

# Matrix Chain Multiplication

Dynamic Programming

# Matrix Chain Multiplication

Dot Product

$$\begin{bmatrix} 1 & 5 & 9 & 7 & 3 & 4 \\ 2 & 1 & 9 & 7 & 2 & 6 \\ 9 & 5 & 2 & 2 & 3 & 5 \\ 6 & 6 & 1 & 3 & 1 & 7 \end{bmatrix} \times \begin{bmatrix} 5 & 1 & 3 \\ 9 & 5 & 1 \\ 8 & 7 & 6 \\ 9 & 6 & 8 \\ 8 & 1 & 3 \\ 2 & 2 & 9 \end{bmatrix} = \begin{bmatrix} 217 & 142 & 163 \\ 182 & 126 & 177 \\ 158 & 73 & 114 \\ 141 & 76 & 120 \end{bmatrix}$$

4x6                      6x3                      4x3

72 Multiplications  
in Total!  
(4x6x3)

## Matrix Chain Multiplication

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & & 10 \times 3 & & 3 \times 12 & & 12 \times 20 & & 20 \times 7 \end{matrix}$$

$$4 \times 10 \times 3 + 4 \times 3 \times 12 + 4 \times 12 \times 20 + 4 \times 20 \times 7 = 1784 \text{ Multiplication Operations}$$

Very Large Computational Times !!

## Matrix Chain Multiplication

$$\begin{array}{ccccccccc} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & & 10 \times 3 & & 3 \times 12 & & 12 \times 20 & & 20 \times 7 \end{array}$$

Goal: Find the optimal way to multiply these matrices to perform the fewest multiplications.

Naïve Approach: Try them all, and pick the most optimal one.

Running time:  $\Omega(4^n/n^{3/2})$  -  $4^n$  dominates! Exponential

# Matrix Chain Multiplication

There is a better way! Dynamic Programming!

Step 1: Check if the problem has Optimal Substructure

If we have an optimal solution for  $A_{i...j}$

Assume the solution has the following parentheses:

$$(\underline{A_{i...k}})(\underline{A_{k+1...j}})$$

If there is a better way to multiply  $(A_{i...k})$ , then we would have a more optimal solution.

This would be a contradiction, as we already stated that we have the optimal solution for  $A_{i...j}$

Therefore this problem has optimal substructure.

## Matrix Chain Multiplication

A matrix series  $A_{i...j}$  can be broken up into a more efficient solution:

$$(A_{i...k})(A_{k+1...j})$$

We want to find out at which 'k' returns the fewest number of multiplications

We need to define our recursive formula:

$M[i,j]$  is the cost of multiplying matrices from  $A_i$  to  $A_j$

# Matrix Chain Multiplication

Now we want to try out a bunch of values for 'k' in order to see what the best one is:

$$M[i,j] = M[i,k] + M[k+1,j] + p_{i-1}p_kp_j$$

100          200          2x3x4

Since we don't know what k is, we try this range of k:

$$\begin{matrix} 100 & 200 \\ (A_{i\dots k})(A_{k+1\dots j}) \\ 2 \times 3 & 3 \times 4 \end{matrix}$$

The minimum returned value is our solution!

$$i \leq k < j$$

Our Final Recursive Formula:

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} \end{cases}$$

## Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

i \ j	1	2	3	4	5
1					
2	x				
3	x	x			
4	x	x	x		
5	x	x	x	x	



## Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

i \ j	1	2	3	4	5
1	0				
2	x	0			
3	x	x	0		
4	x	x	x	0	
5	x	x	x	x	0

Step 1: Fill the table for  $i = j$

## Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

i \ j	1	2	3	4	5
1	0				
2	x	0			
3	x	x	0		
4	x	x	x	0	
5	x	x	x	x	0

Step 2: Fill the table for:

$i=1, j=2$

$i=2, j=3$

$i=3, j=4$

$i=4, j=5$

# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

i \ j	1	2	3	4	5
1	0				
2	x	0			
3	x	x	0		
4	x	x	x	0	
5	x	x	x	x	0

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

$$M[1,2] = \min_{1 \leq k < 2} \{M[1,1] + M[1+1,2] + p_0p_1p_2\}$$

$$M[1,2] = \min_{1 \leq k < 2} \{0 + 0 + 4 \times 10 \times 3\}$$

# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

i \ j	1	2	3	4	5
1	0	120			
2	x	0	360		
3	x	x	0		
4	x	x	x	0	
5	x	x	x	x	0

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

$$M[1,2] = \min_{1 \leq k < 2} \{M[1,1] + M[1+1,2] + p_0p_1p_2\}$$

$$M[1,2] = \min_{1 \leq k < 2} \{0 + 0 + 4 \times 10 \times 3\}$$

$$M[1,2] = 120$$

$$M[2,3] = \min_{2 \leq k < 3} \{M[2,2] + M[2+1,3] + p_1p_2p_3\}$$

$$M[2,3] = \min_{2 \leq k < 3} \{0 + 0 + 10 \times 3 \times 12\}$$

$$M[2,3] = 360$$

# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

i \ j	1	2	3	4	5
1	0	120			
2	x	0	360		
3	x	x	0	720	
4	x	x	x	0	1680
5	x	x	x	x	0

$$M[1,2] = \min_{1 \leq k < 2} \{M[1,1] + M[1+1,2] + p_0p_1p_2\}$$

$$M[1,2] = \min_{1 \leq k < 2} \{0 + 0 + 4 \times 10 \times 3\}$$

$$M[1,2] = 120$$

$$M[2,3] = \min_{2 \leq k < 3} \{M[2,2] + M[2+1,3] + p_1p_2p_3\}$$

$$M[2,3] = \min_{2 \leq k < 3} \{0 + 0 + 10 \times 3 \times 12\}$$

$$M[2,3] = 360$$

$$M[3,4] = \min_{3 \leq k < 4} \{M[3,3] + M[3+1,4] + p_2p_3p_4\}$$

$$M[3,4] = \min_{3 \leq k < 4} \{0 + 0 + 3 \times 12 \times 20\}$$

$$M[3,4] = 720$$

$$M[4,5] = \min_{4 \leq k < 5} \{M[4,4] + M[4+1,5] + p_3p_4p_5\}$$

$$M[4,5] = \min_{4 \leq k < 5} \{0 + 0 + 12 \times 20 \times 7\}$$

$$M[1,5] = 1680$$

# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

i \ j	1	2	3	4	5
1	0	120	264		
2	x	0	360		
3	x	x	0	720	
4	x	x	x	0	1680
5	x	x	x	x	0

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

$$M[1,3] = \min_{1 \leq k < 3}$$

$$k=1$$

$$= M[1,1] + M[1+1,3] + p_0p_1p_3$$

$$= 0 + 360 + 4 \times 10 \times 12$$

$$= 840$$

$$k=2$$

$$= M[1,2] + M[2+1,3] + p_0p_2p_3$$

$$= 120 + 0 + 4 \times 3 \times 12$$

$$= 264 \leftarrow$$



# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

i \ j	1	2	3	4	5
1	0	120	264		
2	x	0	360	1320	
3	x	x	0	720	
4	x	x	x	0	1680
5	x	x	x	x	0

$$M[2,4] = \min_{2 \leq k < 4}$$

$$k=2$$

$$= M[2,2] + M[2+1,4] + p_1p_2p_4$$

$$= 0 + 720 + 10 \times 3 \times 20$$

$$= 1320$$

$$k=3$$

$$= M[2,3] + M[3+1,4] + p_1p_3p_4$$

$$= 360 + 0 + 10 \times 12 \times 20$$

$$= 2760$$

computer science

# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

i \ j	1	2	3	4	5
1	0	120	264		
2	x	0	360	1320	
3	x	x	0	720	1140
4	x	x	x	0	1680
5	x	x	x	x	0

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

$$M[3,5] = \min_{3 \leq k < 5}$$

$$k=3$$

$$= M[3,3] + M[3+1,5] + p_2p_3p_5$$

$$= 0 + 1680 + 3 \times 12 \times 7$$

$$= 1932$$

$$k=4$$

$$= M[3,4] + M[4+1,5] + p_2p_4p_5$$

$$= 720 + 0 + 3 \times 20 \times 7$$

$$= 1140 \leftarrow$$



# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & & 10 \times 3 & & 3 \times 12 & & 12 \times 20 & & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

i \ j	1	2	3	4	5
1	0	120	264	1080	
2	x	0	360	1320	
3	x	x	0	720	1140
4	x	x	x	0	1680
5	x	x	x	x	0

$$M[1,4] = \min_{1 \leq k < 4}$$

$$k=1$$

$$= M[1,1] + M[1+1,4] + p_0p_1p_4$$

$$= 0 + 1320 + 4 \times 10 \times 20$$

$$= 2120$$

$$k=2$$

$$= M[1,2] + M[2+1,4] + p_0p_2p_4$$

$$= 120 + 720 + 4 \times 3 \times 20$$

$$= 1080$$

$$k=3$$

$$= M[1,3] + M[3+1,4] + p_0p_3p_4$$

$$= 264 + 0 + 4 \times 12 \times 20$$

$$= 1224$$

# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & & 10 \times 3 & & 3 \times 12 & & 12 \times 20 & & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

i \ j	1	2	3	4	5
1	0	120	264	1080	
2	x	0	360	1320	1350
3	x	x	0	720	1140
4	x	x	x	0	1680
5	x	x	x	x	0

$$M[2,5] = \min_{2 \leq k < 5}$$

$$k=2$$

$$= M[2,2] + M[2+1,5] + p_1p_2p_5$$

$$= 0 + 1140 + 10 \times 3 \times 7$$

$$= 1350$$

$$k=3$$

$$= M[2,3] + M[3+1,5] + p_1p_3p_5$$

$$= 360 + 1680 + 10 \times 12 \times 7$$

$$= 2880$$

$$k=4$$

$$= M[2,4] + M[4+1,5] + p_1p_4p_5$$

$$= 1320 + 0 + 10 \times 20 \times 7$$

$$= 2720$$

# Matrix Chain Multiplication

We want to start with  $i = j$ , then  $i < j$  starting with a spread of 1, working our way up

i \ j	1	2	3	4	5
1	0	120	264	1080	1344
2	x	0	360	1320	1350
3	x	x	0	720	1140
4	x	x	x	0	1680
5	x	x	x	x	0

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} \end{cases}$$

$$\begin{matrix} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{matrix}$$

$$M[1,5] = \min_{1 \leq k < 5}$$

$$\begin{aligned} k=1 \\ &= M[1,1] + M[1+1,5] + p_0p_1p_5 \\ &= 0 + 1350 + 4 \times 10 \times 7 \\ &= 1630 \end{aligned}$$

$$\begin{aligned} k=2 \\ &= M[1,2] + M[2+1,5] + p_0p_2p_5 \\ &= 120 + 1140 + 4 \times 3 \times 7 \\ &= 1344 \end{aligned}$$

$$\begin{aligned} k=3 \\ &= M[1,3] + M[3+1,5] + p_0p_3p_5 \\ &= 264 + 1680 + 4 \times 12 \times 7 \\ &= 2280 \end{aligned}$$

$$k=4$$

$$\begin{aligned} &= M[1,4] + M[4+1,5] + p_0p_4p_5 \\ &= 1080 + 0 + 4 \times 20 \times 7 \\ &= 1640 \end{aligned}$$

## Matrix Chain Multiplication

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{array}{ccccccccc} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & & 10 \times 3 & & 3 \times 12 & & 12 \times 20 & & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{array}$$

We now know that we can multiply  $A_1$  to  $A_5$   
in as few as 1344 multiplication operations!

But where do we put our brackets?

We must focus on the selected  $k$  values

## Matrix Chain Multiplication

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{array}{ccccccccc} A_1 & \times & A_2 & \times & A_3 & \times & A_4 & \times & A_5 \\ 4 \times 10 & & 10 \times 3 & & 3 \times 12 & & 12 \times 20 & & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{array}$$

$$k=2$$

$$M[1,5] = M[1,2] + M[3,5] + p_0p_2p_5$$

## Matrix Chain Multiplication

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{array}{c} (A_1 \times A_2) (A_3 \times A_4 \times A_5) \\ \begin{array}{ccccc} 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{array} \end{array}$$

↑

$k=2 \leftarrow$

$$M[1,5] = M[1,2] + M[3,5] + p_0p_2p_5$$

## Matrix Chain Multiplication

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

$$\begin{array}{ccccc} (A_1 \times A_2) & ((A_3 \times A_4) & A_5) \\ 4 \times 10 & 10 \times 3 & 3 \times 12 & 12 \times 20 & 20 \times 7 \\ p_0 & p_1 & p_1 & p_2 & p_2 & p_3 & p_3 & p_4 & p_4 & p_5 \end{array}$$

k=2

$$M[1,5] = \underline{M[1,2]} + \underline{M[3,5]} + p_0p_2p_5$$

k=4

$$M[3,5] = M[3,4] + M[5,5] + p_2p_4p_5$$



## Matrix Chain Multiplication

$$M[i,j] = \begin{cases} 0 & \text{if } i=j \\ \min_{i \leq k < j} \{M[i,k] + M[k+1,j] + p_{i-1}p_kp_j\} & \text{otherwise} \end{cases}$$

Check our work!

$$\begin{array}{c} (A_1 \times A_2) ((A_3 \times A_4) A_5) \\ 4 \times 10 \quad 10 \times 3 \quad 3 \times 12 \quad 12 \times 20 \quad 20 \times 7 \\ \left( \begin{array}{c} 120 \\ 4 \times 3 \end{array} \right) \left( \begin{array}{c} 720 \\ 3 \times 20 \quad 20 \times 7 \end{array} \right) \\ \quad \quad \quad 420 \\ \quad \quad \quad 4 \times 3 \quad 3 \times 7 \\ \quad \quad \quad 84 \\ \quad \quad \quad 4 \times 7 \end{array}$$

$$120 + 720 + 420 + 84 = 1344$$



## Matrix Chain Multiplication

**Example:** Given a chain of four matrices  $A_1, A_2, A_3$  and  $A_4$ , with  $p_0 = 5, p_1 = 4, p_2 = 6, p_3 = 2$  and  $p_4 = 7$ . Find  $m[1, 4]$ .