

COMPILER DESIGN PROJECT REPORT - 3

SEMANTIC CHECKER FOR C PROGRAMMING LANGUAGE



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INTRODUCTION

A parser constructs parse trees in the syntax analysis phase. The plain parse-tree constructed in that phase is generally of no use for a compiler, as it does not carry any information of how to evaluate the tree. The productions of context-free grammar, which makes the rules of the language, do not accommodate how to interpret them. For example,

$$E \rightarrow E + T$$

The above CFG production has no semantic rule associated with it, and it cannot help in making any sense of the production. Parsing only verifies that the program consists of tokens arranged in a syntactically valid combination.

SEMANTIC ANALYSIS

Semantics of a language provide meaning to its constructs, like tokens and syntax structure. Semantics help interpret symbols, their types, and their relations with each other. Semantic analysis judges whether the syntax structure constructed in the source program derives any meaning or not.

CFG + semantic rules = Syntax Directed Definitions

For example:

int a = "value";

should not issue an error in lexical and syntax analysis phase, as it is lexically and structurally correct, but it should generate a semantic error as the type of the assignment differs. These rules are set by the grammar of the language and evaluated in semantic analysis.

Semantic analysis is the task of ensuring that the declarations and statements of a program are semantically correct, i.e., that their meaning is clear and consistent with the way in which control structures and data types are supposed to be used.

Semantic analysis is not a separate module within a compiler. It is usually a collection of procedures called at appropriate times by the parser as the grammar requires it. Implementing the

semantic actions is conceptually simpler in recursive descent parsing because they are simply added to the recursive procedures. Implementing the semantic actions in a table - action driven LL(1) parser requires the addition of a third type of variable to the productions and the necessary software routines to process it.

Part of semantic analysis is producing some sort of representation of the program, either object code or an intermediate representation of the program. One - pass compilers will generate object code without using an intermediate representation; code generation is part of the semantic actions performed during parsing. Other compilers will produce an intermediate representation during semantic analysis; most often it will be an abstract syntax tree or quadruples.

Semantic analysis typically involves:

- **Type checking** - Data types are used in a manner that is consistent with their definition (i. e., only with compatible data types, only with operations that are defined for them, etc.)
- **Label Checking** - Labels references in a program must exist.
- **Flow control checks** - Control structures must be used in their proper fashion (no GOTOs into a FORTRAN DO statement, no breaks outside a loop or switch statement, etc.)
- **Array-bound Checking** - Variables being used as an array index should be within the bounds of the array.
- **Scope Resolution** - We need to determine what identifiers are accessible at different points in the program.

Abstract syntax trees have one enormous advantage over other intermediate representations: they can be “ decorated ” , i.e., each node on the AST can have their attributes saved in the AST nodes, which can simplify the task of type checking as the parsing process continues.

ATTRIBUTE GRAMMAR

An attribute is a property whose value is assigned to a grammar symbol. Attribute computation functions (or semantic functions) are associated with the productions of a grammar and are used

to compute the values of an attribute.

An attribute grammar is an extension to a context - free grammar that is used to describe features of a programming language that cannot be described in BNF or can only be described in BNF with great difficulty. Each attribute has well-defined domain of values, such as integer, float, character, string, and expressions. Attribute grammar is a medium to provide semantics to the context-free grammar and it can help specify the syntax and semantics of a programming language. Attribute grammar (when viewed as a parse-tree) can pass values or information among the nodes of a tree.

$$E \rightarrow E + T \{ E.value = E.value + T.value \}$$

The right part of the CFG contains the semantic rules that specify how the grammar should be interpreted. Here, the values of non-terminals E and T are added together and the result is copied to the non-terminal E.

Semantic attributes may be assigned to their values from their domain at the time of parsing and evaluated at the time of assignment or conditions. Based on the way the attributes get their values, they can be broadly divided into two categories :

- **Synthesized attributes** - These attributes get values from the attribute values of their child nodes. Synthesized attributes never take values from their parent nodes or any sibling nodes. For example,

$$S \rightarrow ABC$$

If S is taking values from its child nodes (A,B,C), then it is said to be a synthesized attribute, as the values of ABC are synthesized to S.

- **Inherited attributes** - In contrast to synthesized attributes, inherited attributes can take values from parent and/or siblings. For example,

$$S \rightarrow ABC$$

A can get values from S, B and C. B can take values from S, A, and C. Likewise, C can take values from S, A, and B.

SYNTAX DIRECTED TRANSLATION

Parser uses a CFG(Context-free-Grammar) to validate the input string and produce output for next phase of the compiler. Output could be either a parse tree or abstract syntax tree. Now to interleave semantic analysis with syntax analysis phase of the compiler, we use Syntax Directed Translation.

Syntax Directed Translation are augmented rules to the grammar that facilitate semantic analysis. SDT involves passing information bottom-up and/or top-down the parse tree in form of attributes attached to the nodes. Syntax directed translation rules use 1) lexical values of nodes, 2) constants, 3) attributes associated to the non-terminals in their definitions.

For example,

$E \rightarrow E+T \quad \{ E.val = E.val + T.val \} \quad PR\#1$

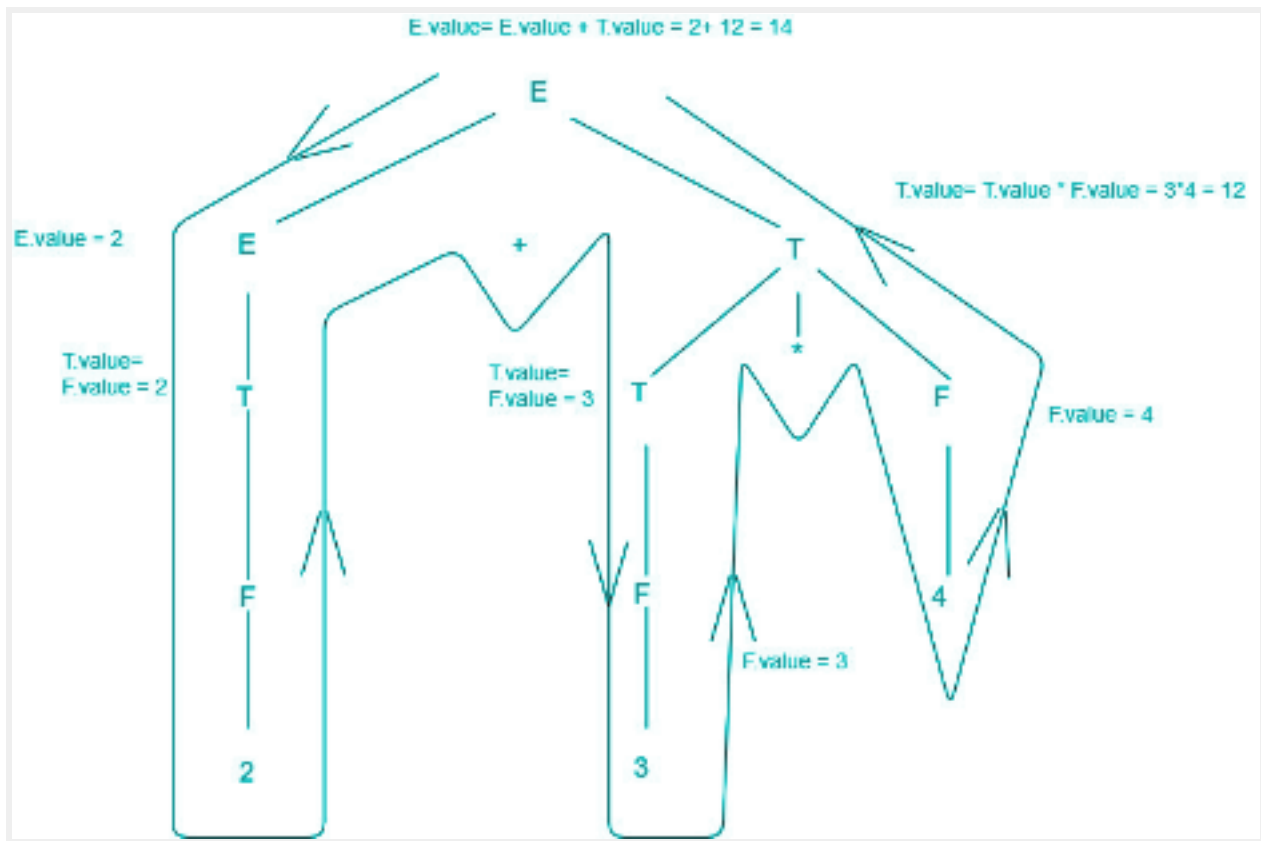
$E \rightarrow T \quad \{ E.val = T.val \} \quad PR\#2$

$T \rightarrow T * F \quad \{ T.val = T.val * F.val \} \quad PR\#3$

$T \rightarrow F \quad \{ T.val = F.val \} \quad PR\#4$

$F \rightarrow INTLIT \quad \{ F.val = INTLIT.lexval \} \quad PR\#5$

Suppose we take the first SDT augmented to [$E \rightarrow E+T$] production rule. The translation rule in consideration has val as attribute for both the non-terminals – E & T. Right hand side of the translation rule corresponds to attribute values of right side nodes of the production rule and vice-versa.



Above diagram shows how semantic analysis could happen. The flow of information happens bottom-up and all the children attributes are computed before parents.

- **S-attributed SDT** - If an SDT uses only synthesized attributes, it is called as S-attributed SDT. These attributes are evaluated using S-attributed SDTs that have their semantic actions written after the production (right hand side). Attributes in S-attributed SDTs are evaluated in bottom-up parsing, as the values of the parent nodes depend upon the values of the child nodes.
- **L-attributed SDT** - This form of SDT uses both synthesized and inherited attributes with restriction of not taking values from right siblings. In L-attributed SDTs, a non-terminal can get values from its parent, child, and sibling nodes. Attributes in L-attributed SDTs are evaluated by depth-first and left-to-right parsing manner.

IMPLEMENTATION

Scope

Implemented using stack. Variable called scope is used which is initially set to 0. Increment and push to stack when opening brace is found. On finding closing brace, pop out the last element from stack. This will assign unique value to each block of code. Stack as a whole contains the scope-id, which is unique. Incase of functions, the passed parameters are also part of the function block, therefore increment the scope variable when you encounter '(' .

Declaration Checking

To check whether an identifier is declared or not, check the symbol table corresponding to the value having scope-id equal to the current stack (which gives scope-id of the matched identifier). If not present, then undeclared.

Type Checking

When you encounter the type, use a global variable called flag which will indicate the type. For example, flag = 1 for int. When you encounter an assignment statement the current flag and the stored type value of the identifier has to match which is found out by consulting the symbol table.

Re-declaration Checking

When inserting a variable, check if the variable with same name exists with same scope-id. This is taken care by hashing the variable name along with the scope-id. Redecclaration is checked every time we insert identifier into the symbol table.

Array Dimension

When you encounter array declaration, pass the array dimension number into the symbol table. This makes sure that it's only positive integer integer.

Symbol Table

It contains fields:

1. Name : identifier name.
2. Token : token is constant or identifier.
3. Type : type of identifier whether int, float, void etc.

4. Scope : Scope could be global, function.
5. Scope-id : Unique value given to each block of code.
6. Function-id: Unique for each function declaration

All the semantic actions are written in **semantic.h** file, including maintaining symbol table .Necessary functions are called from **syntaxChecker.y** and **lexicalAnalyzer.l** for performing semantic actions and to report semantic error if any. Source code for all the files are included below.

SOURCE CODE

Lexical Analyzer

This is the lex program that contains regular expressions and returns whatever tokens are required.

```

1 %{
2 #include <stdio.h>
3 #include "y.tab.h"
4 int lineCount=1;
5 int nestedCommentCount=0;
6 int commentFlag=0;
7 char *tablePtr;
8 void addToken(char*);
9 %}
10
11 digit      [0-9]
12 letter     [a-zA-Z_]
13 hex        [a-fA-F0-9]
14 E          [Ee][+]?[digit]+
15 FS         (f|F|l|letter)
16 IS         (u|U|l|letter)*
17
18 singlelineComment (\V\/*)
19 multilineCommentStart (\V/*)
20 multilineCommentEnd (\V\*/)
21 %x DETECT_COMMENT
22
23 %%
24
25 {singlelineComment}      { lineCount++; }
26 {multilineCommentStart} { BEGIN(DETECT_COMMENT);
27                          nestedCommentCount++;
28                          }
29
30
31
32 <DETECT_COMMENT>{multilineCommentStart} { nestedCommentCount++;
33                                          if(nestedCommentCount>1)
34                                              commentFlag = 1;
35                                          }
36
37 <DETECT_COMMENT>{multilineCommentEnd} { BEGIN(INITIAL); lineCount++;
38                                          if(nestedCommentCount>0)
39                                              nestedCommentCount--;
40                                          if(nestedCommentCount==0)
41                                              BEGIN(INITIAL);
42                                          }
43
44 <DETECT_COMMENT>\n      {lineCount++;}
45 <DETECT_COMMENT>.      {}

```

FIG. 1

```

46
47
48
49
50 "auto"           { return(AUTO); }
51 "break"          { return(BREAK); }
52 "case"           { return(CASE); }
53 "char"           { return(CHAR); }
54 "const"          { return(CONST); }
55 "continue"       { return(CONTINUE); }
56 "default"        { return(DEFAULT); }
57 "do"             { return(DO); }
58 "double"         { return(DOUBLE); }
59 "else"           { return(ELSE); }
60 "enum"           { return(ENUM); }
61 "extern"         { return(EXTERN); }
62 "float"          { return(FLOAT); }
63 "for"            { return(FOR); }
64 "goto"           { return(GOTO); }
65 "if"             { return(IF); }
66 "int"            { return(INT); }
67 "long"           { return(LONG); }
68 "register"       { return(REGISTER); }
69 "return"         { return(RETURN); }
70 "short"          { return(SHORT); }
71 "signed"         { return(SIGNED); }
72 "sizeof"         { return(SIZEOF); }
73 "static"         { return(STATIC); }
74 "struct"         { return(STRUCT); }
75 "switch"         { return(SWITCH); }
76 "typedef"        { return(TYPDEF); }
77 "union"          { return(UNION); }
78 "unsigned"       { return(UNSIGNED); }
79 "void"           { return(VOID); }
80 "volatile"       { return(VOLATILE); }
81 "while"          { return(WHILE); }
82
83 {letter}({letter}){digit}*      { addToken(yytext); return(IDENTIFIER); }
84 {letter}?\"{\\.|\"{^\\|\"}*\"      { addToken(yytext); return(STRING_LITERAL); }
85
86 0[xX](hex)+{IS}?               { addToken(yytext); return(CONSTANT); }
87 0{digit}+{IS}?                 { addToken(yytext); return(CONSTANT); }
88 {hex}-{IS}?                    { addToken(yytext); return(CONSTANT); }
89 {letter}?(\"{\\.|\"{^\\|\"})*'      { addToken(yytext); return(CONSTANT); }
90 {digit}{+|-|*|/}%{FS}?         { addToken(yytext); return(CONSTANT); }

```

FIG. 2

```

91 {digit}*"."{digit}*{E})?(FS)? { addToken(yytext); return(CONSTANT); }
92 {digit}+ "."{digit}*{E})?(FS)? { addToken(yytext); return(CONSTANT); }
93
94 "..." { return(ELLIPSIS); }
95 ">=" { return(RIGHT_ASSIGN); }
96 "<=" { return(LEFT_ASSIGN); }
97 "+=" { return(ADD_ASSIGN); }
98 "-=" { return(SUB_ASSIGN); }
99 "*=" { return(MUL_ASSIGN); }
100 "/=" { return(DIV_ASSIGN); }
101 "%=" { return(MOD_ASSIGN); }
102 "&=" { return(AND_ASSIGN); }
103 "^=" { return(XOR_ASSIGN); }
104 "|=" { return(OR_ASSIGN); }
105 ">>" { return(RIGHT_OP); }
106 "<<" { return(LEFT_OP); }
107 "++" { return(INC_OP); }
108 "--" { return(DEC_OP); }
109 "->" { return(PTR_OP); }
110 "&&" { return(AND_OP); }
111 "||" { return(OR_OP); }
112 "<=" { return(LE_OP); }
113 ">=" { return(GE_OP); }
114 "==" { return(EQ_OP); }
115 "!=" { return(NE_OP); }
116 ";" { return(';'); }
117 "{"|"<%" { return('{'); }
118 "("|"%">" { return('('); }
119 ":" { return(':'); }
120 "." { return('.'); }
121 "=" { return('='); }
122 "(" { return('('); }
123 ")" { return(')'); }
124 "["|"|"<:" { return('['); }
125 "]"|"|">:" { return(']'); }
126 "&" { return('&'); }
127 "&" { return('&'); }
128 "!" { return('!'); }
129 "~" { return('~'); }
130 "-" { return('-'); }
131 "+" { return('+'); }
132 "*" { return('*'); }
133 "/" { return('/'); }
134 "%" { return('%'); }
135 "<" { return('<'); }

```

FIG. 3

```

136 ">"          { return('>'); }
137 "A"          { return('^'); }
138 "|"          { return('|'); }
139 "?"          { return('?'); }
140
141
142 "#include"(...) { lineCount++; }
143 "#define"(...) { lineCount++; }
144
145 [ ]          {}
146 [\t\v\f]    {}
147 [\n]        {lineCount++;}
148 .           {}
149
150 %%
151 yywrap()
152 {
153     return(1);
154 }
155
156 void addToken(char *yytext)
157 {
158     int len = strlen(yytext);
159     tablePtr = (char*)malloc((len+1)*sizeof(char));
160     strcpy(tablePtr, yytext);
161 }
162
163
164
165
166
167
168
169
170
171
172
173
174
175
176
177
178
179
180

```

FIG. 4

Syntax Analyzer

This is the main parser code, a yacc file which contains the declarations, rules and programs and defines the actions which should be taken for various cases.

```

1 %{
2 #include <stdio.h>
3 #include <string.h>
4 #include <stdlib.h>
5 #include "y.tab.h"
6 #include "semantic.h"
7
8 extern FILE *yyin;
9 extern int lineCount;
10 extern char *tablePtr;
11 extern char *tablePtr;
12 extern int nestedCommentCount;
13 extern int commentFlag;
14 extern int arrayIndexErr;
15
16 char *sourceCode=NULL;
17 int errorFlag=0;
18 void makeList(char *,char,int);
19 %}
20
21 %token AUTO BREAK CASE CHAR CONST CONTINUE DEFAULT DO DOUBLE ELSE ENUM
22 %token EXTERN FLOAT FOR GOTO IF INT LONG REGISTER RETURN SHORT SIGNED
23
24 %token SIZEOF STATIC STRUCT SWITCH TYPEDEF UNION UNSIGNED VOID VOLATILE WHILE
25
26 %token IDENTIFIER CONSTANT FLCONSTANT STRING_LITERAL
27
28 %token ELLIPSIS
29
30 %token PTR OP INC_OP DEC_OP LEFT_OP RIGHT_OP LE_OP GE_OP EQ_OP NE_OP
31 %token AND_OP OR_OP MUL_ASSIGN DIV_ASSIGN MOD_ASSIGN ADD_ASSIGN
32 %token SUB_ASSIGN LEFT_ASSIGN RIGHT_ASSIGN AND_ASSIGN
33 %token XOR_ASSIGN OR_ASSIGN TYPE_NAME
34
35 %nonassoc LOWER_THAN_ELSE
36 %nonassoc ELSE
37 %start translation_unit
38 %%
39
40 primary_expression
41 : IDENTIFIER { $$=checkScope(tablePtr,lineCount); }
42 | FLCONSTANT { tempCheckType=4; }
43 | CONSTANT { addConstant(tablePtr, lineCount); }
44 | STRING_LITERAL { makeList(tablePtr, 's', lineCount); }
45 | '(' expression ')' { makeList('(', 'p', lineCount); makeList(")", 'p', lineCount); $$=$2; }

```

FIG. 1

```

46 postfix_expression
47 : primary_expression { $$=$1; }
48 | postfix_expression '[' expression ']' { makeList("[", 'p', lineCount); makeList("]", 'p', lineCount); }
49 | postfix_expression '(' 'p' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
50 | postfix_expression '(' argument_expression_list ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
51 | postfix_expression 'IDENTIFIER' { makeList(tablePtr, 'v', lineCount); }
52 | postfix_expression PTR_OP IDENTIFIER { makeList(tablePtr, 'v', lineCount); }
53 | postfix_expression INC_OP { makeList(tablePtr, 'o', lineCount); }
54 | postfix_expression DEC_OP { makeList(tablePtr, 'o', lineCount); }
55
56
57 argument_expression_list
58 : assignment_expression { $$=$1; }
59 | argument_expression_list ',' assignment_expression { makeList(",", 'p', lineCount); }
60
61 unary_expression
62 : postfix_expression { $$=$1; }
63 | INC_OP unary_expression { makeList("++", 'o', lineCount); }
64 | DEC_OP unary_expression { makeList("--", 'o', lineCount); }
65 | unary_operator cast_expression { makeList("sizeof", 'o', lineCount); }
66 | SIZEOF unary_expression { makeList("sizeof", 'o', lineCount); }
67 | SIZEOF '(' type_name ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
68
69 unary_operator
70 : '&' { makeList("&", 'o', lineCount); }
71 : '*' { makeList("*", 'o', lineCount); }
72 : '+' { makeList("+", 'o', lineCount); }
73 : '-' { makeList("-", 'o', lineCount); }
74 : '~' { makeList("~", 'o', lineCount); }
75 : '!' { makeList("!", 'o', lineCount); }
76
77 cast_expression
78 : unary_expression { $$=$1; }
79 | '(' type_name ')' cast_expression { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
80
81 multiplicative_expression
82 : cast_expression { $$=$1; }
83 | multiplicative_expression '*' cast_expression { makeList("*", 'o', lineCount); }
84 | multiplicative_expression '/' cast_expression { makeList("/", 'o', lineCount); }
85 | multiplicative_expression '%' cast_expression { makeList("%", 'o', lineCount); }
86
87 additive_expression
88 : multiplicative_expression { $$=$1; }
89 | additive_expression '+' multiplicative_expression { makeList("+", 'o', lineCount); }
90 | additive_expression '-' multiplicative_expression { makeList("-", 'o', lineCount); }

```

FIG. 2

```

91 | ;
92 | shift_expression
93 | : additive_expression {$$=$1;}
94 | | shift_expression LEFT_OP additive_expression { makeList("<=", 'o', lineCount); }
95 | | shift_expression RIGHT_OP additive_expression { makeList(">=", 'o', lineCount); }
96 | ;
97 | relational_expression
98 | : shift_expression {$$=$1;}
99 | | relational_expression '<' shift_expression
100 | | relational_expression '>' shift_expression
101 | | relational_expression LE_OP shift_expression { makeList("<=", 'o', lineCount); }
102 | | relational_expression GE_OP shift_expression { makeList(">=", 'o', lineCount); }
103 | ;
104 | equality_expression
105 | : relational_expression {$$=$1;}
106 | | equality_expression EQ_OP relational_expression { makeList("==", 'o', lineCount); }
107 | | equality_expression NE_OP relational_expression { makeList("!=", 'o', lineCount); }
108 | ;
109 | and_expression
110 | : equality_expression {$$=$1;}
111 | | and_expression '&' equality_expression { makeList("&", 'o', lineCount); }
112 | ;
113 | exclusive_or_expression
114 | : and_expression {$$=$1;}
115 | | exclusive_or_expression '^' and_expression { makeList("^", 'o', lineCount); }
116 | ;
117 | inclusive_or_expression
118 | : exclusive_or_expression {$$=$1;}
119 | | inclusive_or_expression '|' exclusive_or_expression { makeList("|", 'o', lineCount); }
120 | ;
121 | logical_and_expression
122 | : inclusive_or_expression {$$=$1;}
123 | | logical_and_expression AND_OP inclusive_or_expression { makeList("&&", 'o', lineCount); }
124 | ;
125 | logical_or_expression
126 | : logical_and_expression {$$=$1;}
127 | | logical_or_expression OR_OP logical_and_expression { makeList("||", 'o', lineCount); }
128 | ;
129 | conditional_expression
130 | : logical_or_expression {$$=$1;}
131 | | logical_or_expression '?' expression ':' conditional_expression { makeList("?:", 'o', lineCount); }
132 | ;
133 | assignment_expression
134 | : conditional_expression {$$=$1;}
135 | | unary_expression assignment_operator assignment_expression {$$=$3; checkType($1,$3,lineCount);}

```

FIG. 3

```

136 | ;
137 | assignment_operator
138 | : '=' { makeList("=", 'o', lineCount); }
139 | | MUL_ASSIGN { makeList("*=", 'o', lineCount); }
140 | | DIV_ASSIGN { makeList("/=", 'o', lineCount); }
141 | | MOD_ASSIGN { makeList("%=", 'o', lineCount); }
142 | | ADD_ASSIGN { makeList("+=", 'o', lineCount); }
143 | | SUB_ASSIGN { makeList("-=", 'o', lineCount); }
144 | | LEFT_ASSIGN { makeList("<<=", 'o', lineCount); }
145 | | RIGHT_ASSIGN { makeList(">>=", 'o', lineCount); }
146 | | AND_ASSIGN { makeList("&=", 'o', lineCount); }
147 | | XOR_ASSIGN { makeList("^=", 'o', lineCount); }
148 | | OR_ASSIGN { makeList("|=", 'o', lineCount); }
149 | ;
150 | expression
151 | : assignment_expression {$$=$1;}
152 | | expression ',' assignment_expression { makeList(",", 'p', lineCount); }
153 | ;
154 | constant_expression
155 | : conditional_expression
156 | ;
157 | declaration
158 | : declaration_specifiers ';' { makeList(";", 'p', lineCount); typeBuffer=' '; }
159 | | declaration_specifiers init_declarator_list ';' { makeList(";", 'p', lineCount); typeBuffer=' '; }
160 | ;
161 | declaration_specifiers
162 | : storage_class_specifier
163 | | storage_class_specifier declaration_specifiers
164 | | type_specifier
165 | | type_specifier declaration_specifiers
166 | | type_qualifier
167 | | type_qualifier declaration_specifiers
168 | ;
169 | init_declarator_list
170 | : init_declarator
171 | | init_declarator_list ',' init_declarator { makeList(",", 'p', lineCount); }
172 | ;
173 | init_declarator
174 | : declarator
175 | | declarator '=' initializer { makeList("=", 'o', lineCount); }
176 | ;
177 | storage_class_specifier
178 | : TYPEDEF { makeList("typedef", 'k', lineCount); }
179 | | EXTERN { makeList("extern", 'k', lineCount); }
180 | | STATIC { makeList("static", 'k', lineCount); }

```

FIG. 4


```

181 | AUTO          { makeList("auto", 'k', lineCount); }
182 | REGISTER      { makeList("register", 'k', lineCount); }
183 | ;
184 type_specifier
185 : VOID          { makeList("void", 'k', lineCount); typeBuffer='v'; }
186 | CHAR          { makeList("char", 'k', lineCount); typeBuffer='c'; }
187 | SHORT         { makeList("short", 'k', lineCount); }
188 | INT           { makeList("int", 'k', lineCount); typeBuffer='l'; }
189 | LONG          { makeList("long", 'k', lineCount); }
190 | FLOAT         { makeList("float", 'k', lineCount); typeBuffer='f'; }
191 | DOUBLE        { makeList("double", 'k', lineCount); }
192 | SIGNED        { makeList("signed", 'k', lineCount); }
193 | UNSIGNED      { makeList("unsigned", 'k', lineCount); }
194 | struct_or_union_specifier
195 | enum_specifier
196 | TYPE_NAME
197 | ;
198 struct_or_union_specifier
199 : struct_or_union IDENTIFIER '{' struct_declaration_list '}'
200 | struct_or_union '{' struct_declaration_list '}'
201 | struct_or_union IDENTIFIER
202 | ;
203 struct_or_union
204 : STRUCT        { makeList("struct", 'k', lineCount); }
205 | UNION         { makeList("union", 'k', lineCount); }
206 | ;
207 struct_declaration_list
208 : struct_declaration
209 | struct_declaration_list struct_declaration
210 | ;
211 struct_declaration
212 : specifier_qualifier_list struct_declarator_list ';' { makeList(";", 'p', lineCount); }
213 | ;
214 specifier_qualifier_list
215 : type_specifier specifier_qualifier_list
216 | type_specifier
217 | type_qualifier specifier_qualifier_list
218 | type_qualifier
219 | ;
220 struct_declarator_list
221 : struct_declarator
222 | struct_declarator_list ',' struct_declarator { makeList(",", 'p', lineCount); }
223 | ;
224 struct_declarator
225 : declarator

```

FIG. 5

```

226 | ':' constant_expression { makeList(":", 'p', lineCount); }
227 | declarator ':' constant_expression { makeList(":", 'p', lineCount); }
228 | ;
229 enum_specifier
230 : ENUM '{' enumerator_list '}' { makeList("enum", 'k', lineCount); }
231 | ENUM IDENTIFIER '{' enumerator_list '}' { makeList("enum", 'k', lineCount); makeList(tablePtr, 'v', lineCount); }
232 | ENUM IDENTIFIER { makeList("enum", 'k', lineCount); makeList(tablePtr, 'v', lineCount); }
233 | ;
234 enumerator_list
235 : enumerator
236 | enumerator_list ',' enumerator { makeList(",", 'p', lineCount); }
237 | ;
238 enumerator
239 : IDENTIFIER { makeList(tablePtr, 'v', lineCount); }
240 | IDENTIFIER '=' constant_expression { makeList("=", 'o', lineCount); makeList("tablePtr", 'v', lineCount); }
241 | ;
242 type_qualifier
243 : CONST { makeList("const", 'k', lineCount); }
244 | VOLATILE { makeList("volatile", 'k', lineCount); }
245 | ;
246 declarator
247 : pointer direct_declarator
248 | direct_declarator
249 | ;
250 direct_declarator
251 : IDENTIFIER { checkDeclaration(tablePtr, lineCount, scopeCount); }
252 | '(' declarator ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
253 | direct_declarator '[' constant_expression ']' { makeList("[", 'p', lineCount); makeList("]", 'p', lineCount); }
254 | direct_declarator '[' ']' { makeList("[", 'p', lineCount); makeList("]", 'p', lineCount); }
255 | direct_declarator '(' parameter_type_list ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
256 | direct_declarator '(' identifier_list ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
257 | direct_declarator '(' ']' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
258 | ;
259 pointer
260 : '*' { makeList("*", 'o', lineCount); }
261 | '*' type_qualifier_list { makeList("*", 'o', lineCount); }
262 | '*' pointer { makeList("*", 'o', lineCount); }
263 | '*' type_qualifier_list pointer { makeList("*", 'o', lineCount); }
264 | ;
265 type_qualifier_list
266 : type_qualifier
267 | type_qualifier_list type_qualifier
268 | ;
269 parameter_type_list
270 : parameter_list

```

FIG. 6

```

271 | parameter_list ',' ELLIPSIS { makeList(" ", 'p', lineCount); makeList(":::", 'o', lineCount); }
272 ;
273 parameter_list
274 : parameter_declaration
275 | parameter_list ',' parameter_declaration { makeList(" ", 'p', lineCount); }
276 ;
277 parameter_declaration
278 : declaration_specifiers declarator
279 | declaration_specifiers abstract_declarator
280 | declaration_specifiers
281 ;
282 identifier_list
283 : IDENTIFIER { checkDeclaration(tablePtr, lineCount, scopeCount); }
284 | identifier_list ',' IDENTIFIER { checkDeclaration(tablePtr, lineCount, scopeCount); makeList(" ", 'p', lineCount); }
285 ;
286 type_name
287 : specifier_qualifier_list
288 | specifier_qualifier_list abstract_declarator
289 ;
290 abstract_declarator
291 : pointer
292 | direct_abstract_declarator
293 | pointer direct_abstract_declarator
294 ;
295 direct_abstract_declarator
296 : '(' abstract_declarator ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
297 | '[' ']' { makeList("[", 'p', lineCount); makeList("]", 'p', lineCount); }
298 | '[' constant_expression ']' { makeList("[", 'p', lineCount); makeList("]", 'p', lineCount); }
299 | direct_abstract_declarator '[' ']' { makeList("[", 'p', lineCount); makeList("]", 'p', lineCount); }
300 | direct_abstract_declarator '[' constant_expression ']' { makeList("[", 'p', lineCount); makeList("]", 'p', lineCount); }
301 | '(' ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
302 | '(' parameter_type_list ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
303 | direct_abstract_declarator '(' ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
304 | direct_abstract_declarator '(' parameter_type_list ')' { makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
305 ;
306 initializer
307 : assignment_expression {$$=$1;}
308 | '{' initializer_list '}'
309 | '{' initializer_list ',' '}'
310 ;
311 initializer_list
312 : initializer
313 | initializer_list ',' initializer { makeList(" ", 'p', lineCount); }
314 ;
315 statement

```

FIG. 7

```

316 : labeled_statement
317 | compound_statement
318 | expression_statement
319 | selection_statement
320 | iteration_statement
321 | jump_statement
322 ;
323 labeled_statement
324 : IDENTIFIER ':' statement { makeList(tablePtr, 'v', lineCount); }
325 | CASE constant_expression ':' statement { makeList(":", 'p', lineCount); makeList("case", 'k', lineCount); }
326 | DEFAULT ':' statement { makeList(":", 'p', lineCount); makeList("default", 'k', lineCount); }
327 ;
328 compound_statement
329 : '(' '}'
330 | '(' statement_list '}'
331 | '(' declaration_list '}'
332 | '(' declaration_list statement_list '}'
333 ;
334 declaration_list
335 : declaration
336 | declaration_list declaration
337 ;
338 statement_list
339 : statement
340 | statement_list statement
341 ;
342 expression_statement
343 : ';' { makeList(";", 'p', lineCount); }
344 | expression ';' { makeList(";", 'p', lineCount); }
345 ;
346 selection_statement
347 : IF '(' expression ')' statement %prec LOWER_THAN_ELSE { makeList("if", 'k', lineCount); makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
348 | IF '(' expression ')' statement ELSE statement { makeList("if", 'k', lineCount); makeList("else", 'k', lineCount); makeList("(", 'p', lineCount); }
349 | SWITCH '(' expression ')' statement { makeList("switch", 'k', lineCount); makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
350 ;
351 iteration_statement
352 : WHILE '(' expression ')' statement { makeList("while", 'k', lineCount); makeList("(", 'p', lineCount); makeList(")", 'p', lineCount); }
353 | DO statement WHILE '(' expression ')' ';' { makeList("do", 'k', lineCount); makeList("while", 'k', lineCount); makeList("(", 'p', lineCount); }
354 ;
355 jump_statement
356 : BREAK { makeList("break", 'k', lineCount); }
357 | CONTINUE { makeList("continue", 'k', lineCount); }
358 | RETURN { makeList("return", 'k', lineCount); }
359 | GOTO IDENTIFIER { makeList("goto", 'k', lineCount); }
360 ;

```

FIG. 8

Semantic Analyzer

This is the semantic phase code , a c header file which contains the function that is called by different production rule in parser phase and taken care of symbol table entry.

```
1 struct tokenList
2 {
3     char *token,type[20],line[100];
4     char *scope[20];
5     int scopeValue;
6     int funcCount;
7     struct tokenList *next;
8 };
9 typedef struct tokenList tokenList;
10 struct funcNode{
11     char funcName[30];
12     int line;
13     char funcReturn[20];
14     struct funcNode *next;
15 };
16 typedef struct funcNode funcNode;
17
18
19 tokenList *symbolPtr = NULL;
20 tokenList *constantPtr = NULL;
21 tokenList *parsedPtr =NULL;
22 extern int functionCount;
23 extern int scopeCount;
24 char typeBuffer=' ';
25 char *sourceCode;
26 int tempCheckType=3;
27
28 int semanticErr=0,lineSemanticCount;
29 int checkScope(char *tempToken,int lineCount)
30 {
31     tokenList *temp=NULL;
32     char type[20];
33     int flag=0,tempFlag=0;
34     for(tokenList *p=symbolPtr;p!=NULL;p=p->next){
35         if(strcmp(tempToken,"printf")==0 || strcmp(tempToken,"scanf")==0){
36             tempFlag=1;
37         }
38         else{
39             if(strcmp(tempToken,p->token)==0){
40                 strcpy(type,p->type);
41                 flag=1;
42                 break;
43             }
44         }
45     }
46 }
```

FIG. 1

```

46     if (flag == 0 && tempFlag == 0 )
47     {
48         printf("\n%s : %d :Undeclared variable \n",sourceCode,lineCount-1);
49         semanticErr++;
50     }
51     else
52     {
53         addSymbol(tempToken,lineCount);
54         if(strcmp(type,"VOID")==0)
55             return(1);
56         if(strcmp(type,"CHAR")==0)
57             return(2);
58         if(strcmp(type,"INT")==0)
59             return(3);
60         if(strcmp(type,"FLOAT")==0)
61             return(4);
62     }
63 }
64 }
65
66 void checkType(int value1,int value2,int lineCount)
67 {
68     lineSemanticCount=lineCount;
69     if(value2 == 0)
70         value2 = tempCheckType;
71     if(value1!=value2)
72     {
73         printf("\n%s : %d :Type Mismatch error \n",sourceCode,lineSemanticCount-1);
74         semanticErr++;
75     }
76     tempCheckType=3;
77 }
78 void checkDeclaration(char *tokenName,int tokenLine,int scopeVal){
79     char type[20];
80     char line[39],lineBuffer[19];
81     snprintf(lineBuffer, 19, "%d", tokenLine);
82     strcpy(line, " ");
83     strcat(line,lineBuffer);
84     switch(typeBuffer){
85         case 'i': strcpy(type,"INT"); break;
86         case 'f': strcpy(type,"FLOAT");break;
87         case 'v': strcpy(type,"VOID");break;
88         case 'c': strcpy(type,"CHAR");break;
89     }
90     for(tokenList *p=symbolPtr;p!=NULL;p=p->next){

```

FIG. 2

```

91     if(strcmp(p->token,tokenName)==0 && p->scopeValue == scopeCount && p->funcCount == functionCount){
92         semanticErr++;
93         if(strcmp(p->type,type)==0){
94             printf("\n%s : %d :Multiple Declaration \n",sourceCode,tokenLine);
95             return;
96         }
97         else{
98             printf("\n%s : %d :Multiple Declaration with Different Type \n",sourceCode,tokenLine);
99             return;
100         }
101     }
102 }
103 addSymbol(tokenName,tokenLine,scopeCount);
104 }
105 }
106 void checkArray(int val,int lineCount){
107     if(val<0){
108         semanticErr++;
109         printf("\n%s : %d :Array Index error\n",sourceCode,lineCount-1);
110     }
111 }
112 void addSymbol(char *tokenName,int tokenLine,int scopeVal){
113     char line[39],lineBuffer[19];
114     snprintf(lineBuffer, 19, "%d", tokenLine);
115     strcpy(line, " ");
116     strcat(line,lineBuffer);
117     char type[20];
118     for(tokenList *p=symbolPtr;p!=NULL;p=p->next)
119         if(strcmp(p->token,tokenName)==0 && p->scopeValue == scopeCount && p->funcCount ==functionCount ){
120             strcat(p->line,line);
121             return;
122         }
123     tokenList *temp=(tokenList *)malloc(sizeof(tokenList));
124     temp->token=(char *)malloc(strlen(tokenName)+1);
125     strcpy(temp->token,tokenName);
126     switch(typeBuffer){
127         case 'i': strcpy(temp->type,"INT"); break;
128         case 'f': strcpy(temp->type,"FLOAT");break;
129         case 'v': strcpy(temp->type,"VOID");break;
130         case 'c': strcpy(temp->type,"CHAR");break;
131     }
132     temp->funcCount=functionCount;
133     if(scopeCount==0){
134         strcpy(temp->scope,"GLOBAL");
135         temp->scopeValue=scopeCount;

```

FIG. 3

```

136     temp->scopeValue=scopeCount;
137 }
138 else{
139     strcpy(temp->scope,"NESTING");
140     temp->scopeValue=scopeCount;
141 }
142 strcpy(temp->line,line);
143 temp->next=NULL;
144 tokenList *p=symbolPtr;
145 if(p==NULL){
146     symbolPtr=temp;
147 }
148 else{
149     while(p->next!=NULL){
150         p=p->next;
151     }
152     p->next=temp;
153 }
154 }
155 }
156 }
157 void addConstant(char *tokenName,int tokenLine){
158     char line[39],lineBuffer[19];
159     snprintf(lineBuffer, 19, "%d", tokenLine);
160     strcpy(line, " ");
161     strcat(line,lineBuffer);
162     for(tokenList *p=constantPtr;p!=NULL;p=p->next)
163         if(strcmp(p->token,tokenName)==0){
164             strcat(p->line,line);
165             return;
166         }
167     tokenList *temp=(tokenList *)malloc(sizeof(tokenList));
168     temp->token=(char *)malloc(strlen(tokenName)+1);
169     strcpy(temp->token,tokenName);
170     strcpy(temp->line,line);
171     temp->next=NULL;
172     tokenList *p=constantPtr;
173     if(p==NULL){
174         constantPtr=temp;
175     }
176     else{
177         while(p->next!=NULL){
178             p=p->next;
179         }
180         p->next=temp;

```

FIG. 4

```

181     p->next=temp;
182 }
183 }
184 }
185 }
186 void makeList(char *tokenName,char tokenType, int tokenLine)
187 {
188     char line[39],lineBuffer[19];
189     snprintf(lineBuffer, 19, "%d", tokenLine);
190     strcpy(line, " ");
191     strcat(line,lineBuffer);
192     char type[20];
193     switch(tokenType)
194     {
195         case 'c':
196             strcpy(type,"Constant");
197             break;
198         case 'v':
199             strcpy(type,"Identifier");
200             break;
201         case 'p':
202             strcpy(type,"Punctuator");
203             break;
204         case 'o':
205             strcpy(type,"Operator");
206             break;
207         case 'k':
208             strcpy(type,"Keyword");
209             break;
210         case 's':
211             strcpy(type,"String Literal");
212             break;
213         case 'd':
214             strcpy(type,"Preprocessor Statement");
215             break;
216     }
217 }
218 }
219 }
220 for(tokenList *p=parsedPtr;p!=NULL;p=p->next)
221     if(strcmp(p->token,tokenName)==0){
222         strcat(p->line,line);
223         return;
224     }
225     tokenList *temp=(tokenList *)malloc(sizeof(tokenList));
226     temp->token=(char *)malloc(strlen(tokenName)+1);

```

FIG. 5

```

210         case 'k':           strcpy(type,"Keyword");
211                             break;
212         case 's':           strcpy(type,"String Literal");
213                             break;
214         case 'd':           strcpy(type,"Preprocessor Statement");
215                             break;
216
217     }
218
219     for(tokenList *p=parsedPtr;p!=NULL;p=p->next)
220     {
221         if(strcmp(p->token,tokenName)==0){
222             strcat(p->line,line);
223             return;
224         }
225         tokenList *temp=(tokenList *)malloc(sizeof(tokenList));
226         temp->token=(char *)malloc(strlen(tokenName)+1);
227         strcpy(temp->token,tokenName);
228         strcpy(temp->type,type);
229         strcpy(temp->line,line);
230         temp->next=NULL;
231
232         tokenList *p=parsedPtr;
233         if(p==NULL){
234             parsedPtr=temp;
235         }
236         else{
237             while(p->next!=NULL){
238                 p=p->next;
239             }
240             p->next=temp;
241         }
242     }
243 }
244
245
246
247
248
249
250
251
252
253
254

```

FIG. 6

EXECUTION OF THE CODE

The following is a shell script to automate the compilation and execution of the mini compiler consist of lexical ,syntax and semantic phases. There is sample input source code given .execution of the same is done and output is shown as sample output and respective symbol table is generated

Shell Script

```

#!/bin/sh
lex lexicalAnalyzer.l
yacc -d syntaxChecker.y
gcc lex.yy.c y.tab.c -w -g
./a.out input.c
rm y.tab.c y.tab.h lex.yy.c

```

Sample Input

```
1#include<stdio.h>
2void func(){
3    int a;
4}
5void main()
6{
7
8    int a=10,b;
9    char c;
10   int f[100][10];
11   a=b;
12   a=b+12;
13   printf("\nHello World");
14   scanf();
15   {
16       int y;
17       int a;
18   }
19
20
21}
22
```

Sample Output

```
aswanth@hp-notebook:~/Projects/CD/Compiler-Design/Semantic Analyzer/src$ ./compile.sh

input.c Parsing Completed
```

Symbol Table

SymbolTable					
Token	Type	Line Number	Scope	Function Number	
func	VOID	2	GLOBAL 0	1	
a	INT	3	NESTING 1	1	
main	VOID	5	GLOBAL 0	2	
a	INT	8 11 12	NESTING 1	2	
b	INT	8 11 12	NESTING 1	2	
c	CHAR	9	NESTING 1	2	
f	INT	10	NESTING 1	2	
printf		13	NESTING 1	2	
scanf		14	NESTING 1	2	
y	INT	16	NESTING 2	2	
a	INT	17	NESTING 2	2	

TEST CASES

Different test cases are added to check implementation of semantic phase. The validation of the same is done below in a table format. Source code for the 6 test cases are added as screenshot below.

Test case Source code

Case1.c

```
1//Testcase to check multiple declaration
2
3#include <stdio.h>
4int main()
5{
6    int a = 5, b = 3;
7    float x;
8    char b;
9
10    int a = b + 10;
11}
```

Case2.c

```
1// Testcase to check undeclared variables
2#include<stdio.h>
3int main()
4{
5    int a=10,b=50,c=10;
6    d = a+b;
7    printf("The sum of %d + %d = %d",a,b,d);
8}
```

Case3.c

```
1// Testcase to check type mismatch
2#include<stdio.h>
3int main()
4{
5    int l=10,a=0,i;
6    float l=5.0;
7    for(i=0;i<l;i++)
8    {
9        printf("\nHello World");
10    }
11
12
13    l = l + 3.14;
14    printf("%d",l);
15}
```

Case4.c

```
1//Testcase to check scope of variables
2#include<stdio.h>
3int func(){
4    int a;
5}
6int main()
7{
8    int a;
9    {
10        int a;
11    }
12    func();
13}
```

Case5.c

```
1// Testcase to check scope of a variable
2
3#include <stdio.h>
4void main()
5{
6    int a = 10, b = 5;
7
8    {
9        printf("Enter a variable");
10        int b;
11        scanf("%d",&b);
12    }
13}
```

Case6.c

```
1// Testcase to check array dimensions
2
3#include <stdio.h>
4int main()
5{
6    int i, arr1[100], arr2[100];
7    for(i=0;i<100;i++)
8    {
9        scanf("%d",&arr1[i]);
10       arr2[i] = arr1[i] * 2;
11    }
12
13    for(i=0;i<100;i++)
14    {
15        printf("%d",arr2[i]);
16    }
17}
```

Test Case Evaluation

File Name	Purpose	Expected Output	Explanation	Status
case1.c	Check Multiple Declaration	case1.c : 9 :Multiple declaration with different Type case1.c : 11 :Multiple Declaration	Multiple declaration of variable a and b	Passed
case2.c	Check Undeclared variable	case2.c : 6 :Undeclared variable case2.c : 6 :Type Mismatch error case2.c : 7 :Undeclared variable	Variable d is undeclared	Passed
case3.c	Check Type mismatch variable	case3.c : 7 :Multiple Declaration with Different Type case3.c : 7 :Type Mismatch error	Variable l is of type <i>int</i> but assignment expression of type float	Passed

case4.c	Check scope of variable	case4.c Parsing Completed	Though multiple declaration of a but all under different scope	Passed
case5.c	Check scope of variable with functions	case5.c Parsing Completed	Multiple declaration under different compound statement	Passed
case6.c	Check array dimension	case6.c : 7 : Invalid array Index	<i>arr2</i> has array dimension negative	Passed

CONCLUSION

After the third phase of this project, we have successfully implemented and added a Semantic Analyzer to the C Compiler. Semantic analysis checks whether the parse tree constructed follows the rules of language. Also, the semantic analyzer keeps track of identifiers, their types and expressions; whether identifiers are declared before use or not, etc. The semantic analyzer produces an annotated syntax tree as an output.