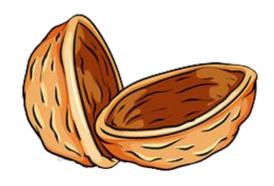
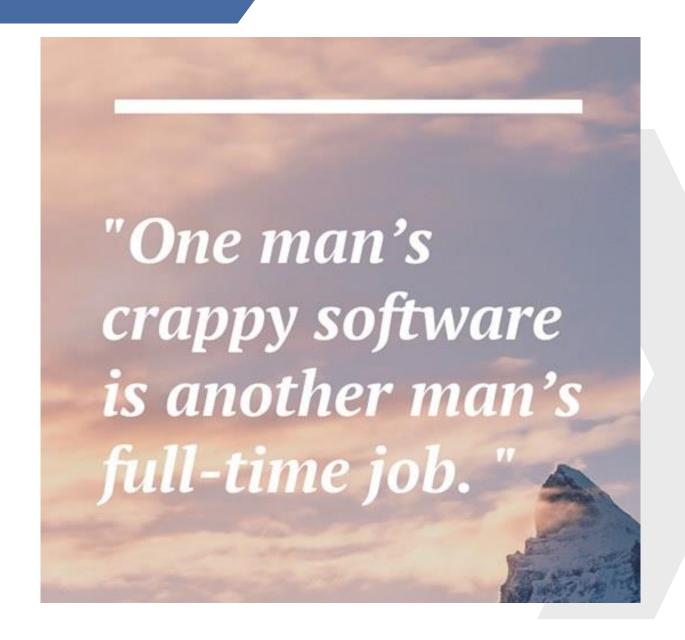
POSIX Threads in a nutshell

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The POSIX IEEE standard

eng.wikipedia.org

POSIX Threads, usually referred to as PThreads, is an execution model that exists independently from a language, as well as a parallel execution model. It allows a program to control multiple different flows of work that overlap in time.

- > Threading API
- > Single process
- > Shared memory space





The POSIX IEEE standard

- Specifies an operating system interface similar to most UNIX systems
 - It extends the C language with primitives that allows the specification of the concurrency
- > POSIX distinguishes between the terms process and thread
 - "A process is an address space with one or more threads executing"
 - "A thread is a single flow of control within a process (a unit of execution)"
- > Every process has at least one thread
 - the "main()" (aka "master") thread; its termination ends the process
 - All the threads share the same address space, and have a private stack



Thread body

> A (P)thread is identified by a C function, called body:

```
void *my_pthread_fn(void *arg)
{
   // Thread body
}
```

- > A thread starts with the first instruction of its body
- > The threads ends when the body function ends
 - It's not the only way a thread can die



Thread creation

> Thread can be created using the primitive

- > pthread t is the type that contains the thread ID
- > pthread_attr_t is the type that contains the parameters of the thread
- > arg is the argument passed to the thread body when it starts



Thread attributes

- > Thread attributes specifies the characteristics of a thread
 - We won't see this; leave empty
- Attributes must be initialized and destroyed always

```
int pthread_attr_init(pthread_attr_t *attr);
int pthread_attr_destroy(pthread_attr_t *attr);
```



Thread termination

A thread can terminate itself calling

```
void pthread_exit(void *retval);
```

- > When the thread body ends after the last "}", pthread_exit() is called implicitly
- > Exception: when main() terminates, exit() is called implicitly



Thread IDs

> Each thread has a unique ID

```
pthread.h

pthread_t pthread_self(void);
```

> The thread ID of the current thread can be obtained using

> Two thread IDs can be compared using



Joining a thread

> A thread can wait the termination of another thread using

- It gets the return value of the thread or PTHREAD_CANCELED if the thread has been killed
- > By default, every thread must be joined
 - The join frees all the internal resources
 - Stack, registers, and so on



Example

```
Let's code!
```

- > Filename: ex create.c
- > The demo explains how to create a thread
 - the main () thread creates another thread (called body ())
 - the body () thread checks the thread lds using pthread_equal () and then ends
 - the main () thread joins the body () thread
- > When compiling under gcc & GNU/Linux, remember
 - the -lpthread option!
 - to add #include "pthread.h"

> Credits to PJ

Semaphores



Semaphores

A semaphore is a counter managed with a set of primitives

It is used for

- > Synchronization
- > Mutual exclusion (critical sections)

POSIX Semaphores can be

- > Unnamed (local to a process)
- > Named (shared between processed through a file descriptor we won't see them)



Unnamed semaphores

Operations permitted:

- initialization /destruction
- > blocking wait / nonblocking wait
 - counter decrement
- > post
 - counter increment
- > counter reading
 - simply returns the counter



Initializing a semaphore

> The sem_t type contains all the semaphore data structures

int sem_init(sem_t *sem, int pshared, unsigned int value);
 - pshared is 0 if sem is not shared between processes

int sem_destroy(sem_t *sem)
 - It destroys the sem semaphore



Semaphore waits

```
int sem_wait(sem_t *sem);
int sem trywait(sem t *sem);
```

- > Under the hood..
- > If the counter is greater than 0 the thread does not block
 - sem trywait never blocks



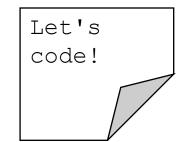
Other semaphore primitives

```
int sem_post(sem_t *sem);
   - It increments the semaphore counter
   - It unblocks a waiting thread

int sem_getvalue(sem_t *sem,int *val);
   - It simply returns the semaphore counter
```



Example(s)



Filename: critical-section.c

> In this example, semaphores are used to implement mutual exclusion in the output of a character in the console

Filename: producer-consumer.c

> In this example, semaphores are used to implement producer-consumer synchronization

Mutexes





What is a (POSIX) mutex?

- > Like a binary semaphore, used for mutual exclusion
 - But.. a mutex can be unlocked only by the thread that locked it
- > Mutexes also support some RT protocols
 - Priority inheritance, priority ceiling...
 - They are not implemented under a lot of UNIX OS
 - Out of scope for this course



Mutex attributes

> Mutex attributes are used to initialize a mutex

```
int pthread_mutexattr_init (pthread_mutexattr_t *attr);
int pthread_mutexattr_destroy (pthread_mutexattr_t
   *attr);
```

> Initialization and destruction of a mutex attribute



Mutex initialization

> Initialize a mutex with a given mutex attribute

Destroys a mutex

```
int pthread_mutex_destroy (pthread_mutex_t *mutex);
```



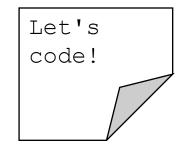
Mutex lock and unlock

- > This primitives implement the blocking lock, the non-blocking lock and the unlock of a mutex
- The mutex lock is NOT a cancellation point

```
int pthread_mutex_lock(pthread_mutex_t *m);
int pthread_mutex_trylock(pthread_mutex_t *m);
int pthread_mutex_unlock(pthread_mutex_t *m);
```



Example



- > Filename: ex mutex.c
- > This is prev. example written using mutexes instead of semaphores.



PThreads scheduling





Scheduling algorithms

- > The POSIX standard specifies in sched.h at least two scheduling strategies which can be used, identified by the symbols SCHED FIFO and SCHED RR
 - Other scheduling policies may be supported by each particular implementation, under the symbol SCHED OTHER

POSIX specifies a Fixed Priority scheduler with at least 32 priorities (0 to 31)

- > Every priority corresponds to a queue, where all the threads with the same priority are inserted
- > The first ready thread in the highest non-empty priority queue is selected for scheduling and becomes the running thread



POSIX and priorities

thread priorities can be specified at creation time into the thread attributes

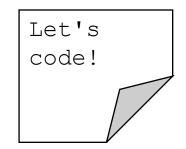


Real-Time and UNIX

- > UNIX systems usually schedule all its threads at low priorities
- > When a RT thread is created, it always preempts all the other applications (i.e. the X server, and all the other demons)
- > For that reason,
 - real-time computations have to be limited
 - only root can use the real-time priorities



Example



- > Filename: ex rr.c
- > The demo explains the behavior of the RT priorities and of the other policies
- > The main () thread creates a high priority thread that activates a low priority thread and two medium priority threads
- > The medium priority threads are scheduled with policies SCHED_RR and SCHED_FIFO
- > When compiling under gcc & GNU/Linux, remember
 - the -lpthread option!
 - to add #include "pthread.h"

> Credits to PJ



How to run the examples



- > Download the Code/ folder from the course website
- Compile
- \$ gcc code.c -o code -lpthread
- > Run (Unix/Linux)
- \$./code
- > Run (Win/Cygwin)
- \$./code.exe



References



Course website

http://hipert.unimore.it/people/paolob/pub/Industrial_Informatics/index.html

My contacts

- > paolo.burgio@unimore.it
- http://hipert.mat.unimore.it/people/paolob/

Resources

- https://computing.llnl.gov/tutorials/pthreads/
- http://man7.org/linux/man-pages/man7/pthreads.7.html
- A "small blog"
 - http://www.google.com