

Red Hat Enterprise Linux 8

Managing, monitoring, and updating the kernel

A guide to managing the Linux kernel on Red Hat Enterprise Linux 8

Last Updated: 2021-08-02

Red Hat Enterprise Linux 8 Managing, monitoring, and updating the kernel

A guide to managing the Linux kernel on Red Hat Enterprise Linux 8

Legal Notice

Copyright © 2021 Red Hat, Inc.

The text of and illustrations in this document are licensed by Red Hat under a Creative Commons Attribution–Share Alike 3.0 Unported license ("CC-BY-SA"). An explanation of CC-BY-SA is available at

http://creativecommons.org/licenses/by-sa/3.0/

. In accordance with CC-BY-SA, if you distribute this document or an adaptation of it, you must provide the URL for the original version.

Red Hat, as the licensor of this document, waives the right to enforce, and agrees not to assert, Section 4d of CC-BY-SA to the fullest extent permitted by applicable law.

Red Hat, Red Hat Enterprise Linux, the Shadowman logo, the Red Hat logo, JBoss, OpenShift, Fedora, the Infinity logo, and RHCE are trademarks of Red Hat, Inc., registered in the United States and other countries.

Linux ® is the registered trademark of Linus Torvalds in the United States and other countries.

Java [®] is a registered trademark of Oracle and/or its affiliates.

XFS [®] is a trademark of Silicon Graphics International Corp. or its subsidiaries in the United States and/or other countries.

MySQL ® is a registered trademark of MySQL AB in the United States, the European Union and other countries.

Node.js ® is an official trademark of Joyent. Red Hat is not formally related to or endorsed by the official Joyent Node.js open source or commercial project.

The OpenStack [®] Word Mark and OpenStack logo are either registered trademarks/service marks or trademarks/service marks of the OpenStack Foundation, in the United States and other countries and are used with the OpenStack Foundation's permission. We are not affiliated with, endorsed or sponsored by the OpenStack Foundation, or the OpenStack community.

All other trademarks are the property of their respective owners.

Abstract

This document provides the users and administrators with necessary information about configuring their workstations on the Linux kernel level. Such adjustments bring performance enhancements, easier troubleshooting or optimized system.

Table of Contents

MAKING OPEN SOURCE MORE INCLUSIVE	6
PROVIDING FEEDBACK ON RED HAT DOCUMENTATION	7
CHAPTER 1. THE LINUX KERNEL RPM 1.1. WHAT AN RPM IS Types of RPM packages 1.2. THE LINUX KERNEL RPM PACKAGE OVERVIEW 1.3. DISPLAYING CONTENTS OF THE KERNEL PACKAGE	8 8 8 9
CHAPTER 2. UPDATING KERNEL WITH YUM 2.1. WHAT IS THE KERNEL 2.2. WHAT IS YUM 2.3. UPDATING THE KERNEL 2.4. INSTALLING THE KERNEL	11 11 11 11 12
3.1. INTRODUCTION TO KERNEL MODULES 3.2. INTRODUCTION TO BOOTLOADER SPECIFICATION 3.3. KERNEL MODULE DEPENDENCIES 3.4. LISTING CURRENTLY LOADED KERNEL MODULES 3.5. LISTING ALL INSTALLED KERNELS 3.6. SETTING A KERNEL AS DEFAULT 3.7. DISPLAYING INFORMATION ABOUT KERNEL MODULES 3.8. LOADING KERNEL MODULES AT SYSTEM RUNTIME 3.9. UNLOADING KERNEL MODULES AT SYSTEM RUNTIME 3.10. UNLOADING KERNEL MODULES AT EARLY STAGES OF THE BOOT PROCESS 3.11. LOADING KERNEL MODULES AUTOMATICALLY AT SYSTEM BOOT TIME 3.12. PREVENTING KERNEL MODULES FROM BEING AUTOMATICALLY LOADED AT SYSTEM BOOT TIME	13 13 14 14 15 16 17 17 18 19 21 22 24
CHAPTER 4. SIGNING KERNEL MODULES FOR SECURE BOOT 4.1. REQUIREMENTS FOR AUTHENTICATING KERNEL MODULES WITH X.509 KEYS 4.2. SOURCES FOR PUBLIC KEYS 4.3. GENERATING A PUBLIC AND PRIVATE KEY PAIR 4.4. EXAMPLE OUTPUT OF SYSTEM KEYRINGS 4.5. ENROLLING PUBLIC KEY ON TARGET SYSTEM BY ADDING THE PUBLIC KEY TO THE MOK LIST 4.6. SIGNING KERNEL MODULES WITH THE PRIVATE KEY	27 28 29 30 31 32 33 34
5.1. UNDERSTANDING KERNEL COMMAND-LINE PARAMETERS5.2. WHAT GRUBBY IS5.3. WHAT BOOT ENTRIES ARE5.4. CHANGING KERNEL COMMAND-LINE PARAMETERS FOR ALL BOOT ENTRIES	36 36 36 36 37 38
6.1. WHAT ARE KERNEL PARAMETERS 6.2. CONFIGURING KERNEL PARAMETERS TEMPORARILY WITH SYSCTL 6.3. CONFIGURING KERNEL PARAMETERS PERMANENTLY WITH SYSCTL	40 40 41 41 42

6.5. CONFIGURING KERNEL PARAMETERS TEMPORARILY THROUGH /PROC/SYS/	43
CHAPTER 7. KEEPING KERNEL PANIC PARAMETERS DISABLED IN VIRTUALIZED ENVIRONMENTS	44
7.1. WHAT IS A SOFT LOCKUP	44
7.2. PARAMETERS CONTROLLING KERNEL PANIC	44
7.3. SPURIOUS SOFT LOCKUPS IN VIRTUALIZED ENVIRONMENTS	45
CHAPTER 8. ADJUSTING KERNEL PARAMETERS FOR DATABASE SERVERS	47
8.1. INTRODUCTION TO DATABASE SERVERS	47
8.2. PARAMETERS AFFECTING PERFORMANCE OF DATABASE APPLICATIONS	47
CHAPTER 9. GETTING STARTED WITH KERNEL LOGGING	49
9.1. WHAT IS THE KERNEL RING BUFFER	49
9.2. ROLE OF PRINTK ON LOG-LEVELS AND KERNEL LOGGING	49
CHAPTER 10. INSTALLING KDUMP	51
10.1. WHAT IS KDUMP	51
10.2. PERFORMING KDUMP INSTALLATION	51
CHAPTER 11. CONFIGURING KDUMP ON THE COMMAND LINE	53
11.1. CONFIGURING KDUMP MEMORY USAGE	53
11.2. CONFIGURING THE KDUMP TARGET	54
11.3. CONFIGURING THE CORE COLLECTOR	57
11.4. CONFIGURING THE KDUMP DEFAULT FAILURE RESPONSES	58
11.5. CONFIGURATION FILE FOR KDUMP	58
11.6. ENABLING AND DISABLING THE KDUMP SERVICE	59
11.7. TESTING THE KDUMP CONFIGURATION	60
11.8. USING KEXEC TO REBOOT THE KERNEL	61
11.9. PREVENTING KERNEL DRIVERS FROM LOADING FOR KDUMP	62
11.10. RUNNING KDUMP ON SYSTEMS WITH ENCRYPTED DISK	63
CHAPTER 12. CONFIGURING KDUMP IN THE WEB CONSOLE	64
12.1. ADDITIONAL RESOURCES	64
12.2. CONFIGURING KDUMP MEMORY USAGE AND TARGET LOCATION IN WEB CONSOLE	64
CHAPTER 13. SUPPORTED KDUMP CONFIGURATIONS AND TARGETS	67
13.1. MEMORY REQUIREMENTS FOR KDUMP	67
13.2. MINIMUM THRESHOLD FOR AUTOMATIC MEMORY RESERVATION	68
13.3. SUPPORTED KDUMP TARGETS	69
13.4. SUPPORTED KDUMP FILTERING LEVELS	70
13.5. SUPPORTED DEFAULT FAILURE RESPONSES	70
13.6. USING FINAL_ACTION PARAMETER	71
13.7. ESTIMATING KDUMP SIZE	71
CHAPTER 14. FIRMWARE ASSISTED DUMP MECHANISMS	73
14.1. FIRMWARE ASSISTED DUMP ON IBM POWERPC HARDWARE	73
14.2. ENABLING FIRMWARE ASSISTED DUMP MECHANISM	73
14.3. FIRMWARE ASSISTED DUMP MECHANISMS ON IBM Z HARDWARE	74
14.4. USING SADUMP ON FUJITSU PRIMEQUEST SYSTEMS	74
CHAPTER 15. ANALYZING A CORE DUMP	76
15.1. INSTALLING THE CRASH UTILITY	76
15.2. RUNNING AND EXITING THE CRASH UTILITY	76
15.3. DISPLAYING VARIOUS INDICATORS IN THE CRASH UTILITY	78
15.4. USING KERNEL OOPS ANALYZER	80
15.5. THE KDUMP HELPER TOOL	81

CHAPTER 16. USING EARLY KDUMP TO CAPTURE BOOT TIME CRASHES	82
16.1. WHAT IS EARLY KDUMP	82
16.2. ENABLING EARLY KDUMP	82
CHAPTER 17. APPLYING PATCHES WITH KERNEL LIVE PATCHING	84
17.1. LIMITATIONS OF KPATCH	84
17.2. SUPPORT FOR THIRD-PARTY LIVE PATCHING	84
17.3. ACCESS TO KERNEL LIVE PATCHES	85
17.4. COMPONENTS OF KERNEL LIVE PATCHING	85
17.5. HOW KERNEL LIVE PATCHING WORKS	85
17.6. SUBSCRIBING THE CURRENTLY INSTALLED KERNELS TO THE LIVE PATCHING STREAM	86
17.7. AUTOMATICALLY SUBSCRIBING ANY FUTURE KERNEL TO THE LIVE PATCHING STREAM	87
17.8. UPDATING KERNEL PATCH MODULES	89
17.9. REMOVING THE LIVE PATCHING PACKAGE	90
17.10. UNINSTALLING THE KERNEL PATCH MODULE	91
17.11. DISABLING KPATCH.SERVICE	92
CHAPTER 18. SETTING LIMITS FOR APPLICATIONS	94
18.1. UNDERSTANDING CONTROL GROUPS	94
18.2. WHAT KERNEL RESOURCE CONTROLLERS ARE	95
18.3. WHAT NAMESPACES ARE	96
18.4. SETTING CPU LIMITS TO APPLICATIONS USING CGROUPS-V1	97
CHAPTER 19. USING CGROUPS-V2 TO SET CPU LIMITS TO APPLICATIONS	101
19.1. MOUNTING CGROUPS-V2	101
19.2. SETTING UP THE CGROUP FOR CPU LIMITING	102
19.3. CONFIGURING CPU LIMITS FOR THE CGROUP	105
CHAPTER 20. USING CONTROL GROUPS VERSION 1 WITH SYSTEMD	108
20.1. ROLE OF SYSTEMD IN CONTROL GROUPS VERSION 1	108
20.2. CREATING TRANSIENT CONTROL GROUPS	109
20.3. CREATING PERSISTENT CONTROL GROUPS	110
20.4. CONFIGURING MEMORY RESOURCE CONTROL SETTINGS ON THE COMMAND-LINE	111
20.5. CONFIGURING MEMORY RESOURCE CONTROL SETTINGS WITH UNIT FILES	111
20.6. REMOVING TRANSIENT CONTROL GROUPS	113
20.7. REMOVING PERSISTENT CONTROL GROUPS	113
20.8. LISTING SYSTEMD UNITS	114
20.9. VIEWING A CONTROL GROUP VERSION 1 HIERARCHY	115
20.10. VIEWING RESOURCE CONTROLLERS	117
20.11. MONITORING RESOURCE CONSUMPTION	117
20.12. CONFIGURING CPUSET CONTROLLER USING SYSTEMD	118
CHAPTER 21. CONFIGURING RESOURCE MANAGEMENT USING CGROUPS VERSION 2 WITH SYSTEMD	
	120
21.1. PREREQUISITES	120
21.2. INTRODUCTION TO RESOURCE DISTRIBUTION MODELS	120
21.3. ALLOCATING CPU RESOURCES USING SYSTEMD	121
21.4. CPU TIME ALLOCATION POLICY OPTIONS FOR SYSTEMD	121
21.5. ALLOCATING MEMORY RESOURCES USING SYSTEMD	122
21.6. MEMORY ALLOCATION CONFIGURATION OPTIONS FOR SYSTEMD	123
21.7. CONFIGURING I/O BANDWIDTH USING SYSTEMD	123
21.8. I/O BANDWIDTH CONFIGURATION OPTIONS FOR SYSTEMD	124
CHAPTER 22. CONFIGURING CPU AFFINITY AND NUMA POLICIES USING SYSTEMD	126
22.1. CONFIGURING CPU AFFINITY USING SYSTEMD	126

22.2. CONFIGURING NUMA USING SYSTEMD	127
22.3. NUMA POLICY CONFIGURATION OPTIONS FOR SYSTEMD	127
CHAPTER 23. ANALYZING SYSTEM PERFORMANCE WITH BPF COMPILER COLLECTION	129
23.1. AN INTRODUCTION TO BCC	129
23.2. INSTALLING THE BCC-TOOLS PACKAGE	129
23.3. USING SELECTED BCC-TOOLS FOR PERFORMANCE ANALYSES	130
Using execsnoop to examine the system processes	130
Using opensnoop to track what files a command opens	131
Using biotop to examine the I/O operations on the disk	131
Using xfsslower to expose unexpectedly slow file system operations	133
CHAPTER 24. ENHANCING SECURITY WITH THE KERNEL INTEGRITY SUBSYSTEM	135
24.1. THE KERNEL INTEGRITY SUBSYSTEM	135
24.2. INTEGRITY MEASUREMENT ARCHITECTURE	136
24.3. EXTENDED VERIFICATION MODULE	136
24.4. TRUSTED AND ENCRYPTED KEYS	136
24.5. WORKING WITH TRUSTED KEYS	137
24.6. WORKING WITH ENCRYPTED KEYS	138
24.7. ENABLING INTEGRITY MEASUREMENT ARCHITECTURE AND EXTENDED VERIFICATION MODULE	
	139
24.8. COLLECTING FILE HASHES WITH INTEGRITY MEASUREMENT ARCHITECTURE	142
24.9. RELATED INFORMATION	143
CHAPTER 25. USING ANSIBLE ROLES TO PERMANENTLY CONFIGURE KERNEL PARAMETERS	144
25.1. INTRODUCTION TO THE KERNEL SETTINGS ROLE	144
25.2. APPLYING SELECTED KERNEL PARAMETERS USING THE KERNEL SETTINGS ROLE	144

MAKING OPEN SOURCE MORE INCLUSIVE

Red Hat is committed to replacing problematic language in our code, documentation, and web properties. We are beginning with these four terms: master, slave, blacklist, and whitelist. Because of the enormity of this endeavor, these changes will be implemented gradually over several upcoming releases. For more details, see our CTO Chris Wright's message.

PROVIDING FEEDBACK ON RED HAT DOCUMENTATION

We appreciate your input on our documentation. Please let us know how we could make it better. To do so:

- For simple comments on specific passages:
 - 1. Make sure you are viewing the documentation in the *Multi-page HTML* format. In addition, ensure you see the **Feedback** button in the upper right corner of the document.
 - 2. Use your mouse cursor to highlight the part of text that you want to comment on.
 - 3. Click the Add Feedback pop-up that appears below the highlighted text.
 - 4. Follow the displayed instructions.
- For submitting more complex feedback, create a Bugzilla ticket:
 - 1. Go to the Bugzilla website.
 - 2. As the Component, use **Documentation**.
 - 3. Fill in the **Description** field with your suggestion for improvement. Include a link to the relevant part(s) of documentation.
 - 4. Click Submit Bug.

CHAPTER 1. THE LINUX KERNEL RPM

The following sections describe the Linux kernel RPM package provided and maintained by Red Hat.

1.1. WHAT AN RPM IS

An RPM package is a file containing other files and their metadata (information about the files that are needed by the system).

Specifically, an RPM package consists of the cpio archive.

The **cpio** archive contains:

- Files
- RPM header (package metadata)
 The rpm package manager uses this metadata to determine dependencies, where to install files, and other information.

Types of RPM packages

There are two types of RPM packages. Both types share the file format and tooling, but have different contents and serve different purposes:

- Source RPM (SRPM)
 An SRPM contains source code and a SPEC file, which describes how to build the source code into a binary RPM. Optionally, the patches to source code are included as well.
- Binary RPM
 A binary RPM contains the binaries built from the sources and patches.

1.2. THE LINUX KERNEL RPM PACKAGE OVERVIEW

The **kernel** RPM is a meta package that does not contain any files, but rather ensures that the following required sub-packages are properly installed:

- **kernel-core** contains the binary image of the kernel, all initramfs-related objects to bootstrap the system, and a minimal number of kernel modules to ensure core functionality. This subpackage alone could be used in virtualized and cloud environments to provide a Red Hat Enterprise Linux 8 kernel with a quick boot time and a small disk size footprint.
- **kernel-modules** contains the remaining kernel modules that are not present in **kernel-core**.

The small set of **kernel** sub-packages above aims to provide a reduced maintenance surface to system administrators especially in virtualized and cloud environments.

Optional kernel packages are for example:

- **kernel-modules-extra** contains kernel modules for rare hardware and modules which loading is disabled by default.
- **kernel-debug** contains a kernel with numerous debugging options enabled for kernel diagnosis, at the expense of reduced performance.
- **kernel-tools** contains tools for manipulating the Linux kernel and supporting documentation.

- **kernel-devel** contains the kernel headers and makefiles sufficient to build modules against the **kernel** package.
- **kernel-abi-whitelists** contains information pertaining to the Red Hat Enterprise Linux kernel ABI, including a list of kernel symbols that are needed by external Linux kernel modules and a **yum** plug-in to aid enforcement.
- **kernel-headers** includes the C header files that specify the interface between the Linux kernel and user-space libraries and programs. The header files define structures and constants that are needed for building most standard programs.

Additional resources

• What are the kernel-core, kernel-modules, and kernel-modules-extras packages?

1.3. DISPLAYING CONTENTS OF THE KERNEL PACKAGE

The following procedure describes how to view the contents of the kernel package and its sub-packages without installing them using the **rpm** command.

Prerequisites

 Obtained kernel, kernel-core, kernel-modules, kernel-modules-extra RPM packages for your CPU architecture

Procedure

• List modules for **kernel**:

```
$ rpm -qlp <kernel_rpm>
(contains no files)
...
```

List modules for kernel-core:

```
$ rpm -qlp < kernel-core_rpm>
...
/lib/modules/4.18.0-80.el8.x86_64/kernel/fs/udf/udf.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/kernel/fs/xfs
/lib/modules/4.18.0-80.el8.x86_64/kernel/fs/xfs.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/kernel/kernel
/lib/modules/4.18.0-80.el8.x86_64/kernel/kernel/trace
/lib/modules/4.18.0-80.el8.x86_64/kernel/kernel/trace/ring_buffer_benchmark.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/kernel/lib
/lib/modules/4.18.0-80.el8.x86_64/kernel/lib/cordic.ko.xz
...
```

List modules for kernel-modules:

```
$ rpm -qlp <kernel-modules_rpm>
...
/lib/modules/4.18.0-80.el8.x86_64/kernel/drivers/infiniband/hw/mlx4/mlx4_ib.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/kernel/drivers/infiniband/hw/mlx5/mlx5_ib.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/kernel/drivers/infiniband/hw/qedr/qedr.ko.xz
```

/lib/modules/4.18.0-80.el8.x86_64/kernel/drivers/infiniband/hw/usnic/usnic_verbs.ko.xz /lib/modules/4.18.0-80.el8.x86_64/kernel/drivers/infiniband/hw/vmw_pvrdma/vmw_pvrdma.ko.xz ...

• List modules for **kernel-modules-extra**:

```
$ rpm -qlp < kernel-modules-extra_rpm>
...
/lib/modules/4.18.0-80.el8.x86_64/extra/net/sched/sch_cbq.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/extra/net/sched/sch_choke.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/extra/net/sched/sch_drr.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/extra/net/sched/sch_dsmark.ko.xz
/lib/modules/4.18.0-80.el8.x86_64/extra/net/sched/sch_gred.ko.xz
...
```

Additional resources

- rpm(8) manual page
- RPM packages

CHAPTER 2. UPDATING KERNEL WITH YUM

The following sections bring information about the Linux kernel provided and maintained by Red Hat (Red Hat kernel), and how to keep the Red Hat kernel updated. As a consequence, the operating system will have all the latest bug fixes, performance enhancements, and patches ensuring compatibility with new hardware.

2.1. WHAT IS THE KERNEL

The kernel is a core part of a Linux operating system, which manages the system resources, and provides interface between hardware and software applications. The Red Hat kernel is a custom-built kernel based on the upstream Linux mainline kernel that Red Hat engineers further develop and harden with a focus on stability and compatibility with the latest technologies and hardware.

Before Red Hat releases a new kernel version, the kernel needs to pass a set of rigorous quality assurance tests.

The Red Hat kernels are packaged in the RPM format so that they are easy to upgrade and verify by the yum package manager.



WARNING

Kernels that have not been compiled by Red Hat are **not** supported by Red Hat.

2.2. WHAT IS YUM

This section refers to description of the **yum** package manager.

Additional resources

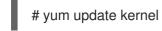
• Configuring basic system settings in Red Hat Enterprise Linux

2.3. UPDATING THE KERNEL

The following procedure describes how to update the kernel using the **yum** package manager.

Procedure

1. To update the kernel, use the following:



This command updates the kernel along with all dependencies to the latest available version.

2. Reboot your system for the changes to take effect.



NOTE

When upgrading from Red Hat Enterprise Linux 7 to Red Hat Enterprise Linux 8, follow relevant sections of the *Upgrading from RHEL 7 to RHEL 8* document.

2.4. INSTALLING THE KERNEL

The following procedure describes how to install new kernels using the **yum** package manager.

Procedure

• To install a specific kernel version, use the following:

yum install kernel-{version}

Additional resources

- Red Hat Code Browser
- Red Hat Enterprise Linux Release Dates

CHAPTER 3. MANAGING KERNEL MODULES

The following sections explain what kernel modules are, how to display their information, and how to perform basic administrative tasks with kernel modules.

3.1. INTRODUCTION TO KERNEL MODULES

The Red Hat Enterprise Linux kernel can be extended with optional, additional pieces of functionality, called kernel modules, without having to reboot the system. On Red Hat Enterprise Linux 8, kernel modules are extra kernel code which is built into compressed **<KERNEL_MODULE_NAME>.ko.xz** object files.

The most common functionality enabled by kernel modules are:

- Device driver which adds support for new hardware
- Support for a file system such as GFS2 or NFS
- System calls

On modern systems, kernel modules are automatically loaded when needed. However, in some cases it is necessary to load or unload modules manually.

Like the kernel itself, the modules can take parameters that customize their behavior if needed.

Tooling is provided to inspect which modules are currently running, which modules are available to load into the kernel and which parameters a module accepts. The tooling also provides a mechanism to load and unload kernel modules into the running kernel.

3.2. INTRODUCTION TO BOOTLOADER SPECIFICATION

The BootLoader Specification (BLS) defines a scheme and the file format to manage the bootloader configuration for each boot option in the drop-in directory without the need to manipulate the bootloader configuration files. Unlike earlier approaches, each boot entry is now represented by a separate configuration file in the drop-in directory. The drop-in directory extends its configuration without having the need to edit or regenerate the configuration files. The BLS extends this concept for the boot menu entries.

Using BLS, you can manage the bootloader menu options by adding, removing, or editing individual boot entry files in a directory. This makes the kernel installation process significantly simpler and consistent across the different architectures.

The **grubby** tool is a thin wrapper script around the BLS and it supports the same **grubby** arguments and options. It runs the **dracut** to create an initial ramdisk image. With this setup, the core bootloader configuration files are static and are not modified after kernel installation.

This premise is particularly relevant in Red Hat Enterprise Linux 8 because the same bootloader is not used in all architectures. **GRUB2** is used in most of them such as the 64-bit ARM, but little-endian variants of IBM Power Systems with Open Power Abstraction Layer (OPAL) uses **Petitboot** and the IBM Z architecture uses **zipl**.

Additional Resources

What grubby is

- What are Boot Entries
- grubby(8) manual page

3.3. KERNEL MODULE DEPENDENCIES

Certain kernel modules sometimes depend on one or more other kernel modules. The /lib/modules/<KERNEL_VERSION>/modules.dep file contains a complete list of kernel module dependencies for the respective kernel version.

The dependency file is generated by the **depmod** program, which is a part of the **kmod** package. Many of the utilities provided by **kmod** take module dependencies into account when performing operations so that **manual** dependency-tracking is rarely necessary.



WARNING

The code of kernel modules is executed in kernel-space in the unrestricted mode. Because of this, you should be mindful of what modules you are loading.

Additional resources

- modules.dep(5) manual page
- depmod(8) manual page

3.4. LISTING CURRENTLY LOADED KERNEL MODULES

The following procedure describes how to view the currently loaded kernel modules.

Prerequisites

• The **kmod** package is installed.

Procedure

• To list all currently loaded kernel modules, execute:

\$ Ismod

```
Size Used by
Module
fuse
            126976 3
           20480 1
uinput
xt_CHECKSUM
                  16384 1
ipt_MASQUERADE
                    16384 1
             16384 1
xt_conntrack
               16384 1
ipt REJECT
             16384 16
nft counter
           16384 0
nf_nat_tftp
nf_conntrack_tftp 16384 1 nf_nat_tftp
tun
            49152 1
```

```
bridge 192512 0
stp 16384 1 bridge
llc 16384 2 bridge,stp
nf_tables_set 32768 5
nft_fib_inet 16384 1
...
```

In the example above:

- The first column provides the **names** of currently loaded modules.
- The second column displays the amount of **memory** per module in kilobytes.
- The last column shows the number, and optionally the names of modules that are **dependent** on a particular module.

Additional resources

- /usr/share/doc/kmod/README file
- Ismod(8) manual page

3.5. LISTING ALL INSTALLED KERNELS

The following procedure describes how to use the command line tool **grubby** to list the **GRUB2** boot entries.

Procedure

To list the boot entries of the kernel:

• To list the boot entries of the kernel, execute:

```
# grubby --info=ALL | grep title
```

The command displays the boot entries of the kernel. The **kernel** field shows the kernel path.

The following procedure describes how to use **grubby** utility to list all installed kernels in their systems using the kernel command line.

As an example, consider listing **grubby-8.40-17**, from the **Grub2** menu on both the BLS and non-BLS installs.

Procedure

To list all installed kernel modules:

• Execute the following command:

```
# grubby --info=ALL | grep title
```

The list of all installed kernels is displayed as follows:

```
title=Red Hat Enterprise Linux (4.18.0-20.el8.x86_64) 8.0 (Ootpa)
title=Red Hat Enterprise Linux (4.18.0-19.el8.x86_64) 8.0 (Ootpa)
title=Red Hat Enterprise Linux (4.18.0-12.el8.x86_64) 8.0 (Ootpa)
title=Red Hat Enterprise Linux (4.18.0) 8.0 (Ootpa)
title=Red Hat Enterprise Linux (0-rescue-2fb13ddde2e24fde9e6a246a942caed1) 8.0 (Ootpa)
```

The above output displays the list of all installed kernels for **grubby-8.40-17**, using the **Grub2** menu.

3.6. SETTING A KERNEL AS DEFAULT

The following procedure describes how to set a specific kernel as default using the **grubby** command-line tool and **GRUB2**.

Procedure

Setting the kernel as default, using the grubby tool

• Execute the following command to set the kernel as default using the **grubby** tool:

grubby --set-default \$kernel path

The command uses a machine ID without the **.conf** suffix as an argument.



NOTE

The machine ID is located in the /boot/loader/entries/ directory.

Setting the kernel as default, using theid argument

• List the boot entries using the id argument and then set an intended kernel as default:

```
# grubby --info ALL | grep id
# grubby --set-default /boot/vmlinuz-<version>.<architecture>
```



NOTE

To list the boot entries using the **title** argument, execute the **# grubby --info=ALL** | **grep title** command.

Setting the default kernel for only the next boot

• Execute the following command to set the default kernel for only the next reboot using the **grub2-reboot** command:

grub2-reboot <index|title|id>



WARNING

Set the default kernel for only the next boot with care. Installing new kernel RPM's, self-built kernels, and manually adding the entries to the /boot/loader/entries/ directory may change the index values.

3.7. DISPLAYING INFORMATION ABOUT KERNEL MODULES

When working with a kernel module, you may want to see further information about that module. This procedure describes how to display extra information about kernel modules.

Prerequisites

• The **kmod** package is installed.

Procedure

• To display information about any kernel module, execute:

\$ modinfo < KERNEL_MODULE_NAME>

For example:

\$ modinfo virtio_net

filename: /lib/modules/4.18.0-94.el8.x86_64/kernel/drivers/net/virtio_net.ko.xz

license: GPL

description: Virtio network driver

rhelversion: 8.1

srcversion: 2E9345B281A898A91319773

alias: virtio:d00000001v* depends: net_failover

intree: Y

name: virtio_net

vermagic: 4.18.0-94.el8.x86_64 SMP mod_unload modversions

• •

parm: napi_weight:int parm: csum:bool parm: gso:bool parm: napi tx:bool

The **modinfo** command displays some detailed information about the specified kernel module. You can query information about all available modules, regardless of whether they are loaded or not. The **parm** entries show parameters the user is able to set for the module, and what type of value they expect.



NOTE

When entering the name of a kernel module, do not append the **.ko.xz** extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

Additional resources

• modinfo(8) manual page

3.8. LOADING KERNEL MODULES AT SYSTEM RUNTIME

The optimal way to expand the functionality of the Linux kernel is by loading kernel modules. The following procedure describes how to use the **modprobe** command to find and load a kernel module into the currently running kernel.

Prerequisites

- Root permissions
- The **kmod** package is installed.
- The respective kernel module is not loaded. To ensure this is the case, list the loaded kernel modules.

Procedure

- Select a kernel module you want to load.
 The modules are located in the /lib/modules/\$(uname -r)/kernel/<SUBSYSTEM>/ directory.
- 2. Load the relevant kernel module:

modprobe < MODULE_NAME>



NOTE

When entering the name of a kernel module, do not append the **.ko.xz** extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

3. Optionally, verify the relevant module was loaded:

\$ Ismod | grep < MODULE_NAME>

If the module was loaded correctly, this command displays the relevant kernel module. For example:

\$ Ismod | grep serio_raw serio_raw 16384 0



IMPORTANT

The changes described in this procedure **will not persist** after rebooting the system. For information on how to load kernel modules to **persist** across system reboots, see Loading kernel modules automatically at system boot time.

Additional resources

modprobe(8) manual page

3.9. UNLOADING KERNEL MODULES AT SYSTEM RUNTIME

At times, you find that you need to unload certain kernel modules from the running kernel. The following procedure describes how to use the **modprobe** command to find and unload a kernel module at system runtime from the currently loaded kernel.

Prerequisites

Root permissions

• The **kmod** package is installed.

Procedure

- Execute the **Ismod** command and select a kernel module you want to unload.
 If a kernel module has dependencies, unload those prior to unloading the kernel module. For details on identifying modules with dependencies, see Listing currently loaded kernel modules and Kernel module dependencies.
- 2. Unload the relevant kernel module:

modprobe -r < MODULE_NAME>

When entering the name of a kernel module, do not append the **.ko.xz** extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.



WARNING

Do not unload kernel modules when they are used by the running system. Doing so can lead to an unstable or non-operational system.

3. Optionally, verify the relevant module was unloaded:

\$ Ismod | grep < MODULE_NAME>

If the module was unloaded successfully, this command does not display any output.



IMPORTANT

After finishing this procedure, the kernel modules that are defined to be automatically loaded on boot, **will not stay unloaded** after rebooting the system. For information on how to counter this outcome, see Preventing kernel modules from being automatically loaded at system boot time.

Additional resources

• modprobe(8) manual page

3.10. UNLOADING KERNEL MODULES AT EARLY STAGES OF THE BOOT PROCESS

In certain situations it is necessary to unload a kernel module very early in the booting process. For example, when the kernel module contains a code, which causes the system to become unresponsive, and the user is not able to reach the stage to permanently disable the rogue kernel module. In that case it is possible to temporarily block the loading of the kernel module using a bootloader.



IMPORTANT

The changes described in this procedure **will not persist** after the next reboot. For information on how to add a kernel module to a denylist so that it will not be automatically loaded during the boot process, see Preventing kernel modules from being automatically loaded at system boot time.

Prerequisites

• You have a loadable kernel module, which you want to prevent from loading for some reason.

Procedure

- Edit the relevant bootloader entry to unload the desired kernel module before the booting sequence continues.
 - Use the cursor keys to highlight the relevant bootloader entry.
 - Press **e** key to edit the entry.

Figure 3.1. Kernel boot menu

```
Red Hat Enterprise Linux (4.18.0-305.10.2.e18_4.x86_64) 8.4 (Ootpa)

Red Hat Enterprise Linux (4.18.0-305.3.1.e18_4.x86_64) 8.4 (Ootpa)

Red Hat Enterprise Linux (4.18.0-240.22.1.e18_3.x86_64) 8.3 (Ootpa)

Red Hat Enterprise Linux (0-rescue-bfeb014063334b2da66b978e95804445) 8.3+

Use the ↑ and ↓ keys to change the selection.

Press 'e' to edit the selected item, or 'c' for a command prompt.
```

- Use the cursor keys to navigate to the line that starts with linux.
- Append **modprobe.blacklist=***module_name* to the end of the line.

Figure 3.2. Kernel boot entry

The **serio_raw** kernel module illustrates a rogue module to be unloaded early in the boot process.

• Press CTRL+x keys to boot using the modified configuration.

Verification

• Once the system fully boots, verify that the relevant kernel module is not loaded.

Ismod | grep serio_raw

Additional resources

Managing kernel modules

3.11. LOADING KERNEL MODULES AUTOMATICALLY AT SYSTEM BOOT TIME

The following procedure describes how to configure a kernel module so that it is loaded automatically during the boot process.

Prerequisites

- Root permissions
- The **kmod** package is installed.

Procedure

- Select a kernel module you want to load during the boot process.
 The modules are located in the /lib/modules/\$(uname -r)/kernel/<SUBSYSTEM>/ directory.
- 2. Create a configuration file for the module:

echo < MODULE_NAME> > /etc/modules-load.d/< MODULE_NAME>.conf



NOTE

When entering the name of a kernel module, do not append the **.ko.xz** extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

3. Optionally, after reboot, verify the relevant module was loaded:

\$ Ismod | grep < MODULE_NAME>

The example command above should succeed and display the relevant kernel module.



IMPORTANT

The changes described in this procedure will persist after rebooting the system.

Additional resources

• modules-load.d(5) manual page

3.12. PREVENTING KERNEL MODULES FROM BEING AUTOMATICALLY LOADED AT SYSTEM BOOT TIME

The following procedure describes how to add a kernel module to a denylist so that it will not be automatically loaded during the boot process.

Prerequisites

- Root permissions
- The **kmod** package is installed.
- Ensure that a kernel module in a denylist is not vital for your current system configuration.

Procedure

1. Select a kernel module that you want to put in a denylist:

\$ Ismod

 Module
 Size
 Used by

 fuse
 126976
 3

 xt_CHECKSUM
 16384
 1

 ipt_MASQUERADE
 16384
 1

 uinput
 20480
 1

 xt_conntrack
 16384
 1

The **Ismod** command displays a list of modules loaded to the currently running kernel.

• Alternatively, identify an unloaded kernel module you want to prevent from potentially loading.

All kernel modules are located in the /lib/modules/<KERNEL_VERSION>/kernel/<SUBSYSTEM>/ directory.

2. Create a configuration file for a denylist:

vim /etc/modprobe.d/blacklist.conf

```
# Blacklists < KERNEL_MODULE_1>
blacklist < MODULE_NAME_1>
install < MODULE_NAME_1> /bin/false
```

Blacklists < KERNEL_MODULE_2> blacklist < MODULE_NAME_2> install < MODULE_NAME_2> /bin/false

Blacklists < KERNEL_MODULE_n> blacklist < MODULE_NAME_n> install < MODULE_NAME_n> /bin/false

The example shows the contents of the **blacklist.conf** file, edited by the **vim** editor. The **blacklist** line ensures that the relevant kernel module will not be automatically loaded during the boot process. The **blacklist** command, however, does not prevent the module from being loaded as a dependency for another kernel module that is not in a denylist. Therefore the **install** line causes the **/bin/false** to run instead of installing a module.

The lines starting with a hash sign are comments to make the file more readable.



NOTE

When entering the name of a kernel module, do not append the **.ko.xz** extension to the end of the name. Kernel module names do not have extensions; their corresponding files do.

3. Create a backup copy of the current initial ramdisk image before rebuilding:

cp /boot/initramfs-\$(uname -r).img /boot/initramfs-\$(uname -r).bak.\$(date +%m-%d-%H%M%S).img

The command above creates a backup **initramfs** image in case the new version has an unexpected problem.

 Alternatively, create a backup copy of other initial ramdisk image which corresponds to the kernel version for which you want to put kernel modules in a denylist:

cp /boot/initramfs-<*SOME_VERSION*>.img /boot/initramfs-<*SOME_VERSION*>.img.bak.\$(date +%m-%d-%H%M%S)

4. Generate a new initial ramdisk image to reflect the changes:

dracut -f -v

• If you are building an initial ramdisk image for a different kernel version than you are currently booted into, specify both target **initramfs** and kernel version:

dracut -f -v /boot/initramfs-<TARGET_VERSION>.img <CORRESPONDING_TARGET_KERNEL_VERSION>

5. Reboot the system:

\$ reboot



IMPORTANT

The changes described in this procedure **will take effect and persist**after rebooting the system. If you improperly put a key kernel module in a denylist, you can face an unstable or non-operational system.

Additional resources

- How do I prevent a kernel module from loading automatically?
- dracut(8) manual page

3.13. COMPILING CUSTOM KERNEL MODULES

You can build a sampling kernel module as requested by various configurations at hardware and software level.

Prerequisites

- You installed the **kernel-devel**, **gcc**, and **elfutils-libelf-devel** packages.
- You have root permissions.
- You created the /root/testmodule/ directory where you compile the custom kernel module.

Procedure

1. Create the /root/testmodule/test.c file with the following content:

```
#include <linux/module.h>
#include <linux/kernel.h>
int init_module(void)
    { printk("Hello World\n This is a test\n"); return 0; }

void cleanup_module(void)
    { printk("Good Bye World"); }
```

The **test.c** file is a source file that provides the main functionality to the kernel module. The file has been created in a dedicated /**root/testmodule**/ directory for organizational purposes. After the module compilation, the /**root/testmodule**/ directory will contain multiple files.

The **test.c** file includes from the system libraries:

• The linux/kernel.h header file is necessary for the printk() function in the example code.

- The linux/module.h file contains function declarations and macro definitions to be shared between several source files written in C programming language.
 Next follow the init_module() and cleanup_module() functions to start and end the kernel logging function printk(), which prints text.
- 2. Create the /root/testmodule/Makefile file with the following content:

```
obj-m := test.o
```

The Makefile contains instructions that the compiler has to produce an object file specifically named **test.o**. The **obj-m** directive specifies that the resulting **test.ko** file is going to be compiled as a loadable kernel module. Alternatively, the **obj-y** directive would instruct to build **test.ko** as a built-in kernel module.

3. Compile the kernel module:

```
# make -C /lib/modules/$(uname -r)/build M=/root/testmodule modules
make: Entering directory '/usr/src/kernels/4.18.0-305.el8.x86_64'
CC [M] /root/testmodule/test.o
Building modules, stage 2.
MODPOST 1 modules
WARNING: modpost: missing MODULE_LICENSE() in /root/testmodule/test.o
see include/linux/module.h for more information
CC /root/testmodule/test.mod.o
LD [M] /root/testmodule/test.ko
make: Leaving directory '/usr/src/kernels/4.18.0-305.el8.x86_64'
```

The compiler creates an object file (**test.o**) for each source file (**test.c**) as an intermediate step before linking them together into the final kernel module (**test.ko**).

After a successful compilation, /root/testmodule/ contains additional files that relate to the compiled custom kernel module. The compiled module itself is represented by the test.ko file.

Verification

1. Optional: check the contents of the /root/testmodule/ directory:

2. Copy the kernel module to the /lib/modules/\$(uname -r)/ directory:

cp /root/testmodule/test.ko /lib/modules/\$(uname -r)/

3. Update the modular dependency list:

depmod -a

4. Load the kernel module:

modprobe -v test insmod /lib/modules/4.18.0-305.el8.x86_64/test.ko

5. Verify that the kernel module was successfully loaded:

Ismod | grep test test 16384 0

6. Read the latest messages from the kernel ring buffer:

dmesg ... [74422.545004] Hello World This is a test

Additional resources

- Managing kernel modules
- Getting started with kernel logging

CHAPTER 4. SIGNING KERNEL MODULES FOR SECURE BOOT

You can enhance the security of your system by using signed kernel modules. The following sections describe how to self-sign privately built kernel modules for use with RHEL 8 on UEFI-based build systems where Secure Boot is enabled. These sections also provide an overview of available options for importing your public key into a target system where you want to deploy your kernel modules.

To sign and load kernel modules, you need to:

If Secure Boot is enabled, the UEFI operating system boot loaders, the Red Hat Enterprise Linux kernel, and all kernel modules have to be signed with a private key and authenticated with the corresponding public key. If they are not signed and authenticated, the system will not be allowed to finish the booting process.

The RHEL 8 distribution includes:

- Signed boot loaders
- Signed kernels
- Signed kernel modules

In addition, the signed first-stage boot loader and the signed kernel include embedded Red Hat public keys. These signed executable binaries and embedded keys enable RHEL 8 to install, boot, and run with the Microsoft UEFI Secure Boot Certification Authority keys that are provided by the UEFI firmware on systems that support UEFI Secure Boot. Note that not all UEFI-based systems include support for Secure Boot.

Prerequisites

To be able to sign externally built kernel modules, install the utilities listed in the following table on the build system.

Table 4.1. Required utilities

Utility	Provided by package	Used on	Purpose
openssl	openssl	Build system	Generates public and private X.509 key pair
sign-file	kernel-devel	Build system	Executable file used to sign a kernel module with the private key
mokutil	mokutil	Target system	Optional utility used to manually enroll the public key
keyctl	keyutils	Target system	Optional utility used to display public keys in the system keyring



NOTE

The build system, where you build and sign your kernel module, does not need to have UEFI Secure Boot enabled and does not even need to be a UEFI-based system.

4.1. REQUIREMENTS FOR AUTHENTICATING KERNEL MODULES WITH X.509 KEYS

In RHEL 8, when a kernel module is loaded, the kernel checks the signature of the module against the public X.509 keys from the kernel system keyring (.builtin_trusted_keys) and the kernel platform keyring (.platform). The .platform keyring contains keys from third-party platform providers and custom public keys. The keys from the kernel system .blacklist keyring are excluded from verification. The following sections provide an overview of sources of keys, keyrings and examples of loaded keys from different sources in the system. Also, you can see how to authenticate a kernel module.

You need to meet certain conditions to load kernel modules on systems with enabled UEFI Secure Boot functionality.

If UEFI Secure Boot is enabled or if the **module.sig_enforce** kernel parameter has been specified:

- You can only load those signed kernel modules whose signatures were authenticated against keys from the system keyring (.builtin_trusted_keys) and the platform keyring (.platform).
- The public key must not be on the system revoked keys keyring (.blacklist).

If UEFI Secure Boot is disabled and the **module.sig_enforce** kernel parameter has not been specified:

• You can load unsigned kernel modules and signed kernel modules without a public key.

If the system is not UEFI-based or if UEFI Secure Boot is disabled:

- Only the keys embedded in the kernel are loaded onto .builtin_trusted_keys and .platform.
- You have no ability to augment that set of keys without rebuilding the kernel.

Table 4.2. Kernel module authentication requirements for loading

Module signed	Public key found and signature valid	UEFI Secure Boot state	sig_enforce	Module load	Kernel tainted
Unsigned	-	Not enabled	Not enabled	Succeeds	Yes
		Not enabled	Enabled	Fails	-
		Enabled	-	Fails	-
Signed	No	Not enabled	Not enabled	Succeeds	Yes
		Not enabled	Enabled	Fails	-
		Enabled	-	Fails	-
Signed	Yes	Not enabled	Not enabled	Succeeds	No
		Not enabled	Enabled	Succeeds	No

Module signed	Public key found and signature valid	UEFI Secure Boot state	sig_enforce	Module load	Kernel tainted
		Enabled	-	Succeeds	No

4.2. SOURCES FOR PUBLIC KEYS

During boot, the kernel loads X.509 keys from a set of persistent key stores into the following keyrings:

- The system keyring (.builtin_trusted_keys)
- The .platform keyring
- The system .blacklist keyring

Table 4.3. Sources for system keyrings

Source of X.509 keys	User can add keys	UEFI Secure Boot state	Keys loaded during boot
Embedded in kernel	No	-	.builtin_trusted_keys
UEFI Secure Boot "db"	Limited	Not enabled	No
		Enabled	.platform
Embedded in shim.efi boot loader	No	Not enabled	No
loudel		Enabled	.builtin_trusted_keys
Machine Owner Key (MOK)	Yes	Not enabled	No
		Enabled	.platform

.builtin_trusted_keys:

- a keyring that is built on boot
- contains trusted public keys
- the **root** privileges are needed to view the keys

.platform:

- a keyring that is built on boot
- contains keys from third-party platform providers and custom public keys
- the **root** privileges are needed to view the keys

.blacklist

- a keyring with X.509 keys which have been revoked
- a module signed by a key from .blacklist will fail authentication even if your public key is in .builtin_trusted_keys

UFFI Secure Boot db:

- a signature database
- stores keys (hashes) of UEFI applications, UEFI drivers, and bootloaders
- the keys can be loaded on the machine

UEFI Secure Boot dbx:

- a revoked signature database
- prevents keys from being loaded
- the revoked keys from this database are added to the .blacklist keyring

4.3. GENERATING A PUBLIC AND PRIVATE KEY PAIR

You need to generate a public and private X.509 key pair to succeed in your efforts of using kernel modules on a Secure Boot-enabled system. You will later use the private key to sign the kernel module. You will also have to add the corresponding public key to the Machine Owner Key (MOK) for Secure Boot to validate the signed module.

Some of the parameters for this key pair generation are best specified with a configuration file.

Procedure

1. Create a configuration file with parameters for the key pair generation:

```
# cat << EOF > configuration_file.config
[req]
default_bits = 4096
distinguished_name = req_distinguished_name
prompt = no
string_mask = utf8only
x509_extensions = myexts
[req distinguished name]
O = Organization
CN = Organization signing key
emailAddress = E-mail address
[myexts]
basicConstraints=critical,CA:FALSE
keyUsage=digitalSignature
subjectKeyIdentifier=hash
authorityKeyIdentifier=keyid
EOF
```

2. Create an X.509 public and private key pair as shown in the following example:

openssl req -x509 -new -nodes -utf8 -sha256 -days 36500 \

- -batch -config configuration_file.config -outform DER \
- -out my_signing_key_pub.der \
- -keyout my_signing_key.priv

The public key will be written to the **my_signing_key_pub.der** file and the private key will be written to the **my_signing_key.priv** file.



IMPORTANT

In RHEL 8, the validity dates of the key pair matter. The key does not expire, but the kernel module must be signed within the validity period of its signing key. For example, a key that is only valid in 2019 can be used to authenticate a kernel module signed in 2019 with that key. However, users cannot use that key to sign a kernel module in 2020.

3. Optionally, you can review the validity dates of your public keys like in the example below:

openssl x509 -inform der -text -noout -in <my_signing_key_pub.der>

Validity

Not Before: Feb 14 16:34:37 2019 GMT Not After: Feb 11 16:34:37 2029 GMT

4. Enroll your public key on all systems where you want to authenticate and load your kernel module.



WARNING

Apply strong security measures and access policies to guard the contents of your private key. In the wrong hands, the key could be used to compromise any system which is authenticated by the corresponding public key.

Additional resources

- openssl(1) manual page
- RHEL Security Guide
- Enrolling public key on target system by adding the public key to the MOK list

4.4. EXAMPLE OUTPUT OF SYSTEM KEYRINGS

You can display information about the keys on the system keyrings using the **keyctl** utility.

The following is a shortened example output of **.builtin_trusted_keys**, **.platform**, and **.blacklist** keyrings from a RHEL 8 system where UEFI Secure Boot is enabled.

keyctl list %:.builtin trusted keys

- 6 keys in keyring:
- ...asymmetric: Red Hat Enterprise Linux Driver Update Program (key 3): bf57f3e87...
- ...asymmetric: Red Hat Secure Boot (CA key 1): 4016841644ce3a810408050766e8f8a29...
- ...asymmetric: Microsoft Corporation UEFI CA 2011: 13adbf4309bd82709c8cd54f316ed...
- ...asymmetric: Microsoft Windows Production PCA 2011: a92902398e16c49778cd90f99e...
- ...asymmetric: Red Hat Enterprise Linux kernel signing key: 4249689eefc77e95880b...
- ...asymmetric: Red Hat Enterprise Linux kpatch signing key: 4d38fd864ebe18c5f0b7...

keyctl list %:.platform

- 4 keys in keyring:
- ...asymmetric: VMware, Inc.: 4ad8da0472073...
- ...asymmetric: Red Hat Secure Boot CA 5: cc6fafe72...
- ...asymmetric: Microsoft Windows Production PCA 2011: a929f298e1...
- ...asymmetric: Microsoft Corporation UEFI CA 2011: 13adbf4e0bd82...

keyctl list %:.blacklist

- 4 keys in keyring:
- ...blacklist: bin:f5ff83a...
- ...blacklist: bin:0dfdbec...
- ...blacklist: bin:38f1d22...
- ...blacklist: bin:51f831f...

The .builtin_trusted_keys keyring above shows the addition of two keys from the UEFI Secure Boot "db" keys as well as the **Red Hat Secure Boot (CA key 1)**, which is embedded in the **shim.efi** boot loader.

The following example shows the kernel console output. The messages identify the keys with an UEFI Secure Boot related source. These include UEFI Secure Boot db, embedded shim, and MOK list.

dmesg | grep 'EFI: Loaded cert'

[5.160660] EFI: Loaded cert 'Microsoft Windows Production PCA 2011: a9290239...

[5.160674] EFI: Loaded cert 'Microsoft Corporation UEFI CA 2011: 13adbf4309b...

[5.165794] EFI: Loaded cert 'Red Hat Secure Boot (CA key 1): 4016841644ce3a8...

Additional resources

keyctl(1), dmesg(1) manual pages

4.5. ENROLLING PUBLIC KEY ON TARGET SYSTEM BY ADDING THE PUBLIC KEY TO THE MOK LIST

When RHEL 8 boots on a UEFI-based system with Secure Boot enabled, the kernel loads onto the system keyring (.builtin_trusted_keys) all public keys that are in the Secure Boot db key database. At the same time the kernel excludes the keys in the dbx database of revoked keys. The sections below describe different ways of importing a public key on a target system so that the system keyring (.builtin_trusted_keys) is able to use the public key to authenticate a kernel module.

The Machine Owner Key (MOK) facility feature can be used to expand the UEFI Secure Boot key database. When RHEL 8 boots on a UEFI-enabled system with Secure Boot enabled, the keys on the MOK list are also added to the system keyring (.builtin_trusted_keys) in addition to the keys from the key database. The MOK list keys are also stored persistently and securely in the same fashion as the Secure Boot database keys, but these are two separate facilities. The MOK facility is supported by shim.efi, MokManager.efi, grubx64.efi, and the mokutil utility.

Enrolling a MOK key requires manual interaction by a user at the UEFI system console on each target system. Nevertheless, the MOK facility provides a convenient method for testing newly generated key pairs and testing kernel modules signed with them.

Procedure

1. Request the addition of your public key to the MOK list:

mokutil --import my_signing_key_pub.der

You will be asked to enter and confirm a password for this MOK enrollment request.

- 2. Reboot the machine.
 - The pending MOK key enrollment request will be noticed by **shim.efi** and it will launch **MokManager.efi** to allow you to complete the enrollment from the UEFI console.
- 3. Enter the password you previously associated with this request and confirm the enrollment. Your public key is added to the MOK list, which is persistent.

Once a key is on the MOK list, it will be automatically propagated to the system keyring on this and subsequent boots when UEFI Secure Boot is enabled.



NOTE

To facilitate authentication of your kernel module on your systems, consider requesting your system vendor to incorporate your public key into the UEFI Secure Boot key database in their factory firmware image.

4.6. SIGNING KERNEL MODULES WITH THE PRIVATE KEY

Users are able to obtain enhanced security benefits on their systems by loading signed kernel modules if the UEFI Secure Boot mechanism is enabled. The following sections describe how to sign kernel modules with the private key.

Prerequisites

- You generated a public and private key pair and know the validity dates of your public keys. For details, see Generating a public and private key pair.
- You enrolled your public key on the target system. For details, see Enrolling public key on target system by adding the public key to the MOK list.
- You have a kernel module in ELF image format available for signing.

Procedure

Execute the sign-file utility with parameters as shown in the example below:

/usr/src/kernels/\$(uname -r)/scripts/sign-file sha256 my_signing_key.priv my_signing_key_pub.der my_module.ko

sign-file computes and appends the signature directly to the ELF image in your kernel module file. The **modinfo** utility can be used to display information about the kernel module's signature, if it is present.



NOTE

The appended signature is not contained in an ELF image section and is not a formal part of the ELF image. Therefore, utilities such as **readelf** will not be able to display the signature on your kernel module.

Your kernel module is now ready for loading. Note that your signed kernel module is also loadable on systems where UEFI Secure Boot is disabled or on a non-UEFI system. That means you do not need to provide both a signed and unsigned version of your kernel module.



IMPORTANT

In RHEL 8, the validity dates of the key pair matter. The key does not expire, but the kernel module must be signed within the validity period of its signing key. The **sign-file** utility will not warn you of this. For example, a key that is only valid in 2019 can be used to authenticate a kernel module signed in 2019 with that key. However, users cannot use that key to sign a kernel module in 2020.

Additional resources

• Displaying information about kernel modules

4.7. LOADING SIGNED KERNEL MODULES

Once your public key is enrolled in the system keyring (.builtin_trusted_keys) and the MOK list, and after you have signed the respective kernel module with your private key, you can finally load your signed kernel module with the the **modprobe** command as described in the following section.

Prerequisites

- You have generated the public and private keypair. For details, see Generating a public and private key pair.
- You have enrolled the public key into the system keyring. For details, see Enrolling public key on target system by adding the public key to the MOK list.
- You have signed a kernel module with the private key. For details, see Signing kernel modules with the private key.

Procedure

1. Verify that your public keys are on the system keyring:

keyctl list %:.builtin_trusted_keys

- 2. Copy the kernel module into the /extra/ directory of the kernel you want:
 - # cp my_module.ko /lib/modules/\$(uname -r)/extra/
- 3. Update the modular dependency list:

depmod -a

4. Load the kernel module and verify that it was successfully loaded:

modprobe -v my_module # Ismod | grep my_module

a. Optionally, to load the module on boot, add it to the /etc/modules-loaded.d/my_module.conf file:

echo "my_module" > /etc/modules-load.d/my_module.conf

Additional resources

Managing kernel modules

CHAPTER 5. CONFIGURING KERNEL COMMAND-LINE PARAMETERS

Kernel command-line parameters are a way to change the behavior of certain aspects of the Red Hat Enterprise Linux kernel at boot time. As a system administrator, you have full control over what options get set at boot. Certain kernel behaviors are only able to be set at boot time, so understanding how to make this changes is a key administration skill.



IMPORTANT

Opting to change the behavior of the system by modifying kernel command-line parameters may have negative effects on your system. You should therefore test changes prior to deploying them in production. For further guidance, contact Red Hat Support.

5.1. UNDERSTANDING KERNEL COMMAND-LINE PARAMETERS

Kernel command-line parameters are used for boot time configuration of:

- The Red Hat Enterprise Linux kernel
- The initial RAM disk
- The user space features

Kernel boot time parameters are often used to overwrite default values and for setting specific hardware settings.

By default, the kernel command-line parameters for systems using the GRUB2 bootloader are defined in the **kernelopts** variable of the /**boot/grub2/grubenv** file for all kernel boot entries.



NOTE

For IBM Z, the kernel command-line parameters are stored in the boot entry configuration file because the zipl bootloader does not support environment variables. Thus, the **kernelopts** environment variable cannot be used.

Additional resources

- kernel-command-line(7), bootparam(7) and dracut.cmdline(7) manual pages
- How to install and boot custom kernels in Red Hat Enterprise Linux 8

5.2. WHAT GRUBBY IS

grubby is a utility for manipulating bootloader-specific configuration files.

You can use **grubby** also for changing the default boot entry, and for adding/removing arguments from a GRUB2 menu entry.

For more details see the **grubby(8)** manual page.

5.3. WHAT BOOT ENTRIES ARE

A boot entry is a collection of options which are stored in a configuration file and tied to a particular kernel version. In practice, you have at least as many boot entries as your system has installed kernels. The boot entry configuration file is located in the /boot/loader/entries/ directory and can look like this:

6f9cc9cb7d7845d49698c9537337cedc-4.18.0-5.el8.x86_64.conf

The file name above consists of a machine ID stored in the /etc/machine-id file, and a kernel version.

The boot entry configuration file contains information about the kernel version, the initial ramdisk image, and the **kernelopts** environment variable, which contains the kernel command-line parameters. The contents of a boot entry config can be seen below:

title Red Hat Enterprise Linux (4.18.0-74.el8.x86_64) 8.0 (Ootpa) version 4.18.0-74.el8.x86_64 linux /vmlinuz-4.18.0-74.el8.x86_64 initrd /initramfs-4.18.0-74.el8.x86_64.img \$tuned_initrd options \$kernelopts \$tuned_params id rhel-20190227183418-4.18.0-74.el8.x86_64 grub_users \$grub_users \$grub_users grub_arg --unrestricted grub_class kernel

The **kernelopts** environment variable is defined in the /boot/grub2/grubenv file.

Additional resources

• How to install and boot custom kernels in Red Hat Enterprise Linux 8

5.4. CHANGING KERNEL COMMAND-LINE PARAMETERS FOR ALL BOOT ENTRIES

This procedure describes how to change kernel command-line parameters for all boot entries on your system.

Prerequisites

• Verify that the **grubby** and **zipl** utilities are installed on your system.

Procedure

To add a parameter:

grubby --update-kernel=ALL --args="<NEW_PARAMETER>"

For systems that use the GRUB2 bootloader, the command updates the /boot/grub2/grubenv file by adding a new kernel parameter to the **kernelopts** variable in that file.

On IBM Z that use the zIPL bootloader, the command adds a new kernel parameter to each /boot/loader/entries/<ENTRY>.conf file.

- On IBM Z, execute the **zipl** command with no options to update the boot menu.
- To remove a parameter:

grubby --update-kernel=ALL --remove-args="<PARAMETER_TO_REMOVE>"

• On IBM Z, execute the **zipl** command with no options to update the boot menu.

Additional resources

- Understanding kernel command-line parameters
- grubby(8) and zipl(8) manual pages
- grubby tool

5.5. CHANGING KERNEL COMMAND-LINE PARAMETERS FOR A SINGLE BOOT ENTRY

This procedure describes how to change kernel command-line parameters for a single boot entry on your system.

Prerequisites

• Verify that the **grubby** and **zipl** utilities are installed on your system.

Procedure

- To add a parameter:
 - # grubby --update-kernel=/boot/vmlinuz-\$(uname -r) --args="<*NEW_PARAMETER*>"
 - On IBM Z, execute the **zipl** command with no options to update the boot menu.
- To remove a parameter use the following:

grubby --update-kernel=/boot/vmlinuz-\$(uname -r) --remove-args=" <PARAMETER_TO_REMOVE>"

• On IBM Z, execute the **zipl** command with no options to update the boot menu.



NOTE

On systems that use the **grub.cfg** file, there is, by default, the **options** parameter for each kernel boot entry, which is set to the **kernelopts** variable. This variable is defined in the **/boot/grub2/grubenv** configuration file.



IMPORTANT

On GRUB2 systems:

- If the kernel command-line parameters are modified for all boot entries, the **grubby** utility updates the **kernelopts** variable in the /boot/grub2/grubenv file.
- If kernel command-line parameters are modified for a single boot entry, the kernelopts variable is expanded, the kernel parameters are modified, and the resulting value is stored in the respective boot entry's /boot/loader/entries/<RELEVANT_KERNEL_BOOT_ENTRY.conf> file.

On zIPL systems:

• **grubby** modifies and stores the kernel command-line parameters of an individual kernel boot entry in the /boot/loader/entries/<*ENTRY*>.conf file.

Additional resources

- Understanding kernel command-line parameters
- grubby(8) and zipl(8) manual pages
- grubby tool

CHAPTER 6. CONFIGURING KERNEL PARAMETERS AT RUNTIME

As a system administrator, you can modify many facets of the Red Hat Enterprise Linux kernel's behavior at runtime. This section describes how to configure kernel parameters at runtime by using the **sysctl** command and by modifying the configuration files in the /etc/sysctl.d/ and /proc/sys/ directories.

6.1. WHAT ARE KERNEL PARAMETERS

Kernel parameters are tunable values which you can adjust while the system is running. There is no requirement to reboot or recompile the kernel for changes to take effect.

It is possible to address the kernel parameters through:

- The **sysctl** command
- The virtual file system mounted at the /proc/sys/ directory
- The configuration files in the /etc/sysctl.d/ directory

Tunables are divided into classes by the kernel subsystem. Red Hat Enterprise Linux has the following tunable classes:

Table 6.1. Table of sysctl classes

Tunable class	Subsystem
abi	Execution domains and personalities
crypto	Cryptographic interfaces
debug	Kernel debugging interfaces
dev	Device-specific information
fs	Global and specific file system tunables
kernel	Global kernel tunables
net	Network tunables
sunrpc	Sun Remote Procedure Call (NFS)
user	User Namespace limits
vm	Tuning and management of memory, buffers, and cache



IMPORTANT

Configuring kernel parameters on a production system requires careful planning. Unplanned changes may render the kernel unstable, requiring a system reboot. Verify that you are using valid options before changing any kernel values.

Additional resources

• sysctl(8), and sysctl.d(5) manual pages

6.2. CONFIGURING KERNEL PARAMETERS TEMPORARILY WITH SYSCTL

The following procedure describes how to use the **sysctl** command to temporarily set kernel parameters at runtime. The command is also useful for listing and filtering tunables.

Prerequisites

Root permissions

Procedure

1. To list all parameters and their values, use the following:





NOTE

The **# sysctl -a** command displays kernel parameters, which can be adjusted at runtime and at boot time.

2. To configure a parameter temporarily, use the command as in the following example:



The sample command above changes the parameter value while the system is running. The changes take effect immediately, without a need for restart.



NOTE

The changes return back to default after your system reboots.

Additional resources

- sysctl(8) manual page
- Configuring kernel parameters permanently with sysctl
- Using configuration files in /etc/sysctl.d/ to adjust kernel parameters

6.3. CONFIGURING KERNEL PARAMETERS PERMANENTLY WITH SYSCTL

The following procedure describes how to use the **sysctl** command to permanently set kernel parameters.

Prerequisites

Root permissions

Procedure

1. To list all parameters, use the following:

sysctl -a

The command displays all kernel parameters that can be configured at runtime.

2. To configure a parameter permanently:

sysctl -w <TUNABLE_CLASS>.<PARAMETER>=<TARGET_VALUE> >> /etc/sysctl.conf

The sample command changes the tunable value and writes it to the /etc/sysctl.conf file, which overrides the default values of kernel parameters. The changes take effect immediately and persistently, without a need for restart.



NOTE

To permanently modify kernel parameters you can also make manual changes to the configuration files in the /etc/sysctl.d/ directory.

Additional resources

- sysctl(8) and sysctl.conf(5) manual pages
- Using configuration files in /etc/sysctl.d/ to adjust kernel parameters

6.4. USING CONFIGURATION FILES IN /ETC/SYSCTL.D/ TO ADJUST KERNEL PARAMETERS

The following procedure describes how to manually modify configuration files in the /etc/sysctl.d/ directory to permanently set kernel parameters.

Prerequisites

Root permissions

Procedure

1. Create a new configuration file in /etc/sysctl.d/:

vim /etc/sysctl.d/<some_file.conf>

2. Include kernel parameters, one per line, as follows:

<TUNABLE_CLASS>.<PARAMETER>=<TARGET_VALUE><TUNABLE_CLASS>.<PARAMETER>=<TARGET_VALUE>

- 3. Save the configuration file.
- 4. Reboot the machine for the changes to take effect.
 - Alternatively, to apply changes without rebooting, execute:

sysctl -p /etc/sysctl.d/<some_file.conf>

The command enables you to read values from the configuration file, which you created earlier.

Additional resources

• sysctl(8), sysctl.d(5) manual pages

6.5. CONFIGURING KERNEL PARAMETERS TEMPORARILY THROUGH /PROC/SYS/

The following procedure describes how to set kernel parameters temporarily through the files in the virtual file system /proc/sys/ directory.

Prerequisites

Root permissions

Procedure

1. Identify a kernel parameter you want to configure:

Is -I /proc/sys/< TUNABLE_CLASS>/

The writable files returned by the command can be used to configure the kernel. The files with read-only permissions provide feedback on the current settings.

2. Assign a target value to the kernel parameter:

echo < TARGET_VALUE> > /proc/sys/< TUNABLE_CLASS>/< PARAMETER>

The command makes configuration changes that will disappear once the system is restarted.

3. Optionally, verify the value of the newly set kernel parameter:

cat /proc/sys/<TUNABLE_CLASS>/<PARAMETER>

Additional resources

- Configuring kernel parameters permanently with sysctl
- Using configuration files in /etc/sysctl.d/ to adjust kernel parameters

CHAPTER 7. KEEPING KERNEL PANIC PARAMETERS DISABLED IN VIRTUALIZED ENVIRONMENTS

When configuring a virtualized environment in Red Hat Enterprise Linux 8 (RHEL 8), you should not enable the **softlockup_panic** and **nmi_watchdog** kernel parameters, as the virtualized environment may trigger a spurious soft lockup that should not require a system panic.

The following sections explain the reasons behind this advice by summarizing:

- What causes a soft lockup.
- Describing the kernel parameters that control a system's behavior on a soft lockup.
- Explaining how soft lockups may be triggered in a virtualized environment.

7.1. WHAT IS A SOFT LOCKUP

A soft lockup is a situation usually caused by a bug, when a task is executing in kernel space on a CPU without rescheduling. The task also does not allow any other task to execute on that particular CPU. As a result, a warning is displayed to a user through the system console. This problem is also referred to as the soft lockup firing.

Additional resources

• What is a CPU soft lockup?

7.2. PARAMETERS CONTROLLING KERNEL PANIC

The following kernel parameters can be set to control a system's behavior when a soft lockup is detected.

softlockup_panic

Controls whether or not the kernel will panic when a soft lockup is detected.

Туре	Value	Effect
Integer	0	kernel does not panic on soft lockup
Integer	1	kernel panics on soft lockup

By default, on RHEL8 this value is 0.

In order to panic, the system needs to detect a hard lockup first. The detection is controlled by the **nmi_watchdog** parameter.

nmi watchdog

Controls whether lockup detection mechanisms (**watchdogs**) are active or not. This parameter is of integer type.

Value	Effect
0	disables lockup detector
1	enables lockup detector

The hard lockup detector monitors each CPU for its ability to respond to interrupts.

watchdog_thresh

Controls frequency of watchdog hrtimer, NMI events, and soft/hard lockup thresholds.

Default threshold	Soft lockup threshold
10 seconds	2 * watchdog_thresh

Setting this parameter to zero disables lockup detection altogether.

Additional resources

- Softlockup detector and hardlockup detector
- Kernel sysctl

7.3. SPURIOUS SOFT LOCKUPS IN VIRTUALIZED ENVIRONMENTS

The soft lockup firing on physical hosts, as described in What is a soft lockup, usually represents a kernel or hardware bug. The same phenomenon happening on guest operating systems in virtualized environments may represent a false warning.

Heavy work-load on a host or high contention over some specific resource such as memory, usually causes a spurious soft lockup firing. This is because the host may schedule out the guest CPU for a period longer than 20 seconds. Then when the guest CPU is again scheduled to run on the host, it experiences a *time jump* which triggers due timers. The timers include also watchdog **hrtimer**, which can consequently report a soft lockup on the guest CPU.

Because a soft lockup in a virtualized environment may be spurious, you should not enable the kernel parameters that would cause a system panic when a soft lockup is reported on a guest CPU.



IMPORTANT

To understand soft lockups in guests, it is essential to know that the host schedules the guest as a task, and the guest then schedules its own tasks.

Additional resources

- What is a soft lockup
- RHEL 8 virtual machine components and their interaction

• Virtual machine reports a "BUG: soft lockup"

CHAPTER 8. ADJUSTING KERNEL PARAMETERS FOR DATABASE SERVERS

There are different sets of kernel parameters which can affect performance of specific database applications. The following sections explain what kernel parameters to configure to secure efficient operation of database servers and databases.

8.1. INTRODUCTION TO DATABASE SERVERS

A database server is a service that provides features of a database management system (DBMS). DBMS provides utilities for database administration and interacts with end users, applications, and databases.

Red Hat Enterprise Linux 8 provides the following database management systems:

- MariaDB 10.3
- MariaDB 10.5 available since RHEL 8.4
- MySQL 8.0
- PostgreSQL 10
- PostgreSQL 9.6
- PostgreSQL 12 available since RHEL 8.1.1
- PostgreSQL 13 available since RHEL 8.4

8.2. PARAMETERS AFFECTING PERFORMANCE OF DATABASE APPLICATIONS

The following kernel parameters affect performance of database applications.

fs.aio-max-nr

Defines the maximum number of asynchronous I/O operations the system can handle on the server.



NOTE

Raising the **fs.aio-max-nr** parameter produces no additional changes beyond increasing the aio limit.

fs.file-max

Defines the maximum number of file handles (temporary file names or IDs assigned to open files) the system supports at any instance.

The kernel dynamically allocates file handles whenever a file handle is requested by an application. The kernel however does not free these file handles when they are released by the application. The kernel recycles these file handles instead. This means that over time the total number of allocated file handles will increase even though the number of currently used file handles may be low.

kernel.shmall

Defines the total number of shared memory pages that can be used system-wide. To use the entire main memory, the value of the **kernel.shmall** parameter should be \leq total main memory size.

kernel.shmmax

Defines the maximum size in bytes of a single shared memory segment that a Linux process can allocate in its virtual address space.

kernel.shmmni

Defines the maximum number of shared memory segments the database server is able to handle.

net.ipv4.ip_local_port_range

Defines the port range the system can use for programs which want to connect to a database server without a specific port number.

net.core.rmem_default

Defines the default receive socket memory through Transmission Control Protocol (TCP).

net.core.rmem_max

Defines the maximum receive socket memory through Transmission Control Protocol (TCP).

net.core.wmem_default

Defines the default send socket memory through Transmission Control Protocol (TCP).

net.core.wmem_max

Defines the maximum send socket memory through Transmission Control Protocol (TCP).

vm.dirty_bytes / vm.dirty_ratio

Defines a threshold in bytes / in percentage of dirty-able memory at which a process generating dirty data is started in the **write()** function.



NOTE

Either vm.dirty_bytes or **vm.dirty_ratio** can be specified at a time.

vm.dirty_background_bytes / vm.dirty_background_ratio

Defines a threshold in bytes / in percentage of dirty-able memory at which the kernel tries to actively write dirty data to hard-disk.



NOTE

Either vm.dirty_background_bytes or vm.dirty_background_ratio can be specified at a time.

vm.dirty_writeback_centisecs

Defines a time interval between periodic wake-ups of the kernel threads responsible for writing dirty data to hard-disk.

This kernel parameters measures in 100th's of a second.

vm.dirty_expire_centisecs

Defines the time after which dirty data is old enough to be written to hard-disk.

This kernel parameters measures in 100th's of a second.

Additional resources

Dirty pagecache writeback and vm.dirty parameters

CHAPTER 9. GETTING STARTED WITH KERNEL LOGGING

Log files are files that contain messages about the system, including the kernel, services, and applications running on it. The logging system in Red Hat Enterprise Linux is based on the built-in **syslog** protocol. Various utilities use this system to record events and organize them into log files. These files are useful when auditing the operating system or troubleshooting problems.

9.1. WHAT IS THE KERNEL RING BUFFER

During the boot process, the console provides a lot of important information about the initial phase of the system startup. To avoid loss of the early messages the kernel utilizes what is called a ring buffer. This buffer stores all messages, including boot messages, generated by the **printk()** function within the kernel code. The messages from the kernel ring buffer are then read and stored in log files on permanent storage, for example, by the **syslog** service.

The buffer mentioned above is a cyclic data structure which has a fixed size, and is hard-coded into the kernel. Users can display data stored in the kernel ring buffer through the **dmesg** command or the /var/log/boot.log file. When the ring buffer is full, the new data overwrites the old.

Additional resources

• syslog(2) and dmesg(1) manual page

9.2. ROLE OF PRINTK ON LOG-LEVELS AND KERNEL LOGGING

Each message the kernel reports has a log-level associated with it that defines the importance of the message. The kernel ring buffer, as described in What is the kernel ring buffer, collects kernel messages of all log-levels. It is the **kernel.printk** parameter that defines what messages from the buffer are printed to the console.

The log-level values break down in this order:

- \bullet 0 Kernel emergency. The system is unusable.
- 1 Kernel alert. Action must be taken immediately.
- 2 Condition of the kernel is considered critical.
- 3 General kernel error condition.
- 4 General kernel warning condition.
- 5 Kernel notice of a normal but significant condition.
- 6 Kernel informational message.
- 7 Kernel debug-level messages.

By default, kernel.printk in RHEL 8 contains the following four values:

sysctl kernel.printk kernel.printk = 7 4 1 7

The four values define the following:

- 1. value. Console log-level, defines the lowest priority of messages printed to the console.
- 2. value. Default log-level for messages without an explicit log-level attached to them.
- 3. value. Sets the lowest possible log-level configuration for the console log-level.
- 4. value. Sets default value for the console log-level at boot time.

 Each of these values above defines a different rule for handling error messages.



IMPORTANT

The default **7 417printk** value allows for better debugging of kernel activity. However, when coupled with a serial console, this **printk** setting is able to cause intense I/O bursts that could lead to a RHEL system becoming temporarily unresponsive. To avoid these situations, setting a **printk** value of **4 417**typically works, but at the expense of losing the extra debugging information.

Also note that certain kernel command line parameters, such as **quiet** or **debug**, change the default **kernel.printk** values.

Additional resources

• syslog(2) manual page

CHAPTER 10. INSTALLING KDUMP

In many cases, the **kdump** service is installed and activated by default on the new Red Hat Enterprise Linux installations. The following sections brings information about **kdump** and how to install **kdump** when it is not enabled by default in some cases.

10.1. WHAT IS KDUMP

kdump is a service providing a crash dumping mechanism. The service enables you to save the contents of the system's memory for later analysis. **kdump** uses the **kexec** system call to boot into the second kernel (a *capture kernel*) without rebooting; and then captures the contents of the crashed kernel's memory (a *crash dump* or a *vmcore*) and saves it. The second kernel resides in a reserved part of the system memory.

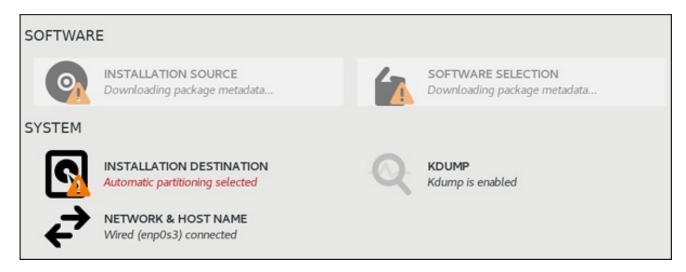


IMPORTANT

A kernel crash dump can be the only information available in the event of a system failure (a critical bug). Therefore, ensuring that **kdump** is operational is important in mission-critical environments. Red Hat advise that system administrators regularly update and test **kexec-tools** in your normal kernel update cycle. This is especially important when new kernel features are implemented.

10.2. PERFORMING KDUMP INSTALLATION

The **Anaconda** installer provides a screen for **kdump** configuration when performing an interactive installation using the graphical or text interface. The installer screen is titled **Kdump** and is available from the main **Installation Summary** screen, and only allows limited configuration – you can only select whether **kdump** is enabled and how much memory is reserved.



Some installation options, such as custom Kickstart installations, in some cases do not install or enable **kdump** by default. If this is the case on your system, follow the procedure below to install **kdump**.

Prerequisites

- An active Red Hat Enterprise Linux subscription
- A repository containing the **kexec-tools** package for your system CPU architecture
- Fulfilled requirements for **kdump** configurations and targets

Procedure

1. Execute the following command to check whether **kdump** is installed on your system:

\$ rpm -q kexec-tools

Output if the package is installed:

kexec-tools-2.0.17-11.el8.x86_64

Output if the package is not installed:

package kexec-tools is not installed

2. Install **kdump** and other necessary packages by:

yum install kexec-tools



IMPORTANT

Starting with Red Hat Enterprise Linux 7.4 (kernel-3.10.0-693.el7) the **Intel IOMMU** driver is supported with **kdump**. For prior versions, Red Hat Enterprise Linux 7.3 (kernel-3.10.0-514[.XYZ].el7) and earlier, it is advised that **Intel IOMMU** support is disabled, otherwise kdump kernel is likely to become unresponsive.

CHAPTER 11. CONFIGURING KDUMP ON THE COMMAND LINE

The memory for **kdump** is reserved during the system boot. The memory size is configured in the system's Grand Unified Bootloader (GRUB) 2 configuration file. The memory size depends on the **crashkernel=** value specified in the configuration file and the size of the system's physical memory.

11.1. CONFIGURING KDUMP MEMORY USAGE

The **crashkernel=** option can be defined in multiple ways. You can specify the **crashkernel=** value or configure the **auto** option. The **crashkernel=auto** boot option, reserves memory automatically, depending on total amount of the system's physical memory. When configured, the kernel will automatically reserve an appropriate amount of required memory for the **kdump** kernel. This helps to prevent Out-of-Memory (OOM) error occurrences.



NOTE

The automatic memory allocation for **kdump** varies based on system hardware architecture and available memory size.

For example, on AMD64 and Intel 64, the **crashkernel=auto** parameter works only when the available memory is more than 1GB and 64-bit ARM architecture and IBM Power Systems have available memory more than 2GB.

If the system has less than the minimum memory threshold for automatic allocation, you can configure the amount of reserved memory manually.

Prerequisites

• Fulfilled **kdump** requirements for configurations and targets

Procedure

- 1. Edit the /etc/default/grub file using the root permissions.
- Set the crashkernel= option to the required value.
 For example, to reserve 128 MB of memory, use the following:

crashkernel=128M

Alternatively, you can set the amount of reserved memory to a variable depending on the total amount of installed memory. The syntax for memory reservation into a variable is **crashkernel=**/**crange1>:**<**size1>**,<**range2>:**<**size2>**. For example:

crashkernel=512M-2G:64M,2G-:128M

The above example reserves 64 MB of memory if the total amount of system memory is 512 MB or higher and lower than 2 GB. If the total amount of memory is more than 2 GB, 128 MB is reserved for **kdump** instead.

Offset the reserved memory.
 Some systems require to reserve memory with a certain fixed offset since crashkernel reservation is very early, and it wants to reserve some area for special usage. If the offset is set, the reserved memory begins there. To offset the reserved memory, use the following syntax:

crashkernel=128M@16M

The example above means that **kdump** reserves 128 MB of memory starting at 16 MB (physical address 0x01000000). If the offset parameter is set to 0 or omitted entirely, **kdump** offsets the reserved memory automatically. This syntax can also be used when setting a variable memory reservation as described above; in this case, the offset is always specified last (for example, **crashkernel=512M-2G:64M,2G-:128M@16M**).

3. Use the following command to update the GRUB2 configuration file:

grub2-mkconfig -o /boot/grub2/grub.cfg



NOTE

The alternative way to configure memory for **kdump** is to append the **crashkernel=** <**SOME_VALUE>** parameter to the **kernelopts** variable with the **grub2-editenv** which will update all of your boot entries. Or you can use the **grubby** utility to update kernel command line parameters of just one entry.

Additional resources

- Memory requirements for kdump
- Configuring kernel command line parameters
- The grub2-mkconfig script silently ignores options in GRUB_CMDLINE_LINUX
- How to manually modify the boot parameter in grub before the system boots
- How to install and boot custom kernels in Red Hat Enterprise Linux 8

11.2. CONFIGURING THE KDUMP TARGET

When a kernel crash is captured, the core dump can be either stored as a file in a local file system, written directly to a device, or sent over a network using the **NFS** (Network File System) or **SSH** (Secure Shell) protocol. Only one of these options can be set at a time, and the default behavior is to store the vmcore file in the /var/crash/ directory of the local file system.

Prerequisites

Fulfilled kdump requirements for configurations and targets.

Procedure

 To store the vmcore file in /var/crash/ directory of the local file system, edit the /etc/kdump.conf file and specify the path:

path /var/crash

The option **path** /**var/crash** represents the path to the file system in which **kdump** saves the **vmcore** file. When you specify a dump target in the /**etc/kdump.conf** file, then the **path** is relative to the specified dump target.

If you do not specify a dump target in the /etc/kdump.conf file, then the path represents the absolute path from the root directory. Depending on what is mounted in the current system, the dump target and the adjusted dump path are taken automatically.



WARNING

kdump saves the **vmcore** file in **/var/crash/var/crash** directory, when the dump target is mounted at **/var/crash** and the option **path** is also set as **/var/crash** in the **/etc/kdump.conf** file. For example, in the following instance, the **ext4** file system is already mounted at **/var/crash** and the **path** are set as **/var/crash**:

grep -v ^# etc/kdump.conf | grep -v ^\$
ext4 /dev/mapper/vg00-varcrashvol
path /var/crash
core_collector makedumpfile -c --message-level 1 -d 31

This results in the /var/crash/var/crash path. To solve this problem, use the option path / instead of path /var/crash

- To change the local directory in which the core dump is to be saved, as **root**, edit the /etc/kdump.conf configuration file as described below.
 - 1. Remove the hash sign ("#") from the beginning of the **#path** /**var/crash** line.
 - 2. Replace the value with the intended directory path. For example:

path /usr/local/cores



IMPORTANT

In Red Hat Enterprise Linux 8, the directory defined as the kdump target using the **path** directive must exist when the **kdump** systemd service is started - otherwise the service fails. This behavior is different from earlier releases of Red Hat Enterprise Linux, where the directory was being created automatically if it did not exist when starting the service.

- To write the file to a different partition, as **root**, edit the /**etc/kdump.conf** configuration file as described below.
 - 1. Remove the hash sign ("#") from the beginning of the **#ext4** line, depending on your choice.
 - device name (the **#ext4** /**dev/vg/lv_kdump** line)
 - file system label (the **#ext4 LABEL=/boot** line)
 - UUID (the #ext4 UUID=03138356-5e61-4ab3-b58e-27507ac41937 line)
 - 2. Change the file system type as well as the device name, label or UUID to the desired values. For example:

ext4 UUID=03138356-5e61-4ab3-b58e-27507ac41937



IMPORTANT

It is recommended to specify storage devices using a **LABEL=** or **UUID=**. Disk device names such as /dev/sda3 are not guaranteed to be consistent across reboot.



IMPORTANT

When dumping to Direct Access Storage Device (DASD) on IBM Z hardware, it is essential that the dump devices are correctly specified in /etc/dasd.conf before proceeding.

- To write the dump directly to a device:
 - 1. Remove the hash sign ("#") from the beginning of the **#raw** /**dev/vg/lv_kdump** line.
 - 2. Replace the value with the intended device name. For example:

raw /dev/sdb1

- To store the dump to a remote machine using the **NFS** protocol:
 - 1. Remove the hash sign ("#") from the beginning of the **#nfs my.server.com:/export/tmp** line.
 - 2. Replace the value with a valid hostname and directory path. For example:

nfs penguin.example.com:/export/cores

- To store the dump to a remote machine using the **SSH** protocol:
 - 1. Remove the hash sign ("#") from the beginning of the **#ssh user@my.server.com** line.
 - 2. Replace the value with a valid username and hostname.
 - 3. Include your **SSH** key in the configuration.
 - Remove the hash sign from the beginning of the **#sshkey** /**root**/.**ssh**/**kdump_id_rsa** line.
 - Change the value to the location of a key valid on the server you are trying to dump to. For example:

ssh john@penguin.example.com sshkey /root/.ssh/mykey

Additional resources

- Supported kdump configurations and targets
- Configuring basic system settings in Red Hat Enterprise Linux

11.3. CONFIGURING THE CORE COLLECTOR

The **kdump** service uses a **core_collector** program to capture the **vmcore** image. In RHEL, the **makedumpfile** utility is the default core collector.

makedumpfile is a dump program that helps to compress the size of a dump file and copy only necessary pages using various dump levels.

makedumpfile makes a small dump file either by compressing dump data or by excluding unnecessary pages or both. It needs the first kernel debug information to understand how first kernel uses the memory. This helps to detect the necessary pages.

Syntax

core_collector makedumpfile -l --message-level 1 -d 31

Options

- -c, -l or -p: specify compress dump file format by each page using either, zlib for -c option, lzo for -l option or snappy for -p option.
- -d (dump_level): excludes pages so that they are not copied to the dump file.
- --message-level: specify the message types. You can restrict outputs printed by specifying message_level with this option. For example, specifying 7 as message_level prints common messages and error messages. The maximum value of message_level is 31

Prerequisites

Fulfilled kdump requirements for configurations and targets.

Procedure

- 1. As **root**, edit the /**etc/kdump.conf** configuration file and remove the hash sign ("#") from the beginning of the **#core_collector makedumpfile -I --message-level 1 -d 31**.
- 2. To enable dump file compression, execute:

core_collector makedumpfile -I --message-level 1 -d 31

where,

- -I specifies the **dump** compressed file format.
- -d specifies dump level as 31.
- --message-level specifies message level as 1.

Also, consider following examples using options **-c** and **-p**:

• To compress dump file using -c:

core_collector makedumpfile -c -d 31 --message-level 1

• To compress dump file using -p:

core collector makedumpfile -p -d 31 --message-level 1

Additional resources

makedumpfile(8) manual page

11.4. CONFIGURING THE KDUMP DEFAULT FAILURE RESPONSES

By default, when **kdump** fails to create a vmcore file at the configured target location, the system reboots and the dump is lost in the process. To change this behavior, follow the procedure below.

Prerequisites

Fulfilled kdump requirements for configurations and targets.

Procedure

- 1. As **root**, remove the hash sign ("#") from the beginning of the **#failure_action** line in the /etc/kdump.conf configuration file.
- 2. Replace the value with a desired action.

failure_action poweroff

Additional resources

Configuring the kdump target

11.5. CONFIGURATION FILE FOR KDUMP

The configuration file for **kdump** kernel is /**etc/sysconfig/kdump**. This file controls the **kdump** kernel command line parameters.

For most configurations, use the default options. However, in some scenarios you might need to modify certain parameters to control the **kdump** kernel behavior. For example, modifying to append the **kdump** kernel command-line to obtain a detailed debugging output.

This section covers information on modifying the KDUMP_COMMANDLINE_REMOVE and KDUMP_COMMANDLINE_REMOVE options for kdump. For information on additional configuration options refer to Documentation/admin-guide/kernel-parameters.txt or the /etc/sysconfig/kdump file.

KDUMP_COMMANDLINE_REMOVE

This option removes arguments from the current **kdump** command line. It removes parameters that may cause **kdump** errors or **kdump** kernel boot failures. These parameters may have been parsed from the previous **KDUMP_COMMANDLINE** process or inherited from the /proc/cmdline file. When this variable is not configured, it inherits all values from the /proc/cmdline file. Configuring this option also provides information that is helpful in debugging an issue.

Example

To remove certain arguments, add them to **KDUMP_COMMANDLINE_REMOVE** as follows:

KDUMP_COMMANDLINE_REMOVE="hugepages hugepagesz slub_debug quiet log_buf_len swiotlb"

KDUMP COMMANDLINE APPEND

This option appends arguments to the current command line. These arguments may have been parsed by the previous **KDUMP_COMMANDLINE_REMOVE** variable.

For the **kdump** kernel, disabling certain modules such as **mce**, **cgroup**, **numa**, **hest_disable** can help prevent kernel errors. These modules may consume a significant portion of the kernel memory reserved for **kdump** or cause **kdump** kernel boot failures.

Example

To disable memory **cgroups** on the **kdump** kernel command line, run the command as follows:

KDUMP_COMMANDLINE_APPEND="cgroup_disable=memory"

Additional resources

- Documentation/admin-guide/kernel-parameters.txt file
- /etc/sysconfig/kdump file

11.6. ENABLING AND DISABLING THE KDUMP SERVICE

To start the **kdump** service at boot time, follow the procedure below.

Prerequisites

- Fulfilled **kdump** requirements for configurations and targets.
- All configurations for installing **kdump** are set up according to your needs.

Procedure

- 1. To enable the **kdump** service, use the following command:
 - # systemctl enable kdump.service

This enables the service for **multi-user.target**.

- 2. To start the service in the current session, use the following command:
 - # systemctl start kdump.service
- 3. To stop the **kdump** service, type the following command:
 - # systemctl stop kdump.service
- 4. To disable the **kdump** service, execute the following command:
 - # systemctl disable kdump.service



WARNING

It is recommended to set **kptr_restrict=1** as default. When **kptr_restrict** is set to (1) as default, the **kdumpctl** service loads the crash kernel even if Kernel Address Space Layout (KASLR) is enabled or not enabled.

Troubleshooting step

When **kptr_restrict** is not set to (1), and if KASLR is enabled, the contents of /**proc/kore** file are generated as all zeros. Consequently, the **kdumpctl** service fails to access the /**proc/kcore** and load the crash kernel.

To work around this problem, the **kexec-kdump-howto.txt** file displays a warning message, which specifies to keep the recommended setting as **kptr_restrict=1**.

To ensure that **kdumpctl** service loads the crash kernel, verify that:

• Kernel **kptr_restrict=1** in the **sysctl.conf** file.

Additional resources

Configuring basic system settings in Red Hat Enterprise Linux

11.7. TESTING THE KDUMP CONFIGURATION

The following procedure describes how to test that the kernel dump process works and is valid before the machine enters production.



WARNING

The commands below cause the kernel to crash. Use caution when following these steps, and never carelessly use them on active production system.

Procedure

- 1. Reboot the system with **kdump** enabled.
- 2. Make sure that **kdump** is running:
 - ~]# systemctl is-active kdump active
- 3. Force the Linux kernel to crash:

echo 1 > /proc/sys/kernel/sysrq echo c > /proc/sysrq-trigger



WARNING

The command above crashes the kernel and a reboot is required.

Once booted again, the *address-YYYY-MM-DD-HH:MM:SS*/vmcore file is created at the location you have specified in /etc/kdump.conf (by default to /var/crash/).



NOTE

In addition to confirming the validity of the configuration, it is possible to use this action to record how long it takes for a crash dump to complete, while a representative load was running.

Additional resources

Configuring the kdump target

11.8. USING KEXEC TO REBOOT THE KERNEL

The **kexec** system call enables loading and booting into another kernel from the currently running kernel, thus performing a function of a boot loader from within the kernel.

The **kexec** utility loads the kernel and the **initramfs** image for the **kexec** system call to boot into another kernel.

The following procedure describes how to manually invoke the **kexec** system call when using the **kexec** utility to reboot into another kernel.

Procedure

1. Execute the **kexec** utility:

kexec -l /boot/vmlinuz-3.10.0-1040.el7.x86_64 --initrd=/boot/initramfs-3.10.0-1040.el7.x86 64.img --reuse-cmdline

The command manually loads the kernel and the initramfs image for the **kexec** system call.

2. Reboot the system:

reboot

The command detects the kernel, shuts down all services and then calls the **kexec** system call to reboot into the kernel you provided in the previous step.



WARNING

When you use the **kexec -e** command to reboot the kernel, the system does not go through the standard shutdown sequence before starting the next kernel, which may cause data loss or an unresponsive system.

11.9. PREVENTING KERNEL DRIVERS FROM LOADING FOR KDUMP

This section explains how to prevent the **kdump** kernel from loading certain kernel drivers using the `/etc/sysconfig/kdump configuration file. Appending the **KDUMP_COMMANDLINE_APPEND=** variable in /etc/sysconfig/kdump file, prevents the **kdump initramfs** from loading the specified kernel module. This helps to prevent **oom killer** or other crash kernel failures.

You can append the **KDUMP_COMMANDLINE_APPEND=** variable using one of the following configuration options:

- rd.driver.blacklist=<modules>
- modprobe.blacklist=<modules>

Procedure

1. Select a kernel module that you intend to block from loading.

\$ Ismod

Module Size Used by
fuse 126976 3

xt_CHECKSUM 16384 1
ipt_MASQUERADE 16384 1
uinput 20480 1

xt_conntrack 16384 1

The **Ismod** command displays a list of modules that are loaded to the currently running kernel.

2. Update the KDUMP_COMMANDLINE_APPEND= variable in the /etc/sysconfig/kdump file.

KDUMP_COMMANDLINE_APPEND="rd.driver.blacklist=hv_vmbus,hv_storvsc,hv_utils,hv_netvsc,hid-hyperv"

Also, consider the following example using the **modprobe.blacklist=<modules>** configuration option.

KDUMP_COMMANDLINE_APPEND="modprobe.blacklist=emcp modprobe.blacklist=bnx2fc modprobe.blacklist=libfcoe modprobe.blacklist=fcoe"

3. Restart the **kdump** service.

\$ systemctl restart kdump

Additional resources

• dracut.cmdline manual page

11.10. RUNNING KDUMP ON SYSTEMS WITH ENCRYPTED DISK

When running a LUKS encrypted partition, systems require certain amount of available memory. If the system has less than the required amount of available memory, the **cryptsetup** utility fails to mount the partition. As a result, capturing the **vmcore** file to an encrypted target location fails in the second kernel (capture kernel).

Using the **kdumpctl estimate** command, you can estimate the amount of memory required for **kdump**. **kdumpctl estimate** prints the **Recommended crashkernel** value, which is the recommended memory size required for **kdump**.

The **Recommended crashkernel** value is calculated based on the current kernel size, kernel module, initramfs, and the LUKS encrypted target memory requirement.

In case you are using the custom **crashkernel=** option, **kdumpctl estimate** prints the **LUKS required size** value, which is the memory size required for LUKS encrypted target.

Procedure

1. Run the following command to print the estimate **crashkernel=** value:

kdumpctl estimate

Encrypted kdump target requires extra memory, assuming using the keyslot with minimum memory requirement

Reserved crashkernel: 256M Recommended crashkernel: 652M

Kernel image size: 47M
Kernel modules size: 8M
Initramfs size: 20M
Runtime reservation: 64M
LUKS required size: 512M
Large modules: <none>

WARNING: Current crashkernel size is lower than recommended size 652M.

- 2. Configure the amount of required memory by increasing the **crashkernel=** value to the desired value.
- 3. Reboot the system.



NOTE

If the **kdump** service still fails to save the dump file to the encrypted target, increase the **crashkernel=** value as required.

CHAPTER 12. CONFIGURING KDUMP IN THE WEB CONSOLE

Setup and test the **kdump** configuration in the RHEL 8 web console.

The web console is part of a default installation of Red Hat Enterprise Linux 8 and enables or disables the **kdump** service at boot time. Further, the web console conveniently enables you to configure the reserved memory for **kdump**; or to select the *vmcore* saving location in an uncompressed or compressed format.

12.1. ADDITIONAL RESOURCES

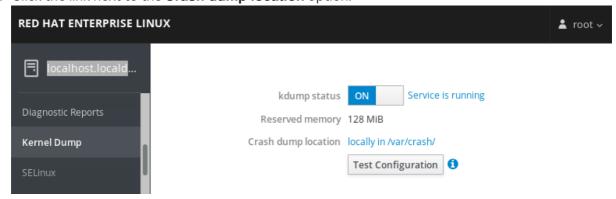
• Getting started using the RHEL web console

12.2. CONFIGURING KDUMP MEMORY USAGE AND TARGET LOCATION IN WEB CONSOLE

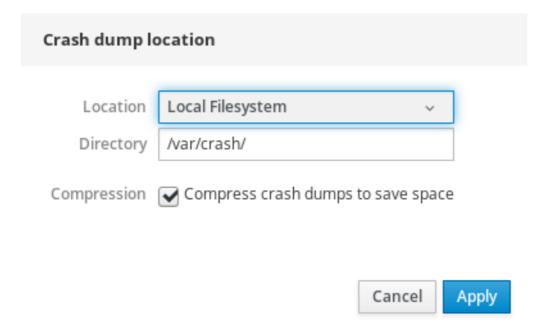
The procedure below shows you how to use the **Kernel Dump** tab in the Red Hat Enterprise Linux web console interface to configure the amount of memory that is reserved for the **kdump** kernel. The procedure also describes how to specify the target location of the **vmcore** dump file and how to test your configuration.

Procedure

- 1. Open the **Kernel Dump** tab and start the **kdump** service.
- 2. Configure the **kdump** memory usage using the command line.
- 3. Click the link next to the **Crash dump location** option.



4. Select the **Local Filesystem** option from the drop-down and specify the directory you want to save the dump in.



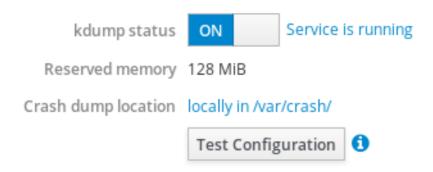
- Alternatively, select the Remote over SSH option from the drop-down to send the vmcore
 to a remote machine using the SSH protocol.
 Fill the Server, ssh key, and Directory fields with the remote machine address, ssh key
 location, and a target directory.
- Another choice is to select the **Remote over NFS** option from the drop-down and fill the **Mount** field to send the vmcore to a remote machine using the NFS protocol.



NOTE

Tick the **Compression** check box to reduce the size of the vmcore file.

5. Test your configuration by crashing the kernel.





WARNING

This step disrupts execution of the kernel and results in a system crash and loss of data.

Additional resources

- Supported kdump targets
- Using secure communications between two systems with OpenSSH

CHAPTER 13. SUPPORTED KDUMP CONFIGURATIONS AND TARGETS

13.1. MEMORY REQUIREMENTS FOR KDUMP

In order for **kdump** to be able to capture a kernel crash dump and save it for further analysis, a part of the system memory has to be permanently reserved for the capture kernel. When reserved, this part of the system memory is not available to the main kernel.

The memory requirements vary based on certain system parameters. One of the major factors is the system's hardware architecture. To find out the exact machine architecture (such as Intel 64 and AMD64, also known as x86_64) and print it to standard output, use the following command:

\$ uname -m

Table 10.1 lists the minimum memory requirements to automatically reserve a memory size for **kdump** on the latest available versions. The size changes according to the system's architecture and total available physical memory.

Table 13.1. Minimum Amount of Reserved Memory Required for kdump

Architecture	Available Memory	Minimum Reserved Memory
AMD64 and Intel 64 (x86_64)	1 GB to 4 GB	160 MB of RAM.
	4 GB to 64 GB	192 MB of RAM.
	64 GB to 1 TB	256 MB of RAM.
	1TB and more	512 MB of RAM.
64-bit ARM architecture (arm64)	2 GB and more	448 MB of RAM.
IBM Power Systems (ppc64le)	2 GB to 4 GB	384 MB of RAM.
	4 GB to 16 GB	512 MB of RAM.
	16 GB to 64 GB	1 GB of RAM.
	64 GB to 128 GB	2 GB of RAM.
	128 GB and more	4 GB of RAM.
IBM Z (s390x)	1 GB to 4 GB	160 MB of RAM.
	4 GB to 64 GB	192 MB of RAM.
	64 GB to 1 TB	256 MB of RAM.

Architecture	Available Memory	Minimum Reserved Memory
	1TB and more	512 MB of RAM.

On many systems, **kdump** is able to estimate the amount of required memory and reserve it automatically. This behavior is enabled by default, but only works on systems that have more than a certain amount of total available memory, which varies based on the system architecture.



IMPORTANT

The automatic configuration of reserved memory based on the total amount of memory in the system is a best effort estimation. The actual required memory may vary due to other factors such as I/O devices. Using not enough of memory might cause that a debug kernel is not able to boot as a capture kernel in case of a kernel panic. To avoid this problem, sufficiently increase the crash kernel memory.

Additional resources

- How has the crashkernel parameter changed between RHEL8 minor releases?
- Technology capabilities and limits tables
- Configuring kdump memory usage
- Configuring kdump memory usage and target location in web console
- Minimum threshold for automatic memory reservation

13.2. MINIMUM THRESHOLD FOR AUTOMATIC MEMORY RESERVATION

On some systems, it is possible to allocate memory for **kdump** automatically, either by using the **crashkernel=auto** parameter in the boot loader configuration file, or by enabling this option in the graphical configuration utility. For this automatic reservation to work, however, a certain amount of total memory needs to be available in the system. The amount differs based on the system's architecture.

Table 10.2 lists the threshold values for automatic memory allocation. If the system has memory less than the specified threshold value, you must configure the memory manually.

Table 13.2. Minimum Amount of Memory Required for Automatic Memory Reservation

Architecture	Required Memory
AMD64 and Intel 64 (x86_64)	2 GB
IBM Power Systems (ppc64le)	2 GB
IBM Z (s390x)	4 GB

Additional resources

- Configuring kdump memory usage
- Configuring kdump memory usage and target location in web console

13.3. SUPPORTED KDUMP TARGETS

When a kernel crash is captured, the vmcore dump file can be either written directly to a device, stored as a file on a local file system, or sent over a network. The table below contains a complete list of dump targets that are currently supported or explicitly unsupported by **kdump**.

Туре	Supported Targets	Unsupported Targets
Raw device	All locally attached raw disks and partitions.	
Local file system	ext2, ext3, ext4, and xfs file systems on directly attached disk drives, hardware RAID logical drives, LVM devices, and mdraid arrays.	Any local file system not explicitly listed as supported in this table, including the auto type (automatic file system detection).
Remote directory	Remote directories accessed using the NFS or SSH protocol over IPv4 .	Remote directories on the rootfs file system accessed using the NFS protocol.
Remote directories accessed using the iSCSI protocol over both hardware and software initiators.	Remote directories accessed using the iSCSI protocol on be2iscsi hardware.	Multipath-based storages.
		Remote directories accessed over IPv6 .
		Remote directories accessed using the SMB or CIFS protocol.
		Remote directories accessed using the FCoE (<i>Fibre Channel over Ethernet</i>) protocol.
		Remote directories accessed using wireless network interfaces.



IMPORTANT

Utilizing firmware assisted dump (**fadump**) to capture a vmcore and store it to a remote machine using SSH or NFS protocol causes renaming of the network interface to **kdump-interface-name>**. The renaming happens if the **interface-name>** is generic, for example *eth#, net#, and so on. This problem occurs because the vmcore capture scripts in the initial RAM disk (**initrd**) add the *kdump-* prefix to the network interface name to secure persistent naming. Since the same **initrd** is used also for a regular boot, the interface name is changed for the production kernel too.

Additional resources

- Configuring kdump target
- Configuring kdump memory usage and target location in web console

13.4. SUPPORTED KDUMP FILTERING LEVELS

To reduce the size of the dump file, **kdump** uses the **makedumpfile** core collector to compress the data and optionally to omit unwanted information. The table below contains a complete list of filtering levels that are currently supported by the **makedumpfile** utility.

Option	Description
1	Zero pages
2	Cache pages
4	Cache private
8	User pages
16	Free pages



NOTE

The **makedumpfile** command supports removal of transparent huge pages and hugetlbfs pages. Consider both these types of hugepages User Pages and remove them using the **8** level.

Additional resources

Configuring the core collector

13.5. SUPPORTED DEFAULT FAILURE RESPONSES

By default, when **kdump** fails to create a core dump, the operating system reboots. You can, however, configure **kdump** to perform a different operation in case it fails to save the core dump to the primary target. The table below lists all default actions that are currently supported.

Option	Description
dump_to_rootfs	Attempt to save the core dump to the root file system. This option is especially useful in combination with a network target: if the network target is unreachable, this option configures kdump to save the core dump locally. The system is rebooted afterwards.
reboot	Reboot the system, losing the core dump in the process.
halt	Halt the system, losing the core dump in the process.
poweroff	Power off the system, losing the core dump in the process.
shell	Run a shell session from within the initramfs, allowing the user to record the core dump manually.
final_action	Enable additional operations such as reboot , halt , and poweroff actions after a successful kdump or when shell or dump_to_rootfs failure action completes. The default final_action option is reboot .

Additional resources

• Configuring the kdump default failure responses

13.6. USING FINAL_ACTION PARAMETER

The **final_action** parameter enables you to use certain additional operations such as **reboot**, **halt**, and **poweroff** actions after a successful **kdump** or when an invoked **failure_response** mechanism using **shell** or **dump_to_rootfs** completes. If the **final_action** option is not specified, it defaults to **reboot**.

Procedure

- 1. Edit the '/etc/kdump.conf file and add the final_action parameter.
 - final_action <reboot | halt | poweroff>
- 2. Restart the kdump service:
 - kdumpctl restart

13.7. ESTIMATING KDUMP SIZE

When planning and building your kdump environment, it is necessary to know how much space is required for the dump file before one is produced.

The **makedumpfile --mem-usage** command provides a useful report about excludable pages, and can be used to determine which dump level you want to assign. Run this command when the system is under representative load, otherwise **makedumpfile --mem-usage** returns a smaller value than is expected in your production environment.

[root@hostname ~]# makedumpfile --mem-usage /proc/kcore

TYPE	PAGES	EXCLUD	ABLE DESCRIPTION
ZERO	501635	yes	Pages filled with zero
		yes	•
CACHE	51657	yes	Cache pages
CACHE_PI	RIVATE 5442	yes	Cache pages + private
USER	16301	yes	User process pages
FREE	77738211	yes	Free pages
KERN_DA	TA 1333192	no	Dumpable kernel data



IMPORTANT

The **makedumpfile --mem-usage** command reports in **pages**. This means that you have to calculate the size of memory in use against the kernel page size. By default the Red Hat Enterprise Linux kernel uses 4 KB sized pages for AMD64 and Intel 64 architectures, and 64 KB sized pages for IBM POWER architectures.

CHAPTER 14. FIRMWARE ASSISTED DUMP MECHANISMS

Firmware assisted dump (fadump) is a dump capturing mechanism, provided as an alternative to the **kdump** mechanism on IBM POWER systems. The **kexec** and **kdump** mechanisms are useful for capturing core dumps on AMD64 and Intel 64 systems. However, some hardware such as mini systems and mainframe computers, leverage the onboard firmware to isolate regions of memory and prevent any accidental overwriting of data that is important to the crash analysis. This section covers **fadump** mechanisms and how they integrate with RHEL. The **fadump** utility is optimized for these expanded dumping features on IBM POWER systems.

14.1. FIRMWARE ASSISTED DUMP ON IBM POWERPC HARDWARE

The **fadump** utility captures the **vmcore** file from a fully-reset system with PCI and I/O devices. This mechanism uses firmware to preserve memory regions during a crash and then reuses the **kdump** userspace scripts to save the **vmcore** file. The memory regions consist of all system memory contents, except the boot memory, system registers, and hardware Page Table Entries (PTEs).

The **fadump** mechanism offers improved reliability over the traditional dump type, by rebooting the partition and using a new kernel to dump the data from the previous kernel crash. The **fadump** requires an IBM POWER6 processor-based or later version hardware platform.

For further details about the **fadump** mechanism, including PowerPC specific methods of resetting hardware, see the **/usr/share/doc/kexec-tools/fadump-howto.txt** file.



NOTE

The area of memory that is not preserved, known as boot memory, is the amount of RAM required to successfully boot the kernel after a crash event. By default, the boot memory size is 256MB or 5% of total system RAM, whichever is larger.

Unlike **kexec-initiated** event, the **fadump** mechanism uses the production kernel to recover a crash dump. When booting after a crash, PowerPC hardware makes the device node /**proc/device-tree/rtas/ibm.kernel-dump** available to the **proc** filesystem (**procfs**). The **fadump-aware kdump** scripts, check for the stored **vmcore**, and then complete the system reboot cleanly.

14.2. ENABLING FIRMWARE ASSISTED DUMP MECHANISM

The crash dumping capabilities of IBM POWER systems can be enhanced by enabling the firmware assisted dump (fadump) mechanism.

Procedure

- 1. Install and configure **kdump**.
- 2. Add fadump=on to the GRUB_CMDLINE_LINUX line in /etc/default/grub file:

GRUB_CMDLINE_LINUX="rd.lvm.lv=rhel/swap crashkernel=auto rd.lvm.lv=rhel/root rhgb quiet fadump=on"

3. (Optional) If you want to specify reserved boot memory instead of using the defaults, configure crashkernel=xxM to GRUB_CMDLINE_LINUX in /etc/default/grub, where xx is the amount of the memory required in megabytes:

GRUB_CMDLINE_LINUX="rd.lvm.lv=rhel/swap crashkernel=xxM rd.lvm.lv=rhel/root rhgb quiet fadump=on"



IMPORTANT

Red Hat recommends to test all boot configuration options before you execute them. If you observe Out of Memory (OOM) errors when booting from the crash kernel, increase the value specified in **crashkernel=** argument until the crash kernel can boot cleanly. Some trial and error may be required in this case.

14.3. FIRMWARE ASSISTED DUMP MECHANISMS ON IBM Z HARDWARE

IBM Z systems support the following firmware assisted dump mechanisms:

- Stand-alone dump (sadump)
- VMDUMP

The **kdump** infrastructure is supported and utilized on IBM Z systems. However, using one of the firmware assisted dump (fadump) methods for IBM Z can provide various benefits:

- The **sadump** mechanism is initiated and controlled from the system console, and is stored on an **IPL** bootable device.
- The **VMDUMP** mechanism is similar to **sadump**. This tool is also initiated from the system console, but retrieves the resulting dump from hardware and copies it to the system for analysis.
- These methods (similarly to other hardware based dump mechanisms) have the ability to capture the state of a machine in the early boot phase, before the **kdump** service starts.
- Although VMDUMP contains a mechanism to receive the dump file into a Red Hat Enterprise
 Linux system, the configuration and control of VMDUMP is managed from the IBM Z Hardware
 console.

IBM discusses **sadump** in detail in the **Stand-alone dump program** article and **VMDUMP** in **Creating dumps on z/VM with VMDUMP** article.

IBM also has a documentation set for using the dump tools on Red Hat Enterprise Linux 7 in the **Using** the Dump Tools on Red Hat Enterprise Linux 7.4 article.

Additional resources

- Stand-alone dump program
- Creating dumps on z/VM with VMDUMP
- Using the Dump Tools on Red Hat Enterprise Linux 7.4

14.4. USING SADUMP ON FUJITSU PRIMEQUEST SYSTEMS

The Fujitsu **sadump** mechanism is designed to provide a **fallback** dump capture in an event when **kdump** is unable to complete successfully. The **sadump** mechanism is invoked manually from the system Management Board (MMB) interface. Using MMB, configure **kdump** like for an Intel 64 or AMD

64 server and then perform the following additional steps to enable **sadump**.

Procedure

1. Add or edit the following lines in the /etc/sysctl.conf file to ensure that kdump starts as expected for sadump:

kernel.panic=0 kernel.unknown_nmi_panic=1



WARNING

In particular, ensure that after **kdump**, the system does not reboot. If the system reboots after **kdump** has fails to save the **vmcore** file, then it is not possible to invoke the **sadump**.

2. Set the **failure_action** parameter in /etc/kdump.conf appropriately as halt or shell.

failure_action shell

Additional resources

• The FUJITSU Server PRIMEQUEST 2000 Series Installation Manual

CHAPTER 15. ANALYZING A CORE DUMP

To determine the cause of the system crash, you can use the **crash** utility, which provides an interactive prompt very similar to the GNU Debugger (GDB). This utility allows you to interactively analyze a core dump created by **kdump**, **netdump**, **diskdump** or **xendump** as well as a running Linux system. Alternatively, you have the option to use Kernel Oops Analyzer or the Kdump Helper tool.

15.1. INSTALLING THE CRASH UTILITY

The following procedure describes how to install the **crash** analyzing tool.

Procedure

- 1. Enable the relevant repositories:
 - # subscription-manager repos --enable baseos repository
 - # subscription-manager repos --enable appstream repository
 - # subscription-manager repos --enable rhel-8-for-x86_64-baseos-debug-rpms
- 2. Install the crash package:
 - # yum install crash
- 3. Install the kernel-debuginfo package:
 - # yum install kernel-debuginfo

The package corresponds to your running kernel and provides the data necessary for the dump analysis.

Additional resources

Configuring basic system settings Red Hat Enterprise Linux

15.2. RUNNING AND FXITING THE CRASH UTILITY

The following procedure describes how to start the crash utility for analyzing the cause of the system crash.

Prerequisites

• Identify the currently running kernel (for example **4.18.0-5.el8.x86_64**).

Procedure

- 1. To start the **crash** utility, two necessary parameters need to be passed to the command:
 - The debug-info (a decompressed vmlinuz image), for example
 /usr/lib/debug/lib/modules/4.18.0-5.el8.x86_64/vmlinux provided through a specific
 kernel-debuginfo package.

• The actual vmcore file, for example /var/crash/127.0.0.1-2018-10-06-14:05:33/vmcore
The resulting crash command then looks like this:

crash /usr/lib/debug/lib/modules/4.18.0-5.el8.x86_64/vmlinux /var/crash/127.0.0.1-2018-10-06-14:05:33/vmcore

Use the same < kernel > version that was captured by kdump.

Example 15.1. Running the crash utility

The following example shows analyzing a core dump created on October 6 2018 at 14:05 PM, using the 4.18.0-5.el8.x86_64 kernel.

WARNING: kernel relocated [202MB]: patching 90160 gdb minimal_symbol values

KERNEL: /usr/lib/debug/lib/modules/4.18.0-5.el8.x86 64/vmlinux

DUMPFILE: /var/crash/127.0.0.1-2018-10-06-14:05:33/vmcore [PARTIAL DUMP]

CPUS: 2

DATE: Sat Oct 6 14:05:16 2018

UPTIME: 01:03:57

LOAD AVERAGE: 0.00, 0.00, 0.00

TASKS: 586

NODENAME: localhost.localdomain RELEASE: 4.18.0-5.el8.x86_64

VERSION: #1 SMP Wed Aug 29 11:51:55 UTC 2018

MACHINE: x86_64 (2904 Mhz)

MEMORY: 2.9 GB

PANIC: "sysrq: SysRq: Trigger a crash"

PID: 10635 COMMAND: "bash"

TASK: ffff8d6c84271800 [THREAD_INFO: ffff8d6c84271800]

CPI I: 1

STATE: TASK_RUNNING (SYSRQ)

crash>

2. To exit the interactive prompt and terminate **crash**, type **exit** or **q**.

Example 15.2. Exiting the crash utility

crash> exit ~]#



NOTE

The **crash** command can also be used as a powerful tool for debugging a live system. However use it with caution so as not to break your system.

Additional resources

A Guide to Unexpected System Restarts

15.3. DISPLAYING VARIOUS INDICATORS IN THE CRASH UTILITY

The following procedures describe how to use the crash utility and display various indicators, such as a kernel message buffer, a backtrace, a process status, virtual memory information and open files.

Displaying the message buffer

• To display the kernel message buffer, type the **log** command at the interactive prompt as displayed in the example below:

crash> log

... several lines omitted ...

EIP: 0060:[<c068124f>] EFLAGS: 00010096 CPU: 2

EIP is at sysrq_handle_crash+0xf/0x20

EAX: 00000063 EBX: 00000063 ECX: c09e1c8c EDX: 00000000 ESI: c0a09ca0 EDI: 00000286 EBP: 00000000 ESP: ef4dbf24

DS: 007b ES: 007b FS: 00d8 GS: 00e0 SS: 0068

Process bash (pid: 5591, ti=ef4da000 task=f196d560 task.ti=ef4da000)

Stack:

c068146b c0960891 c0968653 00000003 00000000 00000002 efade5c0 c06814d0 <0> fffffffb c068150f b7776000 f2600c40 c0569ec4 ef4dbf9c 00000002 b7776000 <0> efade5c0 00000002 b7776000 c0569e60 c051de50 ef4dbf9c f196d560 ef4dbfb4

Call Trace:

[<c068146b>] ? __handle_sysrq+0xfb/0x160

[<c06814d0>] ? write_sysrq_trigger+0x0/0x50

[<c068150f>] ? write_sysrq_trigger+0x3f/0x50

[<c0569ec4>] ? proc_reg_write+0x64/0xa0

[<c0569e60>] ? proc_reg_write+0x0/0xa0

[<c051de50>] ? vfs_write+0xa0/0x190

[<c051e8d1>] ? sys_write+0x41/0x70

[<c0409adc>] ? syscall_call+0x7/0xb

Code: a0 c0 01 0f b6 41 03 19 d2 f7 d2 83 e2 03 83 e0 cf c1 e2 04 09 d0 88 41 03 f3 c3 90 c7 05 c8 1b 9e c0 01 00 00 00 0f ae f8 89 f6 <c6> 05 00 00 00 00 01 c3 89 f6 8d bc 27 00 00 00 00 8d 50 d0 83

EIP: [<c068124f>] sysrq_handle_crash+0xf/0x20 SS:ESP 0068:ef4dbf24

CR2: 00000000000000000

Type **help log** for more information on the command usage.



NOTE

The kernel message buffer includes the most essential information about the system crash and, as such, it is always dumped first in to the **vmcore-dmesg.txt** file. This is useful when an attempt to get the full **vmcore** file failed, for example because of lack of space on the target location. By default, **vmcore-dmesg.txt** is located in the /var/crash/ directory.

Displaying a backtrace

To display the kernel stack trace, use the bt command.

crash> bt

PID: 5591 TASK: f196d560 CPU: 2 COMMAND: "bash"

#0 [ef4dbdcc] crash_kexec at c0494922

#1 [ef4dbe20] oops_end at c080e402

```
#2 [ef4dbe34] no_context at c043089d
#3 [ef4dbe58] bad_area at c0430b26
#4 [ef4dbe6c] do_page_fault at c080fb9b
#5 [ef4dbee4] error code (via page fault) at c080d809
  EAX: 00000063 EBX: 00000063 ECX: c09e1c8c EDX: 00000000 EBP: 00000000
  DS: 007b
              ESI: c0a09ca0 ES: 007b
                                        EDI: 00000286 GS: 00e0
              EIP: c068124f ERR: fffffff EFLAGS: 00010096
  CS: 0060
#6 [ef4dbf18] sysrq_handle_crash at c068124f
#7 [ef4dbf24] __handle_sysrq at c0681469
#8 [ef4dbf48] write sysrq trigger at c068150a
#9 [ef4dbf54] proc reg write at c0569ec2
#10 [ef4dbf74] vfs write at c051de4e
#11 [ef4dbf94] sys_write at c051e8cc
#12 [ef4dbfb0] system call at c0409ad5
  EAX: fffffda EBX: 00000001 ECX: b7776000 EDX: 00000002
  DS: 007b ESI: 00000002 ES: 007b
                                        EDI: b7776000
  SS: 007b ESP: bfcb2088 EBP: bfcb20b4 GS: 0033
              EIP: 00edc416 ERR: 00000004 EFLAGS: 00000246
  CS: 0073
```

Type **bt** <*pid>* to display the backtrace of a specific process or type **help bt** for more information on **bt** usage.

Displaying a process status

• To display the status of processes in the system, use the **ps** command.

```
crash> ps
 PID PPID CPU TASK ST %MEM VSZ RSS COMM
       0 0 c09dc560 RU 0.0 0 0 [swapper]
       0 1 f7072030 RU 0.0
                             0
                                 0 [swapper]
      0 2 f70a3a90 RU 0.0
  0
                            0
                                0 [swapper]
       0 3 f70ac560 RU 0.0 0 0 [swapper]
      0 1 f705ba90 IN 0.0 2828 1424 init
... several lines omitted ...
       1 1 f2592560 IN 0.0 12876 784 auditd
 5566
        1 2 ef427560 IN 0.0 12876 784 auditd
 5567
 5587 5132 0 f196d030 IN 0.0 11064 3184 sshd
> 5591 5587 2 f196d560 RU 0.0 5084 1648 bash
```

Use **ps** *<pid>* to display the status of a single specific process. Use *help ps* for more information on **ps** usage.

Displaying virtual memory information

 To display basic virtual memory information, type the vm command at the interactive prompt.

```
crash> vm
PID: 5591 TASK: f196d560 CPU: 2 COMMAND: "bash"
    MM PGD RSS TOTAL_VM
f19b5900 ef9c6000 1648k 5084k
    VMA START END FLAGS FILE
f1bb0310 242000 260000 8000875 /lib/ld-2.12.so
f26af0b8 260000 261000 8100871 /lib/ld-2.12.so
efbc275c 261000 262000 8100873 /lib/ld-2.12.so
```

```
efbc2a18 268000 3ed000 8000075 /lib/libc-2.12.so
efbc23d8 3ed000 3ee000 8000070 /lib/libc-2.12.so
efbc2888 3ee000 3f0000 8100071 /lib/libc-2.12.so
efbc2cd4 3f0000 3f1000 8100073 /lib/libc-2.12.so
efbc243c 3f1000 3f4000 100073
efbc28ec 3f6000 3f9000 8000075 /lib/libdl-2.12.so
efbc2568 3f9000 3fa000 8100071 /lib/libdl-2.12.so
efbc2f2c 3fa000 3fb000 8100073 /lib/libdl-2.12.so
f26af888 7e6000 7fc000 8000075 /lib/libtinfo.so.5.7
f26aff2c 7fc000 7ff000 8100073 /lib/libtinfo.so.5.7
efbc211c d83000 d8f000 8000075 /lib/libnss_files-2.12.so
efbc2504 d8f000 d90000 8100071 /lib/libnss_files-2.12.so
efbc2950 d90000 d91000 8100073 /lib/libnss_files-2.12.so
f26afe00 edc000 edd000 4040075
f1bb0a18 8047000 8118000 8001875 /bin/bash
f1bb01e4 8118000 811d000 8101873 /bin/bash
f1bb0c70 811d000 8122000 100073
f26afae0 9fd9000 9ffa000 100073
... several lines omitted ...
```

Use **vm** <**pid**> to display information on a single specific process, or use **help vm** for more information on **vm** usage.

Displaying open files

• To display information about open files, use the **files** command.

```
crash> files
PID: 5591 TASK: f196d560 CPU: 2 COMMAND: "bash"
ROOT: / CWD: /root
FD FILE DENTRY INODE TYPE PATH
0 f734f640 eedc2c6c eecd6048 CHR /pts/0
1 efade5c0 eee14090 f00431d4 REG /proc/sysrq-trigger
2 f734f640 eedc2c6c eecd6048 CHR /pts/0
10 f734f640 eedc2c6c eecd6048 CHR /pts/0
255 f734f640 eedc2c6c eecd6048 CHR /pts/0
```

Use **files** < pid> to display files opened by only one selected process, or use help files for more information on files usage.

15.4. USING KERNEL OOPS ANALYZER

The Kernel Oops Analyzer tool analyzes the crash dump by comparing the oops messages with known issues in the knowledge base.

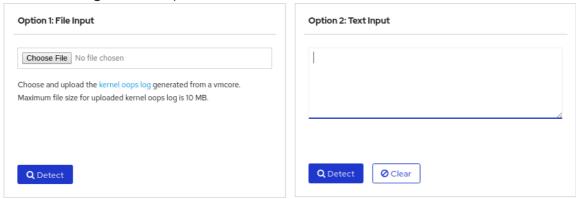
Prerequisites

• Secure an oops message to feed the Kernel Oops Analyzer.

Procedure

- 1. Access the Kernel Oops Analyzer tool.
- 2. To diagnose a kernel crash issue, upload a kernel oops log generated in vmcore.

Alternatively you can also diagnose a kernel crash issue by providing a text message or a
 vmcore-dmesg.txt as an input.



3. Click **DETECT** to compare the oops message based on information from the **makedumpfile** against known solutions.

Additional resources

- Red Hat Customer Portal Labs
- Kernel Oops Analyzer
- A Guide to Unexpected System Restarts

15.5. THE KDUMP HELPER TOOL

The Kdump Helper tool helps to set up the **kdump** using the provided information. Kdump Helper generates a configuration script based on your preferences. Initiating and running the script on your server sets up the **kdump** service.

Additional resources

Kdump Helper

CHAPTER 16. USING EARLY KDUMP TO CAPTURE BOOT TIME CRASHES

As a system administrator, you can utilize the **early kdump** support of the **kdump** service to capture a vmcore file of the crashing kernel during the early stages of the booting process. This section describes what **early kdump** is, how to configure it, and how to check the status of this mechanism.

16.1. WHAT IS EARLY KDUMP

Kernel crashes during the booting phase occur when the **kdump** service is not yet started, and cannot facilitate capturing and saving the contents of the crashed kernel's memory. Therefore, the vital information for troubleshooting is lost.

To address this problem, RHEL 8 introduced the **early kdump** feature as a part of the **kdump** service.

16.2. ENABLING EARLY KDUMP

This section describes how to enable the **early kdump** feature to eliminate the risk of losing information about the early boot kernel crashes.

Prerequisites

- An active Red Hat Enterprise Linux subscription.
- A repository containing the **kexec-tools** package for your system CPU architecture
- Fulfilled **kdump** configuration and targets requirements.

Procedure

1. Verify that the **kdump** service is enabled and active:

systemctl is-enabled kdump.service && systemctl is-active kdump.service enabled active

If **kdump** is not enabled and running, set all required configurations and verify that **kdump** service is enabled.

2. Rebuild the **initramfs** image of the booting kernel with the **early kdump** functionality:

dracut -f --add earlykdump

3. Add the **rd.earlykdump** kernel command line parameter:

grubby --update-kernel=/boot/vmlinuz-\$(uname -r) --args="rd.earlykdump"

4. Reboot the system to reflect the changes

reboot

5. Optionally, verify that **rd.earlykdump** was successfully added and **early kdump** feature was enabled:

-

cat /proc/cmdline

BOOT_IMAGE=(hd0,msdos1)/vmlinuz-4.18.0-187.el8.x86_64 root=/dev/mapper/rhel-root ro crashkernel=auto resume=/dev/mapper/rhel-swap rd.lvm.lv=rhel/root rd.lvm.lv=rhel/swap rhgb quiet rd.earlykdump

journalctl -x | grep early-kdump

Mar 20 15:44:41 redhat dracut-cmdline[304]: early-kdump is enabled.

Mar 20 15:44:42 redhat dracut-cmdline[304]: kexec: loaded early-kdump kernel

Additional resources

- /usr/share/doc/kexec-tools/early-kdump-howto.txt file
- What is early kdump support and how do I configure it?
- Enabling and disabling the kdump service

CHAPTER 17. APPLYING PATCHES WITH KERNEL LIVE PATCHING

You can use the Red Hat Enterprise Linux kernel live patching solution to patch a running kernel without rebooting or restarting any processes.

With this solution, system administrators:

- Can immediately apply critical security patches to the kernel.
- Do not have to wait for long-running tasks to complete, for users to log off, or for scheduled downtime.
- Control the system's uptime more and do not sacrifice security or stability.

Note that not every critical or important CVE will be resolved using the kernel live patching solution. Our goal is to reduce the required reboots for security-related patches, not to eliminate them entirely. For more details about the scope of live patching, see the *Customer Portal Solutions article*.



WARNING

Some incompatibilities exist between kernel live patching and other kernel subcomponents. Read the Limitations of kpatch carefully before using kernel live patching.

17.1. LIMITATIONS OF KPATCH

- The **kpatch** feature is not a general-purpose kernel upgrade mechanism. It is used for applying simple security and bug fix updates when rebooting the system is not immediately possible.
- Do not use the **SystemTap** or **kprobe** tools during or after loading a patch. The patch could fail to take effect until after such probes have been removed.

17.2. SUPPORT FOR THIRD-PARTY LIVE PATCHING

The **kpatch** utility is the only kernel live patching utility supported by Red Hat with the RPM modules provided by Red Hat repositories. Red Hat will not support any live patches which were not provided by Red Hat itself.

If you require support for an issue that arises with a third-party live patch, Red Hat recommends that you open a case with the live patching vendor at the outset of any investigation in which a root cause determination is necessary. This allows the source code to be supplied if the vendor allows, and for their support organization to provide assistance in root cause determination prior to escalating the investigation to Red Hat Support.

For any system running with third-party live patches, Red Hat reserves the right to ask for reproduction with Red Hat shipped and supported software. In the event that this is not possible, we require a similar system and workload be deployed on your test environment without live patches applied, to confirm if the same behavior is observed.

For more information about third-party software support policies, see How does Red Hat Global Support Services handle third-party software, drivers, and/or uncertified hardware/hypervisors or guest operating systems?

17.3. ACCESS TO KERNEL LIVE PATCHES

Kernel live patching capability is implemented as a kernel module (**kmod**) that is delivered as an RPM package.

All customers have access to kernel live patches, which are delivered through the usual channels. However, customers who do not subscribe to an extended support offering will lose access to new patches for the current minor release once the next minor release becomes available. For example, customers with standard subscriptions will only be able to live patch RHEL 8.2 kernel until the RHEL 8.3 kernel is released.

17.4. COMPONENTS OF KERNEL LIVE PATCHING

The components of kernel live patching are as follows:

Kernel patch module

- The delivery mechanism for kernel live patches.
- A kernel module which is built specifically for the kernel being patched.
- The patch module contains the code of the desired fixes for the kernel.
- The patch modules register with the **livepatch** kernel subsystem and provide information about original functions to be replaced, with corresponding pointers to the replacement functions. Kernel patch modules are delivered as RPMs.
- The naming convention is **kpatch_<kernel version>_<kpatch version>_<kpatch release>**. The "kernel version" part of the name has *dots* replaced with *underscores*.

The kpatch utility

A command-line utility for managing patch modules.

The kpatch service

A **systemd** service required by **multiuser.target**. This target loads the kernel patch module at boot time.

The kpatch-dnf package

A DNF plugin delivered in the form of an RPM package. This plugin manages automatic subscription to kernel live patches.

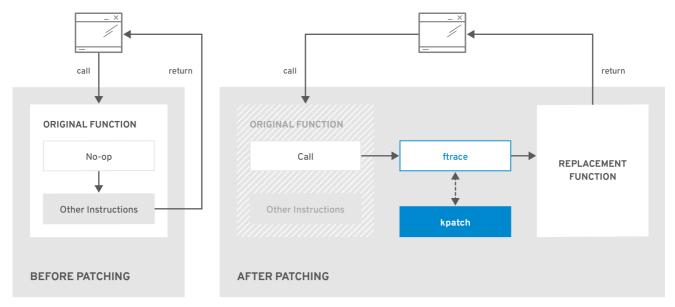
17.5. HOW KERNEL LIVE PATCHING WORKS

The **kpatch** kernel patching solution uses the **livepatch** kernel subsystem to redirect old functions to new ones. When a live kernel patch is applied to a system, the following things happen:

- 1. The kernel patch module is copied to the /var/lib/kpatch/ directory and registered for reapplication to the kernel by **systemd** on next boot.
- 2. The kpatch module is loaded into the running kernel and the new functions are registered to the **ftrace** mechanism with a pointer to the location in memory of the new code.

3. When the kernel accesses the patched function, it is redirected by the **ftrace** mechanism which bypasses the original functions and redirects the kernel to patched version of the function.

Figure 17.1. How kernel live patching works



RHEL 424549 0119

17.6. SUBSCRIBING THE CURRENTLY INSTALLED KERNELS TO THE LIVE PATCHING STREAM

A kernel patch module is delivered in an RPM package, specific to the version of the kernel being patched. Each RPM package will be cumulatively updated over time.

The following procedure explains how to subscribe to all future cumulative live patching updates for a given kernel. Because live patches are cumulative, you cannot select which individual patches are deployed for a given kernel.



WARNING

Red Hat does not support any third party live patches applied to a Red Hat supported system.

Prerequisites

Root permissions

Procedure

1. Optionally, check your kernel version:

uname -r 4.18.0-94.el8.x86_64 2. Search for a live patching package that corresponds to the version of your kernel:

yum search \$(uname -r)

3. Install the live patching package:

yum install "kpatch-patch = \$(uname -r)"

The command above installs and applies the latest cumulative live patches for that specific kernel only.

If the version of a live patching package is 1-1 or higher, the package will contain a patch module. In that case the kernel will be automatically patched during the installation of the live patching package.

The kernel patch module is also installed into the /var/lib/kpatch/ directory to be loaded by the systemd system and service manager during the future reboots.



NOTE

An empty live patching package will be installed when there are no live patches available for a given kernel. An empty live patching package will have a <code>kpatch_version-kpatch_release</code> of O-O, for example <code>kpatch-patch-4_18_0-94-O-O.el8.x86_64.rpm</code>. The installation of the empty RPM subscribes the system to all future live patches for the given kernel.

4. Optionally, verify that the kernel is patched:

```
# kpatch list
Loaded patch modules:
kpatch_4_18_0_94_1_1 [enabled]

Installed patch modules:
kpatch_4_18_0_94_1_1 (4.18.0-94.el8.x86_64)
...
```

The output shows that the kernel patch module has been loaded into the kernel, which is now patched with the latest fixes from the **kpatch-patch-4_18_0-94-1-1.el8.x86_64.rpm** package.

Additional resources

- kpatch(1) manual page
- Configuring basic system settings Red Hat Enterprise Linux

17.7. AUTOMATICALLY SUBSCRIBING ANY FUTURE KERNEL TO THE LIVE PATCHING STREAM

You can use the **kpatch-dnf** DNF plugin to subscribe your system to fixes delivered by the kernel patch module, also known as kernel live patches. The plugin enables **automatic** subscription for any kernel the system currently uses, and also for kernels **to-be-installed in the future**

Prerequisites

Root permissions.

Procedure

1. Optionally, check all installed kernels and the kernel you are currently running:

2. Install the **kpatch-dnf** plugin:

yum install kpatch-dnf

3. Enable automatic subscription to kernel live patches:

This command subscribes all currently installed kernels to receiving kernel live patches. The command also installs and applies the latest cumulative live patches, if any, for all installed kernels.

In the future, when you update the kernel, live patches will automatically be installed during the new kernel installation process.

The kernel patch module is also installed into the /var/lib/kpatch/ directory to be loaded by the system and service manager during future reboots.



NOTE

An empty live patching package will be installed when there are no live patches available for a given kernel. An empty live patching package will have a <code>kpatch_version-kpatch_release</code> of O-O, for example <code>kpatch-patch-4_18_0-240-0-0.el8.x86_64.rpm</code>. The installation of the empty RPM subscribes the system to all future live patches for the given kernel.

4. Optionally, verify that all installed kernels were patched:

```
# kpatch list
Loaded patch modules:
kpatch_4_18_0_240_10_1_0_1 [enabled]

Installed patch modules:
kpatch_4_18_0_240_10_1_0_1 (4.18.0-240.10.1.el8_3.x86_64)
kpatch_4_18_0_240_15_1_0_2 (4.18.0-240.15.1.el8_3.x86_64)
```

The output shows that both the kernel you are running, and the other installed kernel have been patched with fixes from **kpatch-patch-4_18_0-240_10_1-0-1.rpm** and **kpatch-patch-4_18_0-240_15_1-0-1.rpm** packages respectively.

Additional resources

- kpatch(1) and dnf-kpatch(8) manual pages
- Configuring basic system settings in Red Hat Enterprise Linux

17.8. UPDATING KERNEL PATCH MODULES

Since kernel patch modules are delivered and applied through RPM packages, updating a cumulative kernel patch module is like updating any other RPM package.

Prerequisites

• The system is subscribed to the live patching stream, as described in Subscribing the currently installed kernels to the live patching stream.

Procedure

• Update to a new cumulative version for the current kernel:

```
# yum update "kpatch-patch = $(uname -r)"
```

The command above automatically installs and applies any updates that are available for the currently running kernel. Including any future released cumulative live patches.

Alternatively, update all installed kernel patch modules:

yum update "kpatch-patch"



NOTE

When the system reboots into the same kernel, the kernel is automatically live patched again by the **kpatch.service** systemd service.

Additional resources

Configuring basic system settings Red Hat Enterprise Linux

17.9. REMOVING THE LIVE PATCHING PACKAGE

The following procedure describes how to disable the Red Hat Enterprise Linux kernel live patching solution by removing the live patching package.

Prerequisites

- Root permissions
- The live patching package is installed.

Procedure

1. Select the live patching package:

```
# yum list installed | grep kpatch-patch
kpatch-patch-4_18_0-94.x86_64
                                             @@commandline
                                  1-1.el8
```

The example output above lists live patching packages that you installed.

2. Remove the live patching package:

```
# yum remove kpatch-patch-4_18_0-94.x86_64
```

When a live patching package is removed, the kernel remains patched until the next reboot, but the kernel patch module is removed from disk. On future reboot, the corresponding kernel will no longer be patched.

- 3. Reboot your system.
- 4. Verify that the live patching package has been removed:

yum list installed | grep kpatch-patch

The command displays no output if the package has been successfully removed.

5. Optionally, verify that the kernel live patching solution is disabled:

```
# kpatch list
```

Loaded patch modules:

The example output shows that the kernel is not patched and the live patching solution is not active because there are no patch modules that are currently loaded.



IMPORTANT

Currently, Red Hat does not support reverting live patches without rebooting your system. In case of any issues, contact our support team.

Additional resources

- kpatch(1) manual page
- Configuring basic system settings Red Hat Enterprise Linux

17.10. UNINSTALLING THE KERNEL PATCH MODULE

The following procedure describes how to prevent the Red Hat Enterprise Linux kernel live patching solution from applying a kernel patch module on subsequent boots.

Prerequisites

- Root permissions
- A live patching package is installed.
- A kernel patch module is installed and loaded.

Procedure

1. Select a kernel patch module:

```
# kpatch list
Loaded patch modules:
kpatch_4_18_0_94_1_1 [enabled]

Installed patch modules:
kpatch_4_18_0_94_1_1 (4.18.0-94.el8.x86_64)
...
```

2. Uninstall the selected kernel patch module:

```
# kpatch uninstall kpatch_4_18_0_94_1_1
uninstalling kpatch_4_18_0_94_1_1 (4.18.0-94.el8.x86_64)
```

• Note that the uninstalled kernel patch module is still loaded:

```
# kpatch list
Loaded patch modules:
kpatch_4_18_0_94_1_1 [enabled]
Installed patch modules:
</P>
```

When the selected module is uninstalled, the kernel remains patched until the next reboot, but the kernel patch module is removed from disk.

3. Reboot your system.

4. Optionally, verify that the kernel patch module has been uninstalled:

kpatch list

Loaded patch modules:

. . .

The example output above shows no loaded or installed kernel patch modules, therefore the kernel is not patched and the kernel live patching solution is not active.



IMPORTANT

Currently, Red Hat does not support reverting live patches without rebooting your system. In case of any issues, contact our support team.

Additional resources

kpatch(1) manual page

17.11. DISABLING KPATCH.SERVICE

The following procedure describes how to prevent the Red Hat Enterprise Linux kernel live patching solution from applying all kernel patch modules globally on subsequent boots.

Prerequisites

- Root permissions
- A live patching package is installed.
- A kernel patch module is installed and loaded.

Procedure

1. Verify **kpatch.service** is enabled:

systemctl is-enabled kpatch.service enabled

2. Disable kpatch.service:

systemctl disable kpatch.service

Removed / etc/system / multi-user. target. wants / kpatch. service.

• Note that the applied kernel patch module is still loaded:

```
# kpatch list
Loaded patch modules:
kpatch_4_18_0_94_1_1 [enabled]

Installed patch modules:
kpatch_4_18_0_94_1_1 (4.18.0-94.el8.x86_64)
```

3. Reboot your system.

4. Optionally, verify the status of **kpatch.service**:

systemctl status kpatch.service

kpatch.service - "Apply kpatch kernel patches"
 Loaded: loaded (/usr/lib/systemd/system/kpatch.service; disabled; vendor preset: disabled)
 Active: inactive (dead)

The example output testifies that **kpatch.service** has been disabled and is not running. Thereby, the kernel live patching solution is not active.

5. Verify that the kernel patch module has been unloaded:

kpatch list

Loaded patch modules:

<NO RESULT>

Installed patch modules:

kpatch_4_18_0_94_1_1 (4.18.0-94.el8.x86_64)

The example output above shows that a kernel patch module is still installed but the kernel is not patched.



IMPORTANT

Currently, Red Hat does not support reverting live patches without rebooting your system. In case of any issues, contact our support team.

Additional resources

- kpatch(1) manual page
- Configuring basic system settings Red Hat Enterprise Linux

CHAPTER 18. SETTING LIMITS FOR APPLICATIONS

You can use the *control groups* (**cgroups**) kernel functionality to set limits, prioritize or isolate the hardware resources of processes. This allows you to granularly control resource usage of applications to utilize them more efficiently.

18.1. UNDERSTANDING CONTROL GROUPS

Control groups is a Linux kernel feature that enables you to organize processes into hierarchically ordered groups - **cgroups**. The hierarchy (control groups tree) is defined by providing structure to **cgroups** virtual file system, mounted by default on the /sys/fs/cgroup/ directory. It is done manually by creating and removing sub-directories in /sys/fs/cgroup/. Alternatively, by using the system and service manager.

The resource controllers (a kernel component) then modify the behavior of processes in **cgroups** by limiting, prioritizing or allocating system resources, (such as CPU time, memory, network bandwidth, or various combinations) of those processes.

The added value of **cgroups** is process aggregation which enables division of hardware resources among applications and users. Thereby an increase in overall efficiency, stability and security of users' environment can be achieved.

Control groups version 1

Control groups version 1 (cgroups-v1) provide a per-resource controller hierarchy. It means that each resource, such as CPU, memory, I/O, and so on, has its own control group hierarchy. It is possible to combine different control group hierarchies in a way that one controller can coordinate with another one in managing their respective resources. However, the two controllers may belong to different process hierarchies, which does not permit their proper coordination.

The **cgroups-v1** controllers were developed across a large time span and as a result, the behavior and naming of their control files is not uniform.

Control groups version 2

The problems with controller coordination, which stemmed from hierarchy flexibility, led to the development of *control groups version 2*.

Control groups version 2 (cgroups-v2) provides a single control group hierarchy against which all resource controllers are mounted.

The control file behavior and naming is consistent among different controllers.



WARNING

RHEL 8 provides **cgroups-v2** as a technology preview with a limited number of resource controllers. For more information about the relevant resource controllers, see the *cgroups-v2* release note.

This sub-section was based on a Devconf.cz 2019 presentation.[1]

Additional resources

- What kernel resource controllers are
- cgroups(7) manual page
- Role of systemd in control groups

18.2. WHAT KERNEL RESOURCE CONTROLLERS ARE

The functionality of control groups is enabled by kernel resource controllers. RHEL 8 supports various controllers for *control groups version 1* (**cgroups-v1**) and *control groups version 2* (**cgroups-v2**).

A resource controller, also called a control group subsystem, is a kernel subsystem that represents a single resource, such as CPU time, memory, network bandwidth or disk I/O. The Linux kernel provides a range of resource controllers that are mounted automatically by the **systemd** system and service manager. Find a list of currently mounted resource controllers in the /**proc/cgroups** file.

The following controllers are available for **cgroups-v1**:

- **blkio** can set limits on input/output access to and from block devices.
- **cpu** can adjust the parameters of the Completely Fair Scheduler (CFS) scheduler for control group's tasks. It is mounted together with the **cpuacct** controller on the same mount.
- **cpuacct** creates automatic reports on CPU resources used by tasks in a control group. It is mounted together with the **cpu** controller on the same mount.
- **cpuset** can be used to restrict control group tasks to run only on a specified subset of CPUs and to direct the tasks to use memory only on specified memory nodes.
- **devices** can control access to devices for tasks in a control group.
- **freezer** can be used to suspend or resume tasks in a control group.
- **memory** can be used to set limits on memory use by tasks in a control group and generates automatic reports on memory resources used by those tasks.
- net_cls tags network packets with a class identifier (classid) that enables the Linux traffic controller (the tc command) to identify packets that originate from a particular control group task. A subsystem of net_cls, the net_filter (iptables), can also use this tag to perform actions on such packets. The net_filter tags network sockets with a firewall identifier (fwid) that allows the Linux firewall (through iptables command) to identify packets originating from a particular control group task.
- **net prio** sets the priority of network traffic.
- pids can set limits for a number of processes and their children in a control group.
- **perf_event** can group tasks for monitoring by the **perf** performance monitoring and reporting utility.
- **rdma** can set limits on Remote Direct Memory Access/InfiniBand specific resources in a control group.
- **hugetlb** can be used to limit the usage of large size virtual memory pages by tasks in a control group.

The following controllers are available for **cgroups-v2**:

- io A follow-up to blkio of cgroups-v1.
- memory A follow-up to memory of cgroups-v1.
- pids Same as pids in cgroups-v1.
- rdma Same as rdma in cgroups-v1.
- cpu A follow-up to cpu and cpuacct of cgroups-v1.
- **cpuset** Supports only the core functionality (**cpus{,.effective}**, **mems{,.effective}**) with a new partition feature.
- **perf_event** Support is inherent, no explicit control file. You can specify a **v2 cgroup** as a parameter to the **perf** command that will profile all the tasks within that **cgroup**.



IMPORTANT

A resource controller can be used either in a **cgroups-v1** hierarchy or a **cgroups-v2** hierarchy, not simultaneously in both.

Additional resources

- cgroups(7) manual page
- Documentation in /usr/share/doc/kernel-doc-<kernel_version>/Documentation/cgroups-v1/ directory

18.3. WHAT NAMESPACES ARE

Namespaces are one of the most important methods for organizing and identifying software objects.

A namespace wraps a global system resource (for example a mount point, a network device, or a hostname) in an abstraction that makes it appear to processes within the namespace that they have their own isolated instance of the global resource. One of the most common technologies that utilize namespaces are containers.

Changes to a particular global resource are visible only to processes in that namespace and do not affect the rest of the system or other namespaces.

To inspect which namespaces a process is a member of, you can check the symbolic links in the /proc/<*PID*>/ns/ directory.

The following table shows supported namespaces and resources which they isolate:

Namespace	Isolates
Mount	Mount points
UTS	Hostname and NIS domain name
IPC	System V IPC, POSIX message queues

PID	Process IDs
Network	Network devices, stacks, ports, etc
User	User and group IDs
Control groups	Control group root directory

Isolates

Additional resources

Namespace

- namespaces(7) and cgroup_namespaces(7) manual pages
- Understanding control groups

18.4. SETTING CPU LIMITS TO APPLICATIONS USING CGROUPS-V1

Sometimes an application consumes a lot of CPU time, which may negatively impact the overall health of your environment. Use the /sys/fs/ virtual file system to configure CPU limits to an application using control groups version 1 (cgroups-v1).

Prerequisites

- An application whose CPU consumption you want to restrict.
- Verify that the **cgroups-v1** controllers were mounted:

mount -I | grep cgroup

 $tmpfs\ on\ /sys/fs/cgroup\ type\ tmpfs\ (ro,nosuid,nodev,noexec,seclabel,mode=755)$ $cgroup\ on\ /sys/fs/cgroup/systemd\ type\ cgroup$

 $(rw, nosuid, nodev, noexec, relatime, seclabel, xattr, release_agent=/usr/lib/systemd/systemd-cgroups-agent, name=systemd)\\$

cgroup on /sys/fs/cgroup/cpu,cpuacct type cgroup

(rw,nosuid,nodev,noexec,relatime,seclabel,cpu,cpuacct)

cgroup on /sys/fs/cgroup/perf_event type cgroup

(rw,nosuid,nodev,noexec,relatime,seclabel,perf_event)

cgroup on /sys/fs/cgroup/pids type cgroup (rw,nosuid,nodev,noexec,relatime,seclabel,pids)

• • •

Procedure

1. Identify the process ID (PID) of the application you want to restrict in CPU consumption:

top

top - 11:34:09 up 11 min, 1 user, load average: 0.51, 0.27, 0.22 Tasks: 267 total, 3 running, 264 sleeping, 0 stopped, 0 zombie %Cpu(s): 49.0 us, 3.3 sy, 0.0 ni, 47.5 id, 0.0 wa, 0.2 hi, 0.0 si, 0.0 st MiB Mem: 1826.8 total, 303.4 free, 1046.8 used, 476.5 buff/cache

```
MiB Swap: 1536.0 total, 1396.0 free, 140.0 used. 616.4 avail Mem

PID USER PR NI VIRT RES SHR S %CPU %MEM TIME+ COMMAND 6955 root 20 0 228440 1752 1472 R 99.3 0.1 0:32.71 sha1sum 5760 jdoe 20 0 3603868 205188 64196 S 3.7 11.0 0:17.19 gnome-shell 6448 jdoe 20 0 743648 30640 19488 S 0.7 1.6 0:02.73 gnome-terminal-1 root 20 0 245300 6568 4116 S 0.3 0.4 0:01.87 systemd 505 root 20 0 0 0 0 1 0.3 0.0 0:00.75 kworker/u4:4-events_unbound ...
```

The example output of the **top** program reveals that **PID 6955** (illustrative application **sha1sum**) consumes a lot of CPU resources.

2. Create a sub-directory in the **cpu** resource controller directory:

mkdir /sys/fs/cgroup/cpu/Example/

The directory above represents a control group, where you can place specific processes and apply certain CPU limits to the processes. At the same time, some **cgroups-v1** interface files and **cpu** controller-specific files will be created in the directory.

3. Optionally, inspect the newly created control group:

```
# II /sys/fs/cgroup/cpu/Example/
-rw-r----. 1 root root 0 Mar 11 11:42 cgroup.clone_children
-rw-r---- 1 root root 0 Mar 11 11:42 cgroup.procs
-r—r—r--. 1 root root 0 Mar 11 11:42 cpuacct.stat
-rw-r---- 1 root root 0 Mar 11 11:42 cpuacct.usage
-r-r-r--. 1 root root 0 Mar 11 11:42 cpuacct.usage all
-r-r-r--- 1 root root 0 Mar 11 11:42 cpuacct.usage_percpu
-r-r-r--- 1 root root 0 Mar 11 11:42 cpuacct.usage percpu sys
-r-r-r--- 1 root root 0 Mar 11 11:42 cpuacct.usage percpu user
-r-r-r-- 1 root root 0 Mar 11 11:42 cpuacct.usage sys
-r-r-r--. 1 root root 0 Mar 11 11:42 cpuacct.usage_user
-rw-r--- 1 root root 0 Mar 11 11:42 cpu.cfs period us
-rw-r--r--. 1 root root 0 Mar 11 11:42 cpu.cfs_quota_us
-rw-r---- 1 root root 0 Mar 11 11:42 cpu.rt_period_us
-rw-r---- 1 root root 0 Mar 11 11:42 cpu.rt_runtime_us
-rw-r---- 1 root root 0 Mar 11 11:42 cpu.shares
-r-r-r--- 1 root root 0 Mar 11 11:42 cpu.stat
-rw-r---- 1 root root 0 Mar 11 11:42 notify_on_release
-rw-r--r--. 1 root root 0 Mar 11 11:42 tasks
```

The example output shows files, such as **cpuacct.usage**, **cpu.cfs._period_us**, that represent specific configurations and/or limits, which can be set for processes in the **Example** control group. Notice that the respective file names are prefixed with the name of the control group controller to which they belong.

By default, the newly created control group inherits access to the system's entire CPU resources without a limit.

4. Configure CPU limits for the control group:

```
# echo "1000000" > /sys/fs/cgroup/cpu/Example/cpu.cfs_period_us
# echo "200000" > /sys/fs/cgroup/cpu/Example/cpu.cfs_quota_us
```

The **cpu.cfs_period_us** file represents a period of time in microseconds (µs, represented here as "us") for how frequently a control group's access to CPU resources should be reallocated. The upper limit is 1 second and the lower limit is 1000 microseconds.

The **cpu.cfs_quota_us** file represents the total amount of time in microseconds for which all processes collectively in a control group can run during one period (as defined by **cpu.cfs_period_us**). As soon as processes in a control group, during a single period, use up all the time specified by the quota, they are throttled for the remainder of the period and not allowed to run until the next period. The lower limit is 1000 microseconds.

The example commands above set the CPU time limits so that all processes collectively in the **Example** control group will be able to run only for 0.2 seconds (defined by **cpu.cfs_quota_us**) out of every 1 second (defined by **cpu.cfs_period_us**).

5. Optionally, verify the limits:

```
# cat /sys/fs/cgroup/cpu/Example/cpu.cfs_period_us /sys/fs/cgroup/cpu/Example/cpu.cfs_quota_us 1000000 200000
```

6. Add the application's PID to the **Example** control group:

```
# echo "6955" > /sys/fs/cgroup/cpu/Example/cgroup.procs
or
# echo "6955" > /sys/fs/cgroup/cpu/Example/tasks
```

The previous command ensures that a desired application becomes a member of the **Example** control group and hence does not exceed the CPU limits configured for the **Example** control group. The PID should represent an existing process in the system. The **PID 6955** here was assigned to process **sha1sum** /dev/zero &, used to illustrate the use-case of the **cpu** controller.

7. Verify that the application runs in the specified control group:

```
# cat /proc/6955/cgroup

12:cpuset:/

11:hugetlb:/

10:net_cls,net_prio:/

9:memory:/user.slice/user-1000.slice/user@1000.service

8:devices:/user.slice

7:blkio:/

6:freezer:/

5:rdma:/

4:pids:/user.slice/user-1000.slice/user@1000.service

3:perf_event:/

2:cpu,cpuacct:/Example

1:name=systemd:/user.slice/user-1000.slice/user@1000.service/gnome-terminal-server.service
```

The example output above shows that the process of the desired application runs in the **Example** control group, which applies CPU limits to the application's process.

8. Identify the current CPU consumption of your throttled application:

top

top - 12:28:42 up 1:06, 1 user, load average: 1.02, 1.02, 1.00

Tasks: 266 total, 6 running, 260 sleeping, 0 stopped, 0 zombie

%Cpu(s): 11.0 us, 1.2 sy, 0.0 ni, 87.5 id, 0.0 wa, 0.2 hi, 0.0 si, 0.2 st

MiB Mem : 1826.8 total, 287.1 free, 1054.4 used, 485.3 buff/cache

MiB Swap: 1536.0 total, 1396.7 free, 139.2 used. 608.3 avail Mem

PID USER PR NI VIRT RES SHR S %CPU %MEM TIME+ COMMAND

6955 root 20 0 228440 1752 1472 R 20.6 0.1 47:11.43 sha1sum

5760 jdoe 20 0 3604956 208832 65316 R 2.3 11.2 0:43.50 gnome-shell

6448 jdoe 20 0 743836 31736 19488 S 0.7 1.7 0:08.25 gnome-terminal
505 root 20 0 0 0 0 0 0.3 0.0 0:03.39 kworker/u4:4-events_unbound

4217 root 20 0 74192 1612 1320 S 0.3 0.1 0:01.19 spice-vdagentd

Notice that the CPU consumption of the PID 6955 has decreased from 99% to 20%.

Additional resources

- Understanding control groups
- What kernel resource controllers are
- cgroups(7), sysfs(5) manual pages

^[1] Linux Control Group v2 - An Introduction, Devconf.cz 2019 presentation by Waiman Long

CHAPTER 19. USING CGROUPS-V2 TO SET CPU LIMITS TO APPLICATIONS

Some applications use too much CPU time. This can negatively impact the overall health of your environment. You can use *control groups version 2* (**cgroups-v2**) to configure CPU limits to applications and restrict their consumption.

19.1. MOUNTING CGROUPS-V2

During the boot process, RHEL 8 mounts the **cgroup-v1** virtual filesystem by default. To utilize **cgroup-v2** functionality in limiting resources for your applications, manually configure the system.

Prerequisites

• You have root permissions.

Procedure

1. Configure the system to mount **cgroups-v2** by default during system boot by the **systemd** system and service manager:

```
# grubby --update-kernel=/boot/vmlinuz-$(uname -r) -- args="systemd.unified_cgroup_hierarchy=1"
```

This adds the necessary kernel command-line parameter to the current boot entry.

2. Reboot the system for the changes to take effect.

Verification steps

1. Optionally, verify that the **cgroups-v2** filesystem was mounted:

```
# mount -I | grep cgroup
cgroup2 on /sys/fs/cgroup type cgroup2
(rw,nosuid,nodev,noexec,relatime,seclabel,nsdelegate)
```

The cgroups-v2 filesystem was successfully mounted on the /sys/fs/cgroup/ directory.

2. Optionally, inspect the contents of the /sys/fs/cgroup/ directory:

```
-rw-r—r--. 1 root root 0 Apr 29 12:03 io.pressure
-r—r—r--. 1 root root 0 Apr 29 12:03 io.stat
-rw-r—r--. 1 root root 0 Apr 29 12:03 memory.pressure
-r—r--. 1 root root 0 Apr 29 12:03 memory.stat
drwxr-xr-x. 69 root root 0 Apr 29 12:03 system.slice
drwxr-xr-x. 3 root root 0 Apr 29 12:18 user.slice
```

The /sys/fs/cgroup/ directory, also called the *root control group*, by default, contains interface files (starting with cgroup) and controller-specific files such as cpuset.cpus.effective. In addition, there are some directories related to systemd, such as, /sys/fs/cgroup/init.scope, /sys/fs/cgroup/system.slice, and /sys/fs/cgroup/user.slice.

Additional resources

- Understanding control groups
- What kernel resource controllers are
- cgroups(7), sysfs(5) manual pages

19.2. SETTING UP THE CGROUP FOR CPU LIMITING

To restrict CPU consumption of your applications, you need to enable specific CPU controllers and create two levels of child control groups inside the /sys/fs/cgroup/ root control group. The root control group already contains some of the resource controllers by default. Therefore two levels of child control groups are advisable to ensure organizational clarity of cgroup files.

Prerequisites

- You have at least two applications whose CPU consumption you want to restrict.
- You have root permissions.
- You have mounted **cgroups-v2** filesystem.

Procedure

1. Identify the process IDs (PIDs) of applications whose CPU consumption you want to constrict:

```
# top
Tasks: 104 total, 3 running, 101 sleeping, 0 stopped, 0 zombie
%Cpu(s): 17.6 us, 81.6 sy, 0.0 ni, 0.0 id, 0.0 wa, 0.8 hi, 0.0 si, 0.0 st
MiB Mem: 3737.4 total, 3312.7 free, 133.3 used, 291.4 buff/cache
                                   0.0 used. 3376.1 avail Mem
MiB Swap: 4060.0 total, 4060.0 free,
  PID USER
              PR NI VIRT RES SHR S %CPU %MEM
                                                        TIME+ COMMAND
             20 0 18720 1756 1468 R 99.0 0.0 0:31.09 sha1sum
 34578 root
34579 root
             20 0 18720 1772 1480 R 99.0 0.0 0:30.54 sha1sum
   1 root 20 0 186192 13940 9500 S 0.0 0.4 0:01.60 systemd
                   0 0 0 S 0.0 0.0 0:00.01 kthreadd
   2 root
          20 0
   3 root
           0 -20
                       0
                            0 I 0.0 0.0 0:00.00 rcu gp
                   0
   4 root
           0 -20
                            0 I 0.0 0.0 0:00.00 rcu_par_gp
```

The example output reveals that **PID 34578** and **34579** (two illustrative applications of **sha1sum**) consume a lot of resources, namely CPU. Both are example applications used to demonstrate managing the **cgroups-v2** functionality.

2. Enable CPU-related controllers:

```
# echo "+cpu" >> /sys/fs/cgroup/cgroup.subtree_control
# echo "+cpuset" >> /sys/fs/cgroup/cgroup.subtree_control
```

These commands enable the **cpu** and **cpuset** controllers for the immediate children groups of the /**sys/fs/cgroup**/ root control group. A *child group* is where you can specify processes and apply limits to each of the processes based on your criteria.

Users can read the contents of the **cgroup.subtree_control** file at any level to get an idea of what controllers are going to be available for enablement in the immediate child group.



NOTE

By default, the /sys/fs/cgroup/cgroup.subtree_control file in the root control group contains memory and pids controllers.

3. Create the /sys/fs/cgroup/Example/ directory:

mkdir /sys/fs/cgroup/Example/

The /sys/fs/cgroup/Example/ directory defines a child group. Also, the previous step enabled the cpu and cpuset controllers for this child group.

When you create the /sys/fs/cgroup/Example/ directory, some cgroups-v2 interface files and cpu and cpuset controller-specific files are automatically created in the directory. The /sys/fs/cgroup/Example/ directory contains also controller-specific files for the memory and pids controllers.

4. Optionally, inspect the newly created child control group:

```
# II /sys/fs/cgroup/Example/
-r-r-r--. 1 root root 0 Jun 1 10:33 cgroup.controllers
-r-r-r--. 1 root root 0 Jun 1 10:33 cgroup.events
-rw-r---- 1 root root 0 Jun 1 10:33 cgroup.freeze
-rw-r----. 1 root root 0 Jun 1 10:33 cgroup.max.depth
-rw-r—r--. 1 root root 0 Jun 1 10:33 cgroup.max.descendants
-rw-r---- 1 root root 0 Jun 1 10:33 cgroup.procs
-r—r—r--. 1 root root 0 Jun 1 10:33 cgroup.stat
-rw-r—r--. 1 root root 0 Jun 1 10:33 cgroup.subtree_control
-rw-r--r--. 1 root root 0 Jun 1 10:33 cpuset.cpus
-r-r-r--. 1 root root 0 Jun 1 10:33 cpuset.cpus.effective
-rw-r—r--. 1 root root 0 Jun 1 10:33 cpuset.cpus.partition
-rw-r---- 1 root root 0 Jun 1 10:33 cpuset.mems
-r--r--r---. 1 root root 0 Jun 1 10:33 cpuset.mems.effective
-r—r—r--. 1 root root 0 Jun 1 10:33 cpu.stat
-rw-r----. 1 root root 0 Jun 1 10:33 cpu.weight
-rw-r---- 1 root root 0 Jun 1 10:33 cpu.weight.nice
   -r--r--. 1 root root 0 Jun 1 10:33 memory.events.local
```

The example output shows files such as **cpuset.cpus** and **cpu.max**. These files are specific to the **cpuset** and **cpu** controllers. The **cpuset** and **cpu** controllers are manually enabled for the root's (/sys/fs/cgroup/) direct child control groups using the /sys/fs/cgroup/cgroup.subtree_control file.

The directory also includes general **cgroup** control interface files such as **cgroup.procs** or **cgroup.controllers**, which are common to all control groups, regardless of enabled controllers.

The files such as **memory.high** and **pids.max** relate to the **memory** and **pids** controllers, which are in the root control group (/**sys/fs/cgroup**/), and are always enabled by default.

By default, the newly created child group inherits access to all of the system's CPU and memory resources, without any limits.

5. Enable the CPU-related controllers in /sys/fs/cgroup/Example/ to obtain controllers that are relevant only to CPU:

```
# echo "+cpu" >> /sys/fs/cgroup/Example/cgroup.subtree_control
# echo "+cpuset" >> /sys/fs/cgroup/Example/cgroup.subtree_control
```

These commands ensure that the immediate child control group will *only* have controllers relevant to limiting the CPU usage – not to **memory** or **pids** controllers.

6. Create the /sys/fs/cgroup/Example/tasks/ directory:

mkdir /sys/fs/cgroup/Example/tasks/

The /sys/fs/cgroup/Example/tasks/ directory defines a child group with files that relate purely to cpu and cpuset controllers.

7. Optionally, inspect another child control group:

```
# II /sys/fs/cgroup/Example/tasks
-r-r-r--. 1 root root 0 Jun 1 11:45 cgroup.controllers
-r-r-r-- 1 root root 0 Jun 1 11:45 cgroup.events
-rw-r---- 1 root root 0 Jun 1 11:45 cgroup.freeze
-rw-r---- 1 root root 0 Jun 1 11:45 cgroup.max.depth
-rw-r---- 1 root root 0 Jun 1 11:45 cgroup.max.descendants
-rw-r—r--. 1 root root 0 Jun 1 11:45 cgroup.procs
-r-r-r--- 1 root root 0 Jun 1 11:45 cgroup.stat
-rw-r—r--. 1 root root 0 Jun 1 11:45 cgroup.subtree control
-rw-r--- 1 root root 0 Jun 1 11:45 cgroup.threads
-rw-r--- 1 root root 0 Jun 1 11:45 cgroup.type
-rw-r--- 1 root root 0 Jun 1 11:45 cpu.max
-rw-r---- 1 root root 0 Jun 1 11:45 cpu.pressure
-rw-r--r--. 1 root root 0 Jun 1 11:45 cpuset.cpus
-r-r-r--. 1 root root 0 Jun 1 11:45 cpuset.cpus.effective
-rw-r—r--. 1 root root 0 Jun 1 11:45 cpuset.cpus.partition
-rw-r---- 1 root root 0 Jun 1 11:45 cpuset.mems
```

8. Ensure the processes that you want to limit compete for CPU time on the same CPU:

echo "1" > /sys/fs/cgroup/Example/tasks/cpuset.cpus

The previous command ensures that the processes you will place in the **Example/tasks** child control group, compete on the same CPU. This setting is important for the **cpu** controller to activate.



IMPORTANT

The **cpu** controller is only activated if the relevant child control group has at least 2 processes which compete for time on a single CPU.

Additional resources

- Understanding control groups
- What kernel resource controllers are
- Mounting cgroups-v2
- cgroups(7), sysfs(5) manual pages

19.3. CONFIGURING CPU LIMITS FOR THE CGROUP

You need to assign values to specific control files of the **cpu** controller to configure the CPU limits for your applications.

The user has two methods to set the CPU limits to a control group:

- Setting CPU bandwidth (editing the **cpu.max** controller file)
- Setting CPU weight (editing the **cpu.weight** controller file)

Prerequisites

- You have at least two applications whose CPU consumption you want to restrict.
- You have root permissions.
- You have mounted **cgroups-v2** filesystem.
- You have enabled **cpu** and **cpuset** controllers both in a root control group and in a child control group.
- You have created two levels of child control groups inside the root control group.

Procedure

1. Configure CPU limits for the control group:

echo "200000 1000000" > /sys/fs/cgroup/Example/tasks/cpu.max

The first value is the allowed time quota in microseconds for which all processes collectively in a child group can run during one period. The second value specifies the length of the period.

During a single period, when processes in a control group collectively exhaust the time specified by this quota, they are throttled for the remainder of the period and not allowed to run until the next period.

This command sets the CPU time limits so that all processes collectively in the /sys/fs/cgroup/Example/tasks child group can run on the CPU for only 0.2 seconds of every 1 second. That is, one fifth of each second.

2. Optionally, verify the limits:

cat /sys/fs/cgroup/Example/tasks/cpu.max 200000 1000000

3. Add the applications' PIDs to the **Example/tasks** child group:

```
# echo "34578" > /sys/fs/cgroup/Example/tasks/cgroup.procs # echo "34579" > /sys/fs/cgroup/Example/tasks/cgroup.procs
```

The example commands ensure that desired applications become members of the **Example/tasks** child group and do not exceed the CPU limits configured for this child group.

Verification steps

1. Verify that the applications run in the specified control group:

```
# cat /proc/34578/cgroup /proc/34579/cgroup 0::/Example/tasks 0::/Example/tasks
```

The output above shows that the processes of the specified applications that run in the **Example/tasks** child group.

2. Inspect the current CPU consumption of the throttled applications:

top top - 11:13:53 up 23:10, 1 user, load average: 0.26, 1.33, 1.66 Tasks: 104 total, 3 running, 101 sleeping, 0 stopped, 0 zombie %Cpu(s): 3.0 us, 7.0 sy, 0.0 ni, 89.5 id, 0.0 wa, 0.2 hi, 0.2 si, 0.2 st MiB Mem: 3737.4 total, 3312.6 free, 133.4 used, 291.4 buff/cache MiB Swap: 4060.0 total, 4060.0 free, 0.0 used. 3376.0 avail Mem PID USER PR NI VIRT RES SHR S %CPU %MEM TIME+ COMMAND 34578 root 20 0 18720 1756 1468 R **10.0** 0.0 37:36.13 sha1sum 20 0 18720 1772 1480 R **10.0** 0.0 37:41.22 sha1sum 1 root 20 0 186192 13940 9500 S 0.0 0.4 0:01.60 systemd 2 root 20 0 0 S 0.0 0.0 0:00.01 kthreadd 0

```
3 root 0 -20 0 0 0 I 0.0 0.0 0:00.00 rcu_gp
4 root 0 -20 0 0 0 I 0.0 0.0 0:00.00 rcu_par_gp
...
```

Notice that the CPU consumption for the **PID 34578** and **PID 34579** has decreased to 10%. The **Example/tasks** child group limits its processes to 20% of the CPU time collectively. Since there are 2 processes in the control group, each can utilize 10% of the CPU time.

Additional resources

- Understanding control groups
- What kernel resource controllers are
- Mounting cgroups-v2
- Setting up the cgroup for CPU limiting
- cgroups(7), sysfs(5) manual pages

CHAPTER 20. USING CONTROL GROUPS VERSION 1 WITH SYSTEMD

The following sections provide an overview of tasks related to creation, modification and removal of the control groups (**cgroups**). The utilities provided by the **systemd** system and service manager are the preferred way of the **cgroups** management and will be supported in the future.

20.1. ROLE OF SYSTEMD IN CONTROL GROUPS VERSION 1

Red Hat Enterprise Linux 8 moves the resource management settings from the process level to the application level by binding the system of **cgroup** hierarchies with the **systemd** unit tree. Therefore, you can manage the system resources with the **systemct** command, or by modifying the **systemd** unit files.

By default, the **systemd** system and service manager makes use of the **slice**, the **scope** and the **service** units to organize and structure processes in the control groups. The **systemctl** command enables you to further modify this structure by creating custom **slices**. Also, **systemd** automatically mounts hierarchies for important kernel resource controllers in the **/sys/fs/cgroup/** directory.

Three **systemd** unit types are used for resource control:

• **Service** - A process or a group of processes, which **systemd** started according to a unit configuration file. Services encapsulate the specified processes so that they can be started and stopped as one set. Services are named in the following way:

<name>.service

• Scope - A group of externally created processes. Scopes encapsulate processes that are started and stopped by the arbitrary processes through the **fork()** function and then registered by **systemd** at runtime. For example, user sessions, containers, and virtual machines are treated as scopes. Scopes are named as follows:

<name>.scope

• Slice - A group of hierarchically organized units. Slices organize a hierarchy in which scopes and services are placed. The actual processes are contained in scopes or in services. Every name of a slice unit corresponds to the path to a location in the hierarchy. The dash ("-") character acts as a separator of the path components to a slice from the -.slice root slice. In the following example:

<parent-name>.slice

parent-name.slice is a sub-slice of **parent.slice**, which is a sub-slice of the **-.slice** root slice. **parent-name.slice** can have its own sub-slice named **parent-name-name2.slice**, and so on.

The **service**, the **scope**, and the **slice** units directly map to objects in the control group hierarchy. When these units are activated, they map directly to control group paths built from the unit names.

The following is an abbreviated example of a control group hierarchy:

```
Control group /:
-.slice
|--user.slice
| |--user-42.slice
| | |--session-c1.scope
```

```
967 gdm-session-worker [pam/gdm-launch-environment]
    -1035 /usr/libexec/gdm-x-session gnome-session --autostart
/usr/share/gdm/greeter/autostart
-noreset -keeptty -verbose 3
       -1212 /usr/libexec/gnome-session-binary --autostart /usr/share/gdm/greeter/autostart
       -1369 /usr/bin/gnome-shell
       -1732 ibus-daemon --xim --panel disable
       -1752 /usr/libexec/ibus-dconf
       -1762 /usr/libexec/ibus-x11 --kill-daemon
       -1912 /usr/libexec/gsd-xsettings
       -1917 /usr/libexec/gsd-a11y-settings
       -1920 /usr/libexec/gsd-clipboard
  -init.scope
  1 /usr/lib/systemd/systemd --switched-root --system --deserialize 18
  -system.slice
   -rngd.service
   └─800 /sbin/rngd -f
   -systemd-udevd.service
   └─659 /usr/lib/systemd/systemd-udevd
   -chronyd.service
   823 /usr/sbin/chronyd
   -auditd.service
    -761 /sbin/auditd
    -763 /usr/sbin/sedispatch
   -accounts-daemon.service
   876 /usr/libexec/accounts-daemon
   -example.service
     - 929 /bin/bash /home/jdoe/example.sh
     -4902 sleep 1
```

The example above shows that services and scopes contain processes and are placed in slices that do not contain processes of their own.

Additional resources

- Configuring basic system settings in Red Hat Enterprise Linux
- What are kernel resource controllers
- systemd.resource-control(5), cgroups(7), fork(), fork(2) manual pages

20.2. CREATING TRANSIENT CONTROL GROUPS

The transient **cgroups** set limits on resources consumed by a unit (service or scope) during its runtime.

Procedure

• To create a transient control group, use the **systemd-run** command in the following format:

systemd-run --unit=<name> --slice=<name>.slice <command>

This command creates and starts a transient service or a scope unit and runs a custom command in such a unit.

- The **--unit=<name>** option gives a name to the unit. If **--unit** is not specified, the name is generated automatically.
- The --slice=<name>.slice option makes your service or scope unit a member of a specified slice. Replace <name>.slice with the name of an existing slice (as shown in the output of systemctl -t slice), or create a new slice by passing a unique name. By default, services and scopes are created as members of the system.slice.
- Replace *command*> with the command you wish to execute in the service or the scope unit
 - The following message is displayed to confirm that you created and started the service or the scope successfully:
 - # Running as unit < name>.service
- Optionally, keep the unit running after its processes finished to collect run-time information:
 - # systemd-run --unit=<name> --slice=<name>.slice --remain-after-exit <command>

The command creates and starts a transient service unit and runs a custom command in such a unit. The **--remain-after-exit** option ensures that the service keeps running after its processes have finished.

Additional resources

- Understanding control groups
- Role of systemd in control groups
- Configuring basic system settings Red Hat Enterprise Linux
- systemd-run(1) manual page

20.3. CREATING PERSISTENT CONTROL GROUPS

To assign a persistent control group to a service, it is necessary to edit its unit configuration file. The configuration is preserved after the system reboot, so it can be used to manage services that are started automatically.

Procedure

• To create a persistent control group, execute:

systemctl enable < name > . service

The command above automatically creates a unit configuration file into the /usr/lib/systemd/system/ directory and by default, it assigns <name>.service to the system.slice unit.

Additional resources

- Understanding control groups
- Role of systemd in control groups
- Configuring basic system settings Red Hat Enterprise Linux
- systemd-run(1) manual page

20.4. CONFIGURING MEMORY RESOURCE CONTROL SETTINGS ON THE COMMAND-LINE

Executing commands in the command-line interface is one of the ways how to set limits, prioritize, or control access to hardware resources for groups of processes.

Procedure

• To limit the memory usage of a service, run the following:

systemctl set-property example.service MemoryLimit=1500K

The command instantly assigns the memory limit of 1,500 kilobytes to processes executed in a control group the **example.service** service belongs to. The **MemoryLimit** parameter, in this configuration variant, is defined in the /etc/systemd/system.control/example.service.d/50-MemoryLimit.conf file and controls the value of the /sys/fs/cgroup/memory/system.slice/example.service/memory.limit_in_bytes file.

• Optionally, to temporarily limit the memory usage of a service, run:

systemctl set-property --runtime example.service MemoryLimit=1500K

The command instantly assigns the memory limit to the **example.service** service. The **MemoryLimit** parameter is defined until the next reboot in the /run/systemd/system.control/example.service.d/50-MemoryLimit.conf file. With a reboot, the whole /run/systemd/system.control/ directory and **MemoryLimit** are removed.



NOTE

The **50-MemoryLimit.conf** file stores the memory limit as a multiple of 4096 bytes - one kernel page size specific for AMD64 and Intel 64. The actual number of bytes depends on a CPU architecture.

Additional resources

- Understanding control groups
- What kernel resource controllers are
- systemd.resource-control(5) and cgroups(7) manual pages
- Role of systemd in control groups

20.5. CONFIGURING MEMORY RESOURCE CONTROL SETTINGS WITH UNIT FILES

Each persistent unit is supervised by the **systemd** system and service manager, and has a unit configuration file in the /**usr/lib/systemd/system/** directory. To change the resource control settings of the persistent units, modify its unit configuration file either manually in a text editor or from the command-line interface.

Manually modifying unit files is one of the ways how to set limits, prioritize, or control access to hardware resources for groups of processes.

Procedure

1. To limit the memory usage of a service, modify the /usr/lib/systemd/system/example.service file as follows:

... [Service] MemoryLimit=1500K ...

The configuration above places a limit on maximum memory consumption of processes executed in a control group, which **example.service** is part of.



NOTE

Use suffixes K, M, G, or T to identify Kilobyte, Megabyte, Gigabyte, or Terabyte as a unit of measurement.

2. Reload all unit configuration files:

systemctl daemon-reload

3. Restart the service:

systemctl restart example.service

- 4. Reboot the system.
- 5. Optionally, check that the changes took effect:

cat /sys/fs/cgroup/memory/system.slice/example.service/memory.limit_in_bytes 1536000

The example output shows that the memory consumption was limited at around 1,500 Kilobytes.



NOTE

The **memory.limit_in_bytes** file stores the memory limit as a multiple of 4096 bytes - one kernel page size specific for AMD64 and Intel 64. The actual number of bytes depends on a CPU architecture.

Additional resources

Understanding control groups

- What kernel resource controllers are
- systemd.resource-control(5), cgroups(7) manual pages
- Configuring basic system settings in Red Hat Enterprise Linux
- Role of systemd in control groups

20.6. REMOVING TRANSIENT CONTROL GROUPS

You can use the **systemd** system and service manager to remove transient control groups (**cgroups**) if you no longer need to limit, prioritize, or control access to hardware resources for groups of processes.

Transient **cgroups** are automatically released once all the processes that a service or a scope unit contains, finish.

Procedure

- To stop the service unit with all its processes, execute:
 - # systemctl stop <name>.service
- To terminate one or more of the unit processes, execute:
 - # systemctl kill < name > . service -- kill-who=PID, ... -- signal = signal

The command above uses the **--kill-who** option to select process(es) from the control group you wish to terminate. To kill multiple processes at the same time, pass a comma-separated list of PIDs. The **--signal** option determines the type of POSIX signal to be sent to the specified processes. The default signal is *SIGTERM*.

Additional resources

- Understanding control groups
- What kernel resource controllers are
- systemd.resource-control(5), cgroups(7) manual pages
- Role of systemd in control groups
- Configuring basic system settings Red Hat Enterprise Linux

20.7. REMOVING PERSISTENT CONTROL GROUPS

You can use the **systemd** system and service manager to remove persistent control groups (**cgroups**) if you no longer need to limit, prioritize, or control access to hardware resources for groups of processes.

Persistent **cgroups** are released when a service or a scope unit is stopped or disabled and its configuration file is deleted.

Procedure

1. Stop the service unit:

systemctl stop < name > . service

2. Disable the service unit:

systemctl disable < name > .service

3. Remove the relevant unit configuration file:

rm /usr/lib/systemd/system/<name>.service

4. Reload all unit configuration files so that changes take effect:

systemctl daemon-reload

Additional resources

- Understanding control groups
- What kernel resource controllers are
- systemd.resource-control(5), cgroups(7), and systemd.kill(5) manual pages
- Role of systemd in control groups
- Configuring basic system settings Red Hat Enterprise Linux

20.8. LISTING SYSTEMD UNITS

The following procedure describes how to use the **systemd** system and service manager to list its units.

Procedure

• To list all active units on the system, execute the **# systemctl** command and the terminal will return an output similar to the following example:

# systemctl UNIT	LOAD ACTIVE SUB DESCRIPTION
init.scope session-2.scope abrt-ccpp.service abrt-oops.service abrt-vmcore.service abrt-xorg.service	loaded active running System and Service Manager loaded active running Session 2 of user jdoe loaded active exited Install ABRT coredump hook loaded active running ABRT kernel log watcher loaded active exited Harvest vmcores for ABRT loaded active running ABRT Xorg log watcher
slice machine.slice Slice system-getty.slice	loaded active active Root Slice loaded active active Virtual Machine and Container loaded active active
system-getty.slice system-lvm2\x2dpvscan.slice lvm2\x2dpvscan.slice	loaded active active system-
system-sshd\x2dkeygen.slice sshd\x2dkeygen.slice	loaded active active system-

system-systemd\x2dhibernate\x2dresume.slice loaded active active system-systemd\x2dhibernate\x2dresume>
system-user\x2druntime\x2ddir.slice loaded active active system-user\x2druntime\x2ddir.slice
system.slice loaded active active System Slice
user-1000.slice loaded active active User Slice of UID 1000
user-42.slice loaded active active User Slice of UID 42
user.slice loaded active active User and Session Slice
...

- **UNIT** a name of a unit that also reflects the unit position in a control group hierarchy. The units relevant for resource control are a *slice*, a *scope*, and a *service*.
- **LOAD** indicates whether the unit configuration file was properly loaded. If the unit file failed to load, the field contains the state *error* instead of *loaded*. Other unit load states are: *stub*, *merged*, and *masked*.
- ACTIVE the high-level unit activation state, which is a generalization of SUB.
- **SUB** the low-level unit activation state. The range of possible values depends on the unit type.
- **DESCRIPTION** the description of the unit content and functionality.
- To list inactive units, execute:

systemctl --all

• To limit the amount of information in the output, execute:

systemctl --type service,masked

The **--type** option requires a comma-separated list of unit types such as a service and a slice, or unit load states such as loaded and masked.

Additional resources

• Configuring basic system settings Red Hat Enterprise Linux

20.9. VIEWING A CONTROL GROUP VERSION 1 HIERARCHY

The following procedure describes how to display control groups (**cgroups**) hierarchy and processes running in specific **cgroups**.

Procedure

• To display the whole **cgroups** hierarchy on your system, execute **# systemd-cgls**:

```
# systemd-cgls
Control group /:
-.slice
--user.slice
--user-42.slice
--user-42.slice
--user-42.slice
--user-42.slice
--user-42.slice
--user-42.slice
--user-42.slice
--user-42.slice
```

The example output returns the entire **cgroups** hierarchy, where the highest level is formed by *slices*.

To display the cgroups hierarchy filtered by a resource controller, execute # systemd-cgls
 <resource_controller>:

The example output of the above command lists the services that interact with the selected controller.

 To display detailed information about a certain unit and its part of the cgroups hierarchy, execute # systemctl status < system_unit

```
# systemctl status example.service

• example.service - My example service
Loaded: loaded (/usr/lib/systemd/system/example.service; enabled; vendor preset: disabled)
Active: active (running) since Tue 2019-04-16 12:12:39 CEST; 3s ago
Main PID: 17737 (bash)
Tasks: 2 (limit: 11522)
Memory: 496.0K (limit: 1.5M)
CGroup: /system.slice/example.service

—17737 /bin/bash /home/jdoe/example.sh
—17743 sleep 1
```

Apr 16 12:12:39 redhat systemd[1]: Started My example service.

Apr 16 12:12:39 redhat bash[17737]: The current time is Tue Apr 16 12:12:39 CEST 2019

Apr 16 12:12:40 redhat bash[17737]: The current time is Tue Apr 16 12:12:40 CEST 2019

Additional resources

- What kernel resource controllers are
- systemd.resource-control(5), cgroups(7) manual pages

20.10. VIEWING RESOURCE CONTROLLERS

The following procedure describes how to learn which processes use which resource controllers.

Procedure

 To view which resource controllers a process interacts with, execute the # cat proc/<PID>/cgroup command:

cat /proc/11269/cgroup

12:freezer:/

11:cpuset:/

10:devices:/system.slice

9:memory:/system.slice/example.service

8:pids:/system.slice/example.service

7:hugetlb:/

6:rdma:/

5:perf_event:/

4:cpu,cpuacct:/

3:net_cls,net_prio:/

2:blkio:/

1:name=systemd:/system.slice/example.service

The example output relates to a process of interest. In this case, it is a process identified by **PID 11269**, which belongs to the **example.service** unit. You can determine whether the process was placed in a correct control group as defined by the **systemd** unit file specifications.



NOTE

By default, the items and their ordering in the list of resource controllers is the same for all units started by **systemd**, since it automatically mounts all the default resource controllers.

Additional resources

- cgroups(7) manual page
- Documentation in the /usr/share/doc/kernel-doc-<kernel_version>/Documentation/cgroups-v1/ directory

20.11. MONITORING RESOURCE CONSUMPTION

The following procedure describes how to view a list of currently running control groups (**cgroups**) and their resource consumption in real-time.

Procedure

 To see a dynamic account of currently running cgroups, execute the # systemd-cgtop command:

```
# systemd-cgtop
Control Group
                          Tasks %CPU Memory Input/s Output/s
                      607 29.8 1.5G
                          125 - 428.7M
/system.slice
/system.slice/ModemManager.service
                                   3 -
                                            8.6M
/system.slice/NetworkManager.service
                                    3 - 12.8M
/system.slice/accounts-daemon.service
                                   3 - 1.8M
/system.slice/boot.mount
                                   - 48.0K
/system.slice/chronyd.service
                                        2.0M
/system.slice/cockpit.socket
                                     1.3M
/system.slice/colord.service
                                      3.5M
                               3 -
/system.slice/crond.service
                              1 - 1.8M
/system.slice/cups.service
                               1
                                      3.1M
/system.slice/dev-hugepages.mount
                                       - 244.0K
/system.slice/dev-mapper-rhel\x2dswap.swap -
                                         - 912.0K
/system.slice/dev-mqueue.mount
                                       - 48.0K
/system.slice/example.service
                                2 -
                                      2.0M
/system.slice/firewalld.service
                               2 -
                                      28.8M
```

The example output displays currently running **cgroups** ordered by their resource usage (CPU, memory, disk I/O load). The list refreshes every 1 second by default. Therefore, it offers a dynamic insight into the actual resource usage of each control group.

Additional resources

• systemd-cgtop(1) manual page

20.12. CONFIGURING CPUSET CONTROLLER USING SYSTEMD

The **systemd** resource management API allows the user to configure limits on a set of CPUs and NUMA nodes that a service can use. This limit restricts access to system resources utilized by the processes. The requested configuration is written in **cpuset.cpus** and **cpuset.mems**. However, the requested configuration may not be used, as the parent **cgroup** limits either `cpus` or **mems**. To access the current configuration, the **cpuset.cpus.effective** and **cpuset.mems.effective** files are exported to the users.

Procedure

To set AllowedCPUs:

systemctl set-property service_name.service AllowedCPUs=value

For example:

systemctl set-property service_name.service AllowedCPUs=0-5

• To set AllowedMemoryNodes:

systemctl set-property service_name.service AllowedMemoryNodes=value

For example:

systemctl set-property service_name.service AllowedMemoryNodes=0

CHAPTER 21. CONFIGURING RESOURCE MANAGEMENT USING CGROUPS VERSION 2 WITH SYSTEMD

The core of systemd is service management and supervision. systemd ensures that the right services start at the right time and in the correct order during the boot process. When the services are running, they have to run smoothly to use the underlying hardware platform optimally. Therefore, systemd also provides capabilities to define resource management policies and to tune various options, which can improve the performance of the service.

21.1. PREREQUISITES

• Basic knowledge of the Linux cgroup subsystem.

21.2. INTRODUCTION TO RESOURCE DISTRIBUTION MODELS

For resource management, systemd uses the cgroups v2 interface.

Note that RHEL 8 uses cgroups v1 by default. Therefore, you must enable cgroups v2 so that systemd can use the cgroups v2 interface for resource management. For more information on how to enable cgroups v2, see Setting CPU limits to applications using cgroups-v2 .

To modify the distribution of system resources, you can apply one or more of the following resource distribution models:

Weights

The resource is distributed by adding up the weights of all sub-groups and giving each sub-group the fraction matching its ratio against the sum.

For example, if you have 10 cgroups, each with Weight of value 100, the sum is 1000 and each cgroup receives one tenth of the resource.

Weight is usually used to distribute stateless resources. The *CPUWeight=* option is an implementation of this resource distribution model.

Limits

A cgroup can consume up to the configured amount of the resource, but you can also overcommit resources. Therefore, the sum of sub-group limits can exceed the limit of the parent cgroup. The *MemoryMax*= option is an implementation of this resource distribution model.

Protections

A protected amount of a resource can be set up for a cgroup. If the resource usage is below the protection boundary, the kernel will try not to penalize this cgroup in favor of other cgroups that compete for the same resource. An overcommit is also allowed.

The MemoryLow= option is an implementation of this resource distribution model.

Allocations

Exclusive allocations of an absolute amount of a finite resource. An overcommit is not allowed. An example of this resource type in Linux is the real-time budget.

Additional resources

Managing CPU with systemd

- Allocating memory resources using systemd
- Configuring I/O bandwidth with systemd

21.3. ALLOCATING CPU RESOURCES USING SYSTEMD

On a system managed by systemd, each system service is started in its cgroup. By enabling the support for the CPU cgroup controller, the system uses the service-aware distribution of CPU resources instead of the per-process distribution. In the service-aware distribution, each service receives approximately the same amount of CPU time relative to all other services running on the system, regardless of the number of processes that comprise the service.

If a specific service requires more CPU resources, you can grant them by changing the CPU time allocation policy for the service.

Procedure

To set a CPU time allocation policy option when using systemd:

- 1. Check the assigned values of the CPU time allocation policy option in the service of your choice:
 - \$ systemctl show --property <CPU time allocation policy option> <service name>
- 2. Set the required value of the CPU time allocation policy option as a root:
 - # systemctl set-property <service name> <CPU time allocation policy option>=<value>



NOTE

The cgroup properties are applied immediately after they are set. Therefore, the service does not need to be restarted.

The cgroup properties are applied immediately after they are set. Therefore, the service does not need to be restarted.

Verification steps

- To verify whether you successfully changed the required value of the CPU time allocation policy option for your service, run the following command:
 - \$ systemctl show --property <CPU time allocation policy option> <service name>

Additional resources

- CPU time allocation policy options for systemd
- Introduction to resource distribution models.

21.4. CPU TIME ALLOCATION POLICY OPTIONS FOR SYSTEMD

The most frequently used CPU time allocation policy options include:

CPUWeight=

Assigns **higher priority** to a particular service over all other services. You can select a value from the interval 1 – 10,000. The default value is 100.

For example, to give **httpd.service** twice as much CPU as to all other services, set the value to **CPUWeight=200**.

Note that **CPUWeight=** is applied only in cases when the operating system is overloaded.

CPUQuota=

Assigns the **absolute CPU time quota** to a service. The value of this option specifies the maximum percentage of CPU time that a service will receive relative to the total CPU time available, for example **CPUQuota=30%**.

Note that **CPUQuota=** represents the limit value for particular resource distribution models described in Introduction to resource distribution models.

For more information on CPUQuota=, see the systemd.resource-control(5) man page.

Additional resources

- Allocating CPU resources using systemd
- Introduction to resource distribution models

21.5. ALLOCATING MEMORY RESOURCES USING SYSTEMD

This section describes how to use any of the memory configuration options (**MemoryMin**, **MemoryLow**, **MemoryHigh**, **MemoryMax**, **MemorySwapMax**) to allocate memory resources using systemd.

Procedure

To set a memory allocation configuration option when using systemd:

- 1. Check the assigned values of the memory allocation configuration option in the service of your choice:
 - \$ systemctl show --property <memory allocation configuration option> <service name>
- 2. Set the required value of the memory allocation configuration option as a root:
 - # systemctl set-property <service name> <memory allocation configuration option>=<value>



NOTE

The cgroup properties are applied immediately after they are set. Therefore, the service does not need to be restarted.

Verification steps

- To verify whether you successfully changed the required value of the memory allocation configuration option for your service, run the following command:
 - \$ systemctl show --property <memory allocation configuration option> <service name>

۸ ما ما:۱: مسما سمم

Additional resources

- Memory allocation configuration options for systemd
- Introduction to resource distribution models

21.6. MEMORY ALLOCATION CONFIGURATION OPTIONS FOR SYSTEMD

You can use the following options when using systemd to configure system memory allocation

MemoryMin

Hard memory protection. If the memory usage is below the limit, the cgroup memory will not be reclaimed.

MemoryLow

Soft memory protection. If the memory usage is below the limit, the cgroup memory can be reclaimed only if no memory is reclaimed from unprotected cgroups.

MemoryHigh

Memory throttle limit. If the memory usage goes above the limit, the processes in the cgroup are throttled and put under a heavy reclaim pressure.

MemoryMax

Absolute limit for the memory usage. You can use the kilo (K), mega (M), giga (G), tera (T) suffixes, for example MemoryMax=1G.

MemorySwapMax

Hard limit on the swap usage.



NOTE

When you exhaust your memory limit, the Out-of-memory (OOM) killer will stop the running service. To prevent this, lower the **OOMScoreAdjust=** value to increase the memory tolerance.

- Allocating memory resources using systemd
- Introduction to resource distribution models

21.7. CONFIGURING I/O BANDWIDTH USING SYSTEMD

To improve the performance of a specific service in RHEL 8, you can allocate I/O bandwidth resources to that service using systemd.

To do so, you can use the following I/O configuration options:

- IOWeight
- IODeviceWeight
- IOReadBandwidthMax
- IOWriteBandwidthMax
- IOReadIOPSMax

IOWriteIOPSMax

Procedure

To set a **I/O bandwidth configuration** option using systemd:

- Check the assigned values of the I/O bandwidth configuration option in the service of your choice:
 - \$ systemctl show --property <I/O bandwidth configuration option> <service name>
- 2. Set the required value of the I/O bandwidth configuration option as a root:
 - # systemctl set-property <service name> <I/O bandwidth configuration option>=<value>

The cgroup properties are applied immediately after they are set. Therefore, the service does not need to be restarted.

Verification steps

- To verify whether you successfully changed the required value of the I/O bandwidth configuration option for your service, run the following command:
 - \$ systemctl show --property <I/O bandwidth configuration option> <service name>

Additional resources

- I/O bandwidth configuration options for systemd
- Introduction to resource distribution models

21.8. I/O BANDWIDTH CONFIGURATION OPTIONS FOR SYSTEMD

To manage the block layer I/O policies with systemd, the following configuration options are available:

IOWeight

Sets the default I/O weight. The weight value is used as a basis for the calculation of how much of the real I/O bandwidth the service receives in relation to the other services.

IODeviceWeight

Sets the I/O weight for a specific block device.

For example, IODeviceWeight=/dev/disk/by-id/dm-name-rhel-root 200.

IOReadBandwidthMax, IOWriteBandwidthMax

Sets the absolute bandwidth per device or a mount point.

For example, IOWriteBandwith=/var/log 5M.



NOTE

Systemd handles the file-system-to-device translation automatically.

IOReadIOPSMax, IOWriteIOPSMax

A similar option to the previous one: sets the absolute bandwidth in Input/Output Operations Per Second (IOPS).



NOTE

Weight-based options are supported only if the block device is using the CFQ I/O scheduler. No option is supported if the device uses the Multi-Queue Block I/O queuing mechanism.

- Configuring I/O bandwidth with systemd
- Introduction to resource distribution models

CHAPTER 22. CONFIGURING CPU AFFINITY AND NUMA POLICIES USING SYSTEMD

The CPU management, memory management, and I/O bandwidth options deal with partitioning available resources.

22.1. CONFIGURING CPU AFFINITY USING SYSTEMD

CPU affinity settings help you restrict the access of a particular process to some CPUs. Effectively, the CPU scheduler never schedules the process to run on the CPU that is not in the affinity mask of the process.

The default CPU affinity mask applies to all services managed by systemd.

To configure CPU affinity mask for a particular systemd service, systemd provides **CPUAffinity=** both as a unit file option and a manager configuration option in the /etc/systemd/system.conf file.

The **CPUAffinity=** unit file option sets a list of CPUs or CPU ranges that are merged and used as the affinity mask. The **CPUAffinity** option in the /etc/systemd/system.conf file defines an affinity mask for the process identification number (PID) 1 and all processes forked off of PID1. You can then override the **CPUAffinity** on a per-service basis.



NOTE

After configuring CPU affinity mask for a particular systemd service, you must restart the system to apply the changes.

Procedure

To set CPU affinity mask for a particual systemd service using the CPUAffinity unit file option:

- 1. Check the values of the **CPUAffinity** unit file option in the service of your choice:
 - \$ systemctl show --property <CPU affinity configuration option> <service name>
- 2. As a root, set the required value of the **CPUAffinity** unit file option for the CPU ranges used as the affinity mask:
 - # systemctl set-property <service name> CPUAffinity=<value>
- 3. Restart the service to apply the changes.
 - # systemctl restart <service name>

To set CPU affinity mask for a particular systemd service using the manager configuration option:

- 1. Edit the /etc/systemd/system.conf file:
 - # vi /etc/systemd/system.conf
- 2. Search for the **CPUAffinity=** option and set the CPU numbers
- 3. Save the edited file and restart the server to apply the changes.

22.2. CONFIGURING NUMA USING SYSTEMD

Non-uniform memory access (NUMA) is a computer memory subsystem design, in which the memory access time depends on the memory location relative to the processor. Memory close to the CPU has lower latency (local memory) than memory that is local for a different CPU or is shared between a set of CPUs.

In terms of the Linux kernel, NUMA policy governs where (for example, on which NUMA nodes) the kernel allocates physical memory pages for the process.

To configure NUMA, systemd provides the unit file option for **NUMAPolicy** and **NUMAMask**, and the manager configuration option in the /etc/systemd/system.conf file.

Procedure

To set the NUMA memory policy through the **NUMAPolicy unit file** option:

- 1. Check the values of the **NUMAPolicy** unit file option in the service of your choice:
 - \$ systemctl show --property <NUMA policy configuration option> <service name>
- 2. As a root, set the required policy type of the **NUMAPolicy** unit file option:
 - # systemctl set-property <service name> NUMAPolicy=<value>
- 3. Restart the service to apply the changes.
 - # systemctl restart <service name>

To set the **NUMAPolicy** through the **manager configuration** option:

- Edit the /etc/systemd/system.conf file:
 - # vi /etc/systemd/system.conf
- 2. Search for the **NUMAPolicy** option and set the policy type.
- 3. Save the edited file and restart the server to apply the changes.

Additional resources

NUMA policy configuration options for systemd

22.3. NUMA POLICY CONFIGURATION OPTIONS FOR SYSTEMD

systemd provides the following options to configure the NUMA policy:

NUMAPolicy

Controls the NUMA memory policy of the executed processes. The following policy types are possible:

- default
- preferred

- bind
- interleave
- local

NUMAMask

Controls the NUMA node list which is associated with the selected NUMA policy. Note that the **NUMAMask** option is not required to be specified for the following policies:

- default
- local

For the preferred policy, the list specifies only a single NUMA node.

Additional resources

- systemd.resource-control(5), systemd.exec(5), and set_mempolicy(2) manual pages
- Configuring NUMA using systemd

CHAPTER 23. ANALYZING SYSTEM PERFORMANCE WITH BPF COMPILER COLLECTION

As a system administrator, you can use the BPF Compiler Collection (BCC) library to create tools for analyzing the performance of your Linux operating system and gathering information, which could be difficult to obtain through other interfaces.

23.1. AN INTRODUCTION TO BCC

BPF Compiler Collection (BCC) is a library, which facilitates the creation of the extended Berkeley Packet Filter (eBPF) programs. The main utility of eBPF programs is analyzing OS performance and network performance without experiencing overhead or security issues.

BCC removes the need for users to know deep technical details of eBPF, and provides many out-of-the-box starting points, such as the **bcc-tools** package with pre-created eBPF programs.



NOTE

The eBPF programs are triggered on events, such as disk I/O, TCP connections, and process creations. It is unlikely that the programs should cause the kernel to crash, loop or become unresponsive because they run in a safe virtual machine in the kernel.

23.2. INSTALLING THE BCC-TOOLS PACKAGE

This section describes how to install the **bcc-tools** package, which also installs the BPF Compiler Collection (BCC) library as a dependency.

Prerequisites

- An active Red Hat Enterprise Linux subscription
- An enabled repository containing the **bcc-tools** package
- Updated kernel
- Root permissions

Procedure

1. Install bcc-tools:

yum install bcc-tools

The BCC tools are installed in the /usr/share/bcc/tools/ directory.

2. Optionally, inspect the tools:

II /usr/share/bcc/tools/

...

- -rwxr-xr-x. 1 root root 4198 Dec 14 17:53 dcsnoop
- -rwxr-xr-x. 1 root root 3931 Dec 14 17:53 dcstat
- -rwxr-xr-x. 1 root root 20040 Dec 14 17:53 deadlock_detector
- -rw-r--r-. 1 root root 7105 Dec 14 17:53 deadlock_detector.c

```
drwxr-xr-x. 3 root root 8192 Mar 11 10:28 doc
-rwxr-xr-x. 1 root root 7588 Dec 14 17:53 execsnoop
-rwxr-xr-x. 1 root root 6373 Dec 14 17:53 ext4dist
-rwxr-xr-x. 1 root root 10401 Dec 14 17:53 ext4slower
...
```

The **doc** directory in the listing above contains documentation for each tool.

23.3. USING SELECTED BCC-TOOLS FOR PERFORMANCE ANALYSES

This section describes how to use certain pre-created programs from the BPF Compiler Collection (BCC) library to efficiently and securely analyze the system performance on the per-event basis. The set of pre-created programs in the BCC library can serve as examples for creation of additional programs.

Prerequisites

- Installed BCC library
- Root permissions

Using execsnoop to examine the system processes

- 1. Execute the **execsnoop** program in one terminal:
 - # /usr/share/bcc/tools/execsnoop
- 2. In another terminal execute for example:
 - \$ ls /usr/share/bcc/tools/doc/

The above creates a short-lived process of the **Is** command.

3. The terminal running **execsnoop** shows the output similar to the following:

```
PCOMM PID PPID RET ARGS
Is 8382 8287 0 /usr/bin/ls --color=auto /usr/share/bcc/tools/doc/sed 8385 8383 0 /usr/bin/sed s/^ *[0-9]\+ *//
...
```

The **execsnoop** program prints a line of output for each new process, which consumes system resources. It even detects processes of programs that run very shortly, such as **Is**, and most monitoring tools would not register them.

The **execsnoop** output displays the following fields:

- **PCOMM** The parent process name. (**Is**)
- **PID** The process ID. (8382)
- PPID The parent process ID. (8287)
- RET The return value of the exec() system call (0), which loads program code into new processes.

• ARGS - The location of the started program with arguments.

To see more details, examples, and options for **execsnoop**, refer to the /usr/share/bcc/tools/doc/execsnoop_example.txt file.

For more information about **exec()**, see **exec(3)** manual pages.

Using opensnoop to track what files a command opens

1. Execute the **opensnoop** program in one terminal:

/usr/share/bcc/tools/opensnoop -n uname

The above prints output for files, which are opened only by the process of the **uname** command.

2. In another terminal execute:

\$ uname

The command above opens certain files, which are captured in the next step.

3. The terminal running **opensnoop** shows the output similar to the following:

```
PID COMM FD ERR PATH
8596 uname 3 0 /etc/ld.so.cache
8596 uname 3 0 /lib64/libc.so.6
8596 uname 3 0 /usr/lib/locale/locale-archive
...
```

The **opensnoop** program watches the **open()** system call across the whole system, and prints a line of output for each file that **uname** tried to open along the way.

The **opensnoop** output displays the following fields:

- **PID** The process ID. (8596)
- COMM The process name. (uname)
- FD The file descriptor a value that **open()** returns to refer to the open file. (3)
- ERR Any errors.
- PATH The location of files that open() tried to open.
 If a command tries to read a non-existent file, then the FD column returns -1 and the ERR column prints a value corresponding to the relevant error. As a result, opensnoop can help you identify an application that does not behave properly.

To see more details, examples, and options for **opensnoop**, refer to the /usr/share/bcc/tools/doc/opensnoop_example.txt file.

For more information about **open()**, see **open(2)** manual pages.

Using biotop to examine the I/O operations on the disk

1. Execute the **biotop** program in one terminal:

/usr/share/bcc/tools/biotop 30

The command enables you to monitor the top processes, which perform I/O operations on the disk. The argument ensures that the command will produce a 30 second summary.



NOTE

When no argument provided, the output screen by default refreshes every 1 second.

2. In another terminal execute for example:

dd if=/dev/vda of=/dev/zero

The command above reads the content from the local hard disk device and writes the output to the /dev/zero file. This step generates certain I/O traffic to illustrate biotop.

3. The terminal running **biotop** shows the output similar to the following:

```
PID
    COMM
                                   I/O Kbytes
                 D MAJ MIN DISK
                                              AVGms
9568 dd
               R 252 0 vda
                             16294 14440636.0 3.69
                W 252 0 vda
                               1763 120696.0 1.65
    kswapd0
7571 gnome-shell
                R 252 0 vda
                                 834 83612.0
                                              0.33
1891 gnome-shell R 252 0 vda
                                1379 19792.0 0.15
7515 Xorg
               R 252 0 vda
                              280 9940.0
                                         0.28
7579 Ilvmpipe-1
                 R 252 0 vda
                              228 6928.0 0.19
9515 gnome-control-c R 252 0 vda
                                   62 6444.0 0.43
8112 gnome-terminal- R 252 0 vda
                                   67 2572.0
                                              1.54
7807 gnome-software R 252 0 vda
                                   31 2336.0 0.73
9578 awk
               R 252 0 vda
                               17 2228.0
                                          0.66
7578 Ilvmpipe-0
                R 252 0 vda
                               156 2204.0
                                            0.07
9581 pgrep
                R 252 0 vda
                               58 1748.0
                                          0.42
7531 InputThread R 252 0 vda
                                 30 1200.0 0.48
                R 252 0 vda
7504 gdbus
                                3 1164.0 0.30
1983 Ilvmpipe-1
                 R 252 0 vda
                                39 724.0 0.08
                 R 252 0 vda
                                36 652.0 0.06
1982 Ilvmpipe-0
```

The **biotop** output displays the following fields:

- **PID** The process ID. (9568)
- **COMM** The process name. (**dd**)
- **DISK** The disk performing the read operations. (**vda**)
- I/O The number of read operations performed. (16294)
- **Kbytes** The amount of Kbytes reached by the read operations. (14,440,636)
- AVGms The average I/O time of read operations. (3.69)

To see more details, examples, and options for **biotop**, refer to the /usr/share/bcc/tools/doc/biotop_example.txt file.

For more information about **dd**, see **dd(1)** manual pages.

Using xfsslower to expose unexpectedly slow file system operations

1. Execute the **xfsslower** program in one terminal:

/usr/share/bcc/tools/xfsslower 1

The command above measures the time the XFS file system spends in performing read, write, open or sync (**fsync**) operations. The **1** argument ensures that the program shows only the operations that are slower than 1 ms.



NOTE

When no arguments provided, **xfsslower** by default displays operations slower than 10 ms.

2. In another terminal execute, for example, the following:

\$ vim text

The command above creates a text file in the **vim** editor to initiate certain interaction with the XFS file system.

3. The terminal running **xfsslower** shows something similar upon saving the file from the previous step:

```
TIME
       COMM
                    PID TBYTES OFF_KB LAT(ms) FILENAME
                                        7.11 b'vim'
13:07:14 b'bash'
                   4754 R 256
13:07:14 b'vim'
                  4754 R 832
                                       4.03 b'libgpm.so.2.1.0'
                                0
13:07:14 b'vim'
                  4754 R 32
                                       1.04 b'libgpm.so.2.1.0'
                                20
13:07:14 b'vim'
                                        2.30 b'vimrc'
                  4754 R 1982 0
                                        2.52 b'getscriptPlugin.vim'
13:07:14 b'vim'
                  4754 R 1393 0
13:07:45 b'vim'
                  4754 S 0
                                      6.71 b'text'
                               0
                                       5.58 b'text'
13:07:45 b'pool'
                  2588 R 16
                                0
```

Each line above represents an operation in the file system, which took more time than a certain threshold. **xfsslower** is good at exposing possible file system problems, which can take form of unexpectedly slow operations.

The **xfsslower** output displays the following fields:

- **COMM** The process name. (**b'bash'**)
- **T** The operation type. (**R**)
 - Read
 - Write
 - Sync
- OFF_KB The file offset in KB. (0)

• **FILENAME** - The file being read, written, or synced.

To see more details, examples, and options for **xfsslower**, refer to the /usr/share/bcc/tools/doc/xfsslower_example.txt file.

For more information about **fsync**, see **fsync(2)** manual pages.

CHAPTER 24. ENHANCING SECURITY WITH THE KERNEL INTEGRITY SUBSYSTEM

You can increase the protection of your system by utilizing components of the kernel integrity subsystem. The following sections introduce the relevant components and provide guidance on their configuration.

24.1. THE KERNEL INTEGRITY SUBSYSTEM

The integrity subsystem is a part of the kernel which is responsible for maintaining the overall system's data integrity. This subsystem helps to keep the state of a certain system the same from the time it was built thereby it prevents undesired modification on specific system files from users.

The kernel integrity subsystem consists of two major components:

Integrity Measurement Architecture (IMA)

- Measures files' content whenever it is executed or opened. Users can change this behavior by applying custom policies.
- Places the measured values within the kernel's memory space thereby it prevents any modification from the users of the system.
- Allows local and remote parties to verify the measured values.

Extended Verification Module (EVM)

• Protects files' extended attributes (also known as *xattr*) that are related to the system's security, like IMA measurements and SELinux attributes, by cryptographically hashing their corresponding values.

Both IMA and EVM also contain numerous feature extensions that bring additional functionality. For example:

IMA-Appraisal

• Provides local validation of the current file's content against the values previously stored in the measurement file within the kernel memory. This extension forbids any operation to be performed over a specific file in case the current and the previous measure do not match.

EVM Digital Signatures

 Allows digital signatures to be used through cryptographic keys stored into the kernel's keyring.



NOTE

The feature extensions complement each other, but you can configure and use them independently of one another.

The kernel integrity subsystem can harness the Trusted Platform Module (TPM) to harden the system security even more. TPM is a specification by the Trusted Computing Group (TCG) for important cryptographic functions. TPMs are usually built as dedicated hardware that is attached to the platform's

motherboard and prevents software-based attacks by providing cryptographic functions from a protected and tamper-proof area of the hardware chip. Some of the TPM features are:

- Random-number generator
- Generator and secure storage for cryptographic keys
- Hashing generator
- Remote attestation

Additional resources

- Integrity Measurement Architecture (IMA)
- Trusted Computing Group resources

24.2. INTEGRITY MEASUREMENT ARCHITECTURE

Integrity Measurement Architecture (IMA) is a component of the kernel integrity subsystem. IMA aims to maintain the contents of local files. Specifically, IMA measures, stores, and appraises files' hashes before they are accessed, which prevents the reading and execution of unreliable data. Thereby, IMA enhances the security of the system.

24.3. EXTENDED VERIFICATION MODULE

Extended Verification Module (EVM) is a component of the kernel integrity subsystem, which monitors changes in files' extended attributes (xattr). Many security-oriented technologies, including Integrity Measurement Architecture (IMA), store sensitive file information, such as content hashes, in the extended attributes. EVM creates another hash from these extended attributes and from a special key, which is loaded at boot time. The resulting hash is validated every time the extended attribute is used. For example, when IMA appraises the file.

RHEL 8 accepts the special encrypted key under the **evm-key** keyring. The key was created by a *master key* held in the kernel keyrings.

24.4. TRUSTED AND ENCRYPTED KEYS

The following section introduces trusted and encrypted keys as an important part of enhancing system security.

Trusted and encrypted keys are variable-length symmetric keys generated by the kernel that utilize the kernel keyring service. The fact that this type of keys never appear in the user space in an unencrypted form means that their integrity can be verified, which in turn means that they can be used, for example, by the extended verification module (EVM) to verify and confirm the integrity of a running system. User-level programs can only access the keys in the form of encrypted blobs.

Trusted keys need a hardware component: the Trusted Platform Module (TPM) chip, which is used to both create and encrypt (seal) the keys. The TPM seals the keys using a 2048-bit RSA key called the *storage root key* (SRK).



NOTE

To use a TPM 1.2 specification, enable and activate it through a setting in the machine firmware or by using the **tpm_setactive** command from the **tpm-tools** package of utilities. Also, the **TrouSers** software stack needs to be installed and the **tcsd** daemon needs to be running to communicate with the TPM (dedicated hardware). The **tcsd** daemon is part of the **TrouSers** suite, which is available through the **trousers** package. The more recent and backward incompatible TPM 2.0 uses a different software stack, where the **tpm2-tools** or **ibm-tss** utilities provide access to the dedicated hardware.

In addition to that, the user can seal the trusted keys with a specific set of the TPM's platform configuration register (PCR) values. PCR contains a set of integrity-management values that reflect the firmware, boot loader, and operating system. This means that PCR-sealed keys can only be decrypted by the TPM on the same system on which they were encrypted. However, once a PCR-sealed trusted key is loaded (added to a keyring), and thus its associated PCR values are verified, it can be updated with new (or future) PCR values, so that a new kernel, for example, can be booted. A single key can also be saved as multiple blobs, each with different PCR values.

Encrypted keys do not require a TPM, as they use the kernel Advanced Encryption Standard (AES), which makes them faster than trusted keys. Encrypted keys are created using kernel-generated random numbers and encrypted by a *master key* when they are exported into user-space blobs. The master key is either a trusted key or a user key. If the master key is not trusted, the encrypted key is only as secure as the user key used to encrypt it.

24.5. WORKING WITH TRUSTED KEYS

The following section describes how to create, export, load or update trusted keys with the **keyctl** utility to improve the system security.

Prerequisites

- For the 64-bit ARM architecture and IBM Z, the **trusted** kernel module needs to be loaded. For more information on how to load kernel modules, see Managing kernel modules.
- Trusted Platform Module (TPM) needs to be enabled and active. For more information about TPM see, The kernel integrity subsystem and Trusted and encrypted keys.

Procedure

1. To create a trusted key using a TPM, execute:

keyctl add trusted <name> "new <key_length> [options]" <key_ring>

• Based on the syntax, construct an example command as follows:

keyctl add trusted kmk "new 32" @u 642500861

The command creates a trusted key called **kmk** with the length of 32 bytes (256 bits) and places it in the user keyring (**@u**). The keys may have a length of 32 to 128 bytes (256 to 1024 bits).

2. To list the current structure of the kernel keyrings:

```
# keyctl show
Session Keyring
-3 --alswrv 500 500 keyring: _ses
97833714 --alswrv 500 -1 \_ keyring: _uid.1000
642500861 --alswrv 500 500 \ trusted: kmk
```

3. To export the key to a user-space blob, execute:

```
# keyctl pipe 642500861 > kmk.blob
```

The command uses the **pipe** subcommand and the serial number of **kmk**.

4. To load the trusted key from the user-space blob, use the **add** subcommand with the blob as an argument:

```
# keyctl add trusted kmk "load `cat kmk.blob`" @u 268728824
```

5. Create secure encrypted keys based on the TPM-sealed trusted key:

```
# keyctl add encrypted <pass:quotes[name]> "new [format] <pass:quotes[key_type]>:
<pass:quotes[primary_key_name]> <pass:quotes[keylength]>"
<pass:quotes[key_ring]>
```

• Based on the syntax, generate an encrypted key using the already created trusted key:

```
\# keyctl add encrypted encr-key "new trusted:kmk 32" @u 159771175
```

The command uses the TPM-sealed trusted key (**kmk**), produced in the previous step, as a *primary key* for generating encrypted keys.

Additional resources

- **keyctl(1)** manual page
- Trusted and encrypted keys
- Kernel Key Retention Service
- The kernel integrity subsystem

24.6. WORKING WITH ENCRYPTED KEYS

The following section describes managing encrypted keys to improve the system security on systems where a Trusted Platform Module (TPM) is not available.

Prerequisites

• For the 64-bit ARM architecture and IBM Z, the **encrypted-keys** kernel module needs to be loaded. For more information on how to load kernel modules, see Managing kernel modules.

Procedure

1. Use a random sequence of numbers to generate a user key:

keyctl add user kmk-user "\$(dd if=/dev/urandom bs=1 count=32 2>/dev/null)" @u 427069434

The command generates a user key called **kmk-user** which acts as a *primary key* and is used to seal the actual encrypted keys.

2. Generate an encrypted key using the primary key from the previous step:

keyctl add encrypted encr-key "new user:kmk-user 32" @u 1012412758

3. Optionally, list all keys in the specified user keyring:

keyctl list @u

2 keys in keyring:

427069434: --alswrv 1000 1000 user: kmk-user 1012412758: --alswrv 1000 1000 encrypted: encr-key



IMPORTANT

Keep in mind that encrypted keys that are not sealed by a trusted primary key are only as secure as the user primary key (random-number key) that was used to encrypt them. Therefore, the primary user key should be loaded as securely as possible and preferably early during the boot process.

Additional resources

- **keyctl(1)** manual page
- Kernel Key Retention Service

24.7. ENABLING INTEGRITY MEASUREMENT ARCHITECTURE AND EXTENDED VERIFICATION MODULE

Integrity measurement architecture (IMA) and extended verification module (EVM) belong to the kernel integrity subsystem and enhance the system security in various ways. The following section describes how to enable and configure IMA and EVM to improve the security of the operating system.

Prerequisites

Verify that the securityfs filesystem is mounted on the /sys/kernel/security/ directory and the /sys/kernel/security/integrity/ima/ directory exists.

mount

... securityfs on /sys/kernel/security type securityfs (rw,nosuid,nodev,noexec,relatime) ...

 Verify that the **systemd** service manager is already patched to support IMA and EVM on boot time:

```
# dmesg | grep -i -e EVM -e IMA
[ 0.000000] Command line: BOOT IMAGE=(hd0,msdos1)/vmlinuz-4.18.0-167.el8.x86 64
root=/dev/mapper/rhel-root ro crashkernel=auto resume=/dev/mapper/rhel-swap
rd.lvm.lv=rhel/root rd.lvm.lv=rhel/swap rhgb quiet
[ 0.000000] kvm-clock: cpu 0, msr 23601001, primary cpu clock
[ 0.000000] Using crashkernel=auto, the size chosen is a best effort estimation.
[ 0.000000] Kernel command line: BOOT_IMAGE=(hd0,msdos1)/vmlinuz-4.18.0-
167.el8.x86 64 root=/dev/mapper/rhel-root ro crashkernel=auto resume=/dev/mapper/rhel-
swap rd.lvm.lv=rhel/root rd.lvm.lv=rhel/swap rhgb quiet
[ 0.911527] ima: No TPM chip found, activating TPM-bypass!
[ 0.911538] ima: Allocated hash algorithm: sha1
[ 0.911580] evm: Initialising EVM extended attributes:
[ 0.911581] evm: security.selinux
[ 0.911581] evm: security.ima
[ 0.911582] evm: security.capability
[ 0.911582] evm: HMAC attrs: 0x1
[ 1.715151] systemd[1]: systemd 239 running in system mode. (+PAM +AUDIT +SELINUX
+IMA -APPARMOR +SMACK +SYSVINIT +UTMP +LIBCRYPTSETUP +GCRYPT
+GNUTLS +ACL +XZ +LZ4 +SECCOMP +BLKID +ELFUTILS +KMOD +IDN2 -IDN +PCRE2
default-hierarchy=legacy)
[ 3.824198] fbcon: qxldrmfb (fb0) is primary device
[ 4.673457] PM: Image not found (code -22)
[ 6.549966] systemd[1]: systemd 239 running in system mode. (+PAM +AUDIT +SELINUX
+IMA -APPARMOR +SMACK +SYSVINIT +UTMP +LIBCRYPTSETUP +GCRYPT
+GNUTLS +ACL +XZ +LZ4 +SECCOMP +BLKID +ELFUTILS +KMOD +IDN2 -IDN +PCRE2
default-hierarchy=legacy)
```

Procedure

1. Add the following kernel command line parameters:

grubby --update-kernel=/boot/vmlinuz-\$(uname -r) --args="ima_policy=appraise_tcb ima appraise=fix evm=fix"

The command enables IMA and EVM in the *fix* mode for the current boot entry and allows users to gather and update the IMA measurements.

The **ima_policy=appraise_tcb** kernel command line parameter ensures that the kernel uses the default Trusted Computing Base (TCB) measurement policy and the appraisal step. The appraisal part forbids access to files, whose prior and current measures do not match.

- 2. Reboot to make the changes come into effect.
- 3. Optionally, verify that the parameters have been added to the kernel command line:

cat /proc/cmdline

BOOT_IMAGE=(hd0,msdos1)/vmlinuz-4.18.0-167.el8.x86_64 root=/dev/mapper/rhel-root ro crashkernel=auto resume=/dev/mapper/rhel-swap rd.lvm.lv=rhel/root rd.lvm.lv=rhel/swap rhgb quiet ima_policy=appraise_tcb ima_appraise=fix evm=fix

4. Create a kernel master key to protect the EVM key:

keyctl add user kmk "\$(dd if=/dev/urandom bs=1 count=32 2> /dev/null)" @u 748544121

The kernel master key (\mathbf{kmk}) is kept entirely in the kernel space memory. The 32-byte long value of the kernel master key \mathbf{kmk} is generated from random bytes from the $/\mathbf{dev}/\mathbf{urandom}$ file and placed in the user $(@\mathbf{u})$ keyring. The key serial number is on the second line of the previous output.

5. Create an encrypted EVM key based on the **kmk** key:

keyctl add encrypted evm-key "new user:kmk 64" @u 641780271

The command uses **kmk** to generate and encrypt a 64-byte long user key (named **evm-key**) and places it in the user (@u) keyring. The key serial number is on the second line of the previous output.



IMPORTANT

It is necessary to name the user key as **evm-key** because that is the name the EVM subsystem is expecting and is working with.

6. Create a directory for exported keys:

mkdir -p /etc/keys/

7. Search for the **kmk** key and export its value into a file:

keyctl pipe keyctl search @u user kmk > /etc/keys/kmk*

The command places the unencrypted value of the kernel master key (**kmk**) into a file of previously defined location (/**etc/keys**/).

8. Search for the **evm-key** user key and export its value into a file:

keyctl pipe keyctl search @u encrypted evm-key > /etc/keys/evm-key

The command places the encrypted value of the user **evm-key** key into a file of arbitrary location. The **evm-key** has been encrypted by the kernel master key earlier.

9. Optionally, view the newly created keys:

```
# keyctl show
Session Keyring
974575405 --alswrv 0 0 keyring: _ses
299489774 --alswrv 0 65534 \_ keyring: _uid.0
748544121 --alswrv 0 0 \_ user: kmk
641780271 --alswrv 0 0 \_ encrypted: evm-key
```

You should be able to see a similar output.

10. Activate EVM:

echo 1 > /sys/kernel/security/evm

11. Optionally, verify that EVM has been initialized:

dmesg | tail -1

[...] evm: key initialized

Additional resources

- The kernel integrity subsystem
- Integrity measurement architecture
- Extended verification module
- Trusted and encrypted keys

24.8. COLLECTING FILE HASHES WITH INTEGRITY MEASUREMENT ARCHITECTURE

The first level of operation of integrity measurement architecture (IMA) is the *measurement* phase, which allows to create file hashes and store them as extended attributes (xattrs) of those files. The following section describes how to create and inspect the files' hashes.

Prerequisites

- Enable integrity measurement architecture (IMA) and extended verification module (EVM) as described in Enabling integrity measurement architecture and extended verification module.
- Verify that the **ima-evm-utils**, **attr**, and **keyutils** packages are already installed:

yum install ima-evm-utils attr keyutils

Updating Subscription Management repositories.

This system is registered to Red Hat Subscription Management, but is not receiving updates.

You can use subscription-manager to assign subscriptions.

Last metadata expiration check: 0:58:22 ago on Fri 14 Feb 2020 09:58:23 AM CET.

Package ima-evm-utils-1.1-5.el8.x86 64 is already installed.

Package attr-2.4.48-3.el8.x86 64 is already installed.

Package keyutils-1.5.10-7.el8.x86 64 is already installed.

Dependencies resolved.

Nothing to do.

Complete!

Procedure

1. Create a test file:

echo < Test_text> > test_file

IMA and EVM ensure that the example file **test_file** is assigned hash values, which are stored as its extended attributes.

2. Inspect extended attributes of the file:

getfattr -m . -d test_file
file: test_file
security.evm=0sAnDly4VPA0HArpPO/EqiutnNyBql

security.ima=0sAQOEDeuUnWzwwKYk+n66h/vby3eD security.selinux="unconfined_u:object_r:admin_home_t:s0"

The previous example output shows extended attributes related to SELinux and the IMA and EVM hash values. EVM actively adds a **security.evm** extended attribute and detects any offline tampering to xattrs of other files such as **security.ima** that are directly related to content integrity of files. The value of the **security.evm** field is in Hash-based Message Authentication Code (HMAC-SHA1), which was generated with the **evm-key** user key.

Additional resources

- Security hardening
- Integrity measurement architecture
- Extended verification module

24.9. RELATED INFORMATION

- Security hardening
- Integrity Measurement Architecture (IMA)
- Basic and advanced configuration of Security-Enhanced Linux (SELinux)

CHAPTER 25. USING ANSIBLE ROLES TO PERMANENTLY CONFIGURE KERNEL PARAMETERS

As an experienced user with good knowledge of Red Hat Ansible Engine, you can use the **kernel_settings** role to configure kernel parameters on multiple clients at once. This solution:

- Provides a friendly interface with efficient input setting.
- Keeps all intended kernel parameters in one place.

After you run the **kernel_settings** role from the control machine, the kernel parameters are applied to the managed systems immediately and persist across reboots.

25.1. INTRODUCTION TO THE KERNEL SETTINGS ROLE

RHEL System Roles is a collection of roles and modules from Ansible Automation Platform that provide a consistent configuration interface to remotely manage multiple systems.

RHEL System Roles were introduced for automated configurations of the kernel using the **kernel_settings** system role. The **rhel-system-roles** package contains this system role, and also the reference documentation.

To apply the kernel parameters on one or more systems in an automated fashion, use the **kernel_settings** role with one or more of its role variables of your choice in a playbook. A playbook is a list of one or more plays that are human-readable, and are written in the YAML format.

You can use an inventory file to define a set of systems that you want Ansible Engine to configure according to the playbook.

With the **kernel_settings** role you can configure:

- The kernel parameters using the **kernel settings sysctl** role variable
- Various kernel subsystems, hardware devices, and device drivers using the kernel_settings_sysfs role variable
- The CPU affinity for the systemd service manager and processes it forks using the kernel_settings_systemd_cpu_affinity role variable
- The kernel memory subsystem transparent hugepages using the kernel_settings_transparent_hugepages and kernel_settings_transparent_hugepages_defrag role variables

Additional resources

- README.md and README.html files in the /usr/share/doc/rhel-systemroles/kernel settings/ directory
- Working with playbooks
- How to build your inventory

25.2. APPLYING SELECTED KERNEL PARAMETERS USING THE KERNEL SETTINGS ROLE

Follow these steps to prepare and apply an Ansible playbook to remotely configure kernel parameters with persisting effect on multiple managed operating systems.

Prerequisites

- Your Red Hat Ansible Engine subscription is attached to the system, also called control machine, from which you want to run the **kernel_settings** role. See the How do I download and install Red Hat Ansible Engine article for more information.
- Ansible Engine repository is enabled on the control machine.
- Ansible Engine is installed on the control machine.



NOTE

You do not need to have Ansible Engine installed on the systems, also called *managed hosts*, where you want to configure the kernel parameters.

- The **rhel-system-roles** package is installed on the control machine.
- An inventory of managed hosts is present on the control machine and Ansible Engine is able to connect to them.

Procedure

1. Optionally, review the **inventory** file for illustration purposes:

cat /home/jdoe/<ansible_project_name>/inventory
[testingservers]

pdoe@192.168.122.98 fdoe@192.168.122.226

[db-servers] db1.example.com db2.example.com

[webservers] web1.example.com web2.example.com 192.0.2.42

The file defines the **[testingservers]** group and other groups. It allows you to run Ansible Engine more effectively against a specific collection of systems.

- 2. Create a configuration file to set defaults and privilege escalation for Ansible Engine operations.
 - a. Create a new YAML file and open it in a text editor, for example:
 - # vi /home/jdoe/<ansible_project_name>/ansible.cfg
 - b. Insert the following content into the file:

[defaults] inventory = ./inventory [privilege_escalation]
become = true
become_method = sudo
become_user = root
become_ask_pass = true

The **[defaults]** section specifies a path to the inventory file of managed hosts. The **[privilege_escalation]** section defines that user privileges be shifted to **root** on the specified managed hosts. This is necessary for successful configuration of kernel parameters. When Ansible playbook is run, you will be prompted for user password. The user automatically switches to **root** by means of **sudo** after connecting to a managed host.

- 3. Create an Ansible playbook that uses the **kernel_settings** role.
 - a. Create a new YAML file and open it in a text editor, for example:

vi /home/jdoe/<ansible_project_name>/kernel_roles.yml

This file represents a playbook and usually contains an ordered list of tasks, also called *plays*, that are run against specific managed hosts selected from your **inventory** file.

b. Insert the following content into the file:

- name: Configure kernel settings

hosts: testingservers

vars:

kernel settings sysctl:

- name: fs.file-max value: 400000

- name: kernel.threads-max

value: 65536

kernel settings sysfs:

- name: /sys/class/net/lo/mtu

value: 65000

kernel_settings_transparent_hugepages: madvise

roles:

- linux-system-roles.kernel_settings

The **name** key is optional. It associates an arbitrary string with the play as a label and identifies what the play is for. The **hosts** key in the play specifies the hosts against which the play is run. The value or values for this key can be provided as individual names of managed hosts or as groups of hosts as defined in the **inventory** file.

The **vars** section represents a list of variables containing selected kernel parameter names and values to which they have to be set.

The **roles** key specifies what system role is going to configure the parameters and values mentioned in the **vars** section.



NOTE

You can modify the kernel parameters and their values in the playbook to fit your needs.

4. Optionally, verify that the syntax in your play is correct.

ansible-playbook --syntax-check kernel-roles.yml

playbook: kernel-roles.yml

This example shows the successful verification of a playbook.

5. Execute your playbook.

ansible-playbook kernel-roles.yml

BECOME password:

PLAY [Configure kernel settings] ... PLAY RECAP **

fdoe@192.168.122.226 : ok=10 changed=4 unreachable=0 failed=0 skipped=6

rescued=0 ignored=0

pdoe@192.168.122.98 : ok=10 changed=4 unreachable=0 failed=0 skipped=6

rescued=0 ignored=0

Before Ansible Engine runs your playbook, you are going to be prompted for your password and so that a user on managed hosts can be switched to **root**, which is necessary for configuring kernel parameters.

The recap section shows that the play finished successfully (failed=0) for all managed hosts, and that 4 kernel parameters have been applied (changed=4).

6. Restart your managed hosts and check the affected kernel parameters to verify that the changes have been applied and persist across reboots.

Additional resources

- Getting started with RHEL System Roles
- README.html and README.md files in the/usr/share/doc/rhel-systemroles/kernel_settings/ directory
- Working with Inventory
- Configuring Ansible
- Working With Playbooks
- Using Variables
- Roles