ECE 231A HW 3

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Problem 1

(a)

$$\lim_{n \to \infty} [p(X_1, \dots, X_n)]^{\frac{1}{n}} = 2^{\lim_{n \to \infty} \frac{1}{n} \log_2[p(X_1, \dots, X_n)]}$$

$$= 2^{\lim_{n \to \infty} \frac{1}{n} \sum_{i=1}^n \log_2[p(X_i)]}$$

$$= 2^{E[\ln[p(X_i)]]}$$

$$= 2^{-H(x)}$$

(b)

$$E\left[\left(\prod_{i=1}^{n} f(X_{i})\right)^{\frac{1}{n}}\right] = \left(\left(E\left[\left(\prod_{i=1}^{n} f(X_{i})\right)^{\frac{1}{n}}\right]\right)^{n}\right)^{\frac{1}{n}}$$

$$\leq \left(E\left[\prod_{i=1}^{n} f(X_{i})\right]\right)^{\frac{1}{n}}$$

$$= \left(E^{n}[f(X_{1})]\right)^{\frac{1}{n}}$$

$$= E[f(X_{1})]$$

Therefore we have that

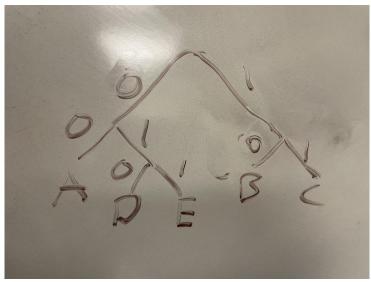
$$E\left[\left(\prod_{i=1}^n f(X_i)\right)^{\frac{1}{n}}\right] \le E[X_i]$$

Problem 2

(a)

$$H(X) = 2.246 \text{ Shannons}$$

(b)



So we have that the average length is $\boxed{2.3}$ bits.

(c)

codeword A is 001 codeword B is 0110 codeword C is 1001 codeword D is 1100 codeword E is 11110 Therefore the SFE codeword average length is $\boxed{3.8\,bits}$.

(d)

We have that the CDF of BAC is

$$P(X_1 = A) + P(X_1 = B, X_2 = A)(P(X_3 = A) + P(X_3 = B) + P(X_3 = C)) = \boxed{0.342}$$

Therefore for the SFE code, we have that

$$\bar{F} = P(X_1 = A) + P(X_1 = B, X_2 = A)(P(X_3 = A) + P(X_3 = B) + P(X_3 = C)) - \frac{1}{2}P(X_1 = B, X_2 = A)$$
 and

$$l = -\lceil \log_2(P(X_1 = B, X_2 = A, X_3 = c)) \rceil + 1 = 8$$

Thus we have that the SFE encoding is 01010110

Problem 3

(a)

$$F(X^{1}) = \begin{cases} 0.2 & \text{if } X^{1} = A \\ 0.5 & \text{if } X^{1} = B \\ 1 & \text{if } X^{1} = C \end{cases}$$

The interval for the first symbol is [0, 0.2).

(b)

$$F(X^{1}X^{2}) = \begin{cases} 0.04 & \text{if } X^{1}X^{2} = AA \\ 0.1 & \text{if } X^{1}X^{2} = AB \\ 0.2 & \text{if } X^{1}X^{2} = AC \\ 0.26 & \text{if } X^{1}X^{2} = BA \\ 0.35 & \text{if } X^{1}X^{2} = BB \\ 0.5 & \text{if } X^{1}X^{2} = BC \\ 0.6 & \text{if } X^{1}X^{2} = CA \\ 0.75 & \text{if } X^{1}X^{2} = CB \\ 1 & \text{if } X^{1}X^{2} = CC \end{cases}$$

Therefore the interval corresponding to AC is [0.1, 0.2)

(c)

For the interval [0,0.2), we have that the midpoint is 0.1 and the length is $-\lceil \log_2(0.2) \rceil + 1 = 4$ bits. Therefore the code word is $\boxed{0001}$ For the interval [0.1,0.2), we have that the midpoint is 0.15 and the length is $-\lceil \log_2(0.01) \rceil + 1 = 8$ bits, so the code word is $\boxed{00100}$

(d)

Therefore the cdf for

$$F(X^{1}X^{2}X^{3}X^{4} = ACCB) = 0.1 + 0.1 * 0.6 = 0.175$$

And thus we have that interval for the ACCB is [0.15, 0.16) so the midpoint is

$$\bar{F} = \frac{0.16 + 0.175}{2} = 0.1675$$

And since $p = 0.2 \cdot 0.5^2 \cdot 0.3 = 0.015$, we have that

$$l = -\lceil \log_2(0.015) \rceil + 1 = 8$$

Thus we have that the SFE code for ACCB is $\boxed{00101010}$

(e)

The binary encoding for 0.16 is 0.001010001111010111 and the binary encoding for 0.1675 is 0.00101011110000101 and the binary encoding for 0.175 is 0.0010110011001101101. Therefore 5 bits can be know if we do not know how ACCB continues.

Problem 4

ABBBCABBBCAAABB