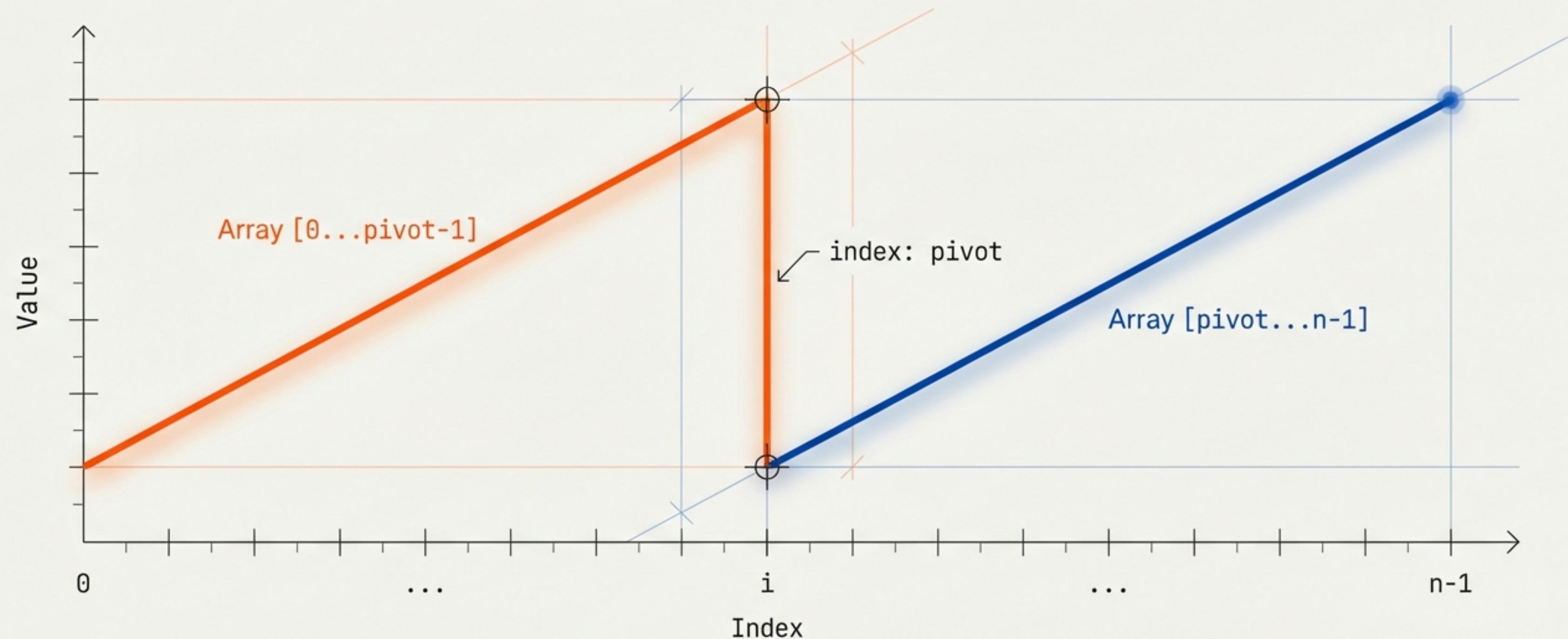


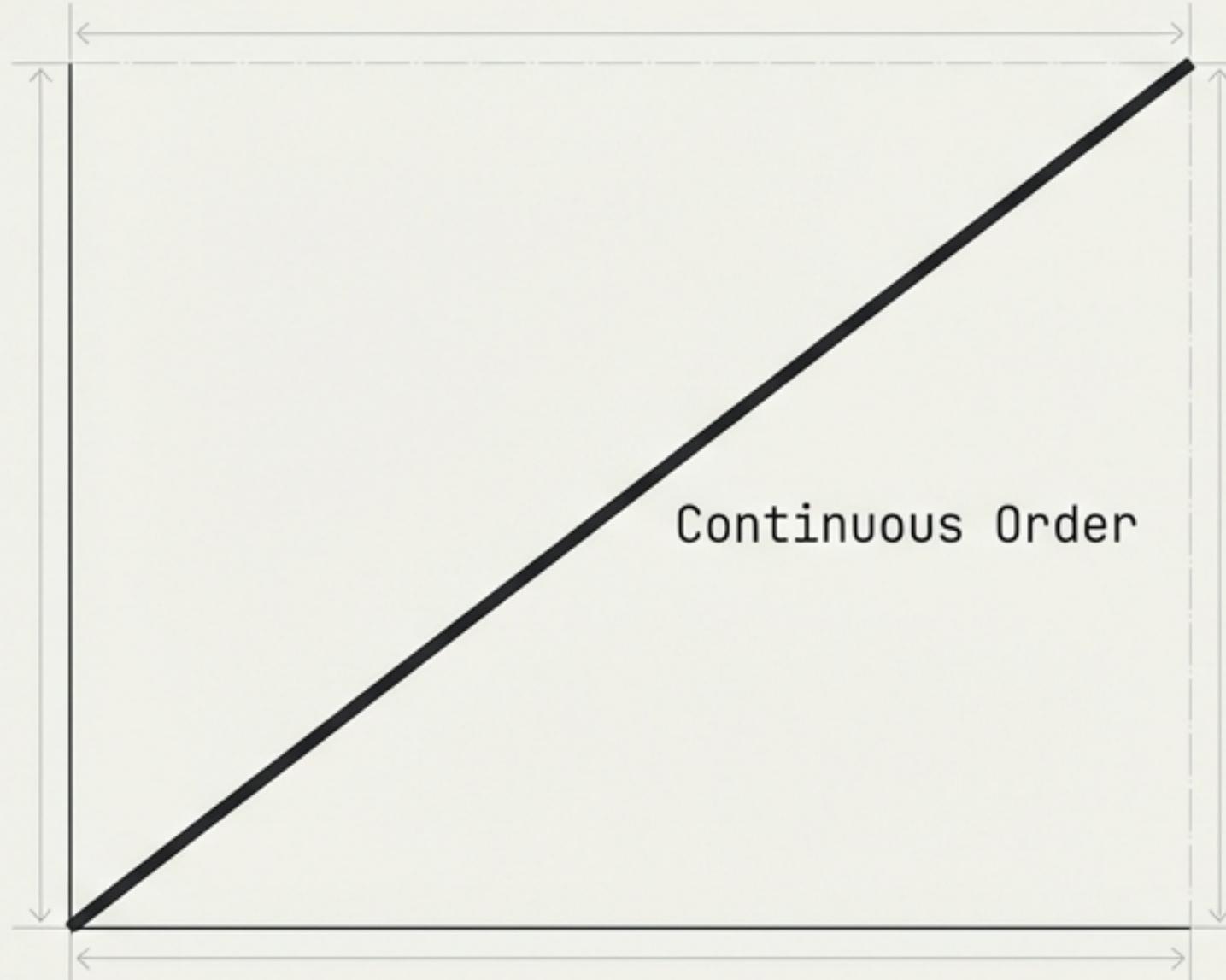
Search in Rotated Sorted Array

Deriving $O(\log n)$ Efficiency through Visual Intuition

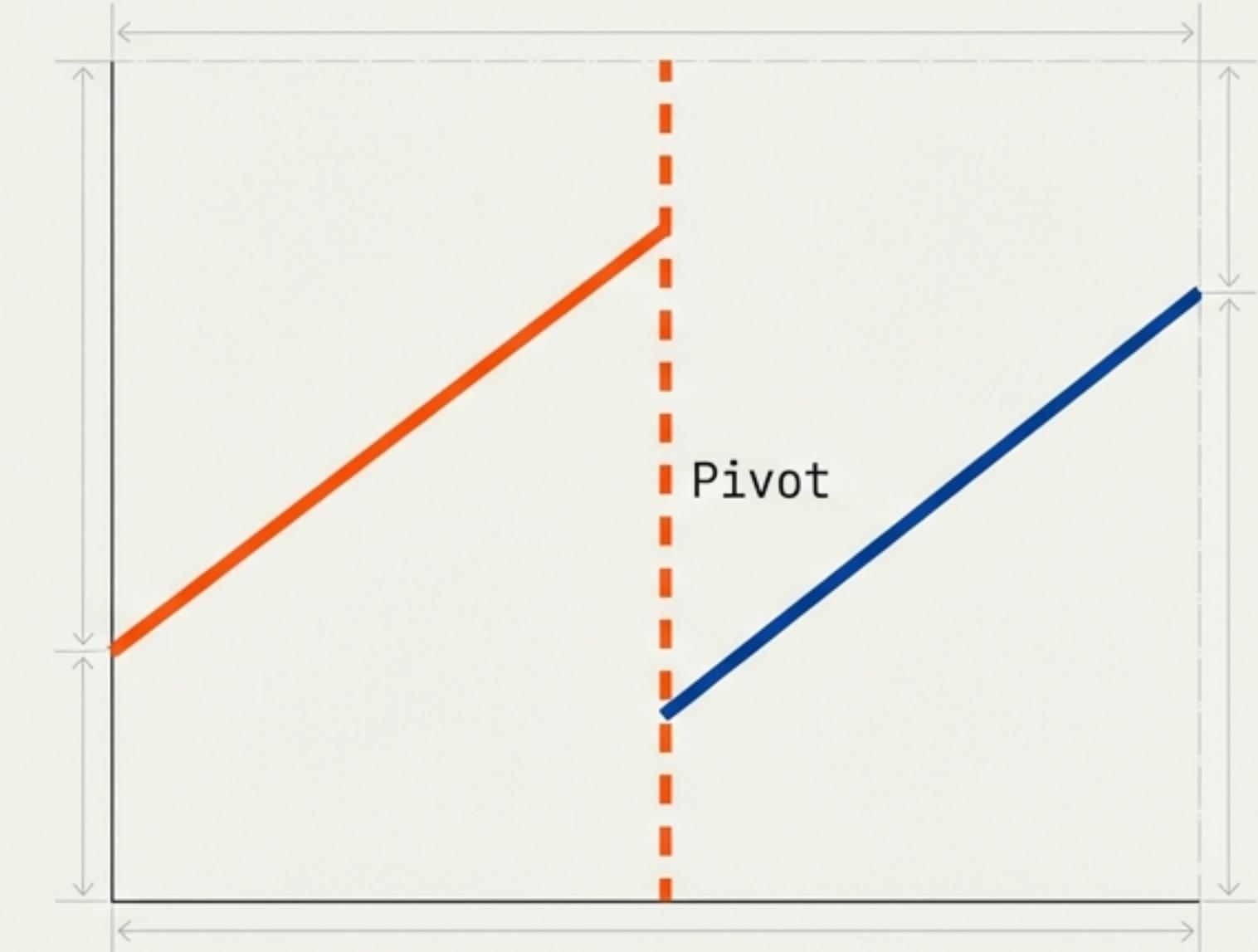


The Geometry of Rotation

Standard Sorted Array

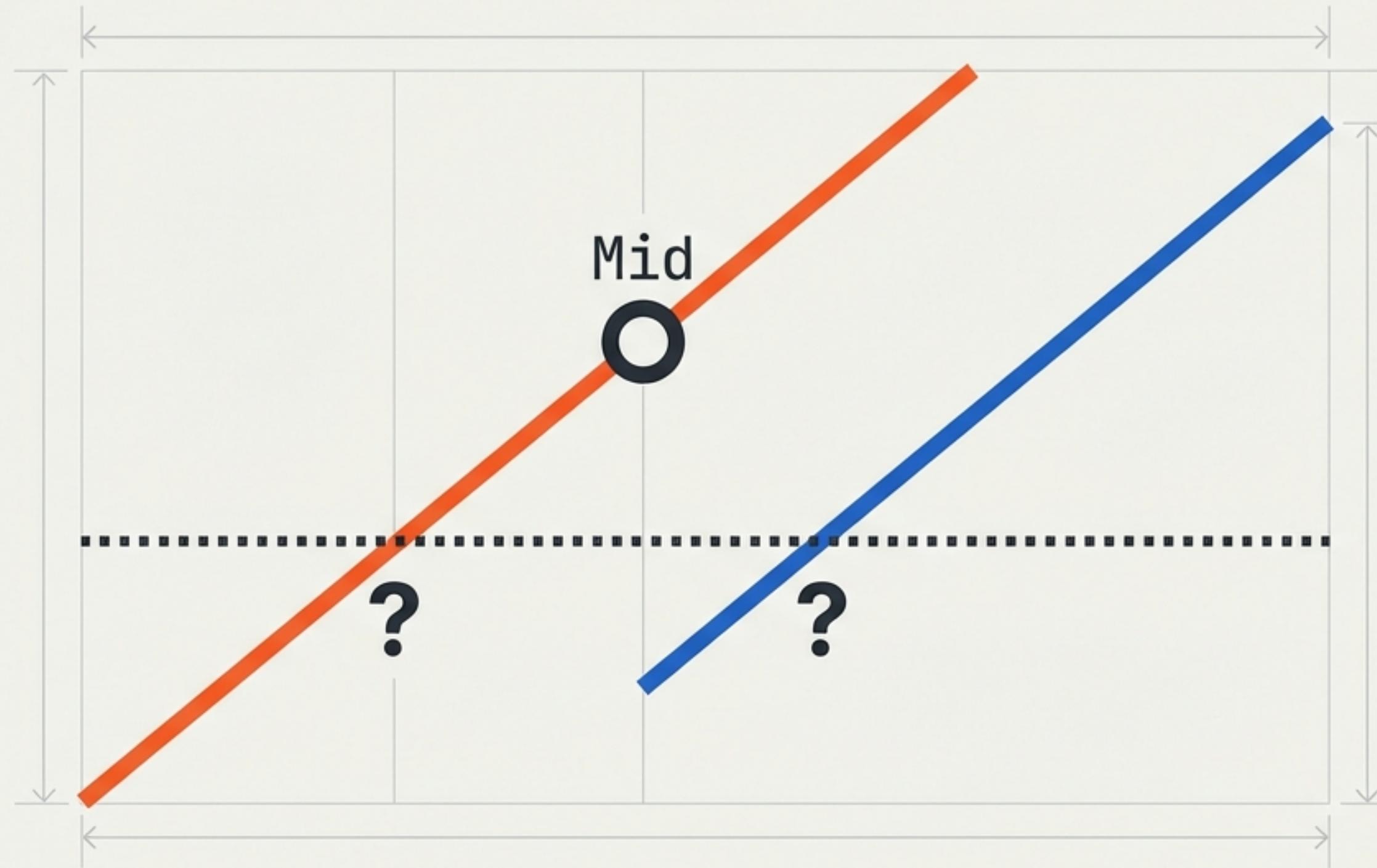


Rotated Sorted Array



We treat the array as a graph. Order exists, but it is discontinuous.

The Binary Search Dilemma



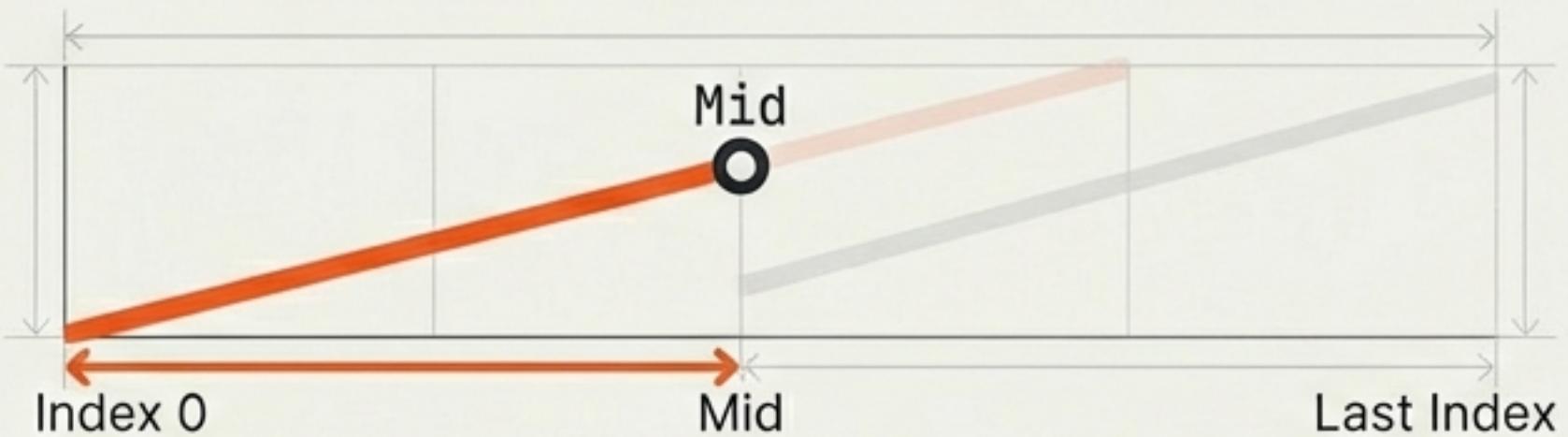
Ambiguity:
If Target < Mid...

Is it here? (**Left**)
Or here? (**Right**)

The Anchor Strategy

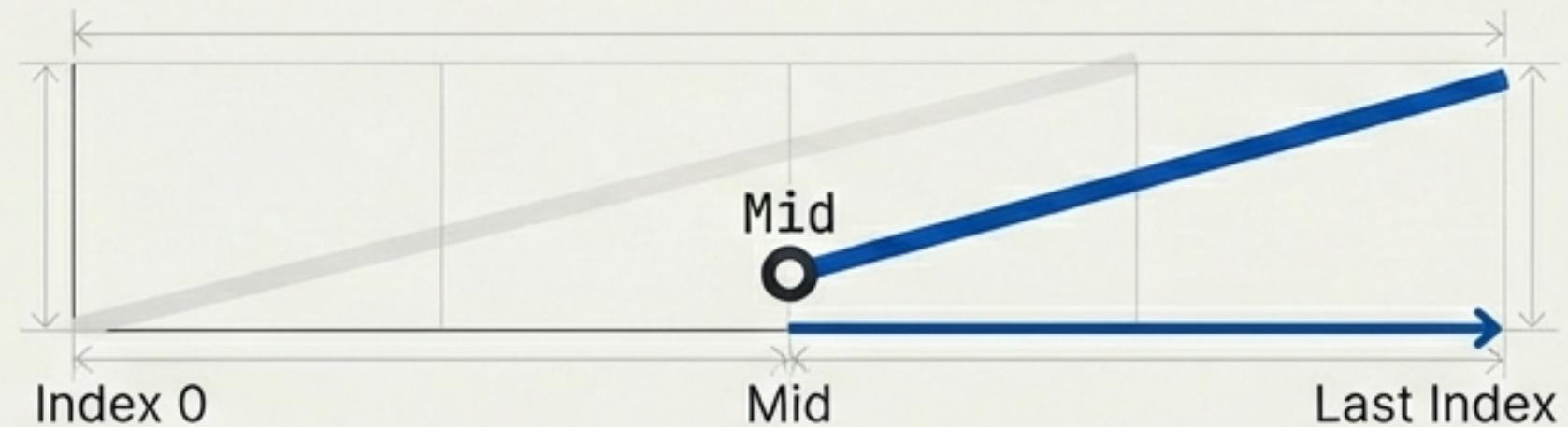
Identifying the sorted half

State A: Left Side Sorted



`if nums[mid] >= nums[left]`

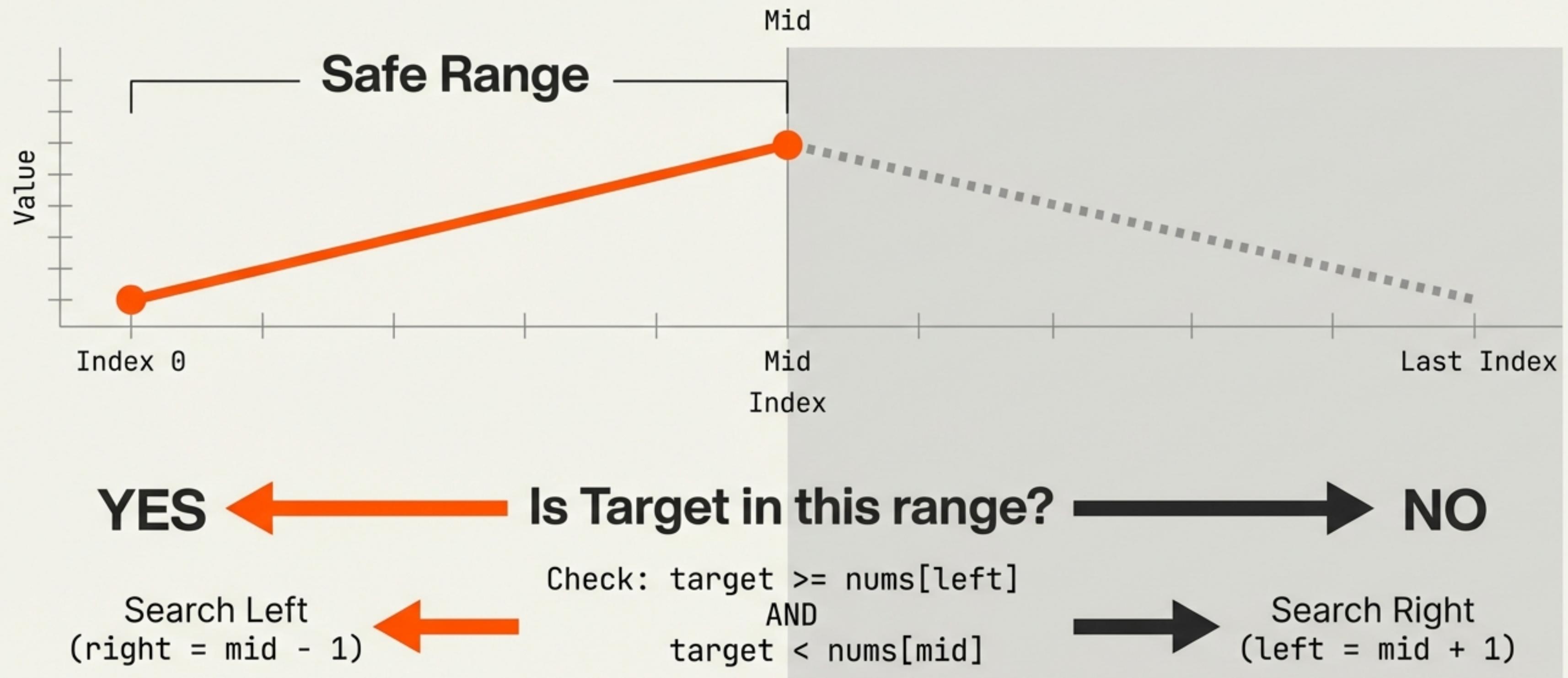
State B: Right Side Sorted



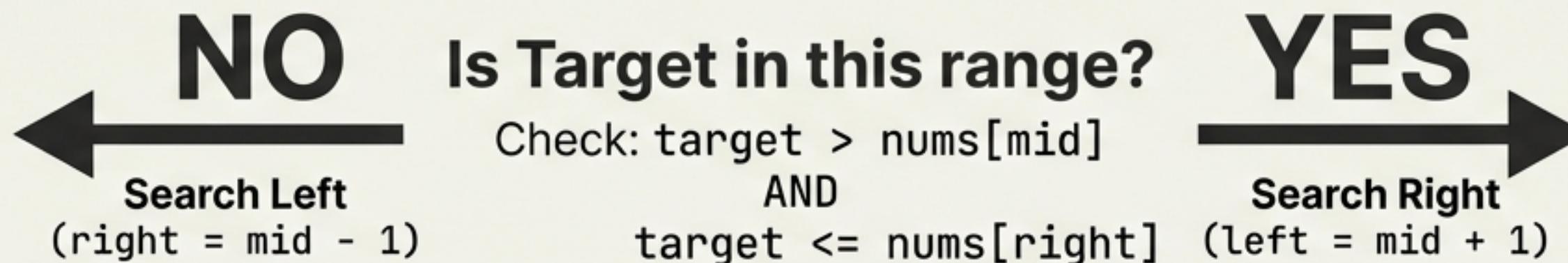
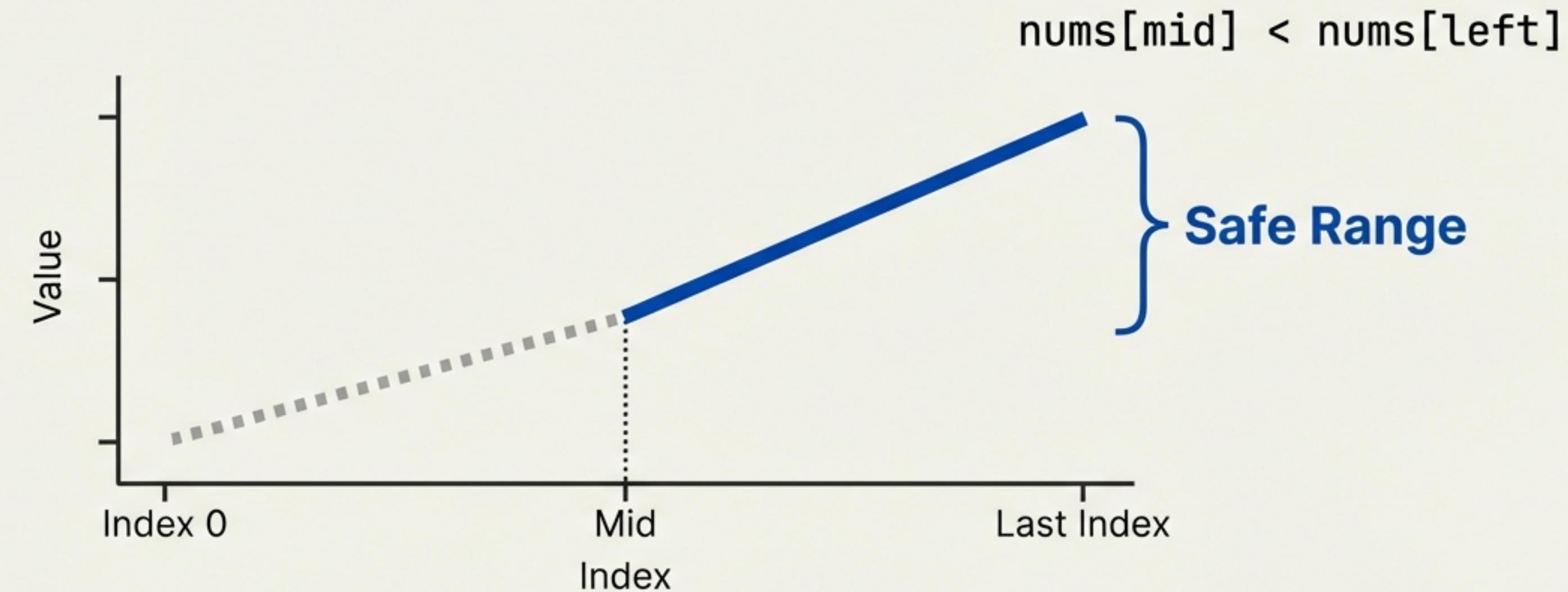
`else`

Case Study 1: Left Sorted Portion

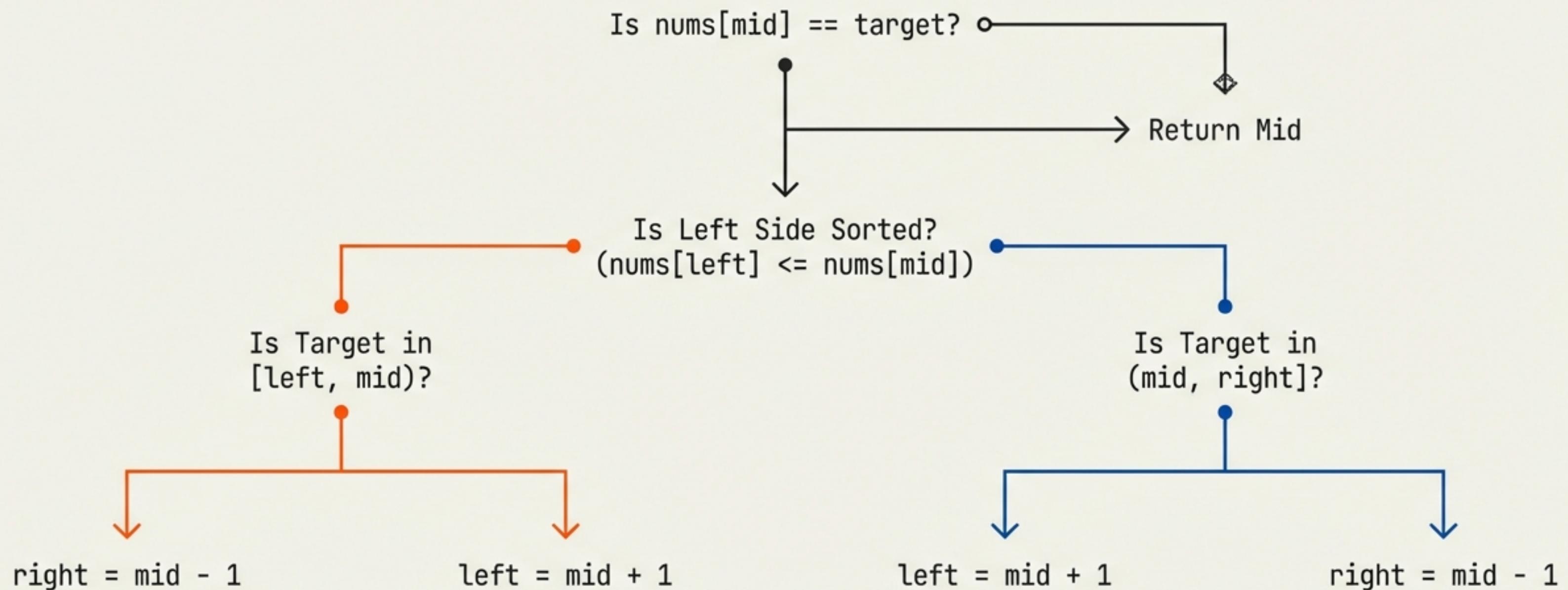
$\text{nums}[\text{mid}] \geq \text{nums}[\text{left}]$



Case Study 2: Right Sorted Portion



The Decision Algorithm



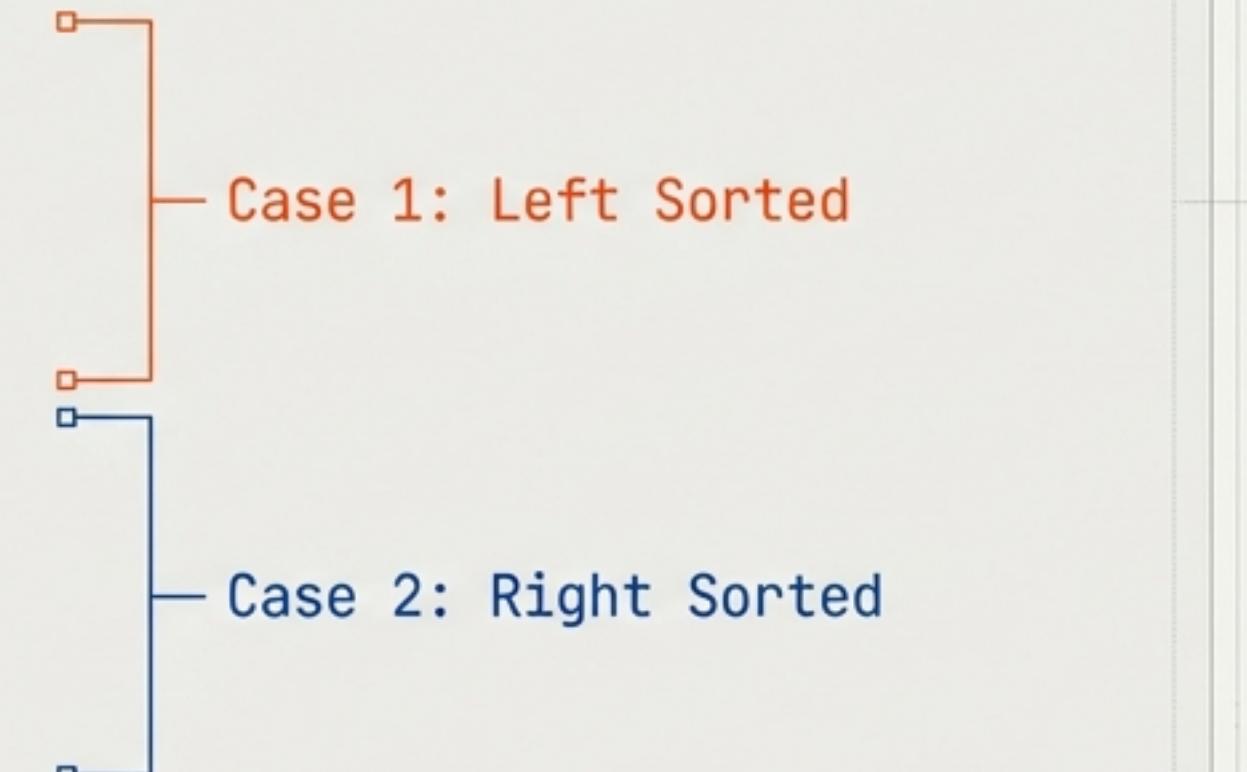
Implementation (Python)

```
def search(nums: List[int], target: int) -> int:
    l, r = 0, len(nums) - 1

    while l <= r:
        mid = (l + r) // 2
        if target == nums[mid]:
            return mid

        # Check for Left Sorted Anchor
        if nums[l] <= nums[mid]:
            if nums[l] <= target < nums[mid]:
                r = mid - 1
            else:
                l = mid + 1
        # Otherwise, Right Sorted Anchor
        else:
            if nums[mid] < target <= nums[r]:
                l = mid + 1
            else:
                r = mid - 1

    return -1
```



Complexity & Edge Cases

Time Complexity

$O(\log n)$

Reason: We eliminate 50% of the search space at every step, regardless of rotation.



Space Complexity

$O(1)$

Reason: Iterative pointers (l , r , mid) only. No recursion stack.

The key is not finding the pivot, but using the sorted half as an anchor.