusive-or might share storage with one of list-1 or list-2.

Examples:

```
 (\text{setq lst1 (list 1 "a" "b")} \\ \text{lst2 (list 1 "A" "b"))} \to (\text{1 "A" "b")} \\ (\text{set-exclusive-or lst1 lst2}) \to (\text{"b" "A" "b" "a")} \\ (\text{set-exclusive-or lst1 lst2 :test $\#'\text{equal}$)} \to (\text{"A" "a")} \\ (\text{set-exclusive-or lst1 lst2 :test $\#'\text{equal}$)} \to \text{NIL} \\ (\text{nset-exclusive-or lst1 lst2}) \to (\text{"a" "b" "A" "b")} \\ (\text{setq lst1 (list (("a" . "b") ("c" . "d") ("e" . "f")))} \\ \to (("a" . "b") ("c" . "d") ("e" . "b") ("d" . "a")))) \\ \to (("c" . "a") ("e" . "b") ("d" . "a")) \\ (\text{nset-exclusive-or lst1 lst2 :test $\#'\text{string= :key $\#'\text{cdr}$)}$} \\ \to (("c" . "d") ("e" . "f") ("c" . "a") ("d" . "a")) \\ (\text{lst1} \to (("a" . "b") ("c" . "d") ("e" . "f")) \\ (\text{lst2} \to (("c" . "a") ("d" . "a")) \\ (\text{lst2} \to (("c" . "a") ("d" . "a")))
```

Side Effects:

nset-exclusive-or is permitted to modify any part, car or cdr, of the list structure of list-1 or list-2.

Exceptional Situations:

Should be prepared to signal an error of type type-error if list-1 and list-2 are not proper lists.

See Also:

Section 3.2.1 (Compiler Terminology), Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

Since the **nset-exclusive-or** side effect is not required, it should not be used in for-effect-only positions in portable code.

subsetp

Syntax:

subsetp list-1 list-2 &key key test test-not ightarrow generalized-boolean

Arguments and Values:

```
list-1—a proper list.
```

list-2—a proper list.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

generalized-boolean—a generalized boolean.

Description:

subsetp returns true if every element of list-1 matches some element of list-2, and false otherwise.

Whether a list element is the same as another list element is determined by the functions specified by the keyword arguments. The first argument to the :test or :test-not function is typically part of an element of *list-1* extracted by the :key function; the second argument is typically part of an element of *list-2* extracted by the :key function.

The argument to the :key function is an element of either *list-1* or *list-2*; the return value is part of the element of the supplied list element. If :key is not supplied or nil, the *list-1* or *list-2* element itself is supplied to the :test or :test-not function.

Examples:

```
(setq cosmos '(1 "a" (1 2))) \to (1 "a" (1 2)) (subsetp '(1) cosmos) \to true (subsetp '((1 2)) cosmos) \to false (subsetp '((1 2)) cosmos :test 'equal) \to true (subsetp '(1 "A") cosmos :test #'equalp) \to true (subsetp '(1) (2)) '((1) (2))) \to false (subsetp '((1) (2)) '((1) (2)) :key #'car) \to true
```

Exceptional Situations:

Should be prepared to signal an error of type type-error if list-1 and list-2 are not proper lists.

See Also:

Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

union, nunion

union, nunion

Function

Syntax:

union list-1 list-2 &key key test test-not \rightarrow result-list nunion list-1 list-2 &key key test test-not \rightarrow result-list

Arguments and Values:

list-1—a proper list.

list-2—a proper list.

test—a designator for a function of two arguments that returns a generalized boolean.

test-not—a designator for a function of two arguments that returns a generalized boolean.

key—a designator for a function of one argument, or nil.

result-list—a list.

Description:

union and nunion return a list that contains every element that occurs in either list-1 or list-2.

For all possible ordered pairs consisting of one element from *list-1* and one element from *list-2*, :test or :test-not is used to determine whether they satisfy the test. The first argument to the :test or :test-not function is the part of the element of *list-1* extracted by the :key function (if supplied); the second argument is the part of the element of *list-2* extracted by the :key function (if supplied).

The argument to the :key function is an element of *list-1* or *list-2*; the return value is part of the supplied element. If :key is not supplied or nil, the element of *list-1* or *list-2* itself is supplied to the :test or :test-not function.

For every matching pair, one of the two elements of the pair will be in the result. Any element from either *list-1* or *list-2* that matches no element of the other will appear in the result.

If there is a duplication between *list-1* and *list-2*, only one of the duplicate instances will be in the result. If either *list-1* or *list-2* has duplicate entries within it, the redundant entries might or might not appear in the result.

The order of elements in the result do not have to reflect the ordering of *list-1* or *list-2* in any way. The result *list* may be eq to either *list-1* or *list-2* if appropriate.

union, nunion

Examples:

Side Effects:

nunion is permitted to modify any part, car or cdr, of the list structure of list-1 or list-2.

Exceptional Situations:

Should be prepared to signal an error of type type-error if list-1 and list-2 are not proper lists.

See Also:

intersection, Section 3.2.1 (Compiler Terminology), Section 3.6 (Traversal Rules and Side Effects)

Notes:

The :test-not parameter is deprecated.

Since the **nunion** side effect is not required, it should not be used in for-effect-only positions in portable code.