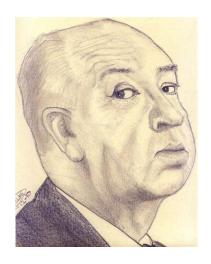
Fundamental mechanism of the quantum computational speedup

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Example of quantum speedup

- Bob, the problem setter, hides a ball in one of four drawers
- Alice, the problem solver, is to locate it by opening drawers
- In the classical case she may need to open up to three drawers, in the quantum case it always takes one
- There is a quantum speed-up

BOB





ALICE



Some jargon

- Drawer and ball problem: oracle problem
- Checking whether the ball is in a drawer: function evaluation (oracle query)
- b = number of the drawer with the ball; checking whether the ball is in drawer a: computing (evaluating) the Kronecker function $\delta(b, a)$
- Quantum speedup: number of function evaluations required by the quantum algorithm vs number required by the best known classical algorithm

Speedup is poorly understood

- Dozens of speedups discovered
- All by means of ingenuity
- No fundamental explanation of the speedup
- No unified explanation of the amount of speedup
- A lacuna of quantum computer science
- Quantum cryptography, the other pillar of quantum information, directly relies on the foundations of quantum mechanics

Representation incompleteness

- Poor understanding of the speedup: the usual representation of quantum algorithms, limited to the process of solving the problem, is physically incomplete
- Drawer number: 01

		meas.
$ 00\rangle_A$	$\Rightarrow U \Rightarrow$	$ 01\rangle_A$

- Quantum register A: contains drawer number a.
- Alice unitarily changes input state it into an output state that encodes the solution, then acquired by a final measurement
- Initial measurement: missing
- Number of the drawer with the ball: not represented physically
- Completing the physical representation explains the speedup

Three steps

- 1) extending the usual representation, limited to the process of solving the problem, to that of setting it
- 2) relativizing the extended representation to Alice, who cannot see the problem setting selected by Bob (should be hidden inside the black box)
- 3) symmetrizing the relativized representation for time-reversal – to represent the reversibility of the computation process

Step 3) provides a quantitative explanation of the speedup

1) extending the representation

	meas. \hat{B}		meas. Â
	$(00\rangle_B + 01\rangle_B + 10\rangle_B + 11\rangle_B) 00\rangle_A$		
	↓		
(6.5)	$ 01\rangle_{\!\scriptscriptstyle B} 00\rangle_{\!\scriptscriptstyle A}$	\Rightarrow U \Rightarrow	$ 01\rangle_{B} 01\rangle_{A}$



- Adding a quantum register B that contains the number of the drawer with the ball
- Unitary transformation between initial and final measurement outcomes
- Extended representation works for Bob (and any external observer)
- Not for Alice
- Input state $|01\rangle_B |00\rangle_A$ would tell her the number of the drawer with the ball before she begins her problem solving action

Relational quantum mechanics

- Quantum state has meaning to an observer
- Rejects the notion of absolute, or observer independent, state of a system
- E. g. a quantum state could be sharp to an observer and a superposition to another

2) relativizing the representation to Alice

meas. B		meas. of \hat{A}
$(00\rangle_B + 01\rangle_B + 10\rangle_B + 11\rangle_B) 00\rangle_A$	\Rightarrow U \Rightarrow	$ 00\rangle_B 00\rangle_A + 01\rangle_B 01\rangle_A + 10\rangle_B 10\rangle_A + 11\rangle_B 11\rangle_A$
		₩.
		$ 01\rangle_{\!\scriptscriptstyle B} 01\rangle_{\!\scriptscriptstyle A}$

- postponing the projection of the quantum state due to the initial measurement at the end of the unitary part of Alice's action
- Throughout it, Alice remains completely ignorant of number of the drawer with the ball selected by Bob
- Legitimate: degree of freedom of quantum description

3) symmetrizing for time-reversal

 1) outcome of initial measurement random, 2) unitary transformation between initial and final measurement outcomes

 Selection of the number of the drawer with the ball by:

1) initial measurement

meas. \hat{B}		
$(00\rangle_B + 01\rangle_B + 10\rangle_B + 11\rangle_B) 00\rangle_A$		
U		
$ 01\rangle_B 00\rangle_A$	\Rightarrow U \Rightarrow	$ 01\rangle_{\!\scriptscriptstyle B} 01\rangle_{\!\scriptscriptstyle A}$

- 2) final measurement
- Which measurement?

		meas. of \hat{A}
$(00\rangle_B + 01\rangle_B + 10\rangle_B + 11\rangle_B) 00\rangle_A$	\Rightarrow U \Rightarrow	$ 00\rangle_B 00\rangle_A + 01\rangle_B 01\rangle_A + 10\rangle_B 10\rangle_A + 11\rangle_B 11\rangle_A$
		Ų
$ 01\rangle_{\!\scriptscriptstyle B} 00\rangle_{\!\scriptscriptstyle A}$	$\leftarrow \mathbf{U}^{\dagger} \leftarrow$	$ 01\rangle_B 01\rangle_A$

- Neither one alone: would introduce preferred direction of time, unjustified in a reversible context
- Share selection evenly between initial and final measurements

3) symmetrizing for time-reversal

Initial and final measurements reduce to partial measurements that evenly and without redundancy contribute to the selection, e. g.:

- initial measurement of \hat{B} , reduced to that of \hat{B}_l , selects 0 of 01; outcome propagated forward in time
- final measurement of \hat{A} , reduced to that of $\hat{A_r}$, selects 1 of 01; outcome propagated backward in time

Performing the two propagations in a sequence time-symmetrizes the representation

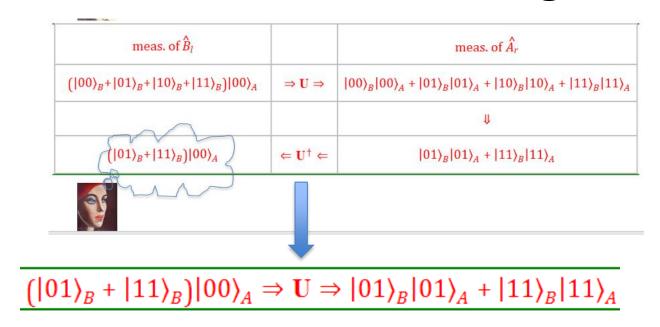
Superposition of all the possible ways of sharing

meas. of \hat{B}_l		meas. of $\hat{A_r}$
$(00\rangle_B + 01\rangle_B + 10\rangle_B + 11\rangle_B) 00\rangle_A$		
1		
$(00\rangle_B + 01\rangle_B) 00\rangle_A$	\Rightarrow U \Rightarrow	$ 00\rangle_B 00\rangle_A + 01\rangle_B 01\rangle_A$
		th .
$ 01\rangle_{\!\scriptscriptstyle B} 00\rangle_{\!\scriptscriptstyle A}$	$\leftarrow \mathbf{U}^{\dagger} \leftarrow$	$ 01\rangle_{_B} 01\rangle_{_A}$



meas. of \hat{B}_l		meas. of $\hat{A_r}$
$(00\rangle_B + 01\rangle_B + 10\rangle_B + 11\rangle_B) 00\rangle_A$	\Rightarrow U \Rightarrow	$ 00\rangle_B 00\rangle_A + 01\rangle_B 01\rangle_A + 10\rangle_B 10\rangle_A + 11\rangle_B 11\rangle_A$
		1
$(01\rangle_B + 11\rangle_B) 00\rangle_A$	$\leftarrow \mathbf{U}^{\dagger} \leftarrow$	$ 01\rangle_{B} 01\rangle_{A}+ 11\rangle_{B} 11\rangle_{A}$

Advanced knowledge



- Computational complexity of the problem: reduced to finding a ball hidden in one of two drawers {01,11}
- Solving the reduced problem classically requires just one function evaluation, a fortiori quantumly

Reduced problem

- Oracle problem can always be solved quantumly with the number of function evaluations required to solve its reduced problem classically
- Reduced problem: original one but for the fact that the problem solver knows in advance a part of the problem setting that corresponds to half solution
- Found an upper bound to the quantum computational complexity of oracle problem
- It holds for any oracle problem and can be computed on the basis of the problem alone

Quantum superposition

 Taking the superposition of all the time-symmetric transformations rebuilds the original transformation with respect to Alice

```
(|00\rangle_B + |01\rangle_B)|00\rangle_A \Rightarrow \mathbf{U} \Rightarrow |00\rangle_B|00\rangle_A + |01\rangle_B|01\rangle_A
+
(|00\rangle_B + |10\rangle_B)|00\rangle_A \Rightarrow \mathbf{U} \Rightarrow |00\rangle_B|00\rangle_A + |10\rangle_B|10\rangle_A
+
+
\vdots
\vdots
(|00\rangle_B + |01\rangle_B + |11\rangle_B)|00\rangle_A \Rightarrow \mathbf{U} \Rightarrow |00\rangle_B|00\rangle_A + |01\rangle_B|01\rangle_A + |10\rangle_B|10\rangle_A + |11\rangle_B|11\rangle_A
```

 Symmetrization for time-reversal partitions the transformation with respect to Alice into a superposition of time-symmetric transformations each solving an instance of the reduced problem

Upper bound checked on the major quantum algorithms

- Deutsch, Deutsch&Jozsa, Simon, Shor, Grover, Abelian hidden subgroup (12 algorithms)
- All optimal in character
- Upper bound always coincides with the number of function evaluations required to solve the problem in an optimal quantum way
- Conjecture: this holds in general, for any oracle problem

Conjecture

```
(|01\rangle_B + |11\rangle_B)|00\rangle_A \Rightarrow \mathbf{U} \Rightarrow |01\rangle_B|01\rangle_A + |11\rangle_B|11\rangle_A
```

- U requires just one function evaluation
- By an optimal quantum algorithm, a non-optimal one would require a higher number
- The quantum description speaks to us in classical logic

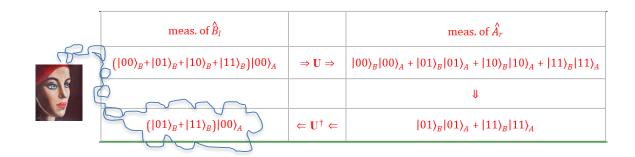
Summarizing

- In all cases, half of the random outcome of the initial measurement (half of the problem setting) selected back in time by final measurement
- This tells nothing to Bob, who knows the outcome of the initial measurement to start with



meas. of \hat{B}_l		meas. of $\hat{A_r}$
$(00\rangle_B + 01\rangle_B + 10\rangle_B + 11\rangle_B) 00\rangle_A$		
Ų		
$(00\rangle_B + 01\rangle_B) 00\rangle_A$	\Rightarrow U \Rightarrow	$ 00\rangle_B 00\rangle_A + 01\rangle_B 01\rangle_A$
		Ų
$ 01\rangle_B 00\rangle_A$	← U [†] ←	$ 01\rangle_{\!\scriptscriptstyle B} 01\rangle_{\!\scriptscriptstyle A}$

To Alice, who is shielded from the outcome of the initial measurement, it tells half of it and thus a corresponding half of the solution



Time-symmetrization by reducing initial and final measurements to complementary partial measurements: inspired by the work of Dolev and Elitzur on the non sequential behavior of the wave function highlighted by partial measurement

Positioning

Two main approaches to the quantum speedup:

- Quantum computer science (quantum complexity classes and their relations to the classical ones)
- Relation between speedup and fundamental quantum features (entanglement/discord)

Positioning – quantum computer science

- Quantum computer science is analytic in character
- The present explanation of the speedup is synthetic, derived from the foundations
- Like in analytic geometry (mathematics on coordinates) and synthetic geometry (derivation of theorems from postulates)
- Analytic counterpart of the upper bound
- Difficulty of deriving it?

Positioning – relation between speedup and the foundations

- The speedup appears to always depend on the exact nature of the problem while the reason for it varies from problem to problem (Vedral, Henderson)
- Present fundamental explanation applies to all oracle problems and is quantitative in character

Conclusion

- Found an upper bound to the computational complexity of any oracle problem, always coinciding with the number of function evaluations required by the optimal quantum algorithm
- Derived in a synthetic way from the foundations of quantum mechanics

Future work

- (1) Checking whether the upper bound always coincides with the number of function evaluations required by the optimal quantum algorithm
- (2) Finding oracle problems liable of interesting speedups, classifying quantum complexity of oracle problems (compare with existing quantum complexity classes)
- (3) Further studying the fundamental implications of an explanation of the speedup that merges time-symmetric quantum mechanics and quantum information

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