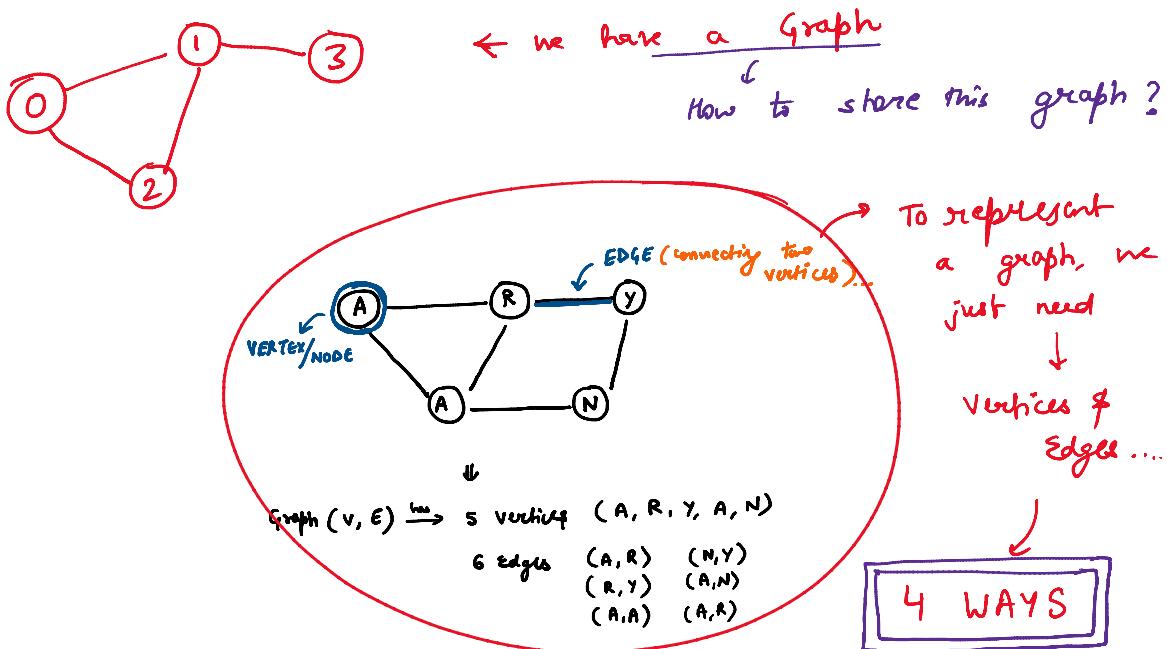


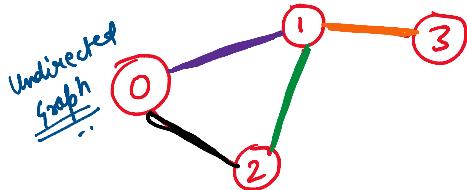
# G-03 Graph Representation || Adjacency Matrix || Adjacency List

|| C++/Java

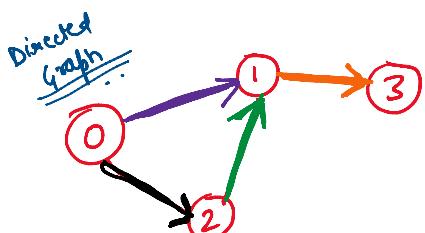
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Method 1: Create a 2-D Array (Adjacency Matrix)



	0	1	2	3
0		1	1	
1	1		1	1
2	1	1		
3		1		



	0	1	2	3
0		1	1	
1				1
2		1		
3				

### Advantages

To know if we have an edge b/w node  $U \neq V$

We can check in  $O(1)$

$\text{matrix}[U][V] == 1 \Rightarrow \text{edge present}$   
 $\neq 1 \Rightarrow \text{edge not present}$

### Dis-advantages

① To store all edges

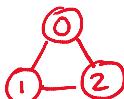
$O(n^2)$  memory

(Even when  $n$  nodes are there)

② To find neighbours of a graph.

We have to iterate over entire row  $i \rightarrow O(n)$  for 1 node

Even if neighbours was just 2 or 3.



No. of Nodes → 3  
No. of Edges → 3  
 $0-1$   
 $1-2$   
 $0-2$

```
#include<bits/stdc++.h>
using namespace std;

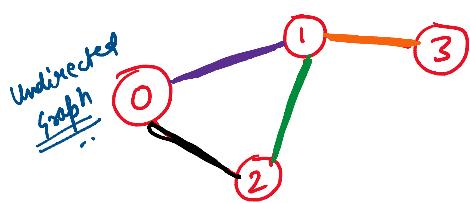
int main()
{
    n = 3, m = 3;
    cin >> n >> m;
    // Graph Nodes are from 0 to n-1
    int gr[n][n];
    for (int i = 0; i < m; i++) {
        int u, v;
        cin >> u >> v;
        gr[u][v] = 1;
        gr[v][u] = 1;
    }
}
```

Just mark the cell  $u-v$  &  $v-u$  to 1

```
import java.io.*;
public class AryanClass {
    public static void main (String[] args) {
        int n = 3, m = 3;
        int gr[][] = new int[n][n];
        // edge 0---1
        gr[0][1] = 1;
        gr[1][0] = 1;
        // Make a 2-D array.
        // edge 1---2
        gr[1][2] = 1;
        gr[2][1] = 1;
        // edge 0---2
        gr[0][2] = 1;
        gr[2][0] = 1;
    }
}
```

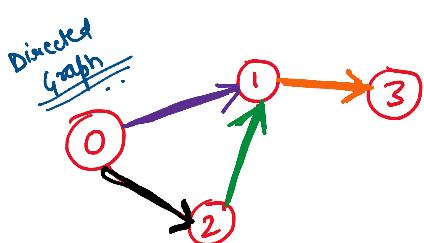
Method 2: Represent via list of Edges & Vertices  $\rightarrow$  (Edge List Representation)

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Vertices = { 0, 1, 2, 3 }

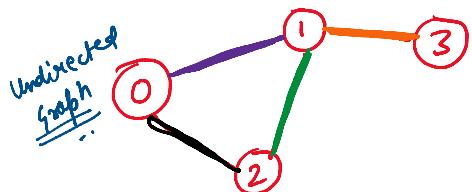
Edges = { (0,1) (1,3) (2,1) (0,2)  
(1,0) (3,1) (1,2) (2,0) }



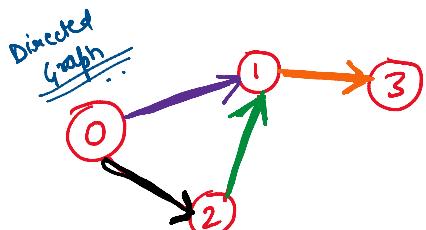
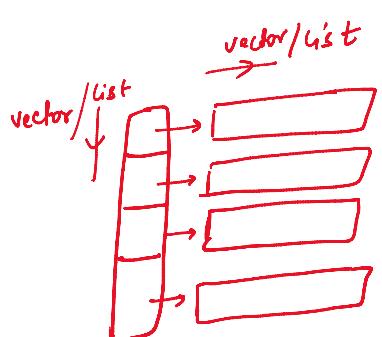
Vertices = { 0, 1, 2, 3 }

Edges = { (0,1) (1,3) (2,1) (0,2) }

\*\*  
Method 4: Adjacency list (MOSTLY USED)



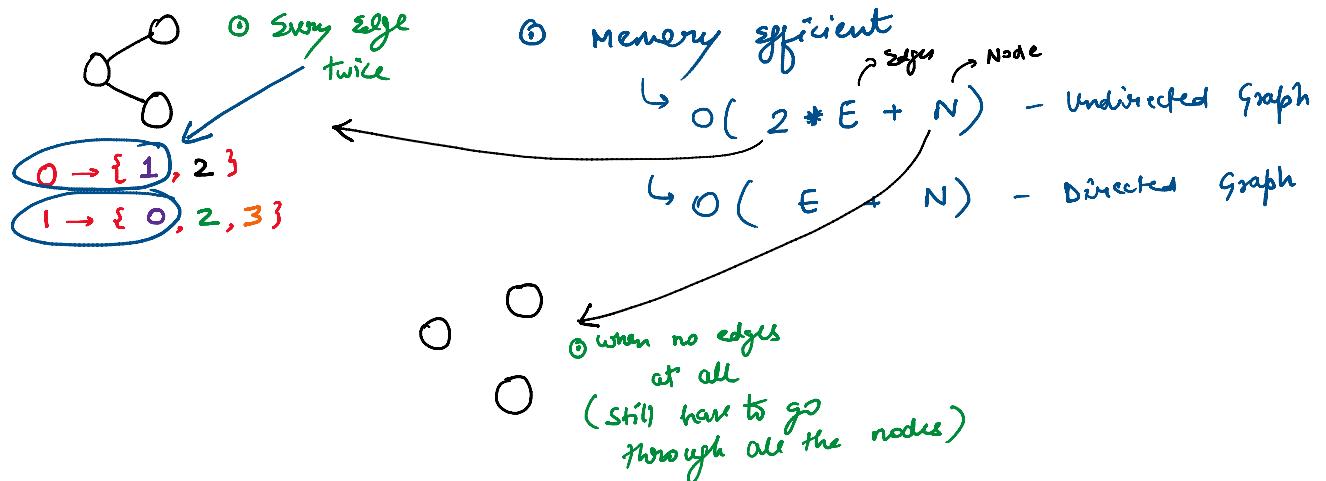
0  $\rightarrow$  { 1, 2 }  
1  $\rightarrow$  { 0, 2, 3 }  
2  $\rightarrow$  { 0, 1 }  
3  $\rightarrow$  { 1 }



0  $\rightarrow$  { 1, 2 }  
1  $\rightarrow$  { 3 }  
2  $\rightarrow$  { 1 }  
3  $\rightarrow$  { }

## Advantages

④ Find neighbours in  $O(\text{neighbours})$  time

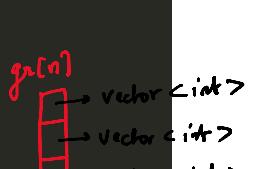


```
#include<bits/stdc++.h>
using namespace std;

int main()
{
    int n, m;
    cin >> n >> m;

    // Graph Nodes are from 0 to n-1
    vector<int> gr[n];
    for (int i = 0; i < m; i++) {
        int u, v;
        cin >> u >> v;

        gr[u].push_back(v);
        gr[v].push_back(u);
    }
}
```



Undirected graph

```
import java.io.*;
import java.util.*;

class AryanClass2 {
    public static void main (String[] args) {
        int n = 3, m = 3;
        ArrayList<ArrayList<Integer>> gr = new ArrayList<>();

        for (int i = 0; i <= n; i++)
            gr.add(new ArrayList<Integer>());

        // edge 0---1
        gr.get(index:0).add(e:1);
        gr.get(index:1).add(e:0);

        // edge 0---2
        gr.get(index:0).add(e:2);
        gr.get(index:2).add(e:0);

        // edge 1---2
        gr.get(index:1).add(e:2);
        gr.get(index:2).add(e:1);
    }
}
```

## Method - 4: Implicit Graph

Given a 2-D array  
Matrix  
↓  
Find connected component/  
groups in the  
graph.

0	0	0	0
1	1	1	0
1	0	0	0
0	0	1	1
0	0	1	0



You don't need to  
convert this to graph  
↓  
It is already a graph  
where cell has neighbours

neighbour:

