

G-07 BFS Problem Type-2 II BFS on Matrix(Implicit Graph)

1091. Shortest Path in Binary Matrix

Medium

4.6K

181



<https://leetcode.com/problems/shortest-path-in-binary-matrix/>

994. Rotting Oranges

Medium

10.1K

340



<https://leetcode.com/problems/rotting-oranges/>

1091. Shortest Path in Binary Matrix

Hint



Medium

4.6K

181



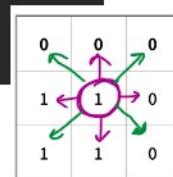
Companies

Given an $n \times n$ binary matrix `grid`, return the length of the shortest **clear path** in the matrix. If there is no clear path, return `-1`.

A **clear path** in a binary matrix is a path from the **top-left** cell (i.e., $(0, 0)$) to the **bottom-right** cell (i.e., $(n - 1, n - 1)$) such that:

- All the visited cells of the path are `0`.
- All the adjacent cells of the path are **8-directionally connected** (i.e., they are different and they share an edge or a corner).

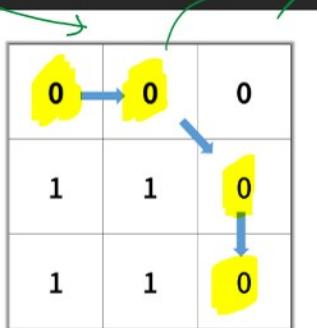
The **length of a clear path** is the number of visited cells of this path.



4 cells
visited
clear path
is 4.

Example 2:

0	0	0
1	1	0
1	1	0



only 0's can be
visited !!

Input: grid = [[0,0,0], [1,1,0], [1,1,0]]
Output: 4

→ explain
→ what algo
BFS → DP
↓
shortest
path...
why not
DP?

Input: grid = [[0,0,0],[1,1,0],[1,1,0]]
Output: 4

Example 3:

Input: grid = [[1,0,0],[1,1,0],[1,1,0]]
Output: -1

0	2	0	0
X	X	0	0
1	1	0	0
1	1	0	0

0	0	3	0
1	X	0	0
1	1	0	0

0	0	0
1	1	0
1	1	0

0	0	0
1	1	0
1	1	0

→ At every unit of time, I am trying to SPREAD & reach to final cell $[n-1][n-1]$

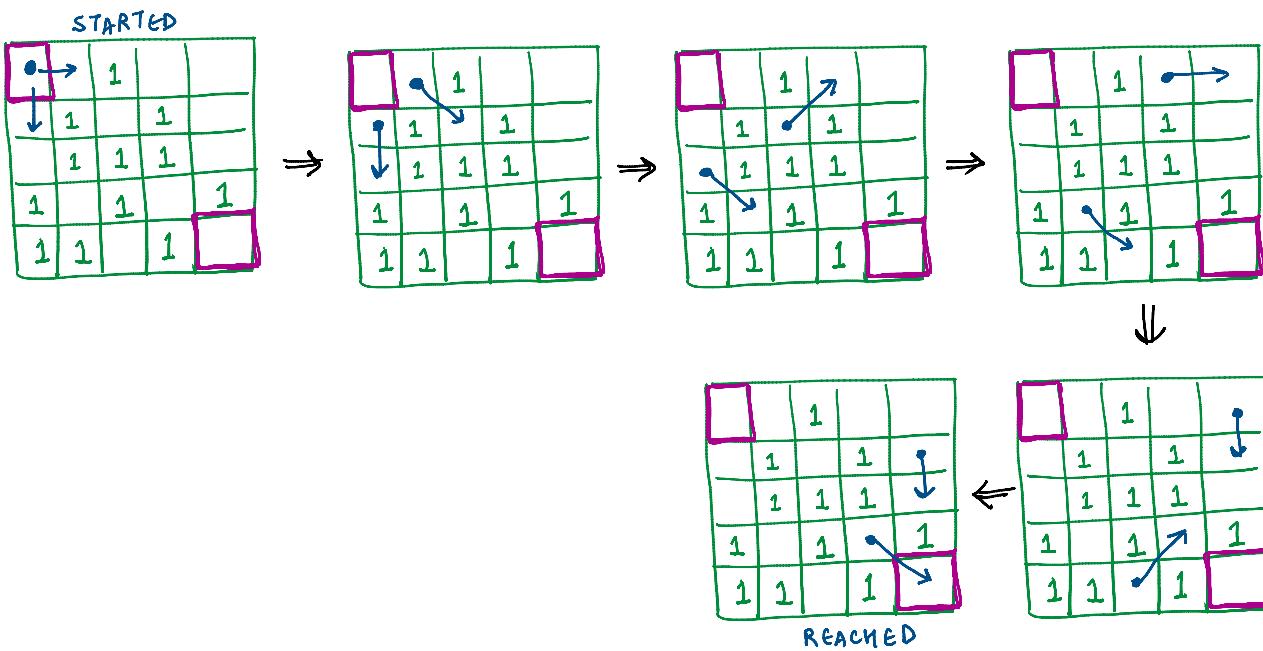
spreading is nothing but going on breadth wise

thus, we can use BFS.

Why??

1	1	1	1
1	1	1	1
1	1	1	1
1	1	1	1
1	1	1	1





"BUT BUT BUT"

Let's confuse you!! You must have seen a same kind of problem in DP also !!

Min Cost Path | DP-6

Difficulty Level : Easy • Last Updated : 17 Feb, 2023

Read Discuss(20+) Courses Practice Video



Given a cost matrix $\text{cost}[]$ and a position (M, N) in $\text{cost}[]$, write a function that returns cost of minimum cost path to reach (M, N) from $(0, 0)$. Each cell of the matrix represents a cost to traverse through that cell. The total cost of a path to reach (M, N) is the sum of all the costs on that path (including both source and destination). You can only traverse down, right and diagonally lower cells from a given cell, i.e., from a given cell (i, j) , cells $(i+1, j)$, $(i, j+1)$, and $(i+1, j+1)$ can be traversed.

Note: You may assume that all costs are positive integers.

Example:

Example:

Input:

1	2	3
-	-	-

only 1 directional i.e. Downwards - Right traversal.
thus DP will fail at these kind of cases ↴

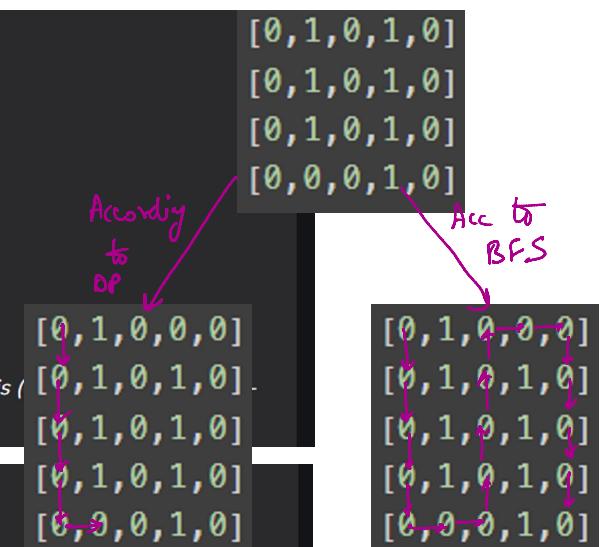
[0, 1, 0, 0, 0]
[0, 1, 0, 1, 0]
[0, 1, 0, 1, 0]

1	2	3
4	8	2
1	5	3

The path with minimum cost is highlighted in the following figure. The path is (2, 2). The cost of the path is 8 (1 + 2 + 2 + 3).

Output:

1	2	3
4	8	2
1	5	3

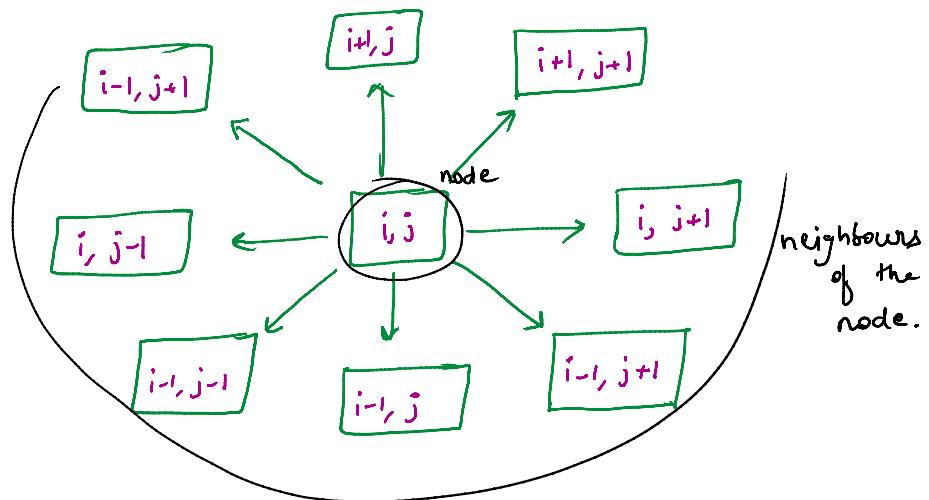


Ans = -1

Ans = 16

Thus, As we can traverse 8-directions, we can't use a DP, we have to use a BFS.

⇒ Now, we know, we traverse Breadth wise thus we are using BFS.
But what is the next level, I have to go on.



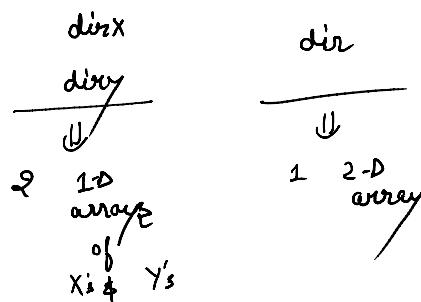
```
int dir[][] = new int[][]{{0,1},{0,-1},{1,0}, {-1,0},{1,-1}, {-1,1}, {-1,-1},{1,1}};
```

To go from node to neighbours we use this dir array

$curr[0] = X$
 $curr[1] = Y$

```
for (int k=0;k<8;k++) {
    int nextX = dir[k][0]+curr[0];
    int nextY = dir[k][1]+curr[1];
```

Neighbour ..



```
class Solution {
public:
    int shortestPathBinaryMatrix(vector<vector<int>>& grid) {
        int m = grid.size(), n = grid[0].size();
        if(grid[0][0]==1 || grid[m-1][n-1]==1) return -1;

        int dir[][2] = {{0,1},{0,-1},{1,0}, {-1,0},{1,-1}, {-1,1}, {-1,-1},{1,1}};
        vector<vector<bool>> visited(m, vector<bool>(n, false));
        visited[0][0] = true;
        queue<vector<int>> q;
        q.push({0,0});

        int ans=0;
        while (!q.empty()) {
            int size = q.size();
            for(int i=0;i<size;i++) {
                auto curr = q.front();
                q.pop();
                if(curr[0]==m-1 && curr[1]==n-1)
                    return ans+1;
            }
            for (int k=0;k<8;k++) {
                int nextX = dir[k][0]+curr[0];
                int nextY = dir[k][1]+curr[1];
                if(nextX<0 && nextX>m && nextY<0 && nextY>n && !visited[nextX][nextY] && grid[nextX][nextY]==0) {
                    q.push({nextX,nextY});
                    visited[nextX][nextY]=true;
                }
            }
            ans++;
        }
        return -1;
    }
};
```

If still it
can't search

Standard
BFS

Check if that is not visited
if 0 Valid location

In bounds of Matrix
should be a value '0'.

currently filled
neighbours coming

Analyze all the nodes at one level, as answer will increase by 1 for these nodes only
when reached end

Going to neighbours

Standard
BFS

dir array to get to the neighbours.

visited to actually visit a cell only once, as if it is visited then it would be the first & shortest path.

If still it
can't reach
cell [n-1][n-1]
then return -1.

time: $O(N^2)$ → BFS in worst case has
to visit all the cells in
grid.

Space: $O(N^2)$
visited Array / Queue

MODIFICATIONS TO MAKE CODE EFFICIENT :-

```
class Solution {
public:
    int shortestPathBinaryMatrix(vector<vector<int>>& grid) {
        int m = grid.size(), n = grid[0].size();
        if(grid[0][0]==1 || grid[m-1][n-1]==1) return -1;

        int dir[8][2] = {{0,1},{0,-1},{1,0},{-1,0},{1,-1},{-1,1},{-1,-1},{1,1}};
        vector<vector<bool>> visited(m, vector<bool>(n, false));
        visited[0][0] = true;
        queue<vector<int>> q;
        q.push({0,0});

        int ans=0;
        while (!q.empty()) {
            int size = q.size();
            for(int i=0;i<size;i++) {
                auto curr = q.front();
                q.pop();
                if(curr[0]==m-1 && curr[1]==n-1) {
                    return ans+1;
                }
                for (int k=0;k<8;k++) {
                    int nextX = dir[k][0]+curr[0];
                    int nextY = dir[k][1]+curr[1];

                    if(nextX>=0 && nextX<m && nextY>=0 && nextY<n && !visited[nextX][nextY] && grid[nextX][nextY]==0) {
                        q.push({nextX,nextY});
                        visited[nextX][nextY]=true;
                    }
                }
            }
            ans++;
        }
    }
}
```

visited array not required, we can modify grid matrix itself.
Not recommended in production code.

Pair is much faster than a vector

```
        }
    }
    ans++;
}
return -1;
};
```

```
i C++ ▾ Auto

1 class Solution {
2 public:
3     int shortestPathBinaryMatrix(vector<vector<int>>& grid) {
4         int m = grid.size(), n = grid[0].size();
5         if(grid[0][0]==1 || grid[m-1][n-1]==1) return -1;
6
7         int dir[8][2] = {{0,1},{0,-1},{1,0}, {-1,0},{1,-1},{-1,1},{-1,-1},{1,1}};
8         vector<vector<bool>> visited(m, vector<bool>(n, false));
9         visited[0][0] = true;
10        queue<vector<int>> q;
11        q.push({0,0});
12
13        int ans=0;
14        while (!q.empty()) {
15            int size = q.size();
16            for(int i=0;i<size;i++) {
17                auto curr = q.front();
18                q.pop();
19                if(curr[0]==m-1 && curr[1]==n-1) {
20                    return ans+1;
21                }
22                for (int k=0;k<8;k++) {
23                    int nextX = dir[k][0]+curr[0];
24                    int nextY = dir[k][1]+curr[1];
25
26                    if(nextX<=0 && nextX>m && nextY<=0 && nextY>n && !visited[nextX][nextY] && grid[nextX][nextY]==0) {
27                        q.push({nextX,nextY});
28                        visited[nextX][nextY]=true;
29                    }
30                }
31            }
32            ans++;
33        }
34        return -1;
35    }
36 };
37
```

```

class Solution {
public:
    int shortestPathBinaryMatrix(vector<vector<int>>& grid) {
        int m = grid.size(), n = grid[0].size();
        if(grid[0][0]==1 || grid[m-1][n-1]==1) return -1;
        int dir[][] = {{0,1},{0,-1},{1,0},{-1,0},{1,-1},{-1,1},{-1,-1},{1,1}};
        vector<vector<bool>> visited(m, vector<bool>(n, false));
        visited[0][0] = true;
        queue<vector<int>> q;
        q.push({0,0});
        int ans=0;
        while (!q.empty()) {
            int size = q.size();
            for(int i=0;i<size;i++) {
                auto curr = q.front();
                q.pop();
                if(curr[0]==m-1 && curr[1]==n-1) {
                    return ans+1;
                }
                for(int j=0;j<8;j++) {
                    int x = curr[0]+dir[j][0];
                    int y = curr[1]+dir[j][1];
                    if(x<0 || x==m || y<0 || y==n || visited[x][y]) continue;
                    visited[x][y] = true;
                    q.push({x,y});
                }
            }
            ans++;
        }
        return -1;
    }
};

```

```

    }
    for (int k=0;k<8;k++) {
        int nextX = dir[k][0]+curr[0];
        int nextY = dir[k][1]+curr[1];
        if(nextX>=0 && nextX<m && nextY>=0 && nextY<n && !visited[nextX][nextY] && grid[nextX][nextY]==0) {
            q.push({nextX,nextY});
            visited[nextX][nextY]=true;
        }
    }
    ans++;
}
return -1;
};

}

```

The screenshot shows a Java code editor with the following code:

```

1 class Solution {
2     public int shortestPathBinaryMatrix(int[][] grid) {
3         int m = grid.length, n = grid[0].length;
4         if(grid[0][0]==1 || grid[m-1][n-1]==1) return -1;
5
6         int dir[][] = new int[][]{{0,1},{0,-1},{1,0},{-1,0},{1,-1},{-1,1},{-1,-1},{1,1}};
7         boolean[][] visited = new boolean[m][n];
8         visited[0][0] = true;
9         Queue<int[]> queue = new LinkedList<>();
10        queue.add(new int[]{0,0});
11
12        int ans=0;
13        while (!queue.isEmpty()) {
14            int size = queue.size();
15            for(int i=0;i<size;i++) {
16                int[] curr = queue.remove();
17                if(curr[0]==m-1 && curr[1]==n-1) {
18                    return ans+1;
19                }
20                for (int k=0;k<8;k++) {
21                    int nextX = dir[k][0]+curr[0];
22                    int nextY = dir[k][1]+curr[1];
23
24                    if(nextX>=0 && nextX<m && nextY>=0 && nextY<n && !visited[nextX][nextY] && grid[nextX][nextY]==0) {
25                        queue.add(new int[]{nextX,nextY});
26                        visited[nextX][nextY]=true;
27                    }
28                }
29            }
30            ans++;
31        }
32        return -1;
33    }
34 }

```

```

class Solution {
    public int shortestPathBinaryMatrix(int[][] grid) {
        int m = grid.length, n = grid[0].length;
        if(grid[0][0]==1 || grid[m-1][n-1]==1) return -1;
        int dir[][] = new int[][]{{0,1},{0,-1},{1,0},{-1,0},{1,-1},{-1,1},{-1,-1},{1,1}};
        boolean[][] visited = new boolean[m][n];
        visited[0][0] = true;
        Queue<int[]> queue = new LinkedList<>();
        queue.add(new int[]{0,0});

```

```
int ans=0;
while (!queue.isEmpty()) {
    int size = queue.size();
    for(int i=0;i<size;i++) {
        int[] curr = queue.remove();
        if(curr[0]==m-1 && curr[1]==n-1) {
            return ans+1;
        }
        for (int k=0;k<8;k++) {
            int nextX = dir[k][0]+curr[0];
            int nextY = dir[k][1]+curr[1];
            if(nextX>=0 && nextX<m && nextY>=0 && nextY<n && !visited[nextX][nextY] && grid[nextX][nextY]==0) {
                queue.add(new int[]{nextX,nextY});
                visited[nextX][nextY]=true;
            }
        }
        ans++;
    }
    return -1;
}
}
```