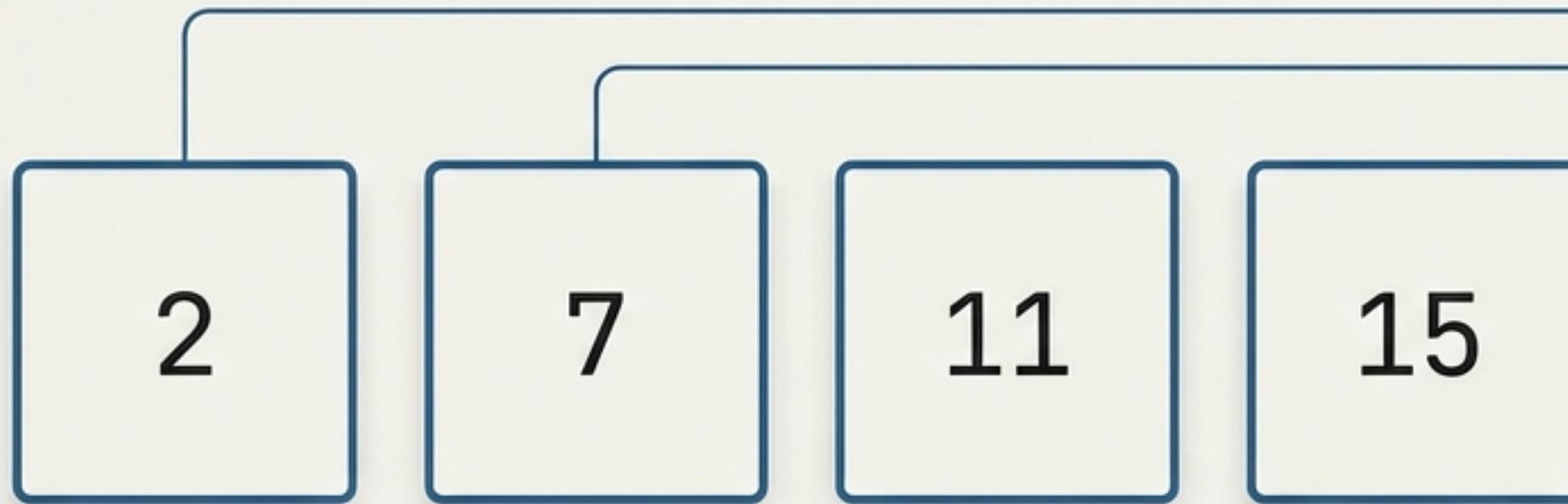


Two Sum: The Art of Algorithmic Trade-offs



Decoding the logic behind the Blind 75's foundational problem.

The Challenge



Input: nums



Target

Constraint: Exactly one solution exists.

Goal: Return indices [0, 1]

Reasoning:

$$\text{nums}[0] + \text{nums}[1] == 9$$

Strategy A: The Unsorted Input

The Complement Concept

Do not scan for a pair sum
($a + b = \text{target}$).

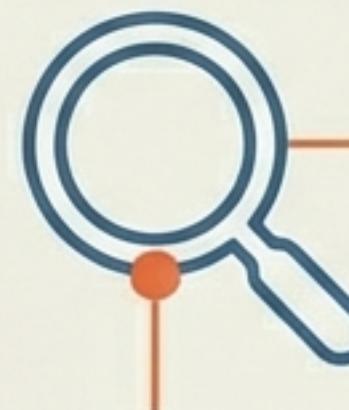
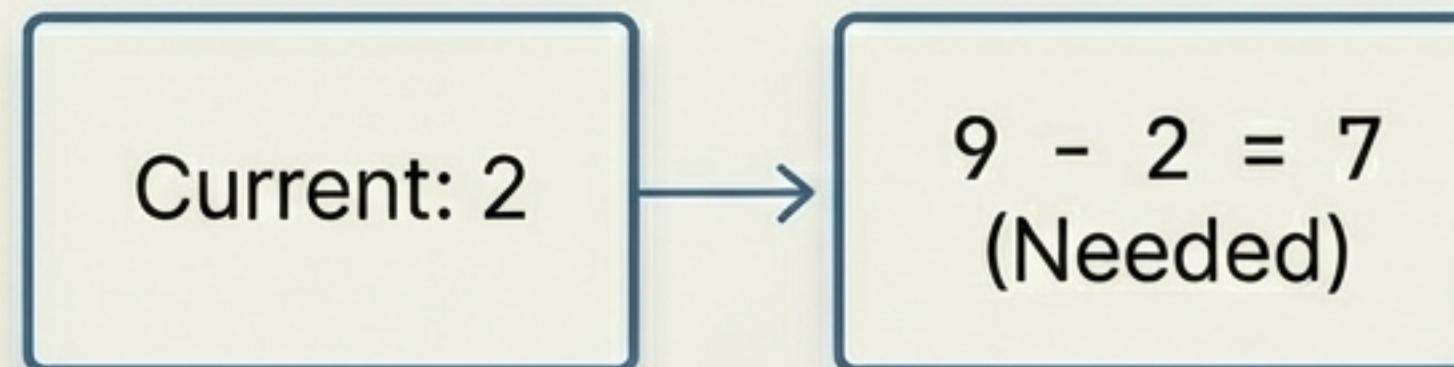
**Iterate once and ask: Have I already seen the difference
($\text{target} - \text{current}$)?**

Key (Value Needed)	Value (Index seen)
	$5 \rightarrow 0$
	$2 \rightarrow 1$
	$-2 \rightarrow 2$
...	...

Tool: Hash Map

Remember visited indices with $O(1)$ lookup time.

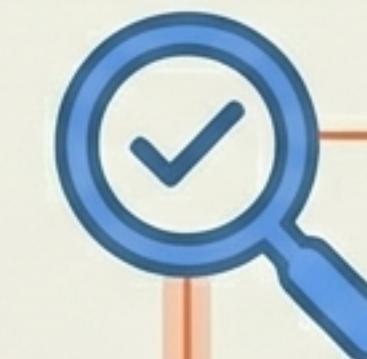
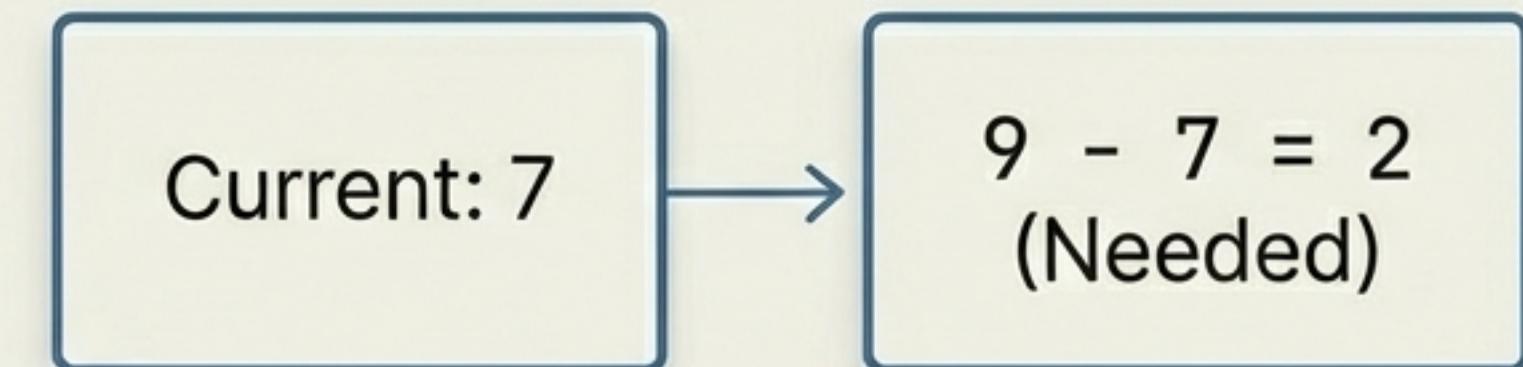
Visualising the One-Pass Hash Map



Not found.
Store {2 : 0}.

Map	
Value	Index

Map is initially empty.



Found! Return
indices [0, 1]

Map	
Value	Index
2	0

Map contains previous
value {2: 0}.

Implementation: Trading Space for Time

```
def twoSum(nums, target):
    prevMap = {} # val : index

    for i, n in enumerate(nums): ←
        diff = target - n
        if diff in prevMap: ←
            return [prevMap[diff], i]
        prevMap[n] = i ←
```

Access index and
value simultaneously

Constant time lookup
 $O(1)$

Build memory bank

Analysis: The Cost of Speed

Time Complexity

O(N)

One pass through the list. Map lookups are constant time.

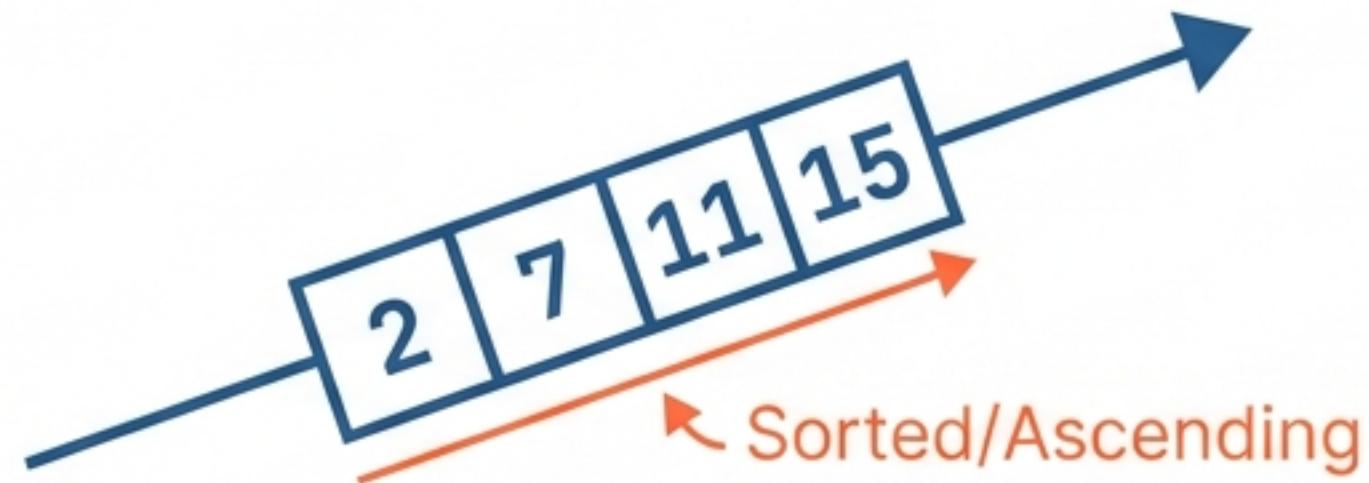
Space Complexity

O(N)

Worst case: store nearly every element in the Map.

Optimal for unsorted data where speed > memory.

The Twist: Sorted Input



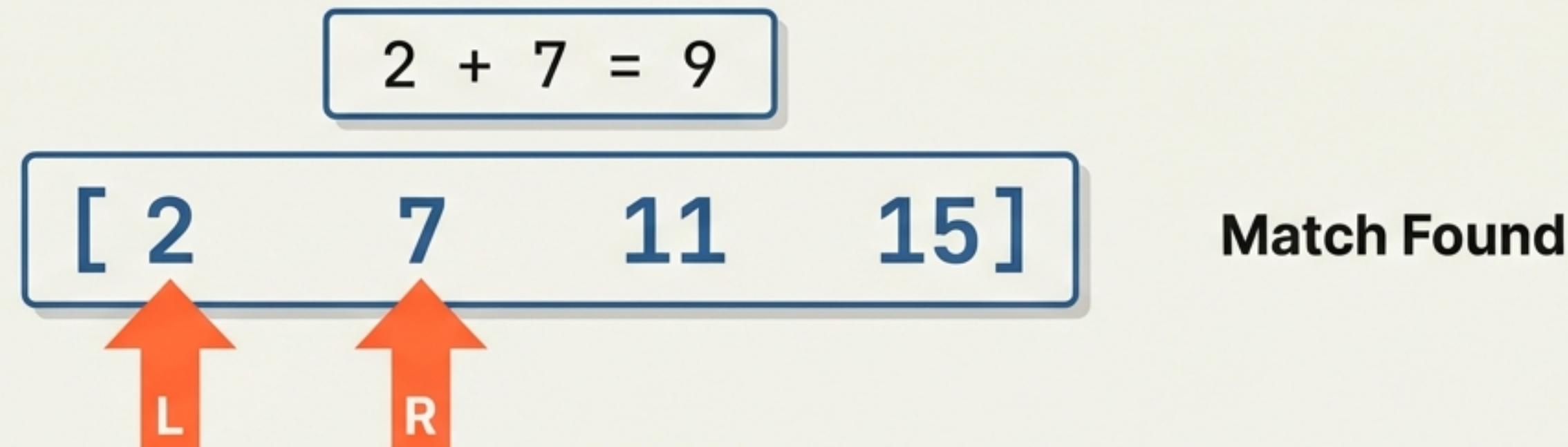
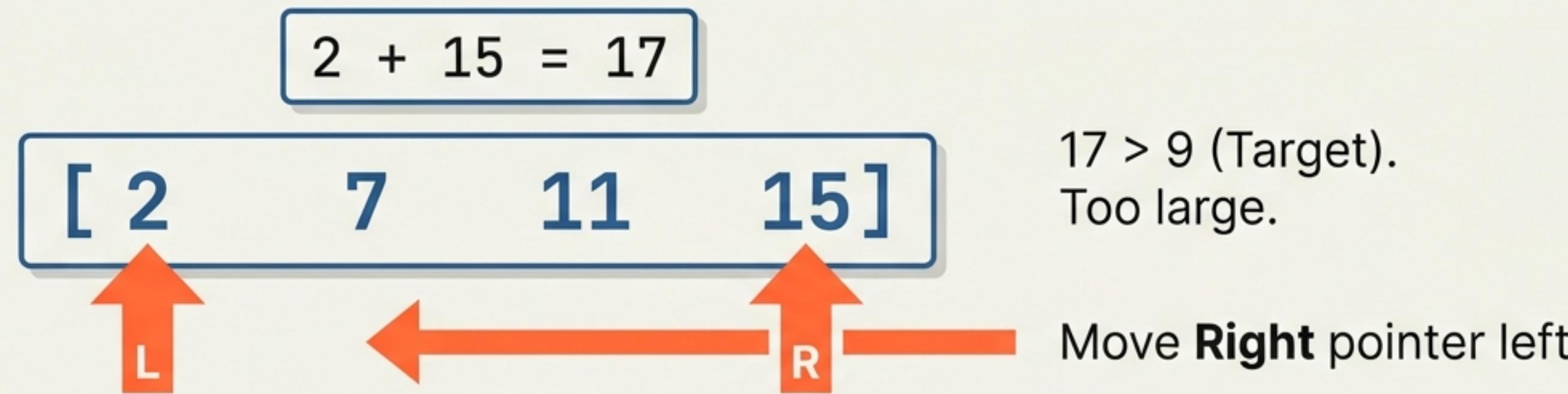
Constraint Change: Input is guaranteed sorted.

Sorted Data + **Directionality** = **No Hash Map Needed**



Tool: Two Pointers

Visualising the Pincer Movement



Implementation: The Constant Space Solution in Heavy Helvetica Now Display with tight tracking

```
l, r = 0, len(nums) - 1  
  
while l < r:  
    curSum = nums[l] + nums[r]  
    if curSum > target:  
        r -= 1  
    elif curSum < target:  
        l += 1  
    else:  
        return [l, r]
```

- Termination condition
- Sum too big?
Decrease right.
- Sum too small?
Increase left.

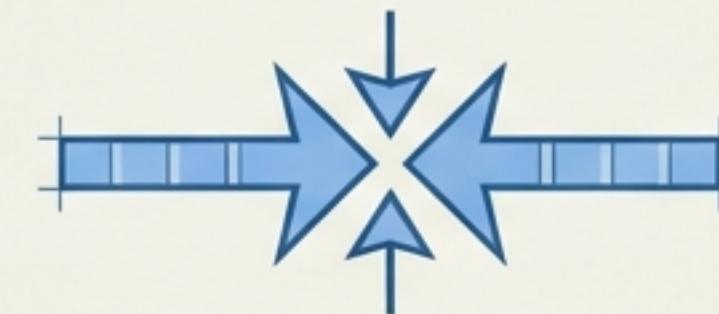
The Verdict: Choosing Your Tool

Unsorted Data



Approach	Hash Map (One Pass)
Time	$O(N)$
Space	$O(N)$
Logic	target - current

Sorted Data



Approach	Two Pointers
Time	$O(N)$
Space	$O(1)$
Logic	Converge from edges

Always check input constraints before selecting your data structure.