

Parallel programming

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I still have no idea



what I'm doing



Definitions

Parallel computing

- › Cooperative
- › Partition computation across different compute engines
- › E.g., PThreads w/Shared mem, but also multi-process on the same machine!

Distributed computing

- › Competitive
 - › Partition computation across different machines
 - › E.g., multiprocess (MPI, MQTT) w/message passing



Why do we need parallel computing?

Increase performance of our machines

› Scale-up

- Solve a "bigger" problem in the same time

› Scale-out (ex: Speedup)

- Solve the same problem «in less time» (or, less power..)



Yes but..

Why (highly) parallel machines...

...and not faster single-core machines?



The answer #1 - Money





The answer #2 – the "hot" one

Moore's law

- › "The number of transistors that we can pack in a given die area doubles every 18 months"

Dennard's scaling

- › "Performance per watt of computing is growing exponentially at roughly the same rate"



The answer #2 – the "hot" one

- › SoC design paradigm



- › Gordon Moore
 - His law is still valid, but...

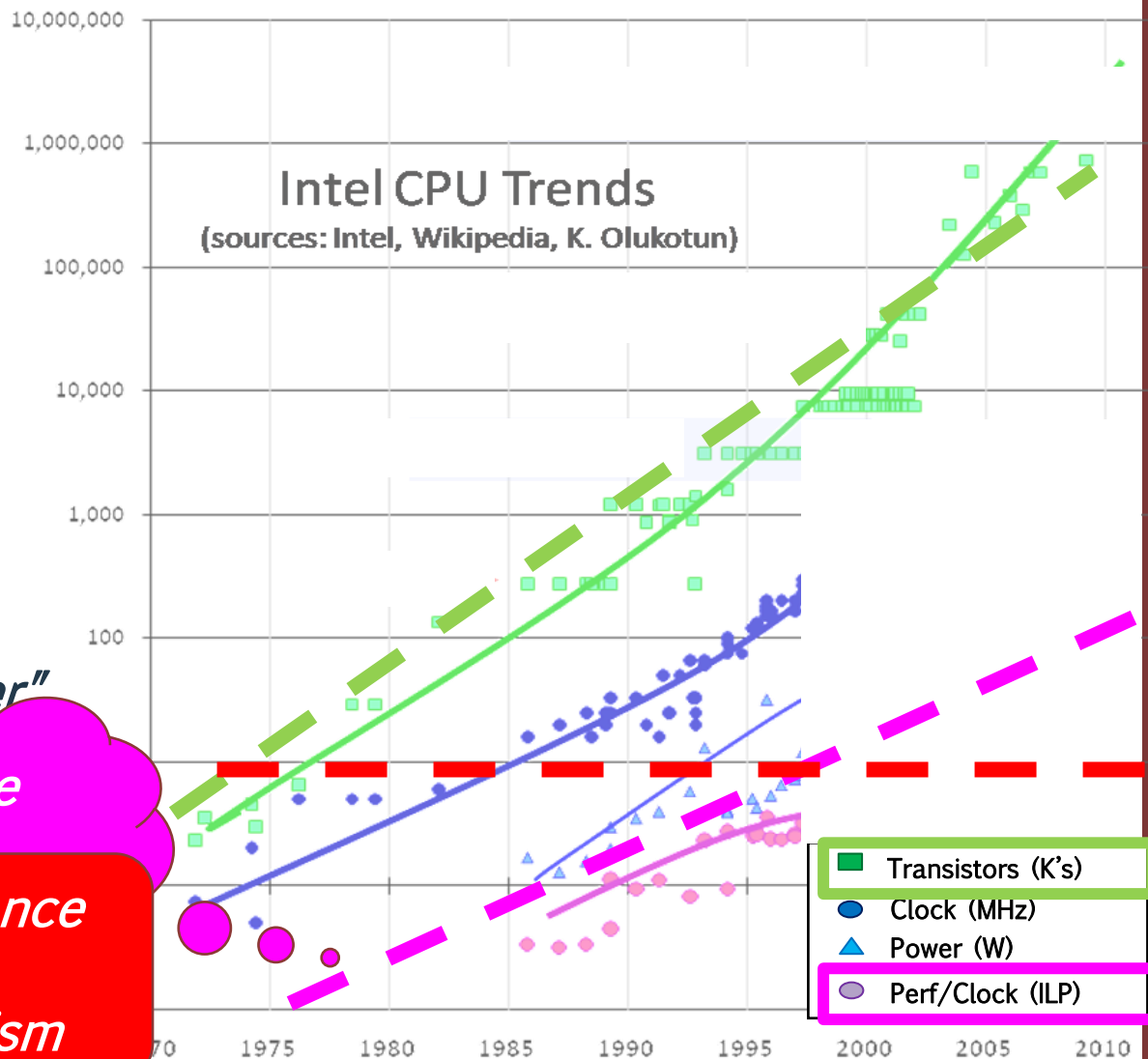
- › *"The free lunch is over"*

– Herb

Performance

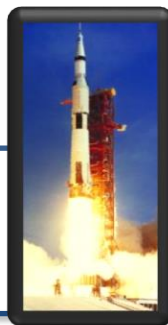
Performance

parallelism

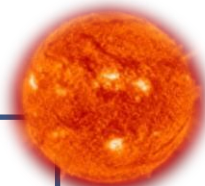




In other words...



Surface of the sun



Rocket nozzle

Nuclear Reactor

Modern computers

First PCs

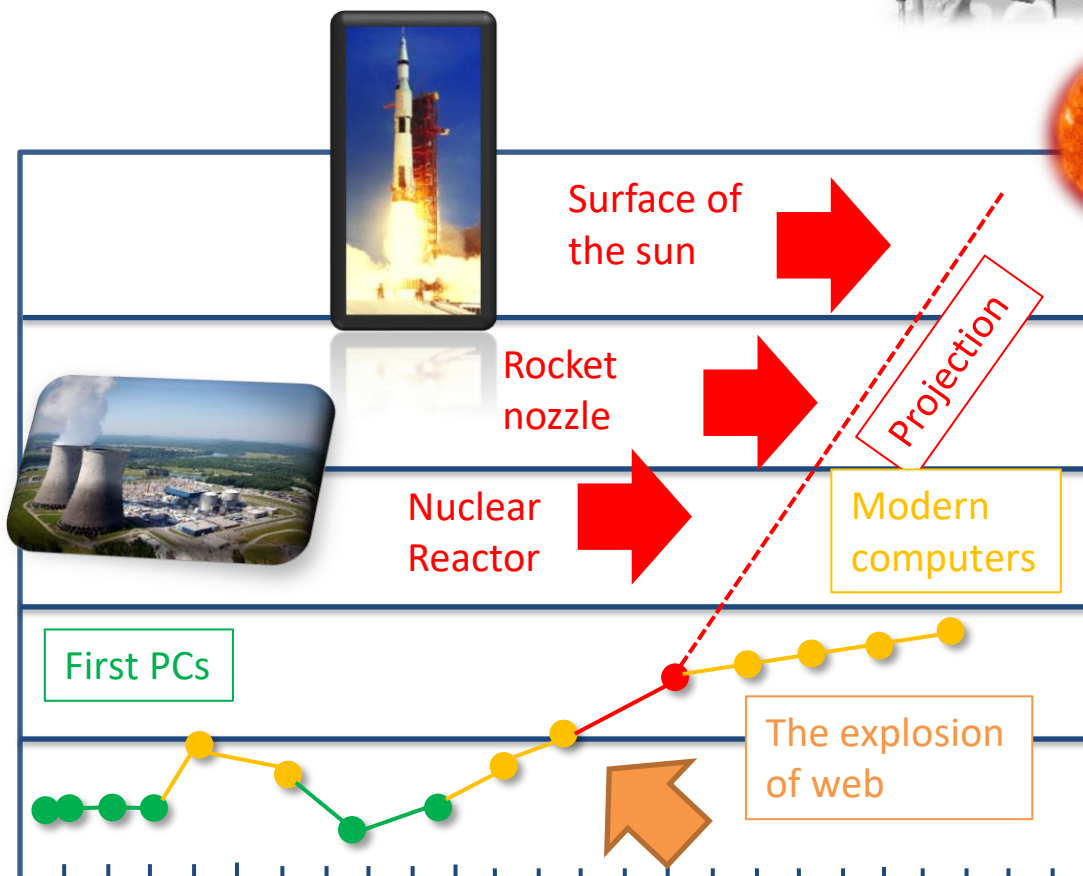
The explosion of web

Hot plate



Summer temperature

1970 1980 1990 2000 2010 2020





Instead of going faster..

- › ..(go faster but through) parallelism!

Problem #1

- › New computer architectures
- › At least, three architectural templates

Problem #2

- › Need to efficiently program them
- › HPC already has this problem!

The problem

- › Programmers **must** know a bit of the architecture!
- › To make parallelization effective
- › "Let's run this on a GPU. It certainly goes faster" (cit.)

The Big problem

- › Effectively programming in parallel is difficult

Brian Kernighan (1942-)

- *Researcher, theory of informatics*
- *Co-authored UNIX and AWK*
- *Wrote "The C Programming Language" book*

"Everyone knows that debugging is twice as hard as writing a program in the first place.

So if you're as clever as you can be when you write it, how will you ever debug it?"





Amdahl's law

- › A sequential program that takes 100 sec to exec
- › Only 95% can run in parallel (it's a lot)
- › And.. you are an extremely good programmer, and you have a machine with 1billion cores, so that part takes 0 sec
- › So,

$$T_{par} = 100_{sec} - 95_{sec} = 5_{sec}$$

$$Speedup = \frac{100_{sec}}{5_{sec}} = 20x$$

...20x, on one billion cores!!!

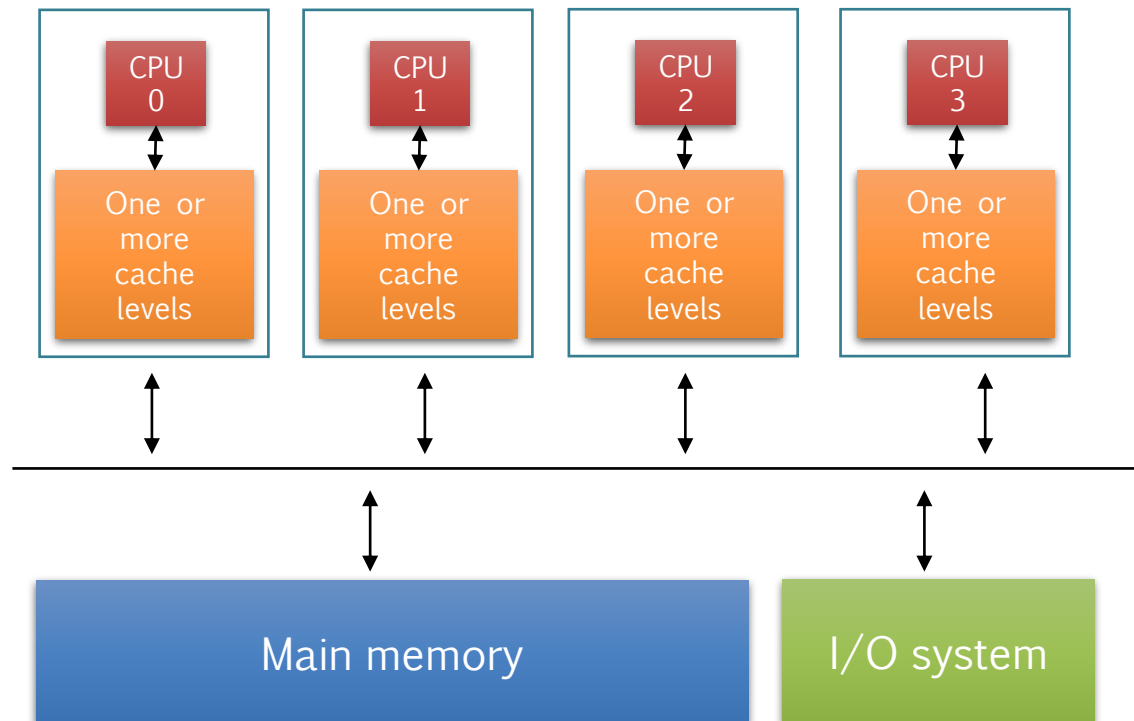




Symmetric multi-processing

- › Memory: centralized with bus interconnect, I/O
- › Typically, **multi**-core (sub)systems
 - Examples: Sun Enterprise 6000, SGI Challenge, Intel (this laptop)

Can be 1 bus, N
busses, or any
network

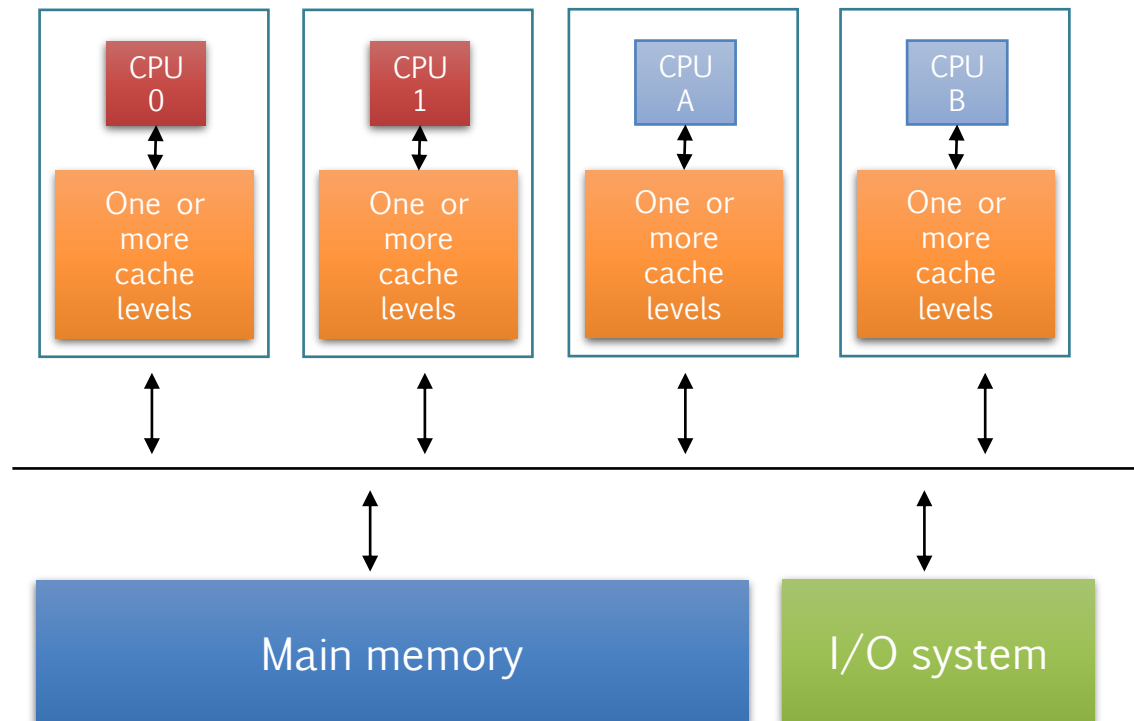




Asymmetric multi-processing

- › Memory: centralized with uniform access time (UMA) and bus interconnect, I/O
- › Typically, multi-core (sub)systems
 - Examples: ARM Big.LITTLE, NVIDIA Tegra X2 (Drive PX)

Can be 1 bus, N
busses, or any
network



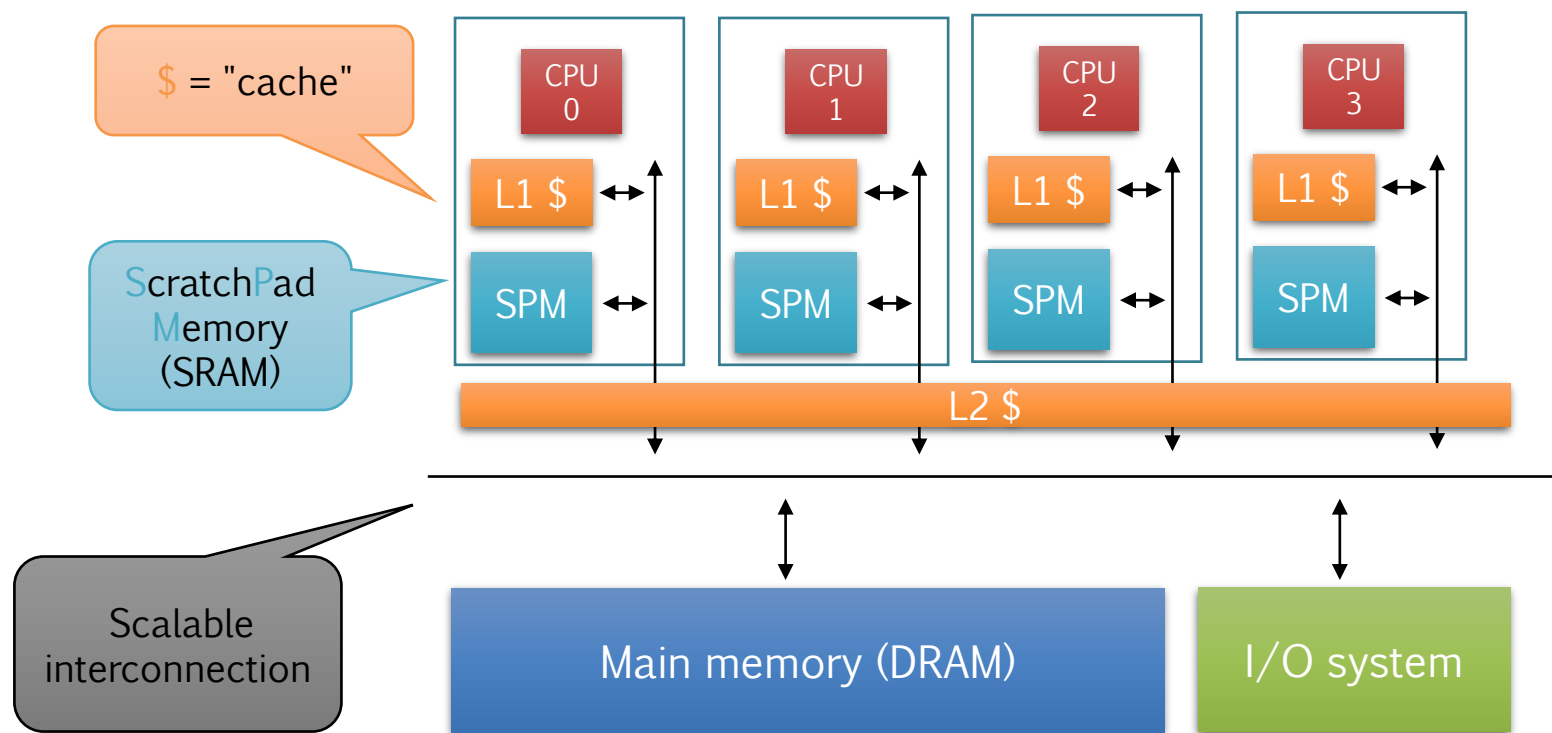


SMP – distributed shared memory

› Non-Uniform Access Time - **NUMA**

› **Scalable** interconnect

- Typically, **many** cores
- Examples: embedded accelerators, GPUs





UMA vs. NUMA

> Shared

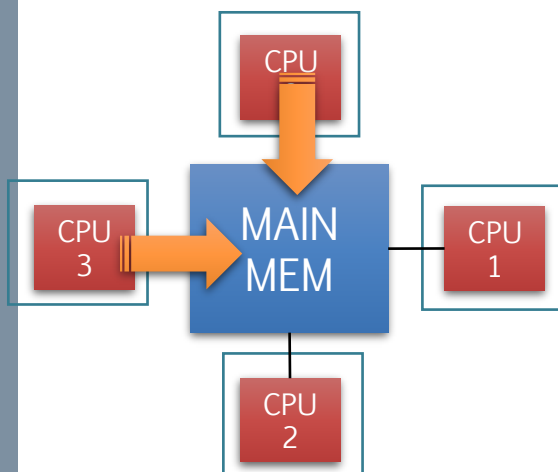
- (

> Uniform

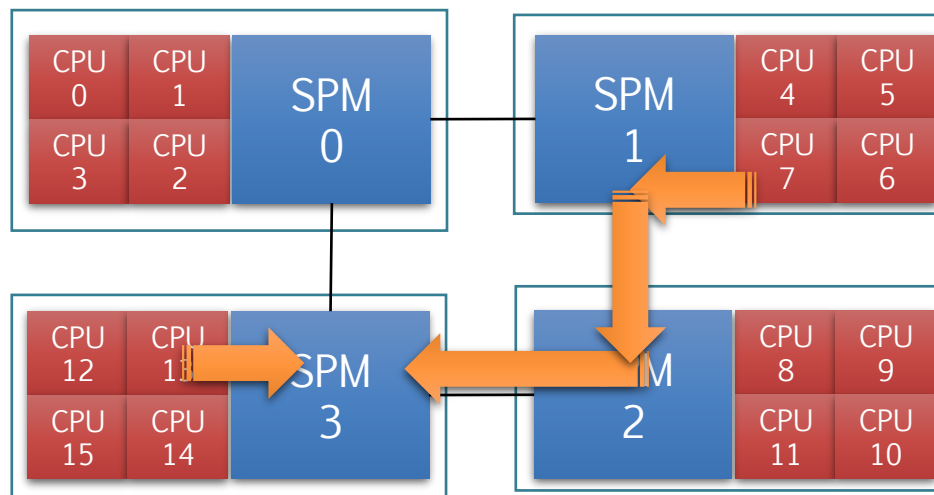
- [

	MEM0	MEM1	MEM2	MEM3
CPU0...3	0 clock	10 clock	20 clock	10 clock
CPU4...7	10 clock	0 clock	10 clock	20 clock
CPU8...11	20 clock	10 clock	0 clock	10 clock
CPU12..15	10 clock	20 clock	10 clock	0 clock

UMA



NUMA





Programming abstractions

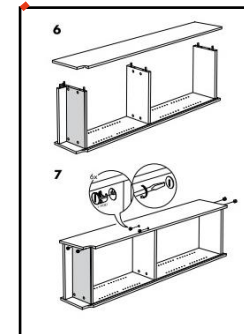
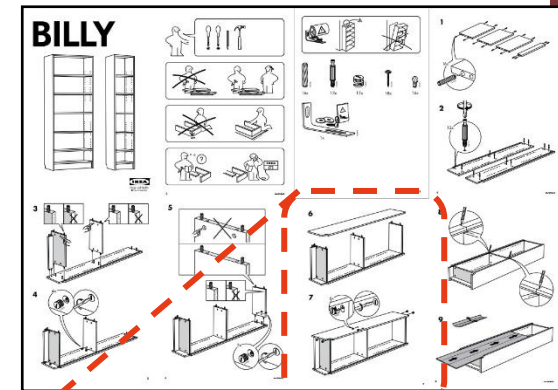


What is...

- › ..a core?
- › ...a program?
- › ...a process?
- › ...a thread?
- › ..a task?

What is...

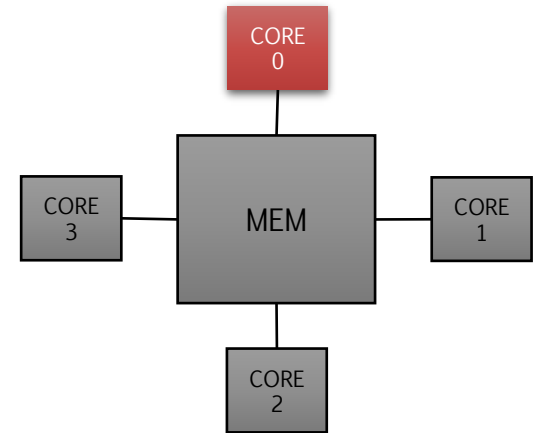
- › ..a core?
 - An electronic circuit to execute instruction (=> programs)
- › ...a program?
 - The implementation of an algorithm
- › ...a process?
 - A program that is executing
- › ...a thread?
 - A unit of execution (of a process)
- › ..a task?
 - A unit of work (of a program)



What is...

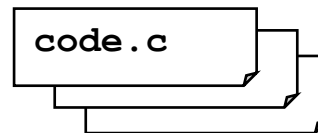
> ..a core?

- An electronic circuit to execute instruction (=> programs)



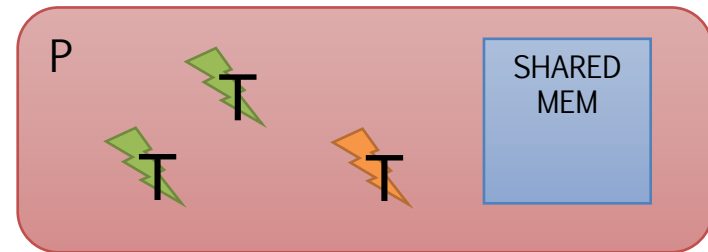
> ...a program?

- The implementation of an algorithm



> ...a process?

- A program that is executing



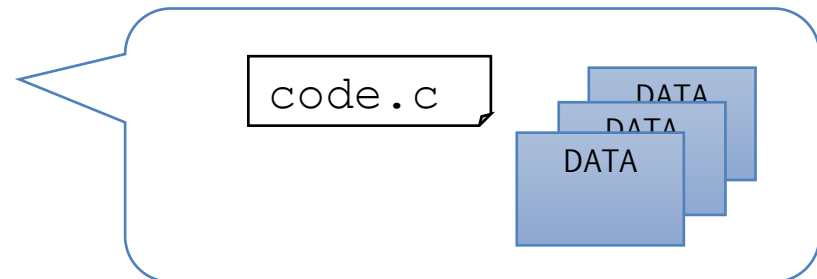
> ...a thread?

- A unit of execution (of a process)



> ..a task?

- A unit of work (of a program)





(Simplest) threading model

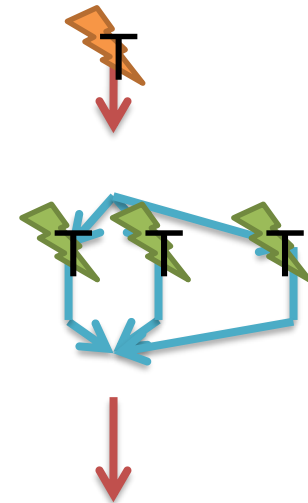
Fork-join execution model

- › The main, single thread spawns a team of **Slave** threads (here, NTHREADS = 3)
- › They all perform computation in parallel
- › At the end, they are joined one by one (aka: barrier)

```
int main()
{
    int err;
    pthread_t mythreads[NTHREADS];
    for (int i=0; i<NTHREADS; i++)
        err = pthread_create (&mythreads[i], // >= PTHREAD_CREATE_JOINABLE
                               &myattr,
                               my_thread_routine,
                               NULL);
    // Here, the main thread can do other things

    for (int i=0; i<NTHREADS; i++)
        pthread_join(mythreads[i], &return_val); // <= JOIN
}
```

Let's see
this in
action





Master-slave threading model

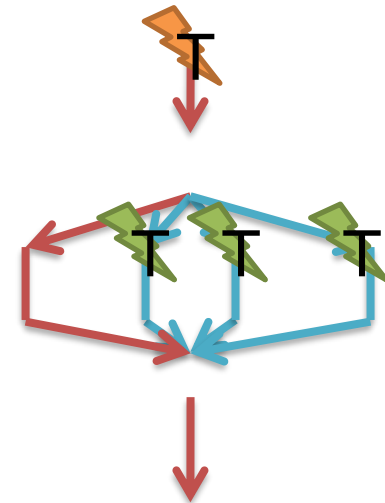
Fork-join execution model

- › The main, **Master** thread spawns a team of **Slave** threads (here, NTHREADS = 3)
- › They all perform computation in parallel
- › At the end, they are joined one by one (aka: barrier)

```
int main()
{
    int err;
    pthread_t mythreads[NTHREADS];
    for (int i=0; i<NTHREADS; i++)
        err = pthread_create (&mythreads[i], // ==> FORK
                               &myattr,
                               my_pthread_fn,
                               NULL);

    // Here, the main thread can do other stuff!
    other_fn();

    for (int i=0; i<NTHREADS; i++)
        pthread_join(mythreads[i], &returnvalue); // <== JOIN
}
```





Work partitioning

Several models, here to cite a few

- › Data parallelism (see also GPGPUS)
 - We're getting there, don't worry...
- › Reduction
- › Task parallelism (aka: work queue)
- › Offloading



Data parallelism

(Aka: data decomposition, loop decomposition, SPMD, SIMD*...)

Parallel threads execute the same operation(s) on multiple data

- › Data is typically an array, a matrix (image)....
 - Note: you typically map the iteration id of the loop to the data index
- › **Partitioning strategy** defines how many iterations (**chunk**) every thread will perform
 - From 1 iteration, to loop size



* Single Program, Multiple Data; Single Instruction, Multiple Data



Exercise

Let's
code!

Create an array of N elements

- › Put inside each array element its index, multiplied by '2'
- › `arr[0] = 0; arr[1] = 2; arr[2] = 4; ...and so on..`

Now, do it in parallel with a team of T PThreads

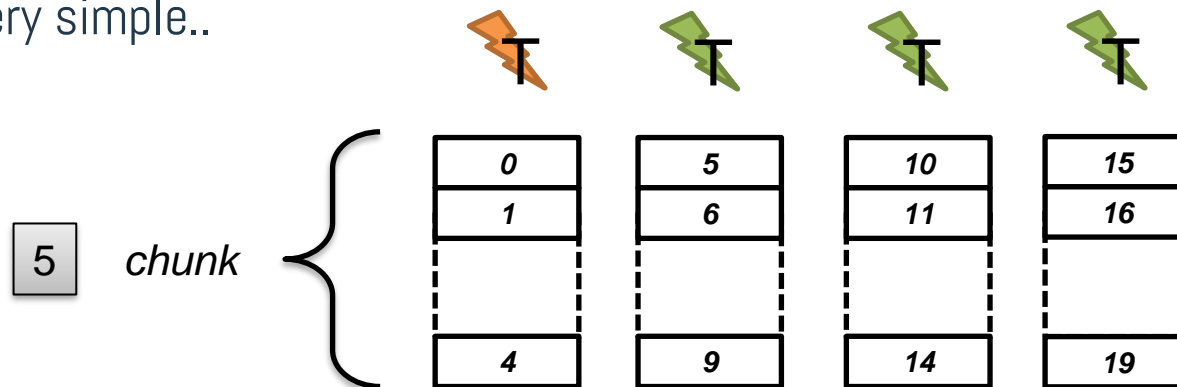
- › Assume N is a multiple of T
- › "Decompose" the `for` construct, so that every thread manages (chunk size is) N/T iterations



Loop partitioning among threads

Let's
code!

- › Case #1: N multiple of T
 - Say, $N = 20, T = 4$
- › $chunk = \#iterations$ for each thread
- › Very simple..



$$chunk = \frac{N}{T} ;$$

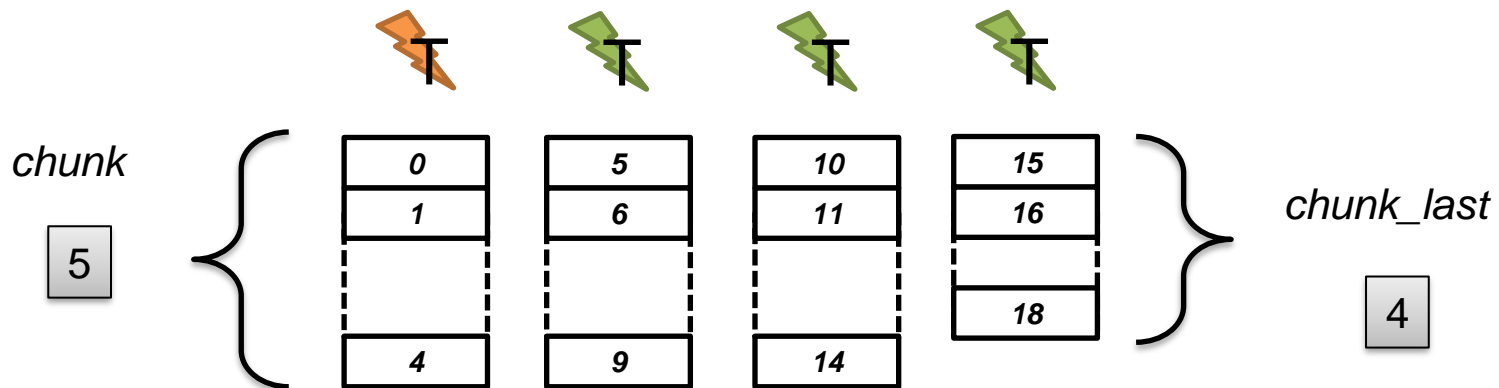
$$i_{start} = thread_{ID} * chunk; \quad i_{end} = i_{start} + chunk$$



Loop partitioning among threads

Let's
code!

- › Case #2: N **not** multiple of T
 - Say, $N = 19, T = 4$
- › $chunk = \#iterations$ for each thread (but last)
 - Last thread has less! ($chunk_{last}$)



$$chunk = \frac{N}{T} + 1; \quad chunk_{last} = N \% chunk$$

$$i_{start} = thread_{ID} * chunk; \quad i_{end} = \begin{cases} i_{start} + chunk & \text{if not last thread} \\ i_{start} + chunk_{last} & \text{if last thread} \end{cases}$$



"Last thread"

Let's
code!

- › Unfortunately, we don't know which thread will be "last" in time
- › But...we don't actually care the order in which iterations are executed
 - If there are not dependencies..
 - And..we do know that

$$0 \leq \text{myid} < \text{NUM_THREADS}$$

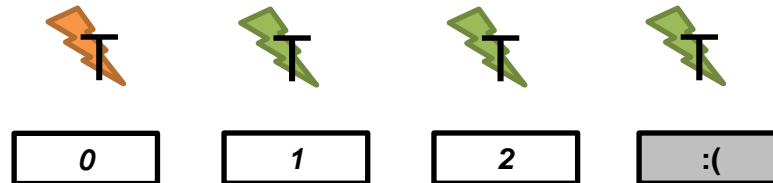
- › We **choose** that last thread as highest number



Loop partitioning among threads

Let's
code!

- › Case #3: N **smaller than** T
 - Say, $N = 3$, $T = 4$
- › Similar to Case #1
 - With some adjustments...





Let's put them together!

Let's
code!

› Case #1 (N multiple of T)

$$chunk = \frac{N}{T} \quad i_{start} = thread_{ID} * chunk; \quad i_{end} = i_{start} + chunk$$

› Case #2 (N not multiple of T)

$$chunk = \frac{N}{T} + 1; \quad chunk_{last} = N \% chunk$$

$$i_{start} = thread_{ID} * chunk; \quad i_{end} = \begin{cases} i_{start} + chunk & \text{if not last thread} \\ i_{start} + chunk_{last} & \text{if last thread} \end{cases}$$



Reduction

wikipedia

*The reduction clause can be used to perform some forms of **recurrence** calculations (involving **mathematically associative and commutative operators**) in parallel. For parallel [...], a **private** copy of each list item is created, one for each implicit task, as if the private clause had been used. [...] The private copy is then initialized as specified above. At the end of the region for which the reduction clause was specified, the original list item is updated by **combining its original value with the final value of each of the private copies**, using the combiner of the specified reduction-identifier.*

E.g., average value of a sequence (array/vector)

- › Create a thread-local copy of a variable
- › Accumulate sums only for the assigned part of the array/vector
- › Then, sum the partial sums (and divide by size)



Exercise

Let's
code!

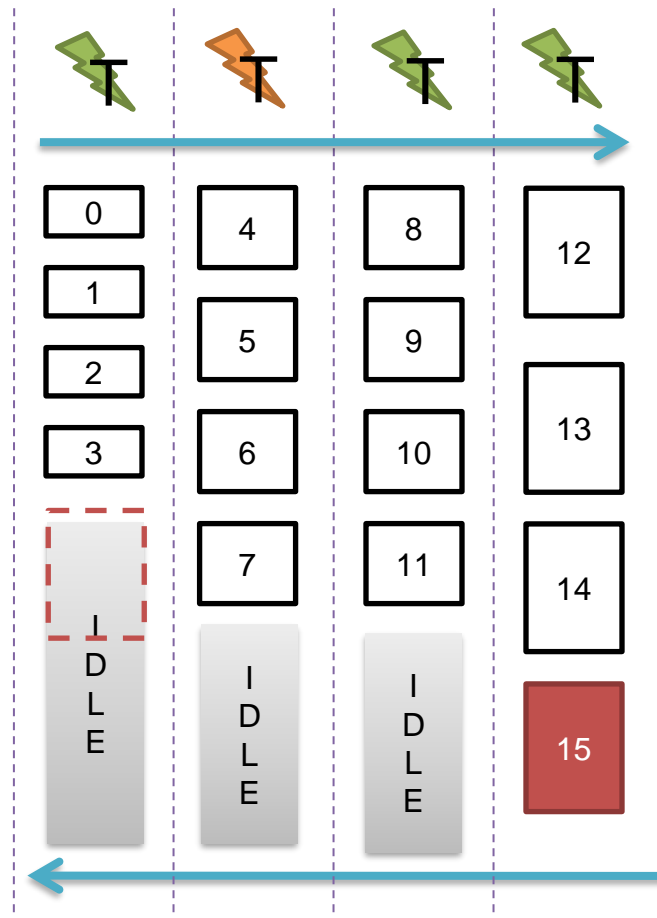
Create an array of N elements

- › ..or a vector
- › Initiate it randomly
- › Now, compute its average value using multi-treading and reduction paradigm



Unbalanced loop partitioning

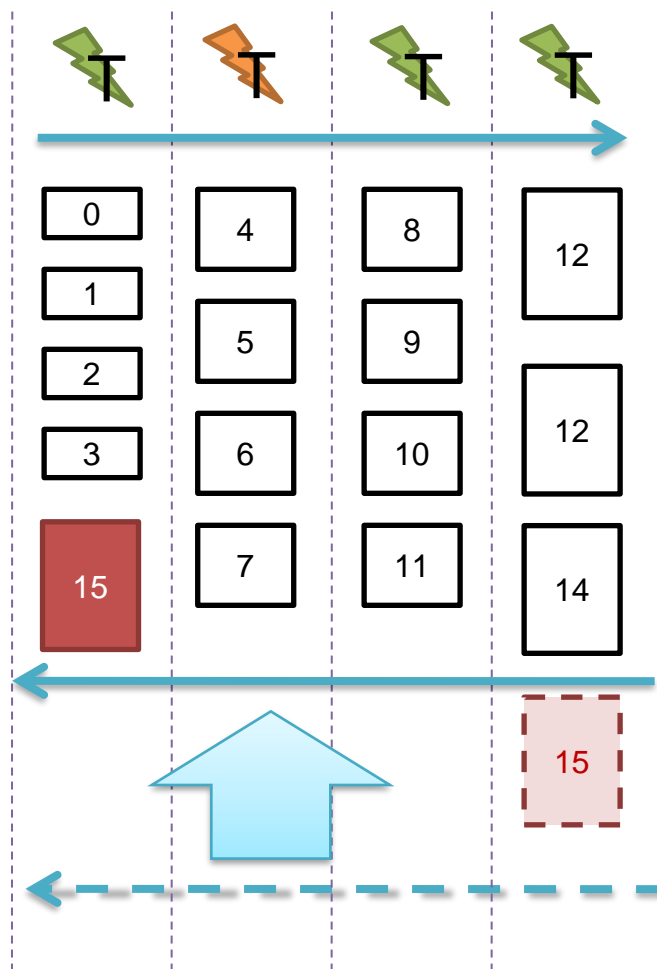
- › So far, we assigned iterations «statically»
 - Might not be effective nor efficient





How can we manage dynamics/irregular workloads?

› We would like something like this..



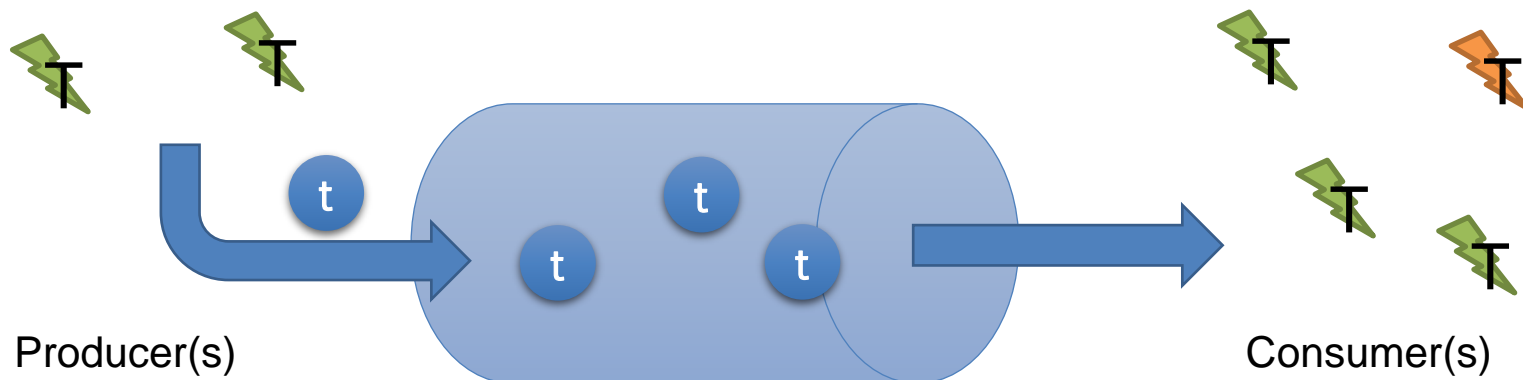


A different parallel paradigm: Tasking

Implements a **producer-consumer paradigm**

Managed by a **task queue**

- › Where units of work (**tasks**)
- › are pushed by threads (**q_push** primitive)
- › and pulled and executed by threads (**q_pop** primitive)



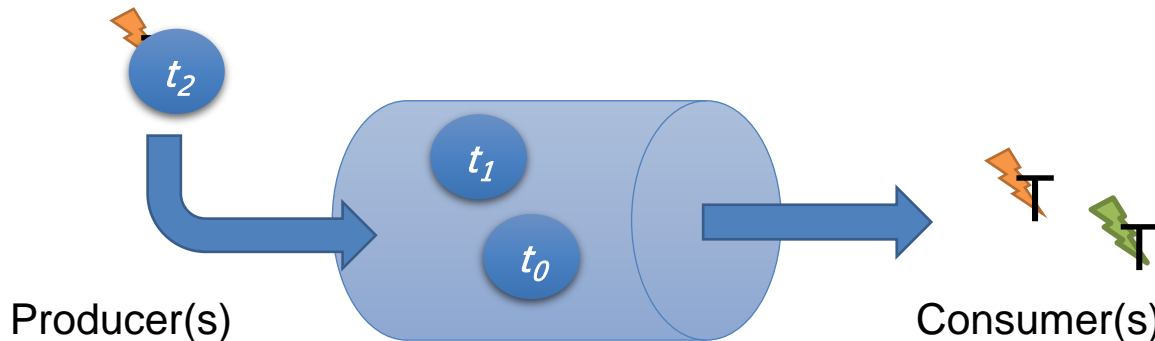


«What» happens «when»?

```
void t0() {  
    // Task 0  
}  
void t1() {  
    // Task 1 pushes t2 in the q  
    q_push(t2());  
}  
void t2() {  
    // Task 2  
}  
  
void *thread_fn(void * args) {  
    // Push t0 and t1  
    q_push(t0());  
    q_push(t1());  
}
```

```
void *other_thread_fn(void * args) {  
    // Pop a task (which one?)  
    t = q_pop();  
    // Execute it  
    t();  
}
```

*pseudo-code





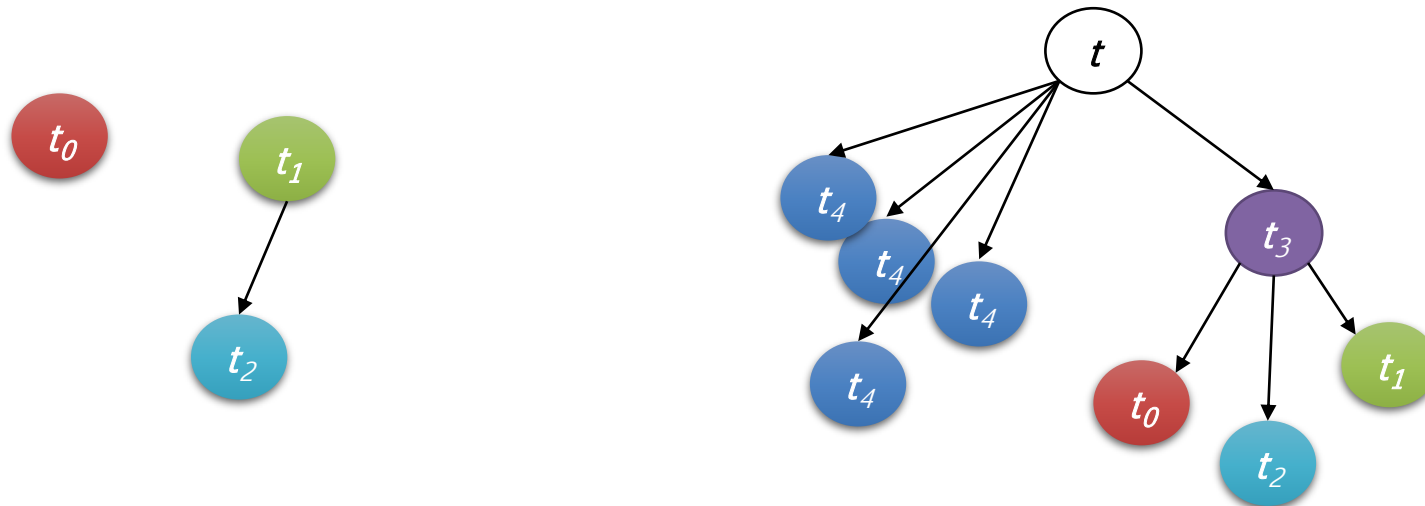
The queue

It's a shared resource

- › Its primitives **q_push** and **q_pop** are *thread-safe*, i.e., internally, the concurrent access to shared data structure is protected by semaphores/mutexes

Typically implemented as a FIFO queue

- › ..but we can also have more complex semantics (e.g., parent-son => DAGs) among tasks
- › We can have queues that are tailored/more efficient for a specific problem/domain/algorithm!





Course website

- › http://hipert.unimore.it/people/paolob/pub/Industrial_Informatics/index.html

My contacts

- › paolo.burgio@unimore.it
- › <http://hipert.mat.unimore.it/people/paolob/>

Resources

- › "Parallel programming" course by "a guy" @UNIMORE
 - https://hipert.unimore.it/people/paolob/pub/Calcolo_Parallelo/index.html
 - <https://github.com/HiPeRT/cp19/>
- › A "small blog"
 - <http://www.google.com>