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Experiment 8

Aim: To implement the Matrix Chain Multiplication (MCM) algorithm using Dynamic Programming, print the DP cost table, and calculate the minimum number of scalar multiplications required to multiply a sequence of matrices.

Problem Statement

Given a chain of matrices $A_1, A_2, A_3, \dots, A_n$ where matrix A_i has dimensions $p[i-1] \times p[i]$, determine the most efficient order to multiply the matrices so that the total number of scalar multiplications is minimized.

Theory

Matrix multiplication is associative; however, different parenthesizations yield different multiplication costs. Dynamic Programming is used because the problem contains optimal substructure and overlapping subproblems.

Algorithm

1. Read number of matrices.
2. Store matrix dimensions in $p[]$.
3. Create DP table $dp[n+1][n+1]$.
4. Initialize $dp[i][i] = 0$.
5. For chain length $L = 2$ to n :

For each valid i :

Compute $j = i + L - 1$

For each k from i to $j-1$:

Compute cost = $dp[i][k] + dp[k+1][j] + p[i-1] \times p[k] \times p[j]$

Choose minimum.

6. Print DP table.

7. Final result = $dp[1][n]$.

Source Code

```
#include <iostream>
#include <vector>
#include <iomanip>
#include <climits>
using namespace std;

void matrixChainOrder(const vector<int>& p) {
    int n = p.size() - 1;
    vector<vector<int>> dp(n + 1, vector<int>(n + 1, 0));

    for (int L = 2; L <= n; L++) {
        for (int i = 1; i <= n - L + 1; i++) {
            int j = i + L - 1;
            dp[i][j] = INT_MAX;
            for (int k = i; k < j; k++) {
                int cost = dp[i][k] + dp[k + 1][j] + p[i - 1] * p[k] * p[j];
                dp[i][j] = min(dp[i][j], cost);
            }
        }
    }

    // Print DP Table in matrix-style format
    cout << "\nMatrix Chain Multiplication DP Table:\n";
    cout << "-----\n\n";

    // Print column headers
    cout << setw(6) << " ";
    for (int j = 1; j <= n; j++)
        cout << setw(6) << j;
    cout << "\n";

    // Print rows
    for (int i = 1; i <= n; i++) {
        cout << setw(3) << i << " "; // row index
        for (int j = 1; j <= n; j++) {
            if (i > j)
                cout << setw(6) << "-";
            else
                cout << setw(6) << dp[i][j];
        }
        cout << "\n";
    }
}
```

```
    cout << "\nMinimum number of multiplications is: " << dp[1][n] << endl;
}

int main() {
    // Example: 6 matrices
    // Dimensions: A1(30x35), A2(35x15), A3(15x5), A4(5x10), A5(10x20), A6(20x25)
    vector<int> p = {30, 35, 15, 5, 10, 20, 25};

    matrixChainOrder(p);

    return 0;
}
```

Output

PS E:\DSA\Experiments\EXP_8_MCM\output> & .\MCM.exe'

Matrix Chain Multiplication DP Table:

	1	2	3	4	5	6
1	0	15750	7875	9375	11875	15125
2	-	0	2625	4375	7125	10500
3	-	-	0	750	2500	5375
4	-	-	-	0	1000	3500
5	-	-	-	-	0	5000
6	-	-	-	-	-	0

Minimum number of multiplications is: 15125

PS E:\DSA\Experiments\EXP_8_MCM\output>

The minimum number of multiplications for the example input is 15125.

Time & Space Complexity

Time Complexity: $O(n^3)$

Space Complexity: $O(n^2)$

Conclusion

Matrix Chain Multiplication demonstrates how Dynamic Programming reduces computation by dividing the problem into smaller overlapping subproblems and combining their solutions efficiently.