

南 开 大 学

计算机学院

编译原理实验报告

预备工作 2: 定义你的编译器、汇编编程 & 熟悉辅助 工具(小组合作)

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摘要

首先,对 SysY 语言进行了上下文无关文法的描述。然后,针对其语言特性,对三个有代表性的程序进行 ARM 汇编语言的翻译。最后,对工作进行总结。

关键字: ARM 汇编; 上下文无关文法

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一、 支持的 SysY 语言特性描述

(一) SysY 语言特性

我们描述了暂时想到的全部 SysY 语言特性,包括 int/float 类型、函数、变量/常量声明、语句、表达式等,参考了实验指导书《SysY2022 语言定义》。

(二) 上下文无关文法描述

1. 终结符集合

 $V_T = \{ \text{ Ident, IntConst, floatConst, const, int, float, void, if, else, while, break, continue, return, (,), ;, [,], +, -, !, *, /, %, <, >, <=, >=, ==, !=, ||, && }$

2. 非终结符集合

 $V_N = \{ \text{BasicUnit, CompUnit, Decl, ConstDel, ConstDefList, BType, ConstDef, ConstL-Val, ConstInitVal, ConstInitValInit, VarDecl, VarDefLsit, VarDef, InitVal, InitValList, FuncDef, FuncType, FuncFParams, FuncFParam, ExpList, Block, BlockItemList, BlockItem, Stmt, Exp, Cond, LVal, PrimaryExp, Number, UnaryExp, FunvRParams, MulExp, AddExp, RelExp, Eq-Exp, LAndExp, LOrExp, ConstExp, UnaryOp, MulOp, AddOp, RelOp, EqOp \}$

3. 开始符号

S = BasicUnit

4. 产生式

```
P = {
BasicUnit \rightarrow Decl \mid FuncDef
CompUnit → BasicUnit | CompUnit BasicUnit
\mathrm{Decl} \to \mathrm{ConstDel} \mid \mathrm{VarDecl}
ConstDel \rightarrow const BType ConstDefList;
ConstDefList \rightarrow ConstDef \mid ConstDef, ConstDefList
BType \rightarrow int \mid float
ConstDef \rightarrow ConstLVal = ConstInitVal
ConstLVal \rightarrow Ident \mid ConstLVal[ConstExp]
ConstInitVal \rightarrow ConstExp \mid |ConstInitValList|
ConstInitValList \rightarrow ConstInitVal \mid ConstInitValList, ConstInitVal
VarDecl \rightarrow BType VarDefList;
VarDefList \rightarrow VarDef \mid VarDefList, VarDef
VarDef \rightarrow ConstLVal \mid ConstLVal = InitVal
InitVal \rightarrow Exp \mid InitValList
\label{eq:initValList} \operatorname{InitValList}, \operatorname{InitVal} \mid \varepsilon
FuncDef \rightarrow FunType Ident(FunctionFParams) Block
FuncType \rightarrow void | int | float
FuncFParams \rightarrow FuncFParam | FuncFParams, FuncFParam | \varepsilon
FuncFParam \rightarrow BType Ident | BType Ident | ExpList
```

```
\text{ExpList} \rightarrow [\text{Exp}] \mid \text{ExpList} [\text{Exp}] \mid \varepsilon
     Block \rightarrow BlockItemList
     BlockItemList \rightarrow BlockItem | BlockItemList BlockItem | \varepsilon
     BlockItem \rightarrow Decl \mid Stmt
     Stmt \rightarrow LVal = Exp; \mid Exp; \mid ; \mid Block \mid if(Cond) Stmt \mid if(Cond) Stmt else Stmt \mid
while (Cond ) Stmt | break; | continue; | return; | return Exp;
     \text{Exp} \to \text{AddExp}
     Cond \rightarrow LOrExp
     LVal \rightarrow Ident \mid LVal[Exp]
     PrimaryExp \rightarrow (Exp) | LVal | Number
     \mathrm{Number} \to \mathbf{IntConst} \mid \mathbf{floatConst}
     UnaryExp → PrimaryExp | Ident (FuncRParams) | Ident () | UnaryOp UnaryExp
     FuncRParams \rightarrow Exp \mid FuncRParams, Exp
     MulExp \rightarrow UnaryExp \mid MulExp MulOp UnaryExp
     AddExp \rightarrow MulExp \mid AddExp AddOp MulExp
     RelExp \rightarrow AddExp \mid RelExp RelOp AddExp
     EqExp \rightarrow RelExp \mid EqExp EqOp RelExp
     LAndExp \rightarrow EqExp \mid LAndExp \&\& EqExp
     LOrExp \rightarrow LAndExp \mid LOrExp \mid LAndExp
     ConstExp \rightarrow AddExp
     UnaryOp \rightarrow + | - | !
     MulOp \rightarrow * | / | \%
     AddOp \rightarrow + | -
     RelOp \rightarrow < | > | <= | >=
     EqOp \rightarrow == |!=
     }
```

5. 对终结符的进一步定义

```
Ident 标识符 Ident \to id-nondigit | Ident id-digit | Ident id-nondigit | id-nondigit = {a, b, c, d, e, f, g, h, i, j, k, l, m, n, o, p, q, r, s, t, u, v, w, x, y, z, A, B, C, D, E, F, G, H, I, J, K, L, M, N, O, P, Q, R, S, T, U, V, W, X, Y, Z, __} id-digit = {0, 1, 2, 3, 4, 5, 6, 7, 8, 9} Ident \to identifier-nondigit | Ident identifier-nondigit | Ident identifier-digit identifier-nondigit \to a | b | c | d | e | f | g | h | i | j | k | l | m | n | o | p | q | r | s | t | u | v | w | x | y | z | A | B | C | D | E | F | G | H | I | J | K | L | M | N | O | P | Q | R | S | T | U | V | W | X | Y | Z identifier-digit \to 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
```

```
IntConst 整形常量 IntConst \rightarrow decimal-const | octal-const | hexadecimal-const decimal-const \rightarrow nonzero-digit | decimal-const digit octal-const \rightarrow 0 | octal-const octal-digit hexadecimal-const \rightarrow hexadecimal-prefix hexadecimal-digit | hexadecimal-const hexadecimal-digit
```

```
\begin{split} & \text{hexadecimal-prefix} \to 0 \le \mid 0 X \\ & \text{nonzero-digit} \to 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \\ & \text{octal-digit} \to 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \\ & \text{hexadecimal-digit} \to 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \mid a \mid b \mid c \mid d \mid e \mid f \mid A \mid B \mid C \mid D \mid E \mid F \end{split}
```

floatConst 浮点数常量 floatConst \rightarrow decimal-floating-constant | hexadecimal-floating-constant decimal-floating-constant \rightarrow fractional-constant exponent-partopt floating-suffixopt | digit-sequence exponent-part floating-suffixopt

$$\label{lem:hexadecimal-floating-constant} \begin{split} & + \text{hexadecimal-prefix hexadecimal-fractional-constant} \mid \text{binary-exponent-part floating-suffixopt} \mid \text{hexadecimal-prefix hexadecimal-digit-sequence} \mid \text{binary-exponent-part floating-suffixopt} \end{split}$$

```
fractional-constant \to digit-sequence
opt . digit-sequence | digit-sequence . exponent-part \to e sign
opt digit-sequence | E signopt digit-sequence sign
\to + | -
```

 ${\it digit}\text{-}{\it sequence}{\rightarrow}{\it digit}|{\it digit}\text{-}{\it sequence}|{\it digit}|$

 $\label{lem:hexadecimal-digit-sequence} hexadecimal-fractional-constant \rightarrow hexadecimal-digit-sequence opt. \ | hexadecimal-digit-sequence | hexadecimal-digit-seq$

 $binary-exponent-part \rightarrow p signopt digit-sequence | P signopt digit-sequence$

 $\mbox{hexadecimal-digit-sequence} \rightarrow \mbox{hexadecimal-digit} \ | \mbox{hexadecimal-digit-sequence} \ \mbox{hexadecimal-digit} | \mbox{hexadecimal-digit-sequence} \ \mbox{hexadecimal-digit} | \mbox{hexadecimal-digit-sequence} \ \mbox{hexadecimal-digit} | \mbox{hexadecimal-digit-sequence} \ \mbox{hexadecimal-digit-sequence} | \mbox{hexadecimal-digit-sequence} \ \mbox{hexadecimal-digit-sequence} | \mbox{hexadecimal-digit-se$

floating-suffix \rightarrow f | l | F | L

二、汇编程序的编写和验证

(一) SysY 程序样例

1. 程序 1

如下所示

p1.c

```
t = b;

b = a + b;

printf("%d", b);

a = t;

i = i + 1;

}
```

p1.c

```
; Function Attrs: noinline nounwind optnone uwtable
        define dso_local i32 @main() #0 {
             \%1 = alloca i32, align 4
             \%2 = alloca i32, align 4
             \%3 = alloca i32, align 4
             \%4 = alloca i32, align 4
             \%5 = alloca i32, align 4
             \%6 = alloca i32, align 4
             store i32 0, i32* %1, align 4
             store i32 0, i32* \%2, align 4
             store i32 1, i32* %3, align 4
             store i32 1, i32* %4, align 4
             \%7 = call \ i32 \ (i8*, \ldots) \ @\_isoc99 \ scanf(i8* getelementptr inbounds ([3 x]))
                        i8], [3 \times i8] * @.str, i64 \ 0, i64 \ 0), i32 * \%6)
             \%8 = load i32, i32* \%2, align 4
             \%9 = \text{call} \ \text{i32} \ (\text{i8*, ...}) \ @printf(\text{i8* getelementptr inbounds} \ ([3 \times \text{i8}], [3 \times \text{i8}]) 
                          i8] * @. str, i64 0, i64 0), i32 %8)
             \%10 = load i32, i32* \%3, align 4
             \%11 = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3]) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3]) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3]) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3]) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3])) = call i32 (i8*, ...) @printf(i8* getelementpt
                       x i8] * @. str, i64 0, i64 0), i32 %10)
             br label %12
18
19
                                                                                                                                             ; preds = \%16, \%0
         12:
20
             \%13 = load i32, i32* \%4, align 4
             \%14 = load i32, i32* \%6, align 4
             %15 = icmp slt i32 %13, %14
             br i1 %15, label %16, label %26
         16:
                                                                                                                                             ; preds = \%12
             \%17 = load i32, i32* \%3, align 4
             store i32 %17, i32* %5, align 4
             \%18 = load i32, i32* \%2, align 4
             \%19 = load i32, i32* \%3, align 4
             \%20 = add \text{ nsw i } 32 \%18, \%19
             store i32 %20, i32* %3, align 4
             \%21 = load i32, i32* \%3, align 4
             \%22 = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([3 x i8], [3
                       x i8]* @.str, i64 0, i64 0), i32 %21)
             \%23 = load i32, i32* \%5, align 4
```

2. 程序 2

如下所示

```
p2.c
```

```
包含的SysY语言特性:
      float 数据类型
      变量声明, 赋值语句
     for循环
      if语句
                                   逻辑运算(逻辑与或等)、关系运算(不
      算术运算(加减乘除、按位与或等)
         等、等于、大于、小于等)
      数组和指针
  #include <stdio.h>
  #include <stdlib.h>
  int main() {
         //数组 指针
         float a [10];
         float* p = (float*)malloc(10 * sizeof(float));
         for (int i = 0; i < 10; i++) {
16
                a[i] = i;
                p[i] = 10-i;
                if (a[i] = p[i]) {
19
                      a[i] = -1;
                      p[i] = -1;
                }
         }
```

p2.ll

```
; 首部参考其自动生成的.11文件
; ModuleID = 'try.c'
source_filename = "try.c"
target datalayout = "e-m:e-p270:32:32-p271:32:32-p272:64:64-i64:64-f80:128-n8:16:32:64-S128"
```

```
target triple = "x86_64-pc-linux-gnu"
   ; Function Attrs: noinline nounwind optnone uwtable
   define dso_local i32 @main() #0 {
     ;数组声明 16字节对齐
     ; float a[10];
    \%1 = alloca [10 x float], align 16
     ; 指针声明 8字节对齐
     ; float * p = (float *) malloc(10 * sizeof(float));
    \%2 = alloca float*, align 8
     ; 调用全局函数@malloc动态分配空间
    \%3 = call noalias i8* @malloc(i64 40) #2
19
     ; The 'bitcast' instruction converts value to type ty2 without changing
        any bits.
    \%4 = bitcast i8* \%3 to float*
     store float * %4, float ** %2, align 8
     ; for 循环中的 int i=0; 4字节对齐
    \%5 = alloca i32, align 4
25
     ; i初始化为0
     store i32 0, i32* %5, align 4
    \%6 = alloca i32, align 4
     br label %7
30
                                                     ; preds = \%42, \%0
    \%8 = load i32, i32* \%5, align 4
     ; 利用slt 判断循环条件i<10是否符合
    \%9 = icmp slt i32 \%8, 10
     ;符合则跳转到label10 否则跳转至label45结束循环
    br i1 %9, label %10, label %45
   10:
                                                     ; preds = \%7
38
        (循环条件符合) 开始循环体部分
    \%11 = load i32, i32* \%5, align 4
40
     ; sitofp 将有符号整型i转换为浮点数
41
    \%12 = sitofp i32 \%11 to float
42
    \%13 = load i32, i32* \%5, align 4
    \%14 = \text{sext i} 32 \%13 \text{ to i} 64
    \%15 = getelementptr inbounds [10 x float], [10 x float]* \%1, i64 0, i64 \%14
     ; a[i]=i
     store float %12, float * %15, align 4
    \%16 = load i32, i32* \%5, align 4
    \%17 = \text{sub nsw i} 32 \ 10, \%16
     \%18 = sitofp i32 \%17 to float
    \%19 = load float*, float** \%2, align 8
```

```
\%20 = load i32, i32* \%5, align 4
     \%21 = \text{sext i} 32 \%20 \text{ to i} 64
     \%22 = getelementptr inbounds float, float* <math>\%19, i64 \%21
54
     store float %18, float* %22, align 4
     \%23 = load i32, i32* \%5, align 4
     \%24 = \text{sext i} 32 \%23 \text{ to i} 64
     %25 = getelementptr inbounds [10 x float], [10 x float] * %1, i64 0, i64 %24
      ; load a[i]到%26
61
     \%26 = load float, float* \%25, align 4
62
     \%27 = load float*, float** \%2, align 8
     \%28 = load i32, i32* \%5, align 4
     \%29 = \text{sext i} 32 \%28 \text{ to i} 64
     %30 = getelementptr inbounds float, float* %27, i64 %29
     ; load p[i]到%26
     \%31 = load float, float* \%30, align 4
     ; if条件判断 a[i]==p[i]
     \%32 = \text{fcmp oeq float } \%26, \%31
     ; 条件成立, 跳转至label%33 否则跳转至label%41
     br i1 %32, label %33, label %41
   33:
                                                            ; preds = \%10
      ; a[i]==p[i]条件成立
     \%34 = load i32, i32* \%5, align 4
     %35 = \text{sext i} 32 \%34 \text{ to i} 64
     %36 = getelementptr inbounds [10 x float], [10 x float] * %1, i64 0, i64 %35
     ; a[i]=-1
     store float -1.0, float* \%36, align 4
     \%37 = load float*, float** \%2, align 8
     \%38 = load i32, i32* \%5, align 4
     %39 = \text{sext i} 32 \%38 \text{ to i} 64
     \%40 = getelementptr inbounds float, float* <math>\%37, i64 \%39
      ; p[i]=-1
     store float -1.0, float * \%40, align 4
     br label %41
                                                            ; preds = \%33, \%10
   41:
89
      ; 分支条件不符合
     br label %42
91
                                                            ; preds = \%41
   42:
93
      ; 取 i
     \%43 = load i32, i32* \%5, align 4
95
      ; i自增 i++
     \%44 = add \text{ nsw i } 32 \%43, 1
     store i32 %44, i32* %5, align 4
     br label %7
```

3. 程序 3

如下所示

```
p3.c
  sysY语言特性:
      int 数据类型
      变量声明, 赋值语句
      while循环
      算术运算(加减乘除、按位与或等)、逻辑运算(逻辑与或等)、关系运算(不
         等、等于、大于、小于等)
      函数
  */
  #include<stdio.h>
  int function(int a, int b){
      return a*b;
  int main()
13
14
         int i, n, f;
         scanf("%d",&n);
         i = 2;
         f = 1;
         while (i \le n)
19
         {
                f = function(f, i);
                i = i + 1;
         printf("%d\n",f);
      return 0;
```

p3.ll

```
; 首部参考其自动生成的.11文件
   ;function函数 输入为整数(i32)
   define dso_local i32 @function(i32 %0, i32 %1) #0 {
    \%3 = \text{mul} \text{ nsw } i32 \%0, \%1 ; \%3 = \%0 * \%1
     ret i32 %3 ; return %3
   }
   ;注意:一切变量名和lable必须依次顺序编号!
   ; main函数
   define dso_local i32 @main() #0 {
    %1 = alloca i32, align 4 ;声明变量i
    %2 = alloca i32, align 4 ;声明变量n
    %3 = alloca i32, align 4 ;声明变量f
    \%4 = call \ i32 \ (i8*, \ldots) \ @\_isoc99\_scanf(i8* getelementptr inbounds \ ([3 x]
        i8], [3 x i8]* @.str, i64 0, i64 0), i32* %2);调用函数scanf("%d",&n);
     store i32 2, i32* \%1, align 4; i = 2;
     store i32 1, i32* \%3, align 4; f =
     br label %5
18
19
   5:
                                                   ; while 循环退出条件
20
    \%6 = load i32, i32* \%1, align 4;
    \%7 = load i32, i32* \%2, align 4; \c \&
    \%8 = load i32, i32* \%3, align 4 
    %9 = icmp sle i32 %6, %7 ; %9为 i <= n 的比较结果
     br i1 %9, label %10, label %13;如果为%9真,那么执行 label%10,否则执行
        label%16
   10:
                                                    ; while 循环体
    %11 = call i32 @function(i32 %6, i32 %8);调用function函数
28
     store i32 %11, i32* %3, align 4;将函数的返回值存入f
    \%12 = add \text{ nsw } i32 \%6, 1
     store i32 %12, i32* %1, align 4; i = i + 1
    br label %5;跳转到退出条件判断
   13:
34
     \%14 = call i32 (i8*, ...) @printf(i8* getelementptr inbounds ([4 x i8], [4])
        x i8]* @.str.1, i64 0, i64 0), i32 %8); 输出f
     ret i32 0 ; main函数返回0
37
   :函数声明
39
   declare dso_local i32 @__isoc99_scanf(i8*, ...) #1
40
   declare dso_local i32 @printf(i8*, ...) #1
42
  ; 尾部参考了其自动生成的.11文件
```

5 . .

(二) 编写等价汇编程序

1. p1.S

如下所示

p1.ll

```
.global main
    main:
                            {r7}
                push
                \mathbf{sub}
                            \mathbf{sp}, \mathbf{sp}, #20
                            r7, sp
                mov
                movs
                            r3, #0
                            r3\;,\;\;[\;r7\;,\;\;\#4]
                \mathbf{str}
                            r3, #1
                movs
                \mathbf{str}
                            r3, [r7, #8]
                            r3\ ,\ \#1
                movs
                \mathbf{str}
                            r3\;,\;\; [\;r7\;,\;\; \#12]
     .while:
                l\,d\,r
                            r2, [r7, #12]
                l\,d\,r
                            r3, [r7, #16]
                            r3, r2
                cmp
                            . \verb|exit_while|
                _{\mathrm{blt}}
                l\,d\,r
                            r3, [r7, #8]
                            r3, [r7, #20]
                \mathbf{str}
                ldr
                            r2, [r7, #8]
                            r3, [r7, #4]
                ldr
                add
                            r3, r3, r2
21
                \mathbf{str}
                            r3, [r7, #8]
                            r3, [r7, #20]
                ldr
                \mathbf{str}
                            r3, [r7, #4]
                            r3, [r7, #12]
                ldr
                            r3\;,\;\;r3\;,\;\;\#1
                adds
                            r3, [r7, #12]
                \mathbf{str}
     .exit_while:
                movs
                            r3, #0
                mov
                            r0, r3
                            r7\;,\;\;r7\;,\;\;\#20
                adds
31
                adds
                           sp, sp, #20
```

编译验证正确。

2. p2.S

如下所示

p2.ll

```
.global main
   main:
             @ args = 0, pretend = 0, frame = 56
             @ frame_needed = 1, uses_anonymous_args = 0
                        {r7, lr}
              push
                       \mathbf{sp}, \mathbf{sp}, \#56
             sub
             add
                       r7, sp, #0
                       r2, .L7
              l\,d\,r
    .LPIC0:
                        r2, pc
             add
              ldr
                        r3, .L7+4
11
              ldr
                       r3, [r2, r3]
              ldr
                       r3, [r3]
              \mathbf{str}
                       r3, [r7, #52]
                        r3,\#0
             mov
                        r0, \#40
             movs
                        malloc (PLT)
              bl
                        r3, r0
             mov
              \mathbf{str}
                        r3, [r7, #8]
             movs
                        r3, #0
22
              \mathbf{str}
                        r3, [r7, #4]
              b
                        .for
25
    .else:
26
    .add i:
27
              ldr
                        r3, [r7, #4]
              adds
                        r3, r3, #1
29
                        r3, [r7, #4]
              \mathbf{str}
30
    .for:
31
              l\,d\,r
                       r3, [r7, #4]
                        r3, #10
             cmp
              bge
                        .exit\_for
              l\,d\,r
                        r3, [r7, #4]
                        s15, r3 @ int
              vmov
37
              vcvt.f32.s32
                               s15, s15
              ldr
                       r3, [r7, #4]
              lsls
                       r3, r3, #2
40
             add
                       r2, r7, #56
41
                       r3, r3, r2
              add
42
                       r3, r3, #44
              \operatorname{subs}
              vstr.32 s15, [r3]
              l\,d\,r
                       r3, [r7, #4]
                       r1, r3, #10
              rsb
                       r3, [r7, #4]
              l\,d\,r
```

```
lsls
                       r3, r3, #2
             ldr
                       r2, [r7, #8]
                       r3, r3, r2
             add
                       s15, r1 @ int
             vmov
             vcvt.f32.s32 \hspace{1.5cm} s15\,,\hspace{1.5cm} s15
             vstr.32 s15, [r3]
             ldr
                       r3, [r7, #4]
                       r3, r3, #2
             lsls
                      r2, r7, #56
             add
                      r3, r3, r2
             add
60
             subs
                      r3, r3, #44
61
             vldr.32 s14, [r3]
62
63
             ldr
                       r3, [r7, #4]
64
                       r3, r3, #2
             lsls
             ldr
                       r2, [r7, #8]
             add
                      r3, r3, r2
             vldr.32 s15, [r3]
             vcmp.f32
                                s14, s15
70
             vmrs
                       APSR_nzcv, FPSCR
71
             bne
                       .else
73
    . if:
74
                       r3, [r7, #4]
             ldr
75
                       r3, r3, #2
             lsls
                       r2\ ,\ \ r7\ ,\ \ \#56
             add
                       r3, r3, r2
             add
                       r3, r3, #44
             subs
                       r2\ ,\ \#0
             mov
                       r2, 49024
             movt
                       r2, [r3]
             \mathbf{str}
                                          @ float
                       r3, [r7, #4]
             l\,d\,r
                       r3, r3, #2
             lsls
85
             ldr
                       r2, [r7, #8]
             add
                       r3, r3, r2
                       r2, #0
             mov
                       r2, 49024
             movt
                                          @ float
             \mathbf{str}
                       r2, [r3]
90
             b .add_i
92
    .exit\_for:
                       r3, #0
             movs
             ldr
                       r1, .L7+8
```

```
.LPIC1:
              add
                        r1, pc
              ldr
                        ^{\rm r2}\;,\;\; .L7{+}4
99
                        r2, [r1, r2]
              ldr
              ldr
                        r1, [r2]
101
                        r2, [r7, #52]
              ldr
                        r1, r2, r1
              eors
103
                        r2, #0
              mov
104
                        .L6
              beq
105
              bl
                        __stack_chk_fail(PLT)
106
    .L6:
                        r0, r3
              mov
              adds
                        r7, r7, #56
              mov
                        sp, r7
                        \{r7, pc\}
              pop
    .L7:
                       _GLOBAL_OFFSET_TABLE_-(.LPIC0+4)
              .word
              .word
                        __stack_chk_guard(GOT)
114
                        _GLOBAL_OFFSET_TABLE_-(.LPIC1+4)
              .word
                        \mathrm{main}\;,\;\;.-\mathrm{main}
              .\,s\,i\,z\,e
116
                       "GCC: (Ubuntu 9.4.0-1ubuntu1~20.04.1) 9.4.0"
              .ident
117
                                 .note.GNU-stack, "", % progbits
              .section
118
```

编译验证正确。

3. p3.S

如下所示

p3.S

```
.arch armv5t
   @ comm section save global variable without initialization
   .comm i , 4 @ 全局变量
   .comm f, 4
   .comm n, 4
   .text
   .align 2<sup>®</sup> 设置对齐
   @ rodata section save constant
   .section .rodata
    \_str0:
    .ascii "%d\0" @ 定义字符串
12
   .align 2
    \_str1:
    .ascii "%d\n"
   .align 2
    . text
```

```
.global function
21
   function: @ function
       str fp, [sp, \#-4]! @ pre-index mode, sp = sp - 4, push fp
       mov fp, sp
       \operatorname{sub} \operatorname{sp}, \operatorname{sp}, \#12 @ allocate space for local variable
       str r0, [fp, #-8] @ r0 = [fp, #-8] = f
       str r1, [fp, \#-12] @ r1 = [fp, \#-12] = i
       mul r2, r0, r1 @计算,函数的参数输入位于r0、r1,结果存于r2寄存器
       mov r0, r2
30
       add sp, fp, \#0
       ldr fp, [sp], #4 @ post-index mode, pop fp, sp = sp + 4
       bx lr @ recover sp fp pc
   .global main
   main:
           @ main函数
37
                   {fp, lr} @ 非叶函数,需要额外存储lr
           push
       add fp, sp, #4
       ldr r1, \_bridge+8 @ r1 = &n
41
       ldr r0, \_bridge+12 @ *r0 = "%d \ 0"
                   isoc99 scanf @ 调用函数读入n, scanf("%d", &n)
           movs
                   r3, #2 @ i=2
45
       ldr r4, bridge
       str r3, [r4]
           movs
                   r3, #1 @ f=1
48
       ldr r4,_bridge+4
49
       str r3, [r4]
                   .L4 @ 无条件跳转, 进入while循环判断
           b
   .L5: @ while循环体
           ldr
                   r3, _bridge
       ldr r0, [r3] @ r0=变量i
54
                   r3, _bridge+4
           ldr
       ldr r1, [r3] @ r1=变量f
                   function @ 参数存入r0r1,调用函数
       ldr r4,_bridge+4
       str r0, [r4] @ 函数返回值存于r0,将其赋给f
                   r3, _bridge @ 将变量i赋给r3
           ldr
60
       ldr r2, [r3]
61
           adds
                   r2, r2, #1 @ r2=i+1
       ldr r4,_bridge
63
       str r2, [r4] @ i=r2
   .L4: @ while循环退出条件
       ldr r3, _bridge
```

三、 总结 编译原理实验报告

```
ldr r0, [r3] @ r0=变量i
                 r3, _bridge+8 @ 变量n
                  r2, [r3] @ 变量n
          ldr
                  r0, r2 @ 比较 i 与 n
          cmp
                  .L5 @ 若 i <= n, 重新进入循环体
          ble
          @ 退出循环
          ldr r3, _bridge+4
      ldr r1, [r3] @ r1=变量f
      ldr r0, \_bridge+16 @ *r0 = "%d\n"
                  printf @ printf("%d\n", f);
      mov r0, #0
      pop {fp, pc} @ return 0
81
   bridge: @ 变量桥梁
82
   . word i
   .word f
   .word n
   .word _str0
   .word _str1
  section .note.GNU-stack, ", % progbits
```

编译验证正确。

综上, 三个程序基本覆盖了主要的 sysY 语言特性。

三、总结

本次作业由林语盈、贺祎昕共同合作完成。

(一) 组内成员分工

对于构造上下文无关文法的任务,终结符、非终结符、终结符定义部分由二人共同完成;产生式部分,按本文中出现的顺序,贺祎昕负责 Block 及之前的定义,林语盈负责其后的定义与全部定义的整体修改。对于汇编程序编写的任务,贺祎昕负责 p1、p2 两个程序的编写和验证,林语盈负责 p3 程序的编写和验证。

(二) 完成工作总述

本次作业能够圆满完成实验要求, 具体完成情况如下:

- 1. 我们确定了编译器支持的语言特性,并参考教材及《SysY2022 语言定义》,用上下文无关文法对其进行详细的描述;
- 2. 我们一共设计了 3 个 SysY 程序,对每个程序给出其 C 语言描述,并对每个程序手动翻译 为汇编程序,再利用汇编器对翻译后的程序生成可执行程序验证正确性;
- 3. 经验证,编写的所有程序均能够运行得到正常结果。

四、 预备工作 3 源码链接

贺祎昕 林语盈

