

Data Structure Algorithm Assignment-3

Question: <https://leetcode.com/problems/climbing-stairs/description/>

Answer: <https://leetcode.com/problems/climbing-stairs/submissions/1628585122>

Description:

Time Complexity = $O(n)$

- You have a for loop that starts at $i = 3$ and runs until $i < n$, so the loop runs approximately $n - 3$ times.
- Each iteration does a constant amount of work (addition and assignment).
- Therefore, time complexity = $O(n)$.

Space Complexity = $O(1)$

- You are only using a fixed number of variables (prev1, prev2, cur), no additional space grows with n .
- Thus, space complexity = $O(1)$ (constant space).

Screenshot:

The screenshot shows a LeetCode problem page for "70. Climbing Stairs". The problem description states: "You are climbing a staircase. It takes n steps to reach the top. Each time you can either climb 1 or 2 steps. In how many distinct ways can you climb to the top?". Example 1 shows input $n = 2$ and output 2, with two ways: 1 step + 1 step, and 2 steps. Example 2 shows input $n = 3$ and output 3, with three ways: 1 step + 1 step + 1 step, 1 step + 2 steps, and 2 steps + 1 step. The constraints are $1 \leq n \leq 45$. The solution is implemented in JavaScript using a for loop starting from 3 and updating prev1 and prev2. The code is as follows:

```
1 /**
2  * @param {number} n
3  * @return {number}
4  */
5 var climbStairs = function(n) {
6   if (n <= 3) return n;
7
8   let prev1 = 3;
9   let prev2 = 2;
10  let cur = 0;
11
12  for (let i = 3; i < n; i++) {
13    cur = prev1 + prev2;
14    prev2 = prev1;
15    prev1 = cur;
16  }
17
18  return cur;
19  return b;
20 };
```

The solution is marked as "Accepted" with a runtime of 0 ms.

Question: <https://leetcode.com/problems/merge-two-sorted-lists/description/>

Answer: <https://leetcode.com/problems/merge-two-sorted-lists/submissions/1628594172>

Description:

Time Complexity: $O(n + m)$

- Each node from list1 and list2 is visited exactly once.
- At each step, you either move list1 or list2 forward by one node.
- If n is the length of list1 and m is the length of list2, you do a total of $n + m$ operations.

Space Complexity: $O(1)$

- You are using a few pointers (Dummy Head, tail, list1, list2), but you are not creating any new nodes — you are just rearranging existing nodes.
- extra space used is constant, regardless of the input size.

Screenshot:

The screenshot shows the LeetCode interface for the problem "21. Merge Two Sorted Lists". The problem description states: "You are given the heads of two sorted linked lists list1 and list2. Merge the two lists into one sorted list. The list should be made by splicing together the nodes of the first two lists. Return the head of the merged linked list." An example is provided: list1 (1, 2, 4) and list2 (1, 3, 4) are merged into a new list (1, 1, 2, 3, 4, 4). The code editor on the right shows a JavaScript solution using a dummy head and a tail pointer to merge the two lists by splicing their nodes. The test results at the bottom show the solution is "Accepted" with a runtime of 0 ms.

```
1 /**
2  * Definition for singly-linked list.
3  * function ListNode(val, next) {
4  *   this.val = (val===undefined ? 0 : val)
5  *   this.next = (next===undefined ? null : next)
6  * }
7  */
8
9  * @param {ListNode} list1
10 * @param {ListNode} list2
11 * @return {ListNode}
12 */
13 var mergeTwoLists = function(list1, list2) {
14   if (!list1) return list2;
15   if (!list2) return list1;
16
17   const dummyhead = { next: null };
18
19   let tail = dummyhead;
20
21   while (list1 && list2) {
22     if (list1.val < list2.val) {
23       tail.next = list1;
24       list1 = list1.next;
25     } else {
26       tail.next = list2;
27       list2 = list2.next;
28     }
29     tail = tail.next;
30   }
31   tail.next = list1 || list2;
32   return dummyhead.next;
33 }
```

Question: <https://leetcode.com/problems/palindrome-linked-list/description/>

Answer: <https://leetcode.com/problems/palindrome-linked-list/submissions/1628597648>

Description:

Time Complexity: $O(n)$

- The while loop (finding middle + reversing first half) goes through about half of the list — $O(n/2)$ operations.
- The second while loop (comparing two halves) also goes through about half — another $O(n/2)$ operations.

Total time = $O(n/2) + O(n/2) = O(n)$.

Space Complexity: $O(1)$

- You're only using a few pointers (slow, fast, prev, next) — all are constant-sized variables.
- No extra space proportional to input size (no array, no recursion stack).
- So, the space complexity is: $O(1)$

Screenshot:

The screenshot shows the LeetCode interface for problem 234, "Palindrome Linked List". The problem description states: "Given the head of a singly linked list, return true if it is a palindrome or false otherwise." Two examples are provided: Example 1 with input [1,2,2,1] and output true, and Example 2 with input [1,2] and output false. The constraints are also listed. On the right, the code editor shows a JavaScript solution. The code defines a ListNode structure and a function isPalindrome that uses a fast-slow pointer technique to find the middle of the list, reverses the first half, and then compares the two halves. The bottom of the screenshot shows the test results, indicating that the solution was "Accepted" with a runtime of 0 ms.

```
1 /**
2  * Definition for singly-linked list.
3  * function ListNode(val, next) {
4  *   this.val = (val===undefined ? 0 : val)
5  *   this.next = (next===undefined ? null : next)
6  * }
7  */
8 /**
9  * @param {ListNode} head
10 * @return {boolean}
11 */
12 var isPalindrome = function(head) {
13   if (!head || !head.next)
14     return true;
15
16   let slow = head;
17   let fast = head;
18   let prev = null;
19
20   // Find the middle of the linked list and reverse the first half
21   while (fast && fast.next) {
22     fast = fast.next.next;
23     let temp = slow.next;
```

Question: <https://leetcode.com/problems/palindrome-linked-list/description/>

Answer: <https://leetcode.com/problems/linked-list-cycle/submissions/1628599176>

Description:

Time Complexity: $O(n)$

- In the worst case, the fast pointer moves through all the nodes of the linked list.
- If there is no cycle, fast will reach the end (null), visiting each node once $\rightarrow O(n)$.
- If there is a cycle, the fast and slow pointers will meet somewhere inside the cycle
- The distance before entering the cycle is at most n , and once inside, the meeting happens in at most n steps $\rightarrow O(n)$

Space Complexity: $O(1)$

- Only two pointers (fast and slow) are used.
- No extra data structures are needed (no arrays, hash maps, etc.).
- Final space complexity: $O(1)$ (constant space).

Screenshot:

141. Linked List Cycle Solved

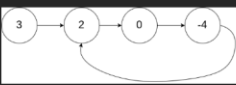
Easy Topics Companies

Given `head`, the head of a linked list, determine if the linked list has a cycle in it.

There is a cycle in a linked list if there is some node in the list that can be reached again by continuously following the `next` pointer. Internally, `pos` is used to denote the index of the node that `tail's next` pointer is connected to. **Note that `pos` is not passed as a parameter.**


Return `true` if there is a cycle in the linked list. Otherwise, return `false`.

Example 1:



Input: `head = [3,2,0,-4]`, `pos = 1`
Output: `true`
Explanation: There is a cycle in the linked list, where the tail connects to the 1st node (0-indexed).

Example 2:



```
JavaScript
function ListNode(val) {
  this.val = val;
  this.next = null;
}

/**
 * @param {ListNode} head
 * @return {boolean}
 */
var hasCycle = function(head) {
  if (!head || !head.next) return false;
  let slow = head;
  let fast = head.next;
  while (slow !== fast) {
    if (!fast || !fast.next) return false;
    slow = slow.next;
    fast = fast.next.next;
  }
  return true;
};
```

Accepted Runtime: 29 ms

Case 1 Case 2 Case 3

Question: <https://leetcode.com/problems/linked-list-cycle/description/>

Answer: <https://leetcode.com/problems/remove-nth-node-from-end-of-list/submissions/1628603281>

Description:

Time Complexity: $O(L)$

$O(L)$ — where L is the length of the linked list.

- The function first advances the head pointer by n steps — $O(n)$.
- Then it continues to traverse the rest of the list along with dummy — up to $O(L - n)$ steps.
- Therefore, the total number of operations is proportional to the length of the list: $O(L)$.

Space Complexity: $O(1)$

$O(1)$ — constant space.

- No extra data structures are used that grow with input size.
- Only a few pointers (res, dummy, and head) are used.

Screenshot:

The screenshot shows a LeetCode problem page for "19. Remove Nth Node From End of List". The problem is marked as "Solved" and "Medium". The description states: "Given the head of a linked list, remove the n^{th} node from the end of the list and return its head." Example 1 shows a linked list [1, 2, 3, 4, 5] where the 4th node (4) is removed, resulting in [1, 2, 3, 5]. Example 2 shows a linked list [1] where the 1st node is removed, resulting in an empty list []. The code editor on the right shows a JavaScript solution using a dummy node and a head pointer to traverse the list and remove the n^{th} node from the end. The test results show "Accepted" with a runtime of 0 ms.

19. Remove Nth Node From End of List

Medium Topics Companies Hint

Given the head of a linked list, remove the n^{th} node from the end of the list and return its head.

Example 1:

Input: head = [1,2,3,4,5], n = 2
Output: [1,2,3,5]

Example 2:

Input: head = [1], n = 1
Output: []

Code

```
JavaScript
7  /**
8   * @param {ListNode} head
9   * @param {number} n
10  * @return {ListNode}
11  */
12
13  var removeNthFromEnd = function(head, n) {
14      let res = new ListNode(0, head);
15      let dummy = res;
16
17      for (let i = 0; i < n; i++) {
18          head = head.next;
19      }
20
21      while (head) {
22          head = head.next;
23          dummy = dummy.next;
24      }
25
26      dummy.next = dummy.next.next;
27
28      return res.next;
29  };

```

Accepted Runtime: 0 ms

Case 1 Case 2 Case 3

Question: <https://leetcode.com/problems/remove-nth-node-from-end-of-list/description/>

Answer: <https://leetcode.com/problems/powx-n/submissions/1628605732>

Description:

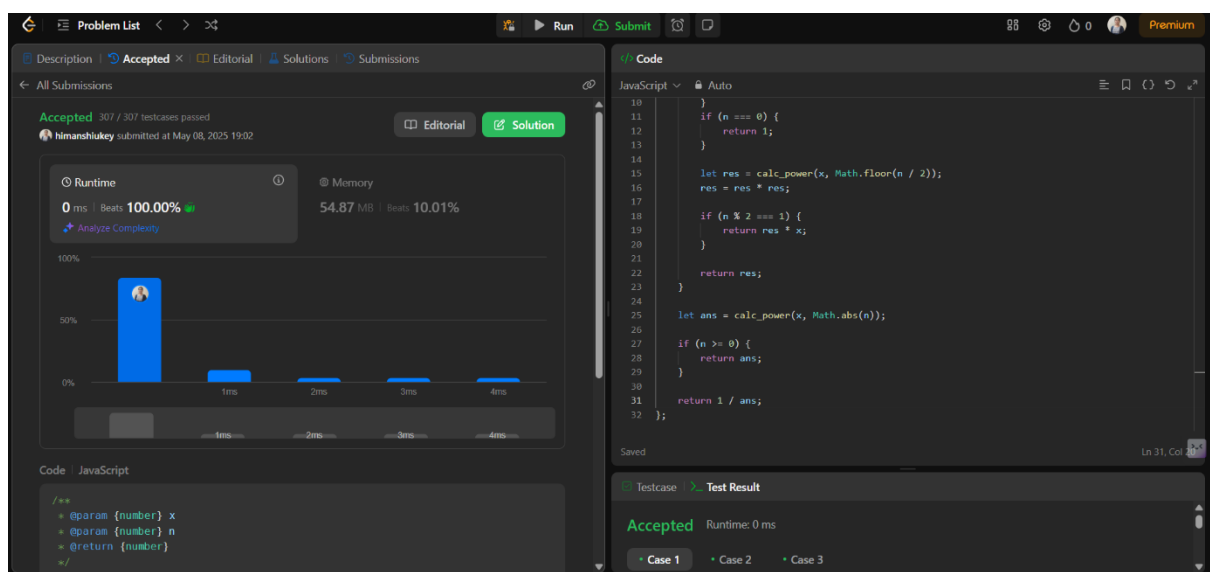
Time Complexity: $O(\log n)$

- The algorithm reduces the exponent n by half in each recursive call (or iteration).
- This is typical of fast exponentiation, where:
- If n is even, it computes $\text{pow}(x * x, n / 2)$
- If n is odd, it computes $x * \text{pow}(x * x, (n - 1) / 2)$
- Hence, the number of steps is proportional to $\log n$.

Space Complexity:

- If implemented recursively: $O(\log n)$ — due to the call stack depth.
- If implemented iteratively: $O(1)$ — only a few variables are used (no recursive stack).

Screenshot:



Question: <https://leetcode.com/problems/powx-n/description/>

Answer: <https://leetcode.com/problems/delete-node-in-a-linked-list/submissions/1628607345>

Description:

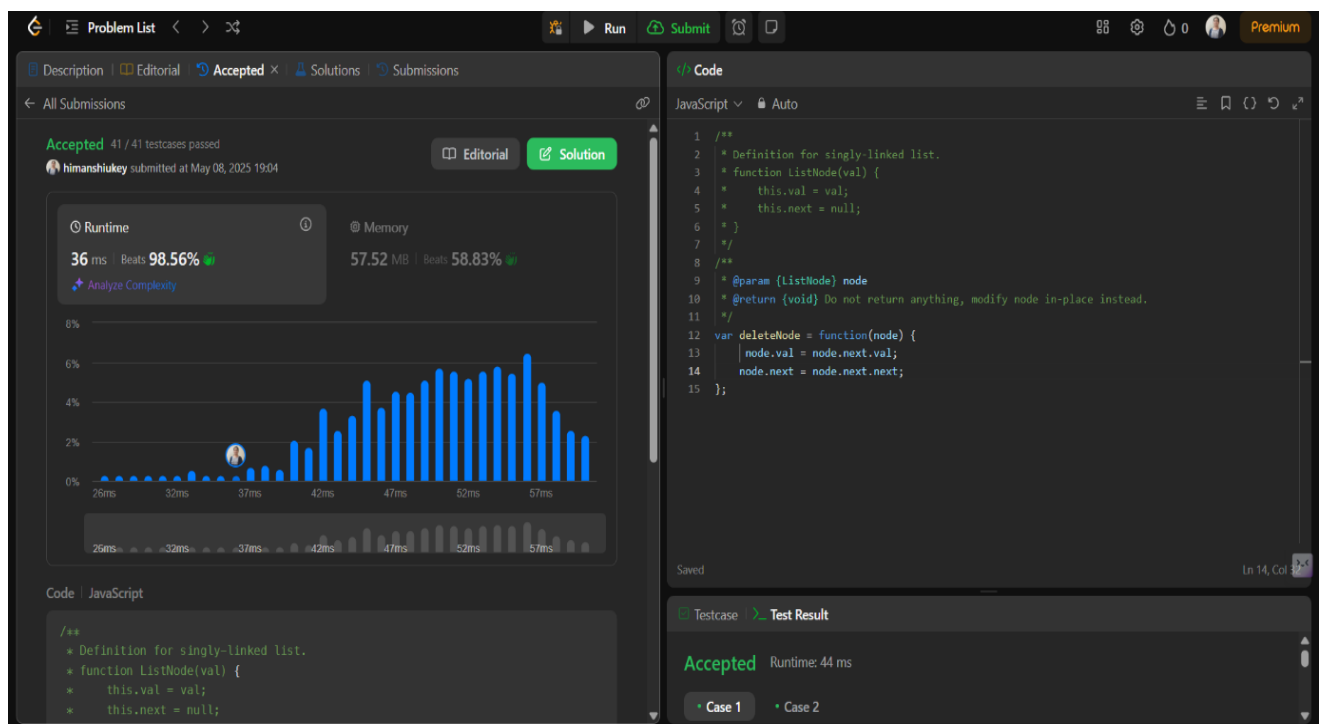
Time Complexity: $O(1)$

- constant time.
- It only accesses and modifies two nodes.

Space Complexity: $O(1)$

- No extra space use

Screenshot:



Thank You For This Assignment