

Mathematical Foundations of Computer Science

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8 Partial Orderings

Recall the definition of a *partial ordering* and an *equivalence relation*. A relation \preceq on a set X is a partial ordering if (1) it is *reflexive*, i.e., $x \preceq x$ for all $x \in X$; (2) *transitive*, i.e., $x \preceq y$ and $y \preceq z$ together imply $x \preceq z$; and (3) *anti-symmetric*, i.e., $x \preceq y$ and $y \preceq x$ only hold if $x = y$. A relation \sim is an equivalence relation if it is (1) reflexive, (2) transitive, and (3) *symmetric*, i.e., $x \sim y$ if and only if $y \sim x$.

8.1 Equivalence Relations as a Partial Ordering

A *partition* \mathcal{P} of V is a set $\{V_1, \dots, V_k\}$ where (1) $V_1 \cup \dots \cup V_k = V$ and (2) the V_i are pairwise disjoint, i.e., $V_i \cap V_j = \emptyset$ for $1 \leq i < j \leq k$. For example, $\{\{1\}, \{2, 3\}, \{4\}\}$ is a partition of $\{1, 2, 3, 4\}$ but $\{\{1\}, \{2, 3\}, \{1, 4\}\}$ is not.

A partition \mathcal{P} of V defines an equivalence relation \sim on V : set $x \sim y$ if and only if x, y are in the same part of \mathcal{P} ; conversely, an equivalence relation \sim defines a partition: throw equivalent elements into the same set; formally,

$$\begin{aligned} E(x) &:= \{y \in V \mid x \sim y\} && \text{(the set of elements equivalent to } x\text{)} \\ \mathcal{P} &:= \{E(x) \mid x \in V\} . \end{aligned}$$

Thus, equivalence relations and partitions are basically the same thing, just represented in a different way. For example, the partition $\{\{1\}, \{2, 3\}, \{4\}\}$ induces an equivalence relation R on $\{1, 2, 3, 4\}$ in which $2 \sim 3$ are equivalent but all other elements are not. Formally, written in set notation, we get

$$R = \{(1, 1), (2, 2), (3, 3), (4, 4), (2, 3), (3, 2)\} .$$

Exercise 8.1. Let E_4 be the set of all equivalence relations on $\{1, 2, 3, 4\}$. Note that E_4 is ordered by set inclusion, i.e.,

$$(E_4, \{(R_1, R_2) \in E_4 \times E_4 \mid R_1 \subseteq R_2\})$$

is a partial ordering.

1. Draw the Hasse diagram of this partial ordering in a nice way.
2. What is the size of the largest chain?
3. What is the size of the largest antichain?

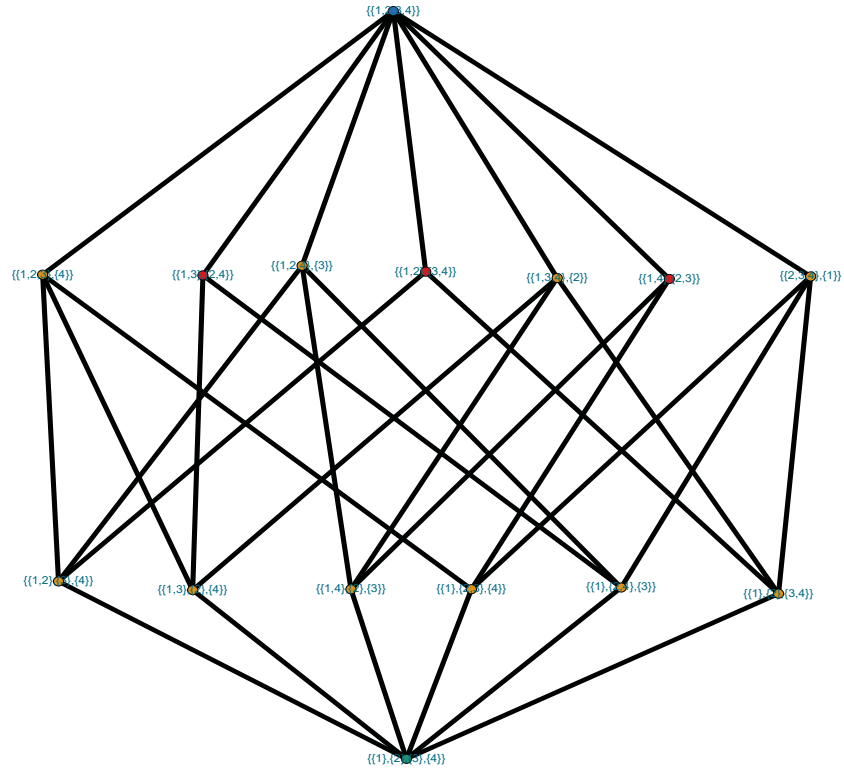


Figure 1: Hasse diagram of the partial ordering

Solution

1. See Figure 1, where relation R is depicted as partition explicitly.
2. As is seen in the figure, the height of the Hasse diagram is 4, and so is the size of the largest chain.
3. As is seen in the figure, the width of the Hasse diagram is 7, and so is the size of the largest antichain.

8.2 Chains and Antichains

Define the partially ordered set (\mathbb{N}_0^n, \leq) as follows: $x \leq y$ if $x_i \leq y_i$ for all $1 \leq i \leq n$. For example, $(2, 5, 4) \leq (2, 6, 6)$ but $(2, 5, 4) \not\leq (3, 1, 1)$.

Exercise 8.2. Consider the infinite partially ordered set (\mathbb{N}_0^n, \leq) .

1. Which elements are minimal? Which are maximal?
2. Is there a minimum? A maximum?
3. Does it have an infinite chain?
4. Does it have arbitrarily large antichains? That is, can you find an antichain A of size $|A| = k$ for every $k \in \mathbb{N}$?

Solution

1. **Claim.** Element x such that for all $1 \leq i \leq n$, $x_i = 0$ is the only minimal element, and there is no maximal element.

Proof. First suppose that there exists some element y , $\exists i, 1 \leq i \leq n$ such that $y_i > 0$. Then we can find an element $y' = (y_1, y_2, \dots, y_i - 1, \dots, y_n) \leq y$, which contradicts with the fact that y is minimal.

Then suppose that there exists some maximal element $z = (z_1, z_2, \dots, z_i, \dots, z_n)$. Then we can find an element $z' = (z_1, z_2, \dots, z_i + 1, \dots, z_n) \geq z$, which contradicts with the fact that z is maximal. \square

Therefore the only minimal element is $(0, 0, \dots, 0)$ and there is no maximal element.

2. **Claim.** $x_0 = (0, 0, \dots, 0)$ is a minimum element, and there is no maximum element.

Proof. For any $x \in \mathbb{N}_0^n$, $\forall i$ in $1 \leq i \leq n$, $0 \leq x_i$. That is, $\forall x \in \mathbb{N}_0^n$, $x_0 \leq x$. Thus x_0 is the minimum element.

Since there is no maximal element, there is no maximum element is easy to see. \square

3. Yes it does have an infinite chain. Or it is better to say it has infinitely many infinite chains. For example, $\{(0, 0, \dots, 0), (1, 0, \dots, 0), (2, 0, \dots, 0), \dots\}$ is one infinite chain. And if we remove finite elements from this chain, still we get an infinite chain. We have infinite ways to remove finite elements, so there are infinitely many infinite chains.
4. Yes it does when $n \geq 2$. But according to the definition of antichain, the smallest size of an antichain is 2 since two elements have to be distinct to be incomparable. So it is better to say that we can find an antichain A of size $|A| = k$ for every $k \in \mathbb{N} \setminus \{1\}$.

Proof. When $n = 1$, every element in \mathbb{N}_0^n is comparable, there is no antichain.

When $n \geq 2$, we can always find such an antichain of size k : $\{(k-1, 0, 0, \dots, 0), (k-2, 1, 0, \dots, 0), \dots, (0, k-1, 0, \dots, 0)\}$ for every $k \in \mathbb{N} \setminus \{1\}$. \square

***Exercise 8.3.** Does every infinite subset $S \subseteq \mathbb{N}_0^n$ contain an infinite chain?

Solution. There are n dimensions in each element. We denote the element as:

$$x = (v_1, v_2, \dots, v_n)$$

Consider a subset of S in which the i^{th} dimension value of its element is a given constant:

$$S_{v_i=c} = \{x \mid x \in S, v_i = c\}$$

Consider the elements number of this subset. If $|S_{v_i=c}|$ is infinite, then we only have to prove $S_{v_i=c} \subseteq \mathbb{N}_0^n$ contain an infinite chain. Therefore, without loss of generalization, we assume for every i and c , $|S_{v_i=c}|$ is not infinite:

$$|S_{v_i=c}| < \infty, i \in \{1, 2, \dots, n\}, c \in \mathbb{N}$$

Because S is an infinite set, then, for any sufficiently large k :

$$\begin{aligned} |S_{v_i \leq k}| &= |\{x \mid x \in S, v_i \leq k\}| < \infty \\ |S_{v_i > k}| &= |S| - |S_{v_i \leq k}| = \infty \end{aligned}$$

Suppose we find a chain of finite size in $S_{v_i \leq k}$, and the maximal element in this chain is $(m_1, m_2, \dots, m_i, \dots, m_n)$, $m_i \leq k$. According to our assumption, for each $j \neq i$:

$$|S_{v_i \geq k, v_j < m_j}| = |\{x \mid x \in S, v_i \geq k, v_j < m_j\}| < \infty$$

Therefore, there are infinite elements larger than the maximal element in the chain:

$$|\{x \mid x \in S, v_i \geq k, v_j \geq m_j \text{ for each } j \neq i\}| = \infty$$

We randomly choose one element $(m'_1, m'_2, \dots, m'_i, \dots, m'_n)$ and add it to the chain. Then we update the value of k , let $k := m'_i$. In this way we can continue this procedure indefinitely and find an infinite chain.

Exercise 8.4. Show that (\mathbb{N}_0^n, \leq) has no infinite antichain. **Hint.** Use the previous exercise.

Proof. Suppose that (\mathbb{N}_0^n, \leq) has an infinite antichain. Then the infinite antichain itself is an infinite subset of \mathbb{N}_0^n , and this infinite subset does not contain an infinite chain. This contradicts with the conclusion of **Exercise 8.3**. \square

Consider the induced ordering on $\{0, 1\}^n$. That is, for $x, y \in \{0, 1\}^n$ we have $x \leq y$ if $x_i \leq y_i$ for every coordinate $i \in [n]$.

Exercise 8.5. Draw the Hasse diagrams of $(\{0, 1\}^n, \leq)$ for $n = 2, 3$.

Proof. The figures are as follows. \square

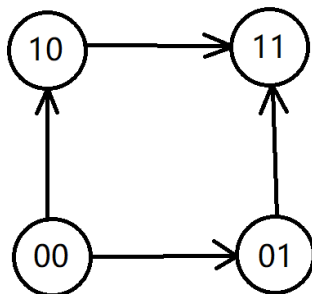


Figure 2: The Hasse diagram of $(\{0, 1\}^n, \leq)$ for $n = 2$

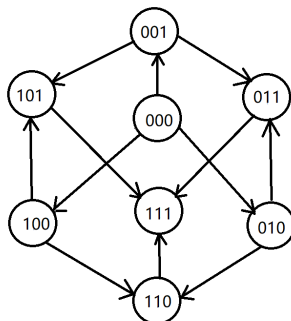


Figure 3: The Hasse diagram of $(\{0, 1\}^n, \leq)$ for $n = 3$

Exercise 8.6. Determine the maximum, minimum, maximal, and minimal elements of $\{0, 1\}^n$.

Proof. maximum: $111 \cdots 1$
 minimum: $000 \cdots 0$
 maximal: $111 \cdots 1$
 minimal: $000 \cdots 0$

□

Exercise 8.7. What is the longest chain of $\{0, 1\}^n$?

Proof. The length of the longest chain is $n + 1$.

There are $n!$ longest chains. Here is an example: $000 \cdots 0 \rightarrow 000 \cdots 01 \rightarrow 000 \cdots 011 \rightarrow \cdots \rightarrow 011 \cdots 11 \rightarrow 111 \cdots 11$. □

****Exercise 8.8.** What is the largest antichain of $\{0, 1\}^n$?

Solution. As shown in the picture, the Hasse diagram of $\{0, 1\}^n$ can be

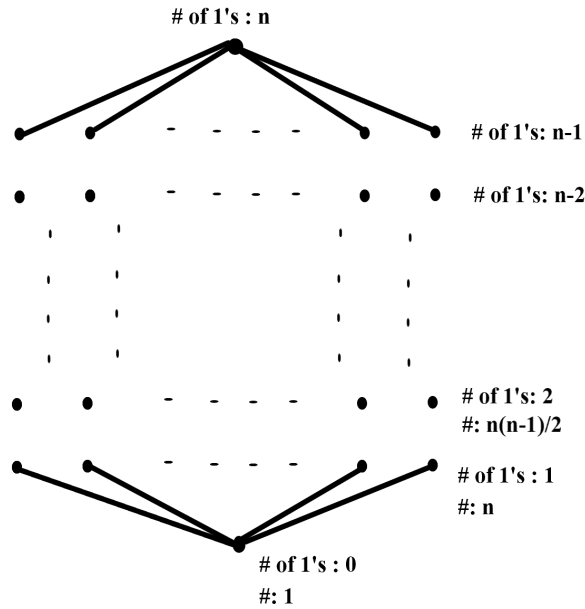


Figure 4: The Hasse diagram of $(\{0, 1\}^n, \leq)$

divided into $n+1$ layers. Since the number of 1's in each layer increases from 0 to n , we have the number of nodes of each layer equal to $\binom{n}{k}$, with $0 \leq k \leq n$. Thus, the largest antichain should be $\{x \mid x \in \{0, 1\}^n, \text{number of 1's in } x \text{ is } \lfloor \frac{n}{2} \rfloor\}$.

Proof. Number all the layers in the Hasse diagram 0 to n from bottom to top. By mathematical induction, we can easily derive that the the number of 1's in the node of i^{th} layer is i . \square