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POSIX Threads

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The POSIX IEEE standard

- ✓ Specifies an **operating system interface similar to most UNIX systems**
 - *It extends the C language with primitives that allows the specification of the concurrency*
- ✓ POSIX distinguishes between the terms process and thread
 - "A **process** is an address space with one or more threads executing in that address space"
 - "A **thread** is a single flow of control within a process (a unit of execution)"
- ✓ Every process has at least one thread
 - the `main()` (aka "**master**") thread; its termination ends the process
 - All the threads **share** the same address space, and have a **private** stack



The PThread library

- ✓ The pthread **primitives** are usually implemented into a pthread **library**
- ✓ All the declarations of the primitives cited in these slides can be found into `sched.h`, `pthread.h` and `semaphore.h`
 - Use `man` to get on-line documentation
- ✓ When compiling under gcc & GNU/Linux, remember the **`-lpthread`** option!

PThread creation, join, end



PThread body

- ✓ A (P)thread is identified by a C function, called body:

```
void *my_pthread_fn(void *arg)
{
    ...
}
```

- ✓ A thread starts with the first instruction of its body
- ✓ The threads ends when the body function ends
 - it's not the only way a thread can die



Thread creation

- ✓ Thread can be created using the primitive

```
int pthread_create ( pthread_t *ID,  
                    pthread_attr_t *attr,  
                    void *(*body)(void *),  
                    void * arg  
                    );
```

- ✓ `pthread_t` is the type that contains the thread ID
- ✓ `pthread_attr_t` is the type that contains the parameters of the thread
- ✓ `arg` is the argument passed to the thread `body` when it starts



Thread attributes

- ✓ Thread attributes specifies the characteristics of a thread
 - Stack size and address
 - Detach state (joinable or detached)
 - Scheduling parameters (priority, ...)

- ✓ Attributes must be initialized and destroyed
 - `int pthread_attr_init(pthread_attr_t *attr);`
 - `int pthread_attr_destroy(pthread_attr_t *attr);`



Thread termination

- ✓ A thread can terminate itself calling

```
void pthread_exit(void *retval);
```

- ✓ When the thread body ends after the last “}”, `pthread_exit()` is called implicitly
- ✓ Exception: when `main()` terminates, `exit()` is called implicitly



Thread IDs

- ✓ Each thread has a unique ID
- ✓ The thread ID of the current thread can be obtained using

```
pthread_t pthread_self(void);
```

- ✓ Two thread IDs can be compared using

```
int pthread_equal(    pthread_t thread1,  
                    pthread_t thread2 );
```



Joining a thread

- ✓ A thread can wait the termination of another thread using

```
int pthread_join ( pthread_t th,  
                  void **thread_return);
```

- ✓ It gets the return value of the thread or `PTHREAD_CANCELED` if the thread has been killed
- ✓ By default, every thread **must** be joined
 - The join frees all the internal resources
 - Stack, registers, and so on



Joining a thread (2)

- ✓ A thread which does not need to be joined has to be declared as **detached**
- ✓ 2 ways to have it:
 - While creating (in father thread) using `pthread_attr_setdetachstate()`
 - The thread itself can become detached calling in its body `pthread_detach()`
- ✓ Joining a detached thread returns **an error**





Example

Let's
code!

- ✓ Filename: `ex_create.c`
- ✓ The demo explains how to create a thread
 - the `main()` thread creates another thread (called `body()`)
 - the `body()` thread checks the thread ids using `pthread_equal()` and then ends
 - the `main()` thread joins the `body()` thread

✓ *Credits to PJ*

Pthread cancellation



Killing a thread

- ✓ A thread can be killed calling

```
int pthread_cancel(pthread_t thread);
```

- ✓ When a thread dies its data structures will be released
 - By the join primitive if the thread is joinable
 - Immediately if the thread is *detached*
 - Why?



PThread cancellation

- ✓ Specifies how to react to a **kill** request

- ✓ There are two different behaviors:

- **deferred** cancellation

- when a kill request arrives to a thread, the thread does not die. The thread will die only when it will execute a primitive that is a **cancellation point**. This is the default behavior of a thread.

- **asynchronous** cancellation

- when a kill request arrives to a thread, the thread dies. The programmer must ensure that all the application data structures are coherent.



Cancellation states and cleanups

- ✓ The user can set the cancellation state of a thread using:

```
int pthread_setcancelstate(int state, int *oldstate);  
int pthread_setcanceltype(int type, int *oldtype);
```

- ✓ The user can protect some regions providing destructors to be executed in case of cancellation

```
int pthread_cleanup_push(void (*routine)(void *),  
                        void *arg);  
int pthread_cleanup_pop(int execute);
```




Cancellation points

- ✓ The **cancellation points** are primitives that can potentially **block** a thread
- ✓ When called, if there is a kill request pending the thread will die
 - `void pthread_testcancel(void);`
 - `sem_wait`, `pthread_cond_wait`, `printf` and all the I/O primitives are cancellation points
 - `pthread_mutex_lock`, is **NOT** a cancellation point
 - Why?
- ✓ A complete list can be found into the POSIX Standard



Cleanup handlers

- ✓ The **user** must **guarantee** that when a thread is killed, the application data remains coherent.
- ✓ The user can **protect** the application code using **cleanup handlers**
 - A cleanup handler is a user function that *cleans up* the application data
 - They are called when the thread **ends** and when it is **killed**



Cleanup handlers (2)

```
void pthread_cleanup_push (void (*routine) (void *),  
                           void *arg);  
void pthread_cleanup_pop (int execute);
```

- They are pushed and popped as in a stack
- If `execute != 0` the cleanup handler is called when popped
- The cleanup handlers are called in LIFO order



Example

Let's
code!

- ✓ Filename: `ex_cancellation.c`
- ✓ Highlights the behavior of:
 - Asynchronous cancellation
 - Deferred cancellation
- ✓ Explains the cleanup handlers usage

POSIX semaphores



Semaphores

- ✓ A semaphore is a counter managed with a set of primitives
- ✓ It is used for
 - Synchronization
 - Mutual exclusion
- ✓ POSIX Semaphores can be
 - Unnamed (local to a process)
 - Named (shared between processes through a file descriptor)



Unnamed semaphores

- ✓ Mainly used with multithread applications
- ✓ Operations permitted:
 - initialization /destruction
 - blocking wait / nonblocking wait
 - counter decrement
 - post
 - counter increment
 - counter reading
 - simply returns the counter



Initializing a semaphore

- ✓ The `sem_t` type contains all the semaphore data structures

```
int sem_init(sem_t *sem, int pshared, unsigned int value);
```

- `pshared` is 0 if `sem` is not shared between processes

```
int sem_destroy(sem_t *sem)
```

- It destroys the `sem` semaphore



Semaphore waits

```
int sem_wait(sem_t *sem);  
int sem_trywait(sem_t *sem);
```

- ✓ Under the hood..
- ✓ If the counter is greater than 0 the thread does not block
 - `sem_trywait` never blocks
- ✓ `sem_wait` is a cancellation point



Other semaphore primitives

```
int sem_post(sem_t *sem);
```

- It increments the semaphore counter
- It unblocks a waiting thread

```
int sem_getvalue(sem_t *sem, int *val);
```

- It simply returns the semaphore counter



Example

Let's
code!

- ✓ Filename: `ex_sem.c`
- ✓ In this example, semaphores are used to implement mutual exclusion in the output of a character in the console.

PThread mutexes



What is a POSIX mutex?

- ✓ Like a **binary semaphore** used for **mutual exclusion**
 - But.. a mutex can be unlocked **only** by the thread that locked it
- ✓ Mutexes also support some RT protocols
 - Priority inheritance
 - Priority ceiling
 - They are not implemented under a lot of UNIX OS
- ✓ Out of scope for this course



Mutex attributes

- ✓ Mutex attributes are used to initialize a mutex

```
int pthread_mutexattr_init (pthread_mutexattr_t *attr);  
int pthread_mutexattr_destroy (pthread_mutexattr_t *attr);
```

- ✓ Initialization and destruction of a mutex attribute



Mutex initialization

- ✓ Initialize a mutex with a given mutex attribute

```
int pthread_mutex_init (pthread_mutex_t *mutex,  
                        const pthread_mutexattr_t *attr);
```

- ✓ Destroys a mutex

```
int pthread_mutex_destroy (pthread_mutex_t *mutex);
```



Mutex lock and unlock

- ✓ This primitives implement the blocking lock, the non-blocking lock and the unlock of a mutex
- ✓ The mutex lock is **NOT** a cancellation point

```
int pthread_mutex_lock(pthread_mutex_t *m);  
int pthread_mutex_trylock(pthread_mutex_t *m);  
int pthread_mutex_unlock(pthread_mutex_t *m);
```




Example

Let's
code!

- ✓ Filename: `ex_mutex.c`
- ✓ This is prev. example written using mutexes instead of semaphores.