ECE 231A HW 1

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Problem 1

If there was no slackness in the Kraft Inequality, then we would have

$$\sum_{i=1}^{m} D^{l_m - l_i} < D^{l_m}$$

For the max length l_m. This means that for some nodes on level l_m that is not a descendent of a codeword, or a codeword. Therefore while this node would be uniquely decodable but not able to form a sentence.

Problem 2

Since a uniquely decodable code is a instantanous code we can use the Kraft Inequality. We have

$$\sum_{i=1}^{6} D^{-l_i} \leq 1$$

The smallest D that satisfies this is $D = 3$, therefore a good lower bound on D would be $\boxed{3}$.

Problem 3

(a)

The entropy of $H(X)$ is

$$H(X) = -\sum_{x \in X} p(x) \log_2 (p(x))$$

and we have that

$$\begin{aligned} H(X|Y) &= \sum_{x \in X, y \in Y} p(y) H(X|Y = y) \\ &= \left(\sum_{x \in S} p(x) \right) H(X|Y = 1) + \left(\sum_{x \notin S} p(x) \right) H(X|Y = 0) \end{aligned}$$

We have

$$\begin{aligned} H(X|Y = 1) &= -\sum_{x \in S} p(x|Y = 1) \log_2 (p(x|Y = 1)) \\ &= -\sum_{x \in S} \frac{p(x)}{\sum_{x \in S} p(x)} \log_2 \left(\frac{p(x)}{\sum_{x \in S} p(x)} \right) \end{aligned}$$

likewise we have

$$\begin{aligned} H(X|Y = 0) &= -\sum_{x \notin S} p(x|Y = 0) \log_2 (p(x|Y = 0)) \\ &= -\sum_{x \notin S} \frac{p(x)}{\sum_{x \notin S} p(x)} \log_2 \left(\frac{p(x)}{\sum_{x \notin S} p(x)} \right) \end{aligned}$$

Therefore we have

$$\begin{aligned} H(X|Y) &= -\sum_{x \in S} p(x) \left(\log_2 (p(x)) - \log_2 \left(\sum_{x \in S} p(x) \right) \right) - \sum_{x \notin S} p(x) \left(\log_2 (p(x)) - \log_2 \left(\sum_{x \notin S} p(x) \right) \right) \\ &= H(X) + \sum_{x \in S} p(x) \log_2 \left(\sum_{x \in S} p(x) \right) + \sum_{x \notin S} p(x) \log_2 \left(\sum_{x \notin S} p(x) \right) \end{aligned}$$

Therefore we have

$$H(X) - H(X|Y) = -\sum_{x \in S} p(x) \log_2 \left(\sum_{x \in S} p(x) \right) - \sum_{x \notin S} p(x) \log_2 \left(\sum_{x \notin S} p(x) \right)$$

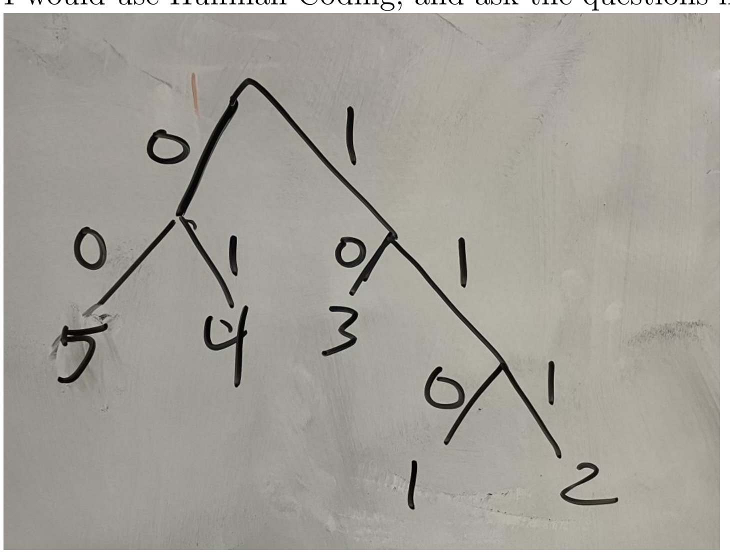
$$H(X) - H(X|Y) = \boxed{-q \log_2(q) - (1-q) \log_2(1-q)}$$

(b)

$H(X) - H(X|Y)$ is maximized when $\sum_{x \in S} p(x) = \sum_{x \notin S} p(x) = \frac{1}{2}$, this is possible when $S = \boxed{2, 5}$ or $S = \boxed{1, 2, 4}$

(d)

I would use Huffman Coding, and ask the questions in the following format,



the average question length would be $\boxed{2.25}$

Problem 4

(a)

If a codeword is l_j long, but if it has to start with $C(i)$, then it would be effectively be concatenating $C(i)$ with a code word from A_{j-i}, ie all the code words with length $l_j - l_i$. Therefore the total number of words of A_j would be the total combinations of A_{j-i}, ie $2^{l_j - l_i}$

Likewise, if a codeword is l_j long, but if it has to end with $C(i)$, then it would be effectively be concatenating a code word from A_{j-i} with $C(i)$, ie all the code words with length $l_j - l_i$. Therefore the total number of words of A_j would be the total combinations of A_{j-i}, ie $2^{l_j - l_i}$

(b)

If we assume that $l_j > l_i$ we would have that the total number of words to remove from A_j would be the toal number of words to remove that start with $C(i)$ plus the total number of words that end with $C(i)$. Therefore we would have that the total number of words to remove would be $2^{l_j - l_i + 1}$. And if $l_j = l_i$ then we would only remove 1 word, $C(i)$.

(c)

We have that for any $1 \leq j \leq k$, in order for the algorithm to not fail, we must have that the number of inital codewords, is greater than the number of removed code words, in otherwords we must have that

$$2^{l_j} > \sum_{i=1}^{j-1} 2^{l_j - l_i + 1}$$

Since, $\sum_{i=1}^{j-1} 2^{l_j - l_i + 1} < \sum_{i=1}^{j} 2^{l_j - l_i + 1}$, we can assure the above inequality will be satisfied by the following inequality:

$$2^{l_j} \geq \sum_{i=1}^{j} 2^{l_j - l_i + 1}$$

This can be generalized to

$$2^{l_k} \geq \sum_{i=1}^{k} 2^{l_k - l_i + 1}$$

Since for $1 \leq j \leq k$:

$$2^{l_k} \geq \sum_{i=1}^{k} 2^{l_k - l_i + 1}$$

$$2^{l_k} 2^{l_j - l_k} \geq 2^{l_j - l_k} \sum_{i=1}^{k} 2^{l_k - l_i + 1}$$

$$2^{l_j} \geq \sum_{i=1}^{k} 2^{l_j - l_i + 1} \geq \sum_{i=1}^{j} 2^{l_j - l_i + 1}$$

Therefore we could rearange the inequality to

$$1 \geq \sum_{i=1}^{j} 2^{-l_i + 1}$$

$$\sum_{i=1}^{j} 2^{-l_i + 1} \leq \frac{1}{2}$$

(d)

Let $E[length(C(U))] = L$, we want to minimize

$$L = \sum_{i=1}^{k} p_i l_i$$

given

$$\sum_{i=1}^{k} 2^{-l_i} \leq \frac{1}{2}$$

Using a langrange multipler, we get

$$J = \sum_{i=1}^{k} p_i l_i + \lambda \sum_{i=1}^{k} 2^{-l_i}$$

$$\frac{\partial J}{\partial l_i} = p_i - \lambda 2^{-l_i} \ln(2)$$

Therefore we get that

$$2^{-l_i} = \frac{p_i}{\lambda \ln(2)}$$

Plugging this into $\sum_{i=1}^{k} 2^{-l_i} \leq \frac{1}{2}$, we get

$$\lambda = \frac{2}{\ln(2)}$$

Therefore we get that

$$2^{-l_i} = \frac{p_i}{2}$$

And thus the optimal code length is

$$l_i^* = -\log_2 p_i + \log_2(2)$$

But since l_i must be an integer we have

$$-\log_2(p_i) \leq l_i \leq -\log_2(p_i) + 2$$

Therefore we have that

$$-\sum_{i=1}^{k} p_i \log_2(p_i) \leq \sum_{i=1}^{k} p_i l_i \leq \left(-\sum_{i=1}^{k} p_i \log_2(p_i)\right) + 2$$

Or in other words:

$$H(U) \leq E[length(C(U))] \leq H(U) + 2$$

Problem 5

(a)

We have that

$$H(X) = -\sum_{i\in\chi_1}(1-\gamma)p(i)\log_2\left((1-\gamma)p(i)\right) - \sum_{i\in\chi_2}\gamma q(i)\log_2\left(\gamma q(i)\right)$$

Likewise for $H(X,Y)$ we have

$$H(X,Y) = -\sum_{y\in\{1,2\}}\sum_{x\in\{1,2,..m\}} P(x,y)\log_2(P(x,y))$$

we have that

$$p(x,1) = \begin{cases}(1-\gamma)p(x) & \text{if } x\in\chi_1\\ 0 & \text{if } x\notin\chi_2\end{cases}$$

$$p(x,2) = \begin{cases}0 & \text{if } x\in\chi_1\\ \gamma q(x) & \text{if } x\notin\chi_2\end{cases}$$

Therefore we have

$$H(X,Y) = -\sum_{i\in\chi_1}(1-\gamma)p(i)\log_2\left((1-\gamma)p(i)\right) - \sum_{i\in\chi_2}\gamma q(i)\log_2\left(\gamma q(i)\right)$$

And thus we have

$$\boxed{H(X,Y) = H(X)}$$

(b)

$$
\begin{aligned}
H(X) &= -\sum_{i\in\chi_1}(1-\gamma)p(i)\log_2\left((1-\gamma)p(i)\right) - \sum_{i\in\chi_2}\gamma q(i)\log_2\left(\gamma q(i)\right) \\
&= -(1-\gamma)\sum_{i\in\chi_1} p(i)\left(\log_2(p(i)) + \log_2(1-\gamma)\right) - \gamma\sum_{i\in\chi_2} q(i)\left(\log_2(q(i)) + \log_2(\gamma)\right) \\
&= -(1-\gamma)\left(\log_2(1-\gamma) + \sum_{i\in\chi_1} p(i)\log_2(p(i))\right) - \gamma\left(\log_2(\gamma) + \sum_{i\in\chi_2} q(i)\log_2(q(i))\right) \\
&= -(1-\gamma)\left(\log_2(1-\gamma) - H(X_1)\right) - \gamma\left(\log_2(\gamma) - H(X_2)\right) \\
&= \boxed{(1-\gamma)\left(H(X_1) - \log_2(1-\gamma)\right) + \gamma\left(H(X_2) - \log_2(\gamma)\right)}
\end{aligned}
$$

(c)

Assuming $H(X_1)$ and $H(X_2)$ are in units of shannons we have, that to maximize $H(x)$ with take the derivative of $H(X)$ with respect to γ we get

$$
\frac{\partial H(X)}{\partial\gamma} = \left(\log_2(1-\gamma) - H(X_1)\right) + \left(H(X_2) - \log_2(\gamma)\right)
$$

This is maximized when

$$
\begin{aligned}
\frac{\partial H(X)}{\partial\gamma} &= 0 \\
\left(\log_2(1-\gamma) - H(X_1)\right) + \left(H(X_2) - \log_2(\gamma)\right) &= 0 \\
H(X_2) - H(X_1) &= \log_2(\gamma) - \log_2(1-\gamma) \\
e^{H(X_2)-H(X_1)} &= \frac{\gamma}{1-\gamma} \\
1 - \gamma e^{H(X_2)-H(X_1)} &= \gamma \\
\gamma &= \boxed{\frac{2^{H(X_2)-H(X_1)}}{1+2^{H(X_2)-H(X_1)}}}
\end{aligned}
$$

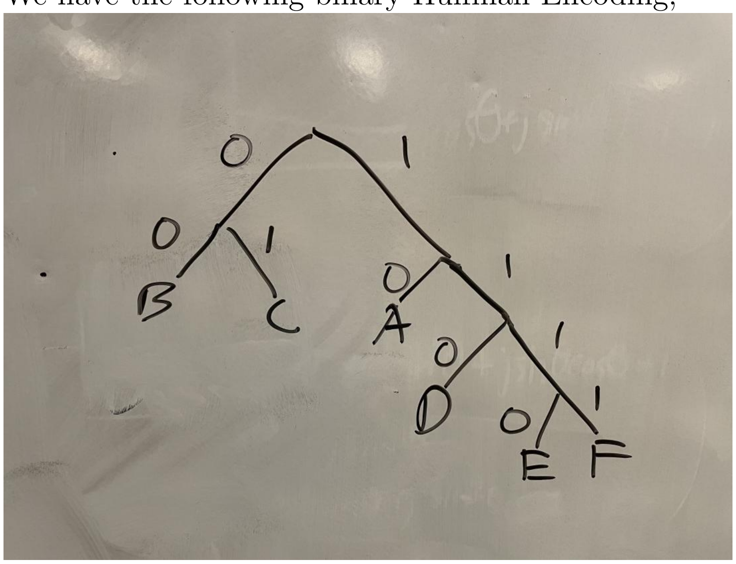
Problem 6

(a)

$$
\begin{aligned}
1, 2H(x) &= -\sum_{x\in X} p(x)\log_2(p(x)) \\
&= p(A)\log_2(p(A)) + p(B)\log_2(p(B)) + p(C)\log_2(p(C)) \\
&+ p(D)\log_2(p(D)) + p(E)\log_2(p(E)) + p(F)\log_2(p(F)) \\
&= \frac{1}{2}\log_2(4) + \frac{1}{4}\log_2(8) + \frac{3}{16}\log_2(\frac{16}{3}) + \frac{1}{16}\log_2(16) \\
&= \boxed{2.452 \text{ shannons}}
\end{aligned}
$$

(b)

We have the following binary Huffman Encoding,



the expected length is $\boxed{2.5}$.