

Median of 2 sorted Arrays :- (self revision)

ex: arr1 = [1, 3, 4, 7, 10, 12]

arr2 = [2, 3, 6, 15]

what we cannot do?

→ we cannot directly apply "BS on real domain"
~~by keeping the check for~~

using the logic:- median is a value
such that $(\# \text{elem} \leq \text{median})$

$= (\# \text{elem} \geq \text{median})$

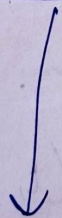
bcz

above median is $\frac{4+6}{2} = 5$

& based on above logic we could be
getting any real no betn (4, 6)

what to do?

sorted array: [1 2 3 4 | 6 7 10 12 15]



if arr = [1 2 3 | 4 5]

So to find

median \Rightarrow we need
the left & right
elem from pt of symmetry
(for even len)
& for odd: left from
symmetry.

bcz that is the only way to find median

⇒ Even len

↳ we know both left & right
half will have equal no
of elem

$$\& \text{ that } n = \left(\frac{n_1 + n_2 + 1}{2} \right)$$

(applies for
both
even & odd)
(L & R) (only in L)

or
$$\left(\frac{n_1 + n_2}{2} \right)$$

* So. we need to print the shortest

$\text{cut} = \left(\frac{n_1 + n_2 + 1}{2} \right)$ elem on the left
side of symmetry

options are: (take from beginning)

take 0 from arr1, cut from arr2

1 " , cut-1 "

2 " , cut-2 "

⋮

cut " > 0 "

{ one of
them
is a
valid
configuration

bcz median
is a specific
value only

& we need ~~the~~ greatest from
left & ~~largest~~ smallest from right

options:

take (# elem from arr1)

$$\geq 0 \rightarrow \min(\text{cnt}, \text{arr1.size}())$$

eg:-
arr1 = [1, 2, 3]
arr2 = [- - -]
median = 5th elem

i.e. max^m elem we
should/can't pick
from arr1.

Now, if we do linearly pick up.

$$0 \rightarrow \min(\text{cnt}, \text{arr1.size}())$$

$$\text{then P.C} = O(N)$$

So can we do it using B.S?

general format to B.S on a Search
space

is check [mid] :- _ 000111 _

But here: (only 1 config is valid)

$$\text{eg:- arr1} = [1, 3, 4, 7, 10, 12]$$

$$\text{arr2} = [2, 3, 6, 15]$$

$$\frac{n_1 + n_2 + 1}{2} = \frac{11}{2} = 5.5 \rightarrow \text{5 elem on left}$$

Options 2

~~can't pick~~

5 elems on left.

so for picking from the smaller array

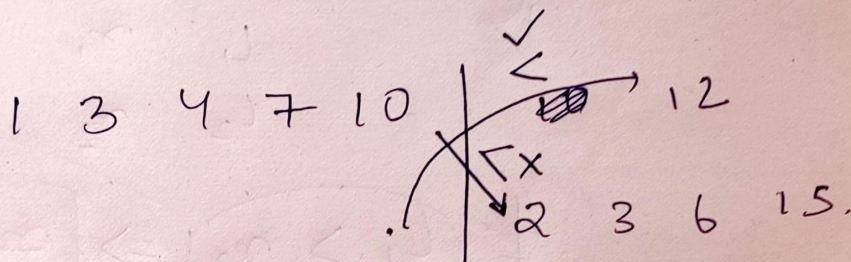
i.e. (0 elem, 1 elem, ... - min(cut, smaller arr. size))

to make 'cut' elems on left + rest from other array)

options:-

0	1	2	3	4
x	x	✓	x	x

1) ① from arr 2.

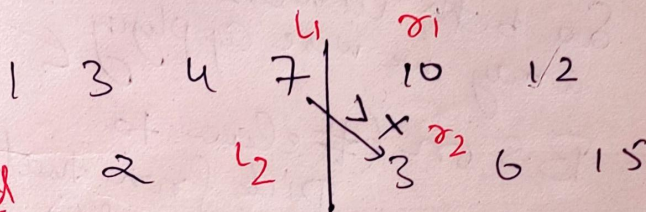


so not valid

2) ①

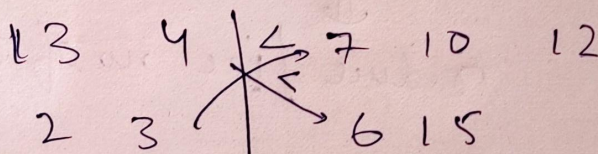
recomparing.

$l_1 < r_2$
 $l_2 < r_1$ for valid



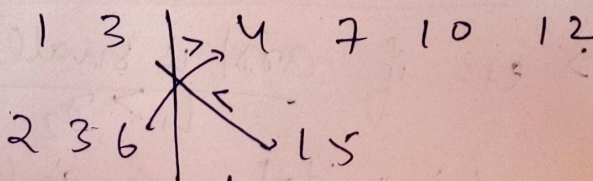
not valid

3) ②



valid config.

4) ③



invalid

options :

ans :

0	1	2	3	4
x	x	✓	x	x

Can we perform a Binary Search on this?

~~only~~



create a check fn

Some Observ'n from how we decide what is ans.



$L_1 \quad r_1$
 $L_2 \quad r_2$

if $(L_i > r_2)$ ⇒

remove bigger
~~reduce~~ n elem
from arr1.

||.

inc n elem on
arr2

So both whichever
array we were applying

BS on (# elem to
pick from that
array for left)



reduce / inc no of elem

⇒ do this way we can go to
go left / right

eg: if $arr1$ is smaller ~~array~~ $arr2$
 $L_1 > r_2$ then use elem from
 $arr1$
 $hi = mid - 1$ ⇒ go left

Note :- There'll never be a case

when $l_1 > r_2$
& $l_2 > r_1$ X

* \rightarrow Sim if arr2 was the smaller array

then for $(l_1 > r_2) \rightarrow$ now dec/~~re~~
remove
bigger elem
from arr1.

so right (# elem
from arr2
for left) \leftarrow

i.e (1st)
elem to arr2

~~arr~~ $l_2 = \text{mid} + 1$

& once $(l_1 < r_2 \text{ \& } l_2 < r_1)$

(mid) rep (no of elem to take from
smaller arr for left)

even sized

$$\text{median} = \frac{\max(l_1, l_2) + \min(r_1, r_2)}{2}$$

Sim for

odd sized array

$$\text{median} = \frac{\max(l_1, l_2)}{2}$$