- 1. Use the Load Effective Address instruction to implement the following arithmetic operations:
 - a. Add 20 to a variable

LEA eax,[eax+20]

b. Multiply a variable by 52

LEA eax,[ebx+ebx*8]//9

LEA eax,[eax+eax*4]//9+36=45

LEA eax,[eax+ebx*4]//45+4=49

LEA eax,[eax+ebx*2]//49+2=51

LEA eax,[eax+ebx]//51+1

c. Multiply a variable by 9

LEA eax,[eax+eax*8]

- 2. Use SHL/SAL, LEA and ADD to implement the following arithmetic operations:
 - a. Multiply a variable by 24

MOV eax,[variable]

MOV ebx,[variable]

//3x ebx

ADD ebx,ebx

ADD ebx,eax

//mul by 3

LEA eax,[ebx]

// mul by 8

SHL eax,3

b. Multiply a variable by 1000

MOV eax,[variable]

MOV ebx,[variable]

MOV edx,[variable]

//shift by 10==1024

SHL eax,10

// 24 times

ADD ebx,ebx

ADD ebx,edx

LEA ebx, [ebx*8]

ADD eax ,-ebx

3. Implement the following while loop in IA-32 assembly language using a post-tested loop.

Why is a post-tested loop implementation possible in this situation?

```
int x = 1;
while (x <= 10)
{
    if (x != 5)
    {
        /* CodeBlock */
        x = x * 2;
    }
    x = x + 1;
}

//int x = 1
    MOV eax, 1

LoopStart:
    // if (x != 5)
    Cmp eax,5</pre>
```

```
JNE IFCondition
//x=x+1
INC eax
//While condition

Cmp eax,10
Jle LoopStart

IFConsition:
//x=x*2
LEA eax, DWORD PTR [eax+eax]
```

Why is a post-tested loop implementation possible in this situation?

Because the condition for the while loop is satisfied, so the post-tested loop is possible.

4. Implement the following C code segment in IA-32 assembly language as efficiently as possible.

```
int j = 5;

for (int i = 0; i < 25; i += 2)

{

        if (i < 3 && j > 23)

        {

            if (i <= 30 && j <= 35 && i%2==0)

            {

                break;

            }

            i = i + 1;

        }

        j += i * 2;
```

```
}
// int j=5
Mov eax,5
ForLoop:
       //int i=0
       Mov ebx,0
       // j<=25
       Cmp eax, 25
       JG AfterLoopCondition
       //if (i < 3 && j > 23)
       Cmp ebx,3
       Jge AfterIfCondition
       Cmp eax,23
       Jle AfterIfCondition
       //i = i + 1;
       INC ebx
      //j += i * 2;
       LEA eax, [eax+ebx*2]
       //i+=2
       LEA ebx,[exa+2]
AfterIfCondition:
       //j += i * 2;
       LEA eax,[eax+ebx*2]
       //i += 2
       LEA ebx,[ebx+2]
       JMP ForLoop
AfterLoopCondition:
```

5. Implement the following C code segment in IA-32 assembly language. Use the cdecl calling convention to implement (Note: treat func1 as both a caller and a callee) int func1(int y)

```
{
int a = 5;
int b = functionOne(a);
int c = functionTwo(a);
int d = b + c + y;
return d;
}
Func1:
    push ebp
    mov ebp, esp

MOV [EBP+8],[y]
```

```
//local variant a=5
sub ebp, 0x4
Mov [EBP-4],5
//int b = functionOne(a)
sub EBP, 0x4
Mov [ESP-8],[b]
//push a
Push [ESP-4]
Call functionOne
//int c = functionTwo(a);
sub EBP, 0x4
Mov [EBP-12],[c]
//push a
Push [ESP-4]
Call functionTwo
//int d = b + c + y;
sub EBP, 0x4
Mov [ESP-16],[d]
//d=b+c+y
LEA [ESP-16] ,[
                   [ESP-8]+[ESP-12] ]
LEA [ESP-16] ,[
                  [ESP-16]+[ESP+8]]
MOV [EBP+4],[ESP-16]
MOV esp, ebp
POP ebp
RET
```

6. Generate the equivalent C code for the following IA-32 assembly code.

Label1:

cmp [var1], 0x12 jle Label2 jmp Label6

Label2:

cmp [var2], 0x27 jg Label3 cmp [var3], 0x19 jg Label4

Label3:

cmp [var2], 0x10 jle Label5 cmp [var3], 0x16 jg Label4 jmp Label5

Label4:

```
; CodeBlock
Label5:
mul [var3], 0x06
add [var2], 0x02
inc [var1]
jmp Label1
Label6:
Void label1(){
       If( var1<= 18) label2();
       else label6();
}
Void label2(){
       If (var2 >39) label3();
       elself (var3 >25) Label4
}
Void label3(){
       If (var2<= 16) label5();
       elseif(var3 >22) label4();
       Else: label5;
Void label4(){
       codeblock;
}
Void Label5(){
       var3= var3*6;
       var2+=2;
       var1+=1;
       label1();
}
```

Void Label6(){

7. Generate the assembly and also evaluate the output of the following program.

NOTE: Please make sure that you don't have leading whitespaces or comma(",") as a seperator for the output. Write down the exact output as your answer

Assume CDECL calling convention and the order of evaluation of arguments from right to left.

```
int main()
int c = 0;
printf("%d\n",c);
C++;
printf("%d\n",c);
++c;
printf("%d\n",c);
printf("%d %d %d %d\n", ++c,++c,c++,c++);
printf("%d ", c++);
return 0;
}
Push ebp
mov ebp, esp
//int c = 0;
Sub esp,0x4
MOV eax,0
//printf("%d\n",c);
Push eax
Call printf
```

```
//c++;
Inc eax
//printf("%d\n",c);
Push eax
Call printf
//++c;
Inc eax
Push eax
//printf("%d\n",c);
Call printf
Push eax
//printf("%d %d %d %d\n", ++c,++c,c++,c++);
Push eax
INC eax
Push eax
Inc eax
Inc eax
Inc eax
Push eax
Push eax
Call prinft
//printf("%d ", c++);
Push eax
Inc eax
Call prinft
//set the return value to the return address
Mov [ebp+4],0
mov esp, ebp
pop ebp
Ret
```

Output:

1