

Lab04-Dynamic Programming

CS214-Algorithm and Complexity, Xiaofeng Gao, Spring 2019.

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1. Given a positive integer n , find the least number of perfect square numbers (e.g., 1, 4, 9, ...) which sum to n .
 - (a) Assume that $OPT(a)$ = the least number of perfect square numbers which sum to a . Please write a recurrence for $OPT(a)$.

$$OPT(a) = \begin{cases} 1, & a = i^2 \\ \min_{j^2 \leq a} \{OPT(a - j^2) + 1\}, & \text{otherwise} \end{cases}$$

- (b) Base on the recurrence, write down your algorithm in the form of *pseudo code*.

Algorithm 1: perfectSquare

Input: An integer n

Output: m (the least number of perfect square numbers which sum to n)

```
1 for  $j \leftarrow 1$  to  $n$  do
2    $M[j] = \text{empty}$ ;
3 Run Find-Solution( $n$ );
4 Function  $M\text{-Find-Solution}(j)$ :
5   if  $M[j]$  is empty then
6     if  $[j^{1/2}]^2 == j$  then
7        $M[j] = 1$ ;
8     else
9        $M[j] = \min_{i^2 \leq j} \{M\text{-Find-Solution}(j - i^2) + 1\}$ ;
10  return  $M[j]$ ;
```

2. Given an input string s (could be empty, and contains only lowercase letters a-z) and a pattern p (could be empty, and contains only lowercase letters a-z and characters like '?' or '*'), please design an algorithm using dynamic programming to determine whether s matches p based on the following rules:

- '?' matches any single character.
- '*' matches any sequence of characters (including the empty sequence).
- The matching should cover the entire input string (not partial).

Assume $m = \text{len}(s)$ and $n = \text{len}(p)$. Output **true** if s matches p , or **false** otherwise.

- (a) Assume that $ANS(i, j)$ means whether the first i ($0 \leq i \leq m$) characters of s match the first j ($0 \leq j \leq n$) characters of p . Please write a recurrence for $ANS(i, j)$.

$$ANS(0, j) = \begin{cases} \text{true}, & j = 0 \\ ANS(0, j - 1), & p[j - 1] = '*' \\ \text{false}, & \text{otherwise} \end{cases}$$

$$ANS(i, 0) = \begin{cases} true, & i = 0 \\ false, & \text{otherwise} \end{cases}$$

for $i, j > 0$,

$$ANS(i, j) = \begin{cases} ANS(i, j-1) \text{ or } ANS(i-1, j-1) \text{ or } ANS(i-1, j), & p[j-1] = '*' \\ ANS(i-1, j-1), & p[j-1] = '?' \\ ANS(i-1, j-1) \text{ and } (s[i-1] == p[j-1]), & \text{otherwise} \end{cases}$$

(b) Base on the recurrence, write down your algorithm in the form of *pseudo code*.

Algorithm 2: stringMatching

Input: an input string s (could be empty, and contains only lowercase letters a-z) and a pattern p (could be empty, and contains only lowercase letters a-z and characters like '?' or '*')

Output: **true** if s matches p , or **false** otherwise

```

1   $m \leftarrow \text{len}(s); n \leftarrow \text{len}(p);$ 
2   $M[0] \leftarrow true;$ 
3  for  $j \leftarrow 1$  to  $n$  do
4      if  $p[j-1] == '*'$  then
5           $M[j] \leftarrow M[j-1];$ 
6      else
7           $M[j] \leftarrow false;$ 
8  for  $i \leftarrow 1$  to  $m$  do
9       $N[0] \leftarrow false;$ 
10     for  $j \leftarrow 1$  to  $n$  do
11         if  $p[j-1] == '*'$  then
12              $N[j] \leftarrow (M[j-1] \text{ or } M[j] \text{ or } N[j-1]);$ 
13         else if  $p[j-1] == '?'$  then
14              $N[j] \leftarrow M[j-1];$ 
15         else
16              $N[j] \leftarrow (M[j-1] \text{ and } (s[i-1] == p[j-1]));$ 
17     swap  $M[]$  and  $N[];$ 
18 output  $N[n];$ 

```

(c) Analyze the time and space complexity of your algorithm.

Solution. Time complexity : The algorithm computes every $ANS(i, j)$ ($0 \leq i \leq m$, $0 \leq j \leq n$) once, so the time complexity is $O(mn)$.

Space complexity : For every $ANS(i, \cdot)$, it needs $ANS(i-1, \cdot)$ and itself, which result in the space complexity of $O(n)$. \square

3. Recall the *String Similarity* problem in class, in which we calculate the edit distance between two strings in a sequence alignment manner.

(a) Implement the algorithm combining dynamic programming and divide-and-conquer strategy in C/C++ with time complexity $O(mn)$ and space complexity $O(m+n)$. (The template [Code-SequenceAlignment.cpp](#) is attached on the course webpage).

- (b) Given $\alpha(x, y) = |\text{ascii}(x) - \text{ascii}(y)|$, where $\text{ascii}(c)$ is the ASCII code of character c , and $\delta = 13$. Find the edit distance between the following two strings.

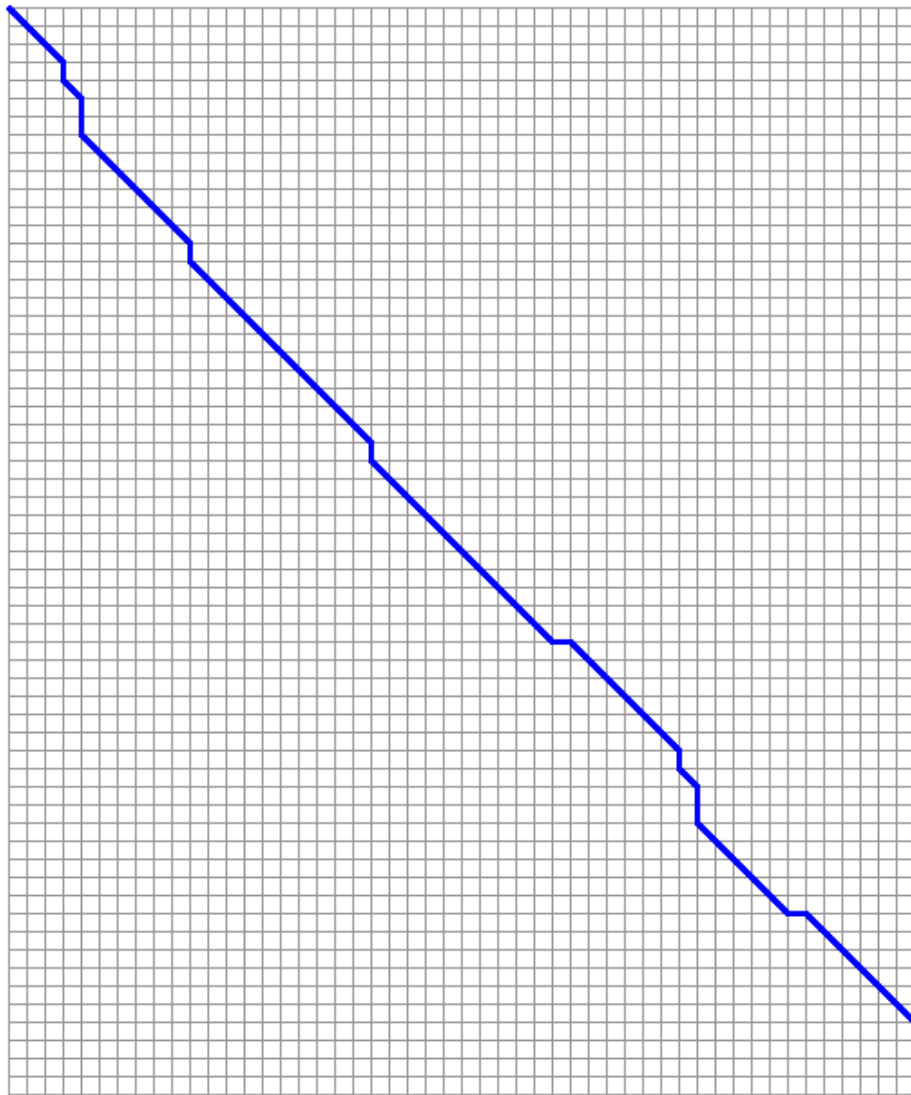
$X[1..60] = PSQAKADIETSJPWUOMZLNLOMOZNLTLQ$
 $CFQHZZRIQOQCOCFPRWOUXXCEMYSWUJ$

$Y[1..50] = SUYLVUSDROFBXUDCOHAAEBKN$
 $AAPNXEVWNLMYUQRPEOCQOCIMZ$

Solution. Run *Code-SequenceAlignment.cpp* in (a) we get that the distance is 439. \square

- (c) **(Bonus)** Visualize the shortest path found in (b) on the corresponding edit distance graph using any tools you like.

Solution. Using **Visual Studio** and **Open CV**, I get the graph as follows. And the code is showed in *VisualizeThePath.cpp*.



\square

Remark: You need to include your .cpp, .pdf and .tex files in your uploaded .rar or .zip file.