Greedy Algorithms*

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Algorithm Course: Shanghai Jiao Tong University

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*Special thanks is given to Prof. Kevin Wayne@Princeton for sharing his lecture notes, and also given to Mr. Hongjian Cao from CS2015@SJTU for producing this slide.

Outline

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1. Basic Methodology
    - Interval Scheduling
    - Interval Partitioning
    - Scheduling to Minimize Lateness
2. More Examples
    - Optimal Caching
    - Coin Changing

Interval Scheduling: An Introductory Example

- Job j starts at s_j and finishes at f_j.
- Two jobs are compatible if they don't overlap.

Goal: find maximum subset of mutually compatible jobs.

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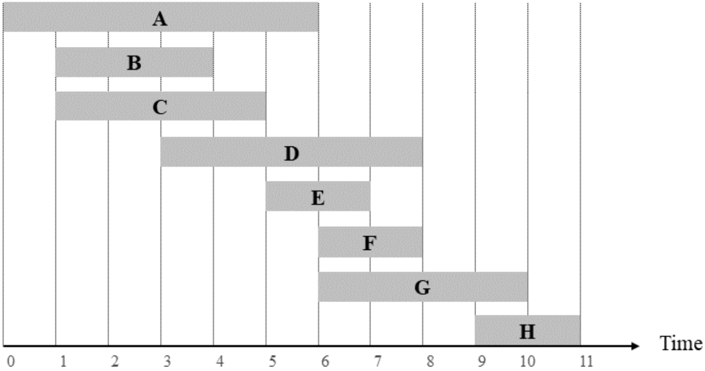

Greedy Strategy

Optimization Problem: Given a problem Π with domain $\mathbf{X}$, choose a subset or determine a sequence according to some maximization or minimization objective. (Each $X \in \mathbf{X}$ is an instance of Π)

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Note: it makes a *locally optimal* choice in hope that this choice will lead to a *globally optimal* solution.

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Note: it makes a *locally optimal* choice in hope that this choice will lead to a *globally optimal* solution.

Interval Scheduling Problem: Consider jobs in some natural order. Take each job provided to judge its compatibility with the ones already taken.

Attempts

[Earliest start time] Consider jobs in ascending order of s_j.

[Earliest finish time] Consider jobs in ascending order of f_j.

[Shortest interval] Consider jobs in ascending order of $f_j - s_j$.

[Fewest conflicts] For each job j, count the number of conflicting jobs c_j. Schedule in ascending order of c_j.

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Counter Example:

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Greedy Interval Scheduling Algorithm

Algorithm 1: Greedy Interval Scheduling

```
Sort jobs by finish times so that $f_1 \leq f_2 \leq \ldots \leq f_n$;
$A \leftarrow \emptyset$ ;                    // set of jobs selected
for $j = 1$ to $n$ do
    if job $j$ is compatible with $A$ then
        $A \leftarrow A \cup \{j\}$;
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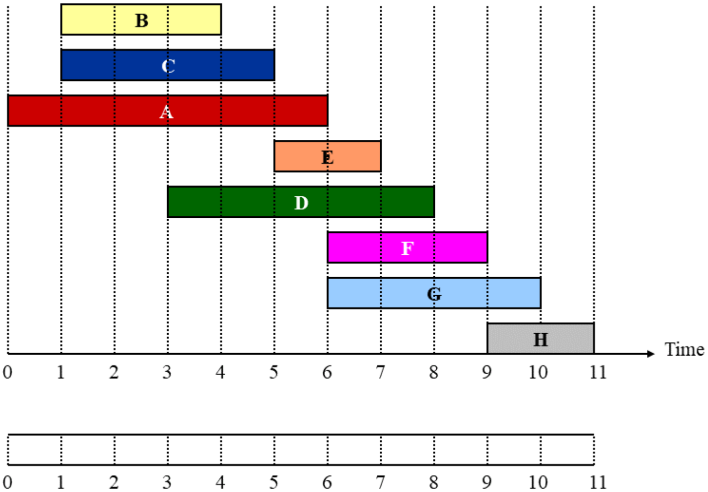
Implementation: $O(n \log n)$.

- After each iteration, set job j^* that was added last to A.
- Job j is compatible with A if $s_j \geq f_{j^*}$.

Demo



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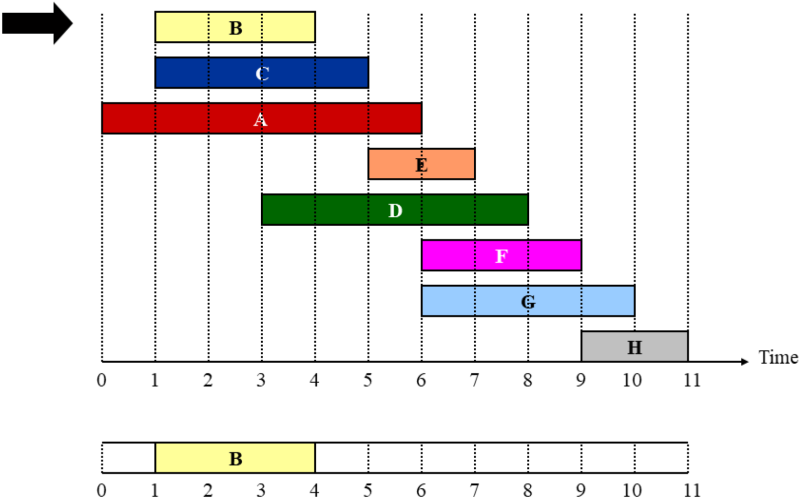

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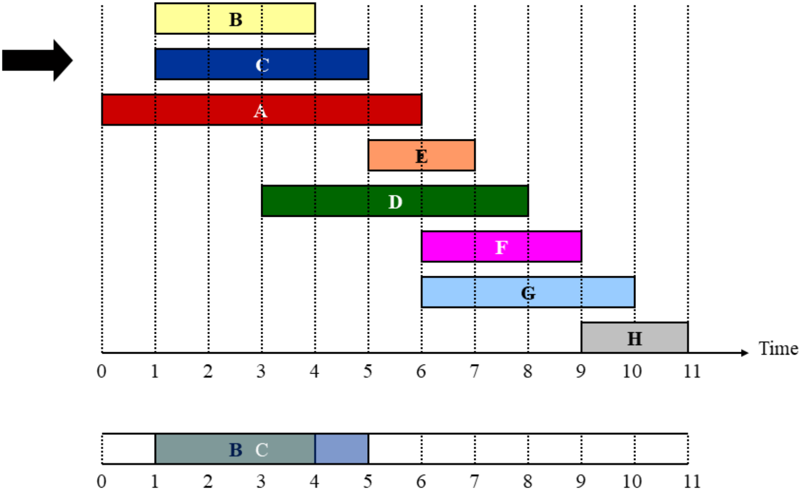

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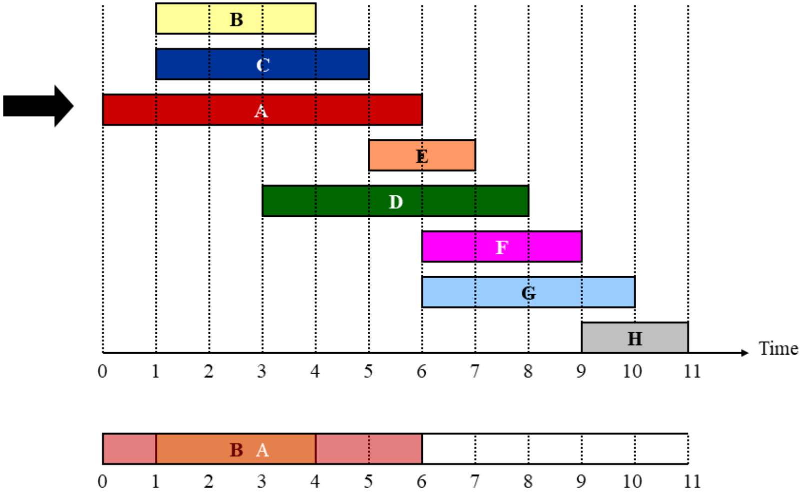

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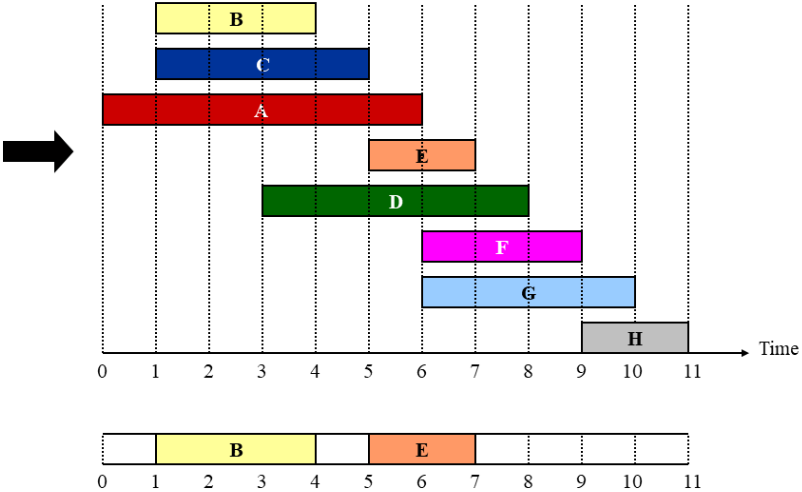

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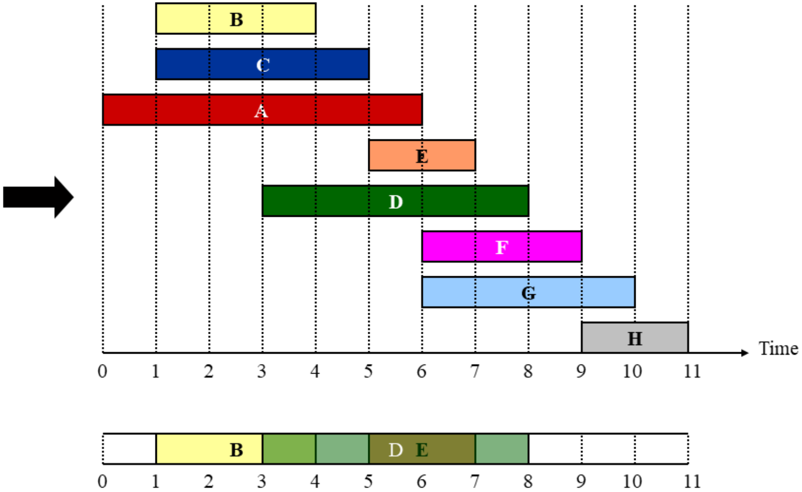

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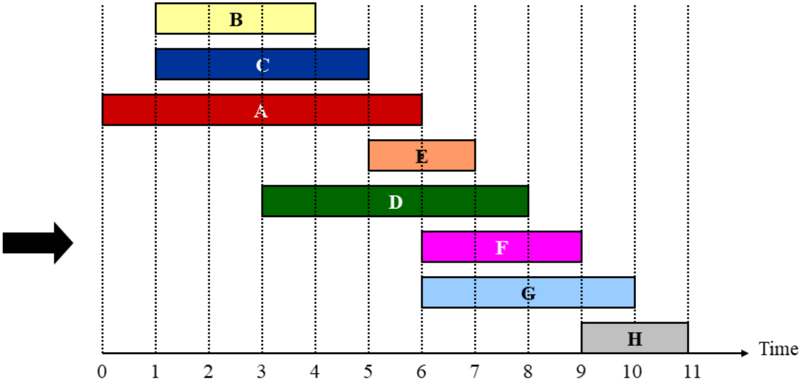

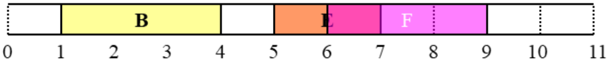

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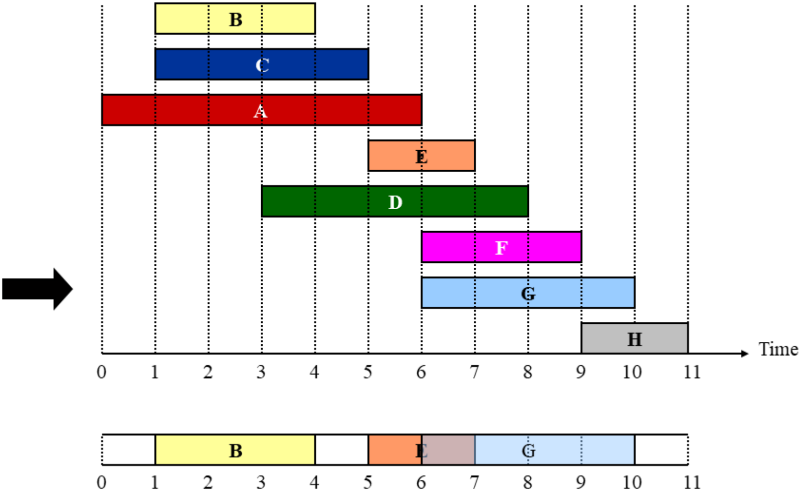

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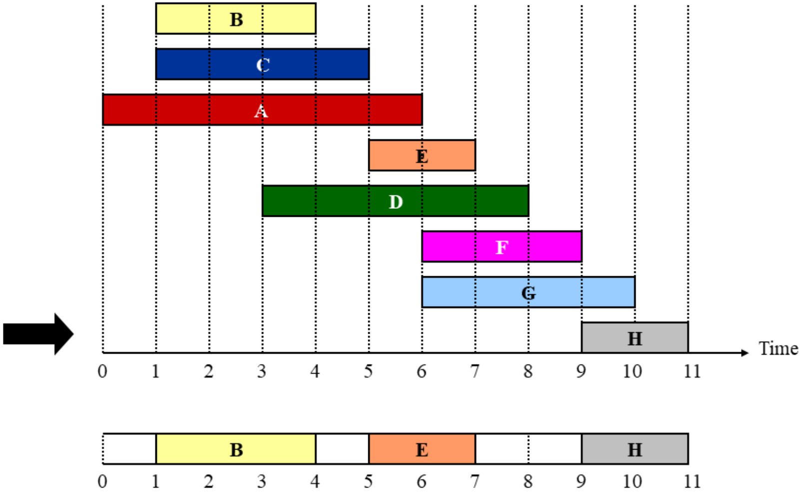

Notation

Greedy Solution: $\{B, E, H\}$

Optimal Solutions: (not necessarily unique)

$\{A, F, H\}, \{B, E, H\}, \{B, F, H\}, \{C, E, H\}, \{C, F, H\}$

Feasible Solutions: (can work, but may not be the best)

$\emptyset, \{A\}, \{B\}, \{C\}, \{D\}, \{E\}, \{F\}, \{G\}, \{H\};$
$\{A, F\}, \{A, G\}, \{A, H\}, \{B, E\}, \{B, F\}, \{B, G\}, \{B, H\},$
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Let $i_1, i_2, \cdots, i_k$ denote set of jobs selected by greedy.

Let $j_1, j_2, \cdots, j_m$ denote set of jobs in an optimal solution with $i_1 = j_1$, $i_2 = j_2, \cdots, i_r = j_r$ for the largest possible value of r.

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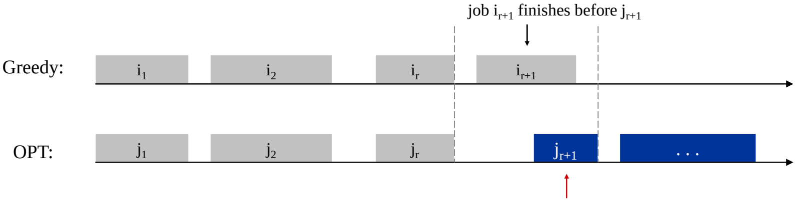

why not replace job j_{r+1}
with job i_{r+1}?

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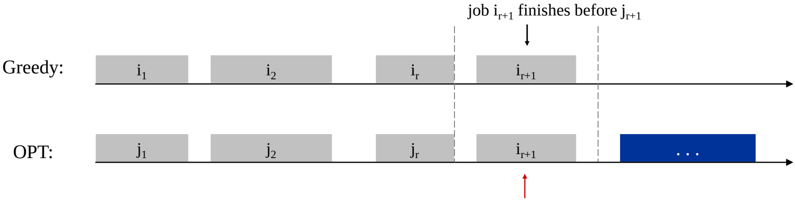

solution still feasible and optimal,
but contradicts the maximality of r.

Outline

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Lecture j starts at s_j and finishes at f_j.

Goal: find minimum number of classrooms to schedule all lectures so that no two occur at the same time in the same room.

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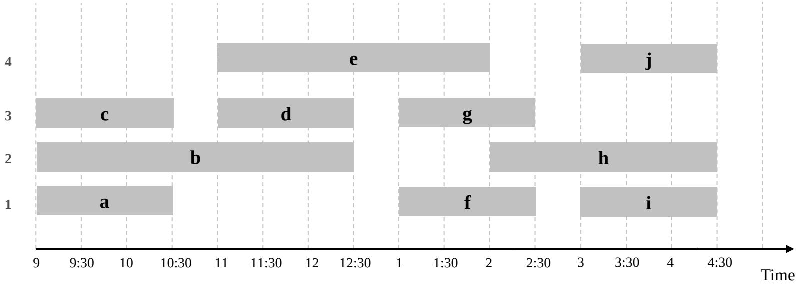
Example: This schedule uses 4 classrooms to schedule 10 lectures.



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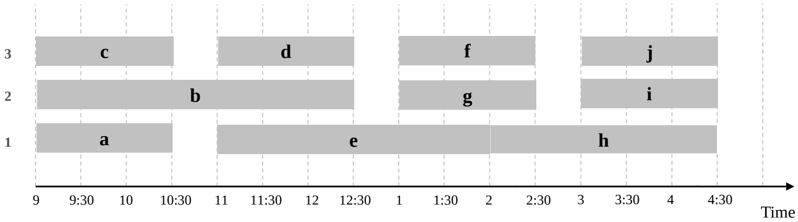
Example: This schedule uses only 3.



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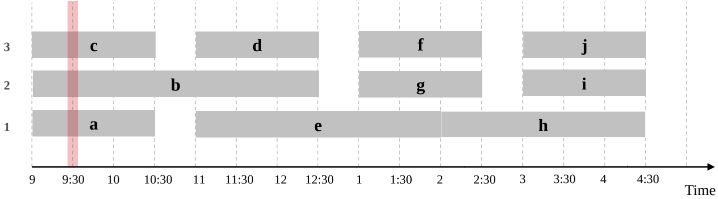

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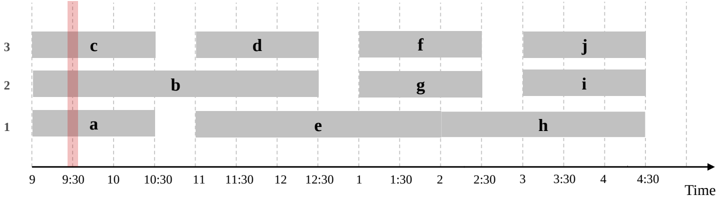

Question. Does there always exist a schedule equal to depth of intervals?

Greedy Interval Partitioning Algorithm

Algorithm 2: Interval Partitioning Greedy Algorithm

```
Sort intervals by starting time so that $s_1 \leq s_2 \leq ... \leq s_n$;
$d \leftarrow 0$ ;              // number of allocated classrooms
for $j = 1$ to $n$ do
    if lecture j is compatible with some classroom k then
        schedule lecture $j$ in classroom $k$;
    else
        allocate a new classroom $d + 1$;
        schedule lecture $j$ in classroom $d + 1$;
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return $A$;
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Implementation: $O(n \log n)$.

- For classroom k, maintain the finish time of the last job added.
- Keep the classrooms in a priority queue.

Correctness Proof

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Key observation $\Rightarrow$ all schedules use $\geq d$ classrooms.

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   - Scheduling to Minimize Lateness
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Scheduling to Minimize Lateness

- Single resource processes one job at a time.
- Job j requires t_j units of processing time and is due at time d_j.
- If j starts at time s_j, it finishes at time $f_j = s_j + t_j$.
- Lateness: $l_j = \max\{0, f_j - d_j\}$.

Goal: schedule all jobs to minimize maximum lateness $L = \max l_j$.

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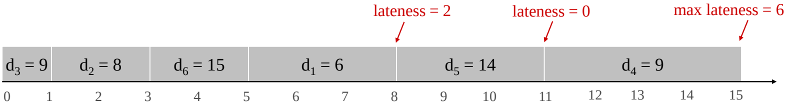
	1	2	3	4	5	6
t_j	3	2	1	4	3	2
d_j	6	8	9	9	14	15



Attempt: Consider jobs in ascending order by some strategy

[Shortest processing time first] Sort by processing time t_j.

[Earliest deadline first] Sort by deadline d_j.

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counterexample

	1	2
t_j	1	10
d_j	100	10

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d_j	2	10

A Greedy Algorithm: Earliest Deadline First

Algorithm 3: Greedy Minimizing Lateness

```
Sort $n$ jobs by deadline so that $d_1 \leq d_2 \leq ... \leq d_n$;
$t \leftarrow 0$;
for $j = 1$ to $n$ do
    Assign job $j$ to interval $[t, t + t_j]$;
    $s_j \leftarrow t, f_j \leftarrow t + t_j$;
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return intervals $[s_j, f_j]$;
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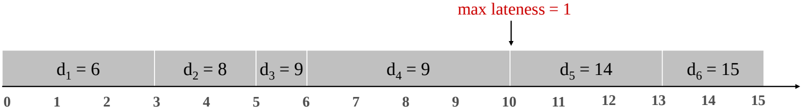
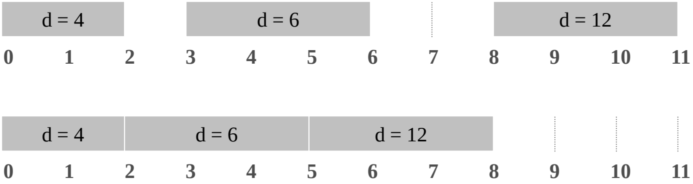
Implementation: $O(n \log n)$.



Correctness Proof: Reduce Optimal Solution

Observation. There exists an optimal schedule with no idle time.



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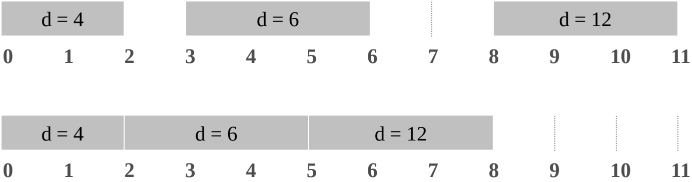
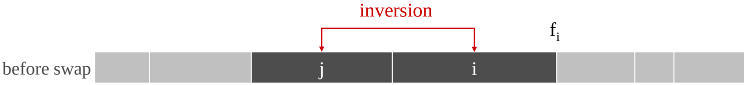


Observation. The greedy schedule has no idle time.

Correctness Proof: Optimal Solution vs Algorithm Solution

Definition. Given a schedule S, an inversion is a pair of jobs i and j such that: $i < j$ but j scheduled before i.



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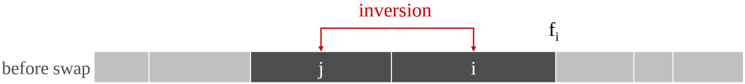

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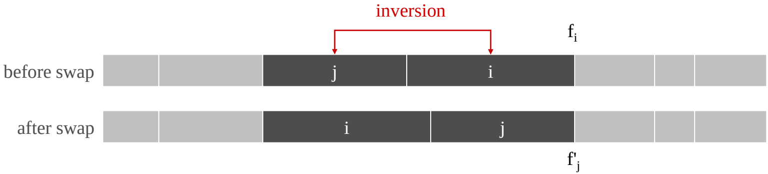
Observation. Greedy schedule has no inversions.

Observation. If a schedule (with no idle time) has an inversion, it has one with a pair of inverted jobs scheduled consecutively.

Correctness Proof: Inversions

Claim. Swapping two consecutive, inverted jobs reduces the number of inversions by one and does not increase the max lateness.



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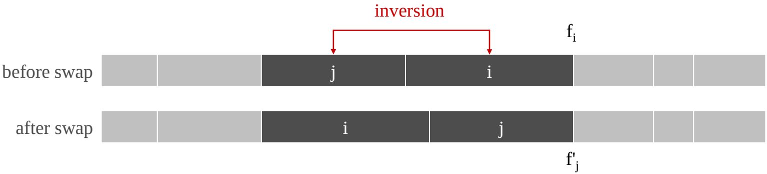

Proof. Let l be the lateness before the swap, and let l' be it afterwards.

- $l'_k = l_k$ for all $k \neq i, j$
- $l'_i \leq l_i$
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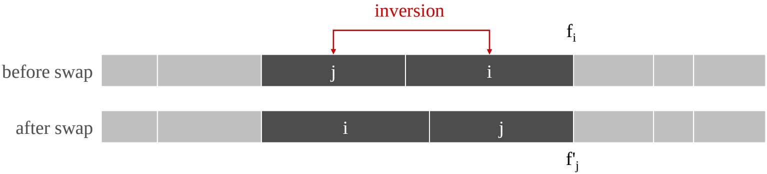

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$$\begin{aligned} l'_j &= f'_j - d_j && \text{(definition)} \\ &= f_i - d_j && \text{(j finishes at time f_i)} \\ &\leq f_i - d_i && (i < j) \\ &\leq l_i && \text{(definition)} \end{aligned}$$

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- ◦ If S^* has no inversions, then $S = S^*$.

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- If S^* has no inversions, then $S = S^*$.
- If S^* has an inversion, let $i - j$ be an adjacent inversion.

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- If S^* has no inversions, then $S = S^*$.
- If S^* has an inversion, let $i - j$ be an adjacent inversion.

  Swapping i and j does not increase the maximum lateness and strictly decreases the number of inversions. This contradicts definition of S^*. □

Greedy Analysis Strategies

Greedy algorithm stays ahead. Show that after each step of the greedy algorithm, its solution is at least as good as any other algorithm's.

Structural. Discover a simple "structural" bound asserting that every possible solution must have a certain value. Then show that your algorithm always achieves this bound.

Exchange argument. Gradually transform any solution to the one found by the greedy algorithm without hurting its quality.

Other greedy algorithms. Kruskal, Prim, Dijkstra, Huffman, ...

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Optimal Offline Caching

- Cache with capacity to store k items.
- Sequence of m item requests $d_1, d_2, ..., d_m$.
- Cache hit: item already in cache when requested.
- Cache miss: item not in cache when requested: must bring requested item into cache, and evict some existing item, if full.

Goal. Eviction schedule that minimizes number of cache misses.

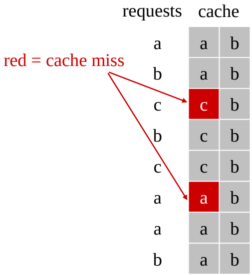
Optimal Offline Caching

- Cache with capacity to store k items.
- Sequence of m item requests $d_1, d_2, \ldots, d_m$.
- Cache hit: item already in cache when requested.
- Cache miss: item not in cache when requested: must bring requested item into cache, and evict some existing item, if full.

Goal. Eviction schedule that minimizes number of cache misses.

Example. $k = 2$, initial cache = ab,

requests: a, b, c, b, c, a, a, b.



red = cache miss

requests	cache	
a	a	b
b	a	b
c	c	b
b	c	b
c	c	b
a	a	b
a	a	b
b	a	b

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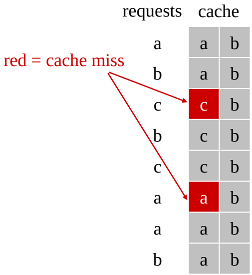
Example. $k = 2$, initial cache $= ab$,

requests: a, b, c, b, c, a, a, b.

Optimal eviction schedule: 2 cache misses.



red = cache miss

requests	cache	
a	a	b
b	a	b
c	c	b
b	c	b
c	c	b
a	a	b
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Farthest-in-future. Evict item in the cache that is not requested until farthest in the future.

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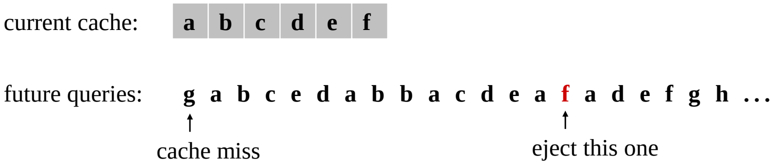

current cache:

a	**b**	**c**	**d**	**e**	**f**

future queries: **g a b c e d a b b a c d e a f a d e f g h . . .**

cache miss

eject this one

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↑ cache miss (g)

↑ eject this one (f)

Theorem. [Bellady, 1960s] FF is an optimal eviction schedule.

Proof. Algorithm and theorem are intuitive; proof is subtle.

Reduced Eviction Schedules

Definition. A reduced schedule is a schedule that only inserts an item into the cache in a step in which that item is requested.

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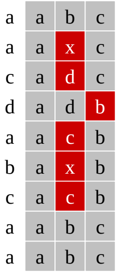
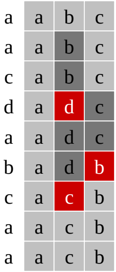
Intuition. Can transform an unreduced schedule into a reduced one with no more cache misses (means no more cache insertion/replacement here).



a	a	b	c
a	a	x	c
c	a	d	c
d	a	d	b
a	a	c	b
b	a	x	b
c	a	c	b
a	a	b	c
a	a	b	c

an unreduced schedule



a	a	b	c
a	a	b	c
c	a	b	c
d	a	d	c
a	a	d	c
b	a	d	b
c	a	c	b
a	a	c	b
a	a	c	b

a reduced schedule

Reduced Eviction Schedules

Claim. Given any unreduced schedule S, we can transform it into a reduced schedule S' with no more cache replacement (insertion).

[†] doesn't enter cache at requested time

Reduced Eviction Schedules

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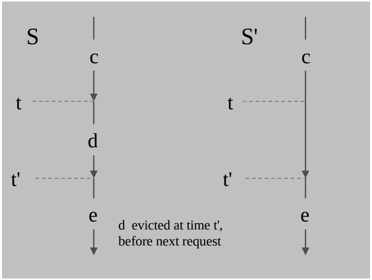
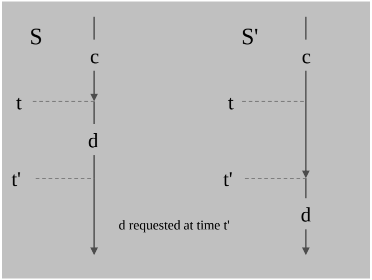
Proof. (by induction on number of unreduced[†] items)
Suppose S brings d into the cache at time t, without a request. Let c be the item S evicts when it brings d into the cache.



Case 1



Case 2

[†] doesn't enter cache at requested time

Theorem. FF is an optimal eviction algorithm

Proof. (by induction on number of requests j)

Invariant: There exists an optimal reduced schedule S that makes the same eviction schedule as S_{FF} through the first $j + 1$ requests.

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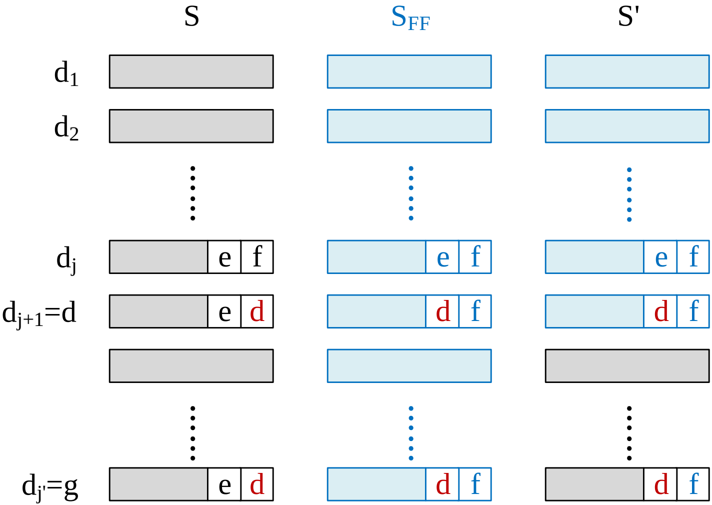
Consider $(j + 1)^{th}$ request $d = d_{j+1}$. Since S and S_{FF} have agreed up until now, they have the same cache contents before request $j + 1$.

- Case 1: (d is already in the cache) $S' = S$ satisfies invariant.
- Case 2: (d is not in the cache; S, S_{FF} evict same element) $S' = S$ satisfies invariant.
- Case 3: (d is not in the cache; S_{FF} evicts e; S evicts $f \neq e$) Let S' agree with S_{FF} at the $j + 1$ requests; we show that having element f in cache is no worse than having element e.

An Illustration of Case 3



Correctness Proof (Continued)

Let j' be the first time after $j + 1$ that S and S' take a different action (must involve e or f or both), and let g be item requested at time j'.

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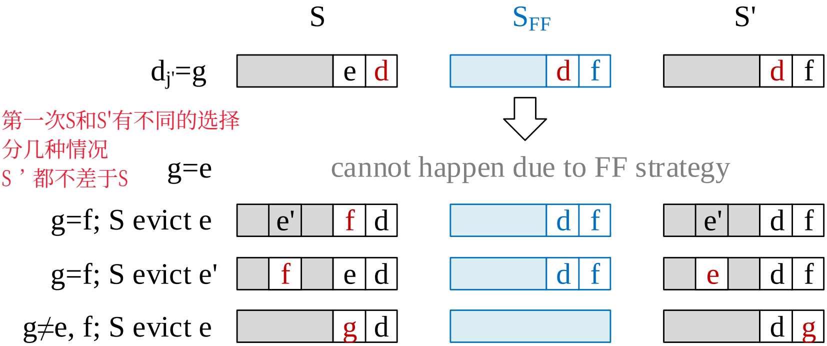

第一次S和S'有不同的选择
分几种情况
S’都不差于S

Correctness Proof (Continued)

Case 3a: $g = e$. Can't happen with Farthest-In-Future since there must be a request for f before e.

Case 3b: $g = f$. Element f can't be in cache of S, so let e' be the element that S evicts.

- if $e' = e$, S' accesses f from cache; now S and S' have same cache
- if $e' \neq e$, S' evicts e' and brings e into the cache; now S and S' have the same cache.

Case 3c: $g \neq e, f$. S must evict e (otherwise S' would take the same action). Make S' evict f; now S and S' have the same cache. $\square$

Caching Perspective

Online vs. Offline algorithms.

- Offline: full sequence of requests is known a priori.
- Online (reality): requests are not known in advance.
- Caching is among most fundamental online problems in CS.

---

‡FF with direction of time reversed!

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LIFO. Evict page brought in most recently.
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Theorem. FF is optimal offline eviction algorithm.

- Provides basis for understanding and analyzing online algorithms.
- LRU is k-competitive. [Section 13.8 in Cornell Book]
- LIFO is arbitrarily bad.

[‡] FF with direction of time reversed!

Outline

1. Basic Methodology
   - Interval Scheduling
   - Interval Partitioning
   - Scheduling to Minimize Lateness

2. More Examples
   - Optimal Caching
   - Coin Changing

Coin Changing

Goal. Given US currency denominations:

1 (cent), 5 (nickel), 10 (dime), 25 (quarter), 100 (dollar),

devise a changing method using fewest number of coins.

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Example. $2.89.

Cashier's Algorithm

Algorithm 4: Cashier's Algorithm

```
Sort coins denominations by value: $c_1 < c_2 < ... < c_n$;
$S \leftarrow \emptyset$ ;                                  // coins selected
while $x \neq 0$ do
    let $k$ be largest integer such that $c_k < x$;
    if $k = 0$ then
        return "no solution found";
    $x \leftarrow x - c_k$;
    $S \leftarrow S \cup \{k\}$;
return $S$;
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Question. Is cashier's algorithm optimal?

Properties of Optimal Solution

Property. Number of pennies ≤ 4.
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Property.Number of nickels + Number of dimes ≤ 2.
Proof.

- Replace 3 dimes and 0 nickels with 1 quarter and 1 nickel;
- Replace 2 dimes and 1 nickel with 1 quarter.
- Recall: at most 1 nickel.

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If not, it needs enough coins of type $c_1, \ldots, c_{k-1}$ to add up to x.

k	c_k	All optimal solutions must satisfy	Max value of coins 1, 2, …, k-1 in any OPT
1	1	P ≤ 4	-
2	5	N ≤ 1	4
3	10	N + D ≤ 2	4 + 5 = 9
4	25	Q ≤ 3	20 + 4 = 24
5	100	no limit	75 + 24 = 99

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Problem reduces to coin-changing $x - c_k$ cents, which, by induction, is optimally solved by greedy algorithm.

Is Cashier's Algorithm Work for Any Denominations?

Observation 1. Greedy is sub-optimal for US postal denominations:

$$1, 10, 21, 34, 70, 100, 350, 1225, 1500.$$

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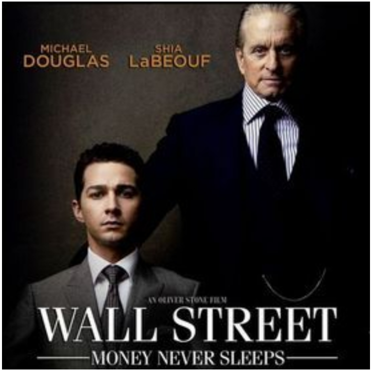
Observation 2. Even no feasible solution with system ¢={7, 8, 9}.

- Cashier's algorithm: 15¢= 9+???
- Optimal: 15¢= 7 + 8.

Movie: Wall Street (1987)



§

Greed is good.

Greed is right.

Greed works.

Greed clarifies, cuts through, and captures the essence of the evolutionary spirit.

— Gordon Gecko
(Michael Douglas)

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§ Watch the movie segment at the class webpage.