PCB = info about process -pid, state, loaction

Ready queue = queue in which process waits for CPU

IO-Wait-State = Process waits for IO by DMA

Process address space = process space = code, data, stack,heap/extra

Mirco kernel = windows

Kernel space and User space

Kernel mode and User mode

Wait time

Throughput

ISR

**IVT** 

Ex Maskable interrupt = RST interrupts , INTR

Non Maskable interrupt = TRAP

RTOS = LinuxRT, ADA,

Real Time OS = the processes with highest priority are those which has earliest deadline!!

CPU Schedular is the part of the kernel that will schedule a process.

Schedule a process = it will select process from ready and allocate CPU to it for a certain duration

At this time it must consider the factors like Wt, Ta, Throughput, response time!!!

Basic scheduling techniques ---

1. FCFS = FIFO

Which process is selected from the ready Queue = the frontmost process
How long will process get the CPU? = ideally until termination
Practically = till the IO instruction occurs

OR till interrupt occurs

Adv: simple, all processes will get a change to run sooner or later

Disadv: no priority, throughput of system may reduce

Ta of low Tcpu process may be high as it is waiting after the high Tcpu process time

Operating System - Silberschatz, galvin

Process	Arrival Time	Тсри
P1	2	5ms
P2	0	3ms
Р3	1	6ms

Calculate Avg Ta and Avg Wt using FIFO scheduling algo ---

GANTT CHART / TIMELINE

		•		
P2	P3	P1		

03	39	914		

Process	Wt (Start - arrival)	Ta (Wt + Tcpu)
P1	9-2 = 7	7 + 5 = 12
P2	0-0 = 0	0+3 = 3
Р3	3 - 1=2	2 + 6=8

Avg Wt = 
$$(7+0+2)/3 = 3$$
  
Avg Ta =  $(12+3+8)/3 = 7.6667$ 

#### Ex2---

Process	Arrival Time	Tcpu
P1	0	3
P2	1	3
Р3	2	4

P1	P2	Р3
03	36	610

Process	Wt	Та
P1	0-0=0	0+3=3
P2	3-1=2	2 +3=5
Р3	6-2=4	4+4=8

### 2 . SJF = Shortest job first

Which process is selected from the ready Queue = processes with lowest Tcpu is selected How long will process get the CPU? = ideally until termination

Practically = till the IO instruction occurs

OR till interrupt occurs

Adv - Throughput improves, as Ta reduces as small Tcpu process waits for less time Disadv - 1. Practical implementation is not possible (theoretical)

Scanf("%d",&val)

for(i=0;i<val;i++) }}}} Kernel does not know val ? -- kernel cannot say how many times loop will run --can it predict Tcpu before the code runs ?

2. STARVATION = a process with high Tcpu MAY never get a chance to run WHY? As lower Tcpu processes keep coming in the queue

Process	Arrival time	Tcpu
---------	--------------	------

P1	0	6
P2	2	3
Р3	0	4

Calculate avg Wt and avg Ta using SJF scheduling algo .

Ready Queue R<-----P1--->F

# Gantt chart/timeline

Р3	P2	P1
04	47	713

Process	Wt ( start-arrival )	Ta (Wt +Tcpu )
P1	7-0=7	7+6 = 13
P2	4-2=2	2 + 3=5
P2	0-0=0	0+4 = 4
Avg	(7+2+0)/3 = 3	(13+5+4)/3 = 7.3333

Process	Arrival time	Тсри
P1	0	1
P2	0	5
Р3	0	2
P4	4	1

# Calculate avg Wt and avg Ta using Non Preemptive SJF

Ready Queue	R <p3-p2-p1>F</p3-p2-p1>	P1 is selected
	R <p3-p2>F</p3-p2>	P3 us selected
	R <p2>F</p2>	P2 is selected
	R<>F	

P1	Р3	P2	P4
01	13	38	89

Process	Wt	Та
P1	0-0	0+1=1
P2	3-0=3	3+5=8
Р3	1-0=1	1+2=3
P4	8-4=4	4+1=5
Avg	(0+3+1+4)/4=2	(1+8+3+5)/4 = 4.25

# Preemptive and Non Preemptive SJF

To preempt = Apply force to get the resource Who can preempt? Process with higher priority can preempt a process with lower priority

Process	Arrival time	Тсри
P1	0	1
P2	0	5
Р3	0	2
P4	4	1

# Calculate avg Wt and avg Ta using Preemptive SJF

Ready Queue	R <p3-p2-p1>F</p3-p2-p1>	P1 is selected to use the CPU
	R <p3-p2>F</p3-p2>	P3 us selected
	R <p2>F</p2>	P2 is selected
	R<>F	P4 preempts P2
	R <p2(4ms remains)="">F</p2(4ms>	

# P2 remaining Tcpu(4) Incoming P4 Tcpu (1)

P1	Р3	P2	P4 preempts P2	P2 resumes
01	13	34	45	59

Process	Wt (start -arrival) ( resume -preempt)	Та
P1	0-0	0+1=1
P2	3-0=3 5-4=1 (3+1)=4	4+5=9
Р3	1-0=1	1+2=3
P4	<mark>4-4=0</mark>	0+1=1
Avg	(0+4+1+0)/4=1.25	(1+9+3+1)/4 = 3.5

\_\_\_\_\_

Windows OS ---- WSL = Windows Subsystem for Linux = it allows us to use Linux OS through windows !!!

We do not have to install Linux separately for practice

Linux distribution -----

Linux based OS ----- Flavors of Linux 1. Red Hat Linux 2. Ubuntu 3. Fedora Core 4. Suse 5. Mint 6. Mandriva 7. BOSS 8. Debian 9. Kali 10. ... 11. .... 12. ... 13. ... 14. ... 15. .... 16. ... Once the WSL is set up !!!! Go to start ----Ubuntu ----->open = The terminal is opened !!! You see a prompt !!! ======== THIS IS THE Ubuntu CLI (interface between System and end-user) EndUser ------System Once you see the **prompt**, enter See the prompt is repeated on every line many times Type a command clear See that the prompts and cleared and we are at the top This will show the contents of the current Type a command Is folder(directory) Type a command pwd It will show the Present Working Directory = current folder /home/unix = What does it mean ? / = topmost folder ----- sub folder ---- home ----sub folder ----unix Windows topmost folder = many top folders = C: , D: , E: etc But Linux has ONLY one topmost folder = / (this folder is denoted by forward slash) Go to the topmost folder Type command cd cd / Cd = change space forward slash directory (folder) Type pwd and observe current folder name

Ex1

Ex2

Ex3

Ex4

Ex5

Ex6

Ex7

Ex8

Go to the earlier folder

Again type pwd and observe output

cd /home/unix




