ASSIGNMENT # 03	1
DESIGN AND ANALYSIS OF	
ALGORITHMS : CS 302	
Sequential of the sequence of	
Name of Group members:	
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Section: F	
Due Date: 16th Dec 2020	
Group question allocated: 9	
	書
PROBLEM STATEMENT: Using dynamic programming	
understand and solve partition-problem	
$S = \{3, 1, 1, 2, 2, 1\}$	
Identifying and explaining the problem	
statement furiner:	
Divide the guen array into two subset	
that the sum of both subsets is equal.	
1> now the partition can be of anysize	
is further more elements in subset must be	
unique i e vie element in one subset	
shouldn't belong to the other subject as nell.	3
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And the state of t	

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Ly	This is basically an extension of subset sum	1
	problem in which the sum is already	
	given and we need to partition the array into	
	subset such that both subset equal to that sum.	
	the state of the s	
	- Similarity of this problem with	
	knapsack problem solution:	
V 3-	If we recall what knapsack was, we	
	were given an dem array union was duided	
	into weight array and value array.	
	And a weight was given unich was equal	
	to the weight of the knapsack.	
	Here we would say that the guin entay is	
	weight array, and we will see further	
	trat by even lodd sum of array technique	
	we will get the sum and hence	
	Array guen is analogous to reight array	
	sun is analogous to weight.	
	Furthermore, in knapsack we do is	
	on each dem we check whether to pick	
	that dem and put it in the knap sack or	
	not, in short words yes or no / True or	
	false which is very much similar to our	
	"Equal subset problem" in which we	
	first see we pick elemed or not if sum equal	
	we return true.	

	Solution approach explained:	3
	The array quen	
	arr={3,1,1,2,2,13.	
	The odd-eur sum approach.	
	The sum of the given array	
-	will be:	
	sum = 3 + 1 + 1 + 2 + 2 + 1	
	sum = 10	
	sum = 10 = 5.	
	2 2	_
	Now here notice the sum of the array	
	is even so naturally we get the assurance	
	that this array can be diricted into	
	or partitioned into subsets equally	
1	so that both of the subsets have equal	
	sum is possible	
	a counter example:	,
	if the given array would be like	
	arr= {3,1,1,2,2,1,13	
	The sum of this array with be:	
	sum = 3 + 1 + 1 + 2 + 2 + 1 + 1	
	sum = 11	
	sim = 11 = 5.5	
	2 2	
_	This means it is odd hence this array CANNOT	
	COPY	

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	be partitioned such that subsets have	- 10 3
	equal sum.	
	In this case False will be returned	
	Now for soming this the approach we	
	going to use is "Topdown Knapsack 01"	
	godig w use is repaired	
	approach.	
	The code / Algorithm of the	
	equal Partition-Problem or simply	
	Partition-Problem.	
9	Code /Algo for array sum:	
	int sum=0; « size of away.	
	for lunt i=0; icsuize; i++)	
	S 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	
	sum = sum + arr [i];	
	7	
4	Main code/Algo	
8	if (sum 1.2 != 0)	
	{ -> This wondution means if	
	return false; array sem is odd.	
,		
		1
	The state of the s	
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else if (sum 7.2 == 0)	the evension	5
{		7 Posize of an	ay
retur	n subset sum (a	rt, sum/2, n.)	
	£ 1	. (
func	tion which the q	vien array	
700 1575	implement	2	
top	down approach.	- 170 - 170	
3	B1 10	3	
Subsetsun	fretion:		
	Three or false rel	wrned	
1	+ code (main)	Note: we want	
	; i L n+1; i++)	one subset which	
1.78	; j L size+1 ; z +1)	equals to the even	
		sum automatically	
ifli	==0)	it implies the	
t[i]	[j] = False;	remaining elemts	
if (j=	=0)	will gum up to	
	[j]=true;	the same sem	
00 i= size 0	efanes	nence both equal	
j= Etem		sum.	
=> we make	t[n+1][W+1]		ų.
	<i>></i>		
	sizl sum. Janay	A Section of the sect	
t[s1+	1][sum+1] arrau	y/ Idarray (matis	
		nade.	

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n+1 sum+1

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	0	1	2	3	ч	5	cum	7	
0	ī	F	F	F	Ŧ	F	ow: {3,1,1,2,2,1		
1	T	F	F	τ	F	F	sum:		
2	T	τ	F	F	Т	F	8 3 4		
3	T	Т	T	F	Т	Т			
4	T	F	Т	T	T	T			
5	T,	F	T	Т	Т	T			
6	T	Τ	T	Т	T	τ			
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		Consider our [0][0] arr [7:2 elents.							
	190	mean an: empty sum:0.							
	sum: O possible.								
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	l=1,j=1	8
	ar[0]=3 L=1 x	
	t[1][1] = t[0][1] =	
	t[1][1] = F	
J	i=1, j=2	
	ar[0]=31=2 x	
	t[1][2] = t[0][2]	
	t[1][2] = F	,U 18
_	i=1, j=3	
	an [0] = 3 L= 3 V	
	t[1][3] = t[1][3-3] OR t[0][3]	
	+[1][3] = +[1][0] OR +[0][3]	
	t[1][3] = T 1) F	
	t[1][3] = T	
L	i=1, j=4	
	arr[0] = 31 = 24	
	t[1][4] = t[1][4-3] OR t[0][9]	
	= t[1][1] 11 t[0][4]	
100	= FIIF	
	= F	