

## COMP 330 Homework 2

Ryan Sowa, ID: 260886668

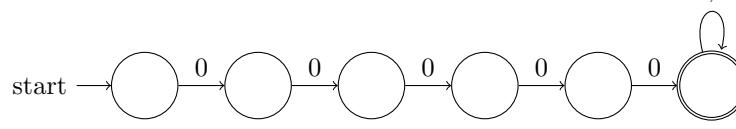
1. (35 points) For each of the following languages give a proof that it is or is not regular

(a)

$$\{0^m 1^n \mid m \geq 5 \text{ and } n \geq 0\}.$$



*Proof.* This language is regular, because the following NFA can be constructed for this language:



(b)

$$\{0^m 1^n \mid m \geq n^2\}.$$



*Proof.* Let us call this language  $L$ . Suppose by contradiction that  $L$  is regular. This implies that there must exist a DFA  $M$  for  $L$ . Let  $x$  be the number of states in  $M$ . Since we only have  $x$  states and we have  $0^x \in L$ , we must have to revisit a state. In other words, there exists some  $j, k \in \mathbb{N}$  where  $0 \leq j < j + k \leq x$  such that  $0^j$  and  $0^{j+k}$  end up in the same state. However, suppose  $0^{j+k}1^{\sqrt{j+k}} \in L$ . From the reasoning above,  $0^j1^{\sqrt{j+k}}$  must also be in  $L$ . However, since

$$j \not\geq j + k$$

, ( $k > 0$ ) we have found a contradiction  $\Rightarrow \Leftarrow$ , and we can conclude that  $L$  is NOT a regular language.

□

- (c) The set of strings in  $\{0, 1\}^*$  which are not of the form  $ww$  for some  $w \in \{0, 1\}^*$ .

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*Proof.* Let us call this language  $L$ . Suppose by contradiction that  $L$  is regular. This implies that there must exist a DFA  $M$  for  $L$ . Let  $x$  be the number of states in  $M$ . Since we only have  $x$  states and we have  $0^x \in L$ , we must have to revisit a state. In other words, there exists some  $j, k \in \mathbb{N}$  where  $0 \leq j < j + k \leq x$  such that  $0^j$  and  $0^{j+k}$  end up in the same state. However, suppose  $0^{j+k}10^j1 \in L$ . From the reasoning above,  $0^j10^j1$  must also be in  $L$ . However, this expression is of the form  $ww$ , so it is not in  $L \Rightarrow \text{contradiction}$ , and we can conclude that  $L$  is NOT a regular language.

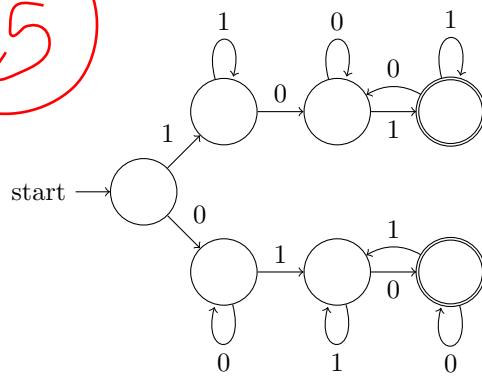
□

- (d) Over the alphabet  $\Sigma = \{0, 1\}$ :

$$L = \{x \mid x \text{ contains the same number of } 01\text{'s and } 10\text{'s as substrings}\}$$

*Proof.* This language is regular, since the following DFA can be constructed for it:

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□

- (e) Over the alphabet  $\Sigma = \{0, 1, 2\}$ :

$$L = \{x \mid x \text{ contains the same number of } 01\text{'s and } 10\text{'s as substrings}\}$$

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*Proof.* Let us call this language  $L$ . Suppose by contradiction that  $L$  is regular. This implies  $L$  must have a pumping length,  $P$  such that any string  $S \in L$ , where  $|S| \geq P$ , can be divided into 3 parts,  $x, y$  and  $z$ . Without loss of generality, suppose  $S = 1020^P1^P$  and let  $P = 3$ . We get the string  $S = 102000111$ , and we divide the string

into  $x = \varepsilon$ ,  $y = 102$ , and  $z = 000111$ . By the Pumping Lemma,  $xy^i z$  must also be in  $L$  if  $L$  is regular. However, take  $i = 2$ . Then, we get the string  $102102000111$ , which does not contain the same number of 01's and 10's as substrings  $\Rightarrow \Leftarrow$ . Therefore, we can conclude by the Pumping Lemma that this language is NOT regular.

□

- (f) The first two Fibonacci numbers are 0 and 1, and each subsequent number is the sum of the previous two: 0, 1, 1, 2, 3, 5, 8, 13, . . . Now the language in question is

$$\{0^n \mid n \text{ is a Fibonacci number}\}.$$

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*Proof.* Let us call this language  $L$ . Suppose by contradiction that  $L$  is regular. This implies that there must exist a DFA  $M$  for  $L$ . Let  $x$  be the number of states in  $M$ . Since we only have  $x$  states and we have  $0^x \in L$ , we must have to revisit a state. In other words, there exists some  $j, k \in \mathbb{N}$  where  $0 \leq j < j+k \leq x$  such that  $0^j$  and  $0^{j+k}$  end up in the same state. However, suppose  $0^j \in L$ . From the reasoning above,  $0^{j+nk}$  where  $n \in \mathbb{Z}$  must also be in  $L$ . However, this implies that each term in the Fibonacci sequence differs by exactly  $k$ , which is trivially false  $\Rightarrow \Leftarrow$ . Therefore, we can conclude that  $L$  is NOT a regular language.

□

- (g) The set of strings in  $\{0,1\}^*$  which are not palindromes:

$$\{w \in \{0,1\}^* \mid w \neq w^R\}.$$

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*Proof.* Let us call this language  $L$ . Suppose by contradiction that  $L$  is regular. This implies  $L$  must have a pumping length,  $P$  such that any string  $S \in L$ , where  $|S| \geq P$ , can be divided into 3 parts,  $x, y$  and  $z$ . Without loss of generality, suppose  $S = 0^P 100^P$  and let  $P = 3$ . We get the string  $S = 00010000$ , and we divide the string into  $x = 00$ ,  $y = 0$ , and  $z = 10000$ . By the Pumping Lemma,  $xy^i z$  must also be in  $L$  if  $L$  is regular. However, take  $i = 2$ . Then, we get the string  $000010000$ , which is a palindrome (not in  $L$ )  $\Rightarrow \Leftarrow$ . Therefore, we can conclude by the Pumping Lemma that this language is NOT regular.

□

2. (a) (3 points) Find a left-most derivation for  $aaabbabbba$  in the following context-free grammar:

$$S \rightarrow aB \mid bA$$

$$A \rightarrow a \mid aS \mid bAA$$

$$B \rightarrow b \mid bS \mid aBB$$

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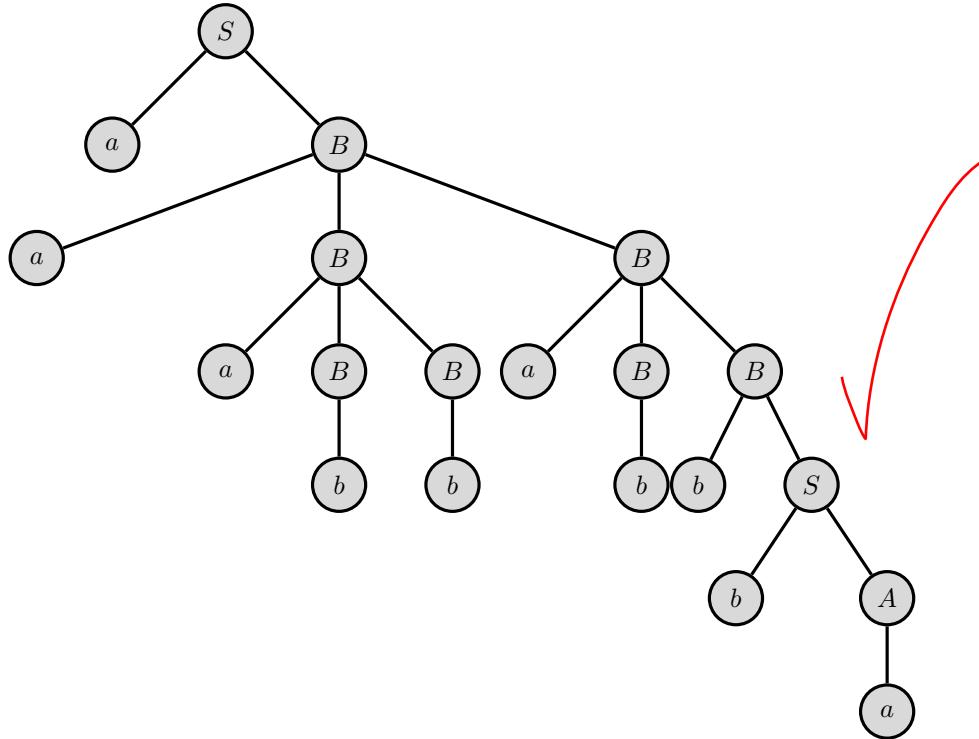
The problem here is that we don't know if  $j$  and  $j+k$  are CONSECUTIVE Fibonacci numbers.

The left-most derivation for  $aaabbabbba$  in the following context-free grammar is shown below:

(2)

$$\begin{aligned} S &\Rightarrow aB \Rightarrow aaBB \Rightarrow aaaBBB \Rightarrow aaabBB \Rightarrow aaabbB \Rightarrow aaabbaBB \\ &\Rightarrow aaabbabB \Rightarrow aaabbabbS \Rightarrow aaabbabbbA \Rightarrow aaabbabbba \end{aligned}$$

- (b) (2 points) Draw the corresponding parse-tree of your left-most derivation.

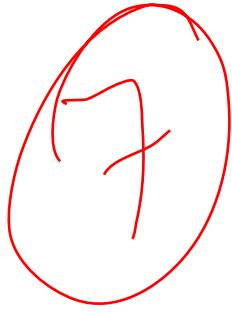


3. (10 points) Show that the language of the grammar  $S \rightarrow 0S1 \mid 1S0 \mid SS \mid \varepsilon$  is

$$\{w \in \{0, 1\}^* \mid w \text{ contains the same number of zeros and ones}\}.$$

*Proof.* Let  $G$  be the context-free grammar above. Proving the statement above is equivalent to proving the following proposition:

$P(n)$ : For any string in  $\{0, 1\}^*$  of length  $n$  such that  $n \in \mathbb{N}$ ,  $G$  can generate all possible strings of length  $n$  which contain the same number of zeros and ones.



### Base case:

We want to prove  $P(0)$  or the following statement: For any string of length 0,  $G$  can generate all possible strings of length 0 which contain the same number of zeros and ones. For a string of length 0, there is only one string containing an equal number of zeros and ones: the empty string.  $G$  can generate this string using the rule  $S \rightarrow \epsilon$ .

### Induction Hypothesis:

Assume  $P(k)$  is true for some  $n = k$ .

### Inductive Step:

We want to prove  $P(k+2)$  or the following statement: For any string of length  $k+2$ ,  $G$  can generate all possible strings of length  $k+2$  which contain the same number of zeros and ones. Since  $P(k)$  is true by the induction hypothesis, we need to analyse all the cases where we add a zero and a one (since we are adding 2 and we must have the same number of 0's and 1's) to our string of length  $k$  to our string. We can add a zero or a one in one of the following ways: 1. Add a zero at the beginning of the string and a one at the end of the string, 2. Add a one at the beginning of the string and a zero at the end of the string, 3. Add 01 or 10 to the beginning or end of our string.

Case 1: We can achieve this using the:  $S \rightarrow 0S1$  (let us call this rule 1).

Case 2: We can achieve this using:  $S \rightarrow 1S0$  (let us call this rule 2).

Case 3: We apply the rule:  $S \rightarrow SS$ . We can then apply rule 1 or rule 2 on the first "S" to add 01 or 10, respectively at the beginning of the string. Similarly, we can then apply rule 1 or rule 2 on the second "S" to add 01 or 10, respectively at the end of the string.

When we have reached the amount of 1's and 0's desired, we can replace all  $S$ 's with  $\epsilon$ 's.

Since we have proved all the cases, we can conclude that the proposition is true.

□

4. (20 points) Construct a context free grammar for the set of all words  $w$  over the alphabet  $\{0, 1\}$  such that each prefix of  $w$  has at least as many

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0's as 1's. You have to prove that (i) every such word can be generated with your grammar, and (ii) every word generated by your grammar has the desired property.

$$S \rightarrow 0S \mid 0S1 \mid SS \mid \varepsilon$$

(i)

*Proof.* Let the context free grammar we constructed above be denoted by  $G$  and the language above  $L$ .

$P(n)$ : All strings of length  $n$ , such that  $n \in \mathbb{N}$ , in  $L$  can be generated by  $G$ .

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**Base Case:** We want to prove  $P(0)$  or the statement: All strings of length 0 in  $L$  can be generated by  $G$ . The only possible string in  $L$  of length 0 is the empty string, and  $G$  can generate the empty string with the rule  $S \rightarrow \varepsilon$ .

**Induction Hypothesis:** Assume  $P(k)$  is true for some  $n = k$ .

**Induction Step:** We want to prove  $P(k+1)$  or the statement: All strings of length  $k+1$  in  $L$  can be generated by  $G$ . By our induction hypothesis, we know that all strings of length  $k$  in  $L$  can be generated by  $G$ . Let us take some string of length  $k$  in  $L$  and call it  $S$ . To reach a string of length  $k+1$  we can either 1. add a 0 to  $S$  or 2. add a 1 to  $S$ .

Case 1: Adding a 0 to  $S$  must be in  $L$  since the number of 0's in  $S$  must be still greater than or equal to the number of 1's. Our context free grammar can generate this string using the rules  $S \rightarrow SS \rightarrow S0S \rightarrow S0$ .

Case 2: According to our language, if we are adding a 1, we must have a greater number of 0's than 1's in  $S$ , since the total number of 0's in our string of  $k+1$  must be greater than or equal to the total number of 1's. Let  $m$  be the amount 0's in  $S$  minus the amount of 1's in  $S$ . Since  $m > 0$ , we can pair an extra 0 up with this added 1 using the rule  $S \rightarrow 0S1$ . Therefore, adding a 1 can also be generated by  $G$ .

When we have reached the amount of 1's and 0's desired, we can replace all  $S$ 's with  $\varepsilon$ 's.

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Since  $P(k+1)$  is true, we can conclude that all strings of length  $n$  in  $L$  can be generated by  $G$ .

□

(ii)

*Proof.* Proposition: Every word generated by our grammar,  $G$  has the desired property.

**Base Case:** The empty string can be generated by  $G$ . This is clearly true using the rule  $S \rightarrow \epsilon$ .

**Induction Hypothesis:** Suppose  $S$  is generated by our grammar and is in the language above,  $L$ .

**Induction Step:** We want to prove that applying any of the rules of  $G$ , our new string,  $T$  will remain in  $L$ . By our induction hypothesis, we know that  $S$  will contain at least as many 0's as 1's.

Since  $S \rightarrow 0S$  is simply adding extra 0's to  $S$ , we can confirm that  $T$  is still in  $L$  as the number of 0's in  $T$  will still be greater than or equal to the number of 1's in  $T$  (it will be in  $L$ ). For the rule  $S \rightarrow 0S1$ , we are adding an equal amount of 0's and 1's to  $S$ , so know that the number of 0's in  $T$  will still be greater than or equal to the number of 1's in  $T$  (it will be in  $L$ ). For the rule  $S \rightarrow SS$ , since we know that  $S$  contains at least as many 0's as 1's, if we multiply the number of 0's and the number of 1's in  $S$  by 2,  $T$  will trivially still contain at least as many 0's as 1's (it will be in  $L$ ). Lastly, for the rule  $S \rightarrow \epsilon$ , we are not changing the number of 0's or 1's in  $S$ , so  $T$  will still be in  $L$ .

□

5. (20 points) For each of the following languages construct a context-free grammar that generates that language:

(a)

$$\text{(S)} \quad \{0, 1\}^*. \quad S \rightarrow 0S \mid 1S \mid \epsilon$$

(b)

$$\text{(S)} \quad \{0^m 1^n \mid m \geq n \text{ and } m - n \text{ is even}\}. \quad S \rightarrow 0S1 \mid 00S \mid \epsilon$$

(c) The complement of  $\{0^n 1^n \mid n \geq 0\}$  over the alphabet  $\{0, 1\}$ .

$$\text{(S)} \quad S \rightarrow A0 \mid 1A \mid 0S1 \\ A \rightarrow 0A \mid 1A \mid \epsilon$$

(d) The set of strings in  $\{0, 1\}^*$  which are not palindromes:

$$\{w \in \{0, 1\}^* \mid w \neq w^R\}.$$

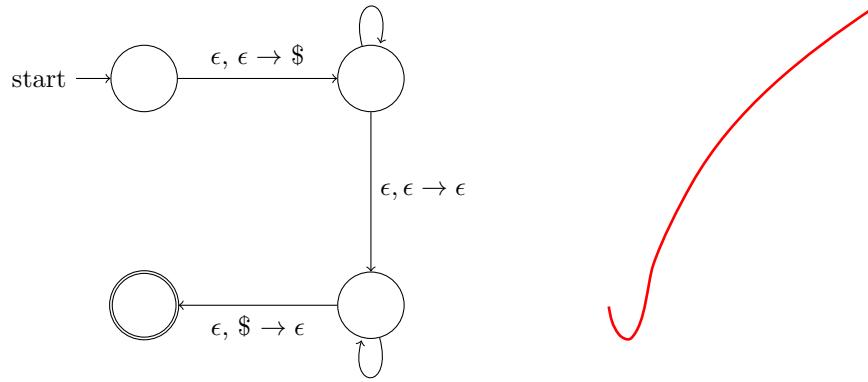
$$\text{(S)} \quad S \rightarrow 0A1 \mid 1A0 \mid 0S0 \mid 1S1 \\ A \rightarrow 0A \mid 1A \mid \epsilon$$

6. (10 points) Use the equivalence of context-free grammars and push-down automata to show that if  $A$  and  $B$  are regular languages, then  $\{xy \mid x \in A, y \in B, |x| = |y|\}$  is context-free.

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Since  $A$  and  $B$  are regular, the following push-down automata will accept this language:

Next letter in string accepted by  $A$ ,  $\epsilon \rightarrow$  Next letter in string accepted by  $A$



Next letter in (different) string accepted by  $B$ , Every letter in the alphabet  $\rightarrow \epsilon$

In this push-down automata, we only accept a string  $x$  which is in  $A$  followed by a string  $y$  which is in  $B$  such that  $|x| = |y|$ . Since push-down automatas are equivalent to context free languages, a context-free language can also be constructed for this language, and hence the language is context-free.