COMPILER DESIGN PROJECT REPORT - 4

INTERMEDIATE CODE GENERATION FOR C LANGUAGE



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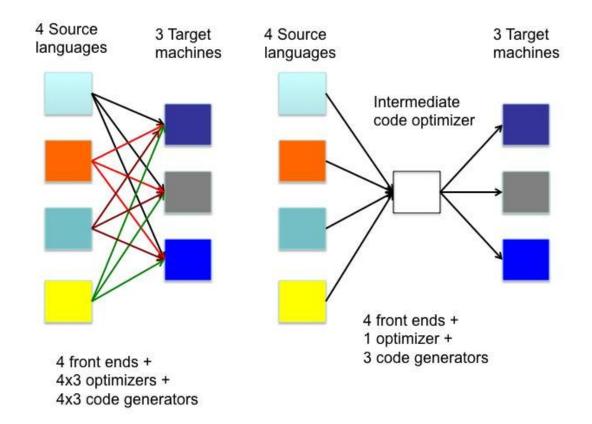
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INTRODUCTION

Compilers generate machine code, whereas interpreters interpret intermediate code. Interpreters are easier to write and can provide better error messages (symbol table is still available). However, they are at least 5 times slower than machine code generated by compilers and also require much more memory than machine code generated by compilers.

While generating machine code directly from source code is possible, it entails two problems:

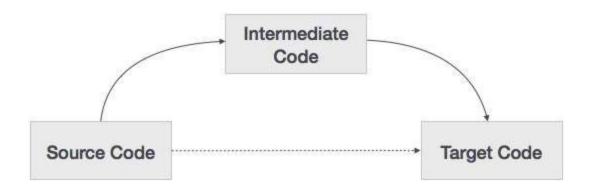
- 1) With **m** languages and n target machines, we need to write **m** front ends, $\mathbf{m} \times \mathbf{n}$ optimizers, and $\mathbf{m} \times \mathbf{n}$ code generators
- 2) The code optimizer which is one of the largest and very-difficult-to-write components of a compiler, cannot be reused.



By converting source code to an intermediate code, a machine-independent code optimizer may be written. This means just **m** front ends, **n** code generators and **1** optimizer.

A source code can directly be translated into its target machine code, but we need to translate the source code into an intermediate code which is then translated to its target code. The main reasons for this are:

- 1) If a compiler translates the source language to its target machine language without having the option for generating intermediate code, then for each new machine, a full native compiler is required.
- 2) Intermediate code eliminates the need of a new full compiler for every unique machine by keeping the analysis portion same for all the compilers.
- 3) The second part of compiler, synthesis, is changed according to the target machine.
- 4) It becomes easier to apply the source code modifications to improve code performance by applying code optimization techniques on the intermediate code.



Intermediate code can be either language specific (e.g., Bytecode for Java) or language independent (three-address code).

INTERMEDIATE REPRESENTATION

The aim of Intermediate Code Generation is to translate the program into a format expected by the compiler back-end. Intermediate code must be easy to produce and easy to translate to machine code.

Why use an intermediate representation?

- 1) It's easy to change the source or the target language by adapting only the front-end or back-end (portability).
- 2) It makes optimization easier: one needs to write optimization methods only for the intermediate representation.
- 3) The intermediate representation can be directly interpreted.

Intermediate codes can be represented in a variety of ways and they have their own benefits.

- 1) **High Level IR** High-level intermediate code representation is very close to the source language itself. They can be easily generated from the source code and we can easily apply code modifications to enhance performance. But for target machine optimization, it is less preferred.
- 2) Low Level IR This one is close to the target machine, which makes it suitable for register and memory allocation, instruction set selection, etc. It is good for machine-dependent optimizations.

Some general forms of intermediate representation are:

- 1) Graphical IR (parse tree, abstract syntax trees, DAG. . .)
- 2) Linear IR (ie., non graphical)
- 3) Three Address Code (TAC): instructions of the form "result=op1 operator op2"
- 4) Static single assignment (SSA) form: each variable is assigned once
- 5) Continuation-passing style (CPS): general form of IR for functional language

THREE-ADDRESS CODE

Intermediate code generator receives input from its predecessor phase, semantic analyzer, in the form of an annotated syntax tree. That syntax tree then can be converted into a linear representation, e.g., postfix notation. Intermediate code tends to be machine independent code. Therefore, code generator assumes to have unlimited number of memory storage (register) to

generate code. The intermediate code generator will try to divide this expression into sub-expressions and then generate the corresponding code.

A statement involving no more than three references(two for operands and one for result) is known as three address statement. A sequence of three address statements is known as three address code. Three address statement is of the form x = y op z, here x, y, z will have address (memory location). Sometimes a statement might contain less than three references but it is still called three address statement.

Example – The three address code for the expression a + b * c + d:

- T 1 = b * c
- T 2 = a + T 1
- T 3 = T 2 + d

where T 1, T 2, T 3 are temporary variables.

A three-address code has at most three address locations to calculate the expression. A three-address code can be represented in two forms: quadruples and triples.

1) Quadruples - Each instruction in quadruples presentation is divided into four fields: operator, arg1, arg2, and result. Consider the example, a = b + c * d; It is represented below in quadruples format:

Ор	arg1	arg2	result
*	С	d	r1
+	b	r1	r2
+	r2	r1	r3
=	r3		a

2) Triples - Each instruction in triples presentation has three fields: op, arg1, and arg2. The results of respective sub-expressions are denoted by the position of expression. Triples represent similarity with DAG and syntax tree. They are equivalent to DAG while representing expressions.

Ор	arg1	arg2
*	c	d
+	b	(0)
+	(1)	(0)
=	(2)	

Triples face the problem of code immovability while optimization, as the results are positional and changing the order or position of an expression may cause problems.

3) Indirect Triples - This representation is an enhancement over triples representation. It uses pointers instead of position to store results. This enables the optimizers to freely re-position the sub-expression to produce an optimized code.

IMPLEMENTATION

When generating IR at this level, do not need to worry about optimizing it. It's okay to generate IR that has lots of unnecessary assignments, redundant computations, etc.

All the lexical analyzer part are done in lexicalAnalyzer.l which returns each token along with its value or type based on whether it is variable or constant or keyword. Syntax and semantic check has been done in syntaxChecker.y. All the productions and respective semantic actions are written here

The functions that are generating intermediate code has been written in icg{.h,.c}. The main functions in icg files are

- threeAddressCode* appendCode(char *code)
 Append the three address code provided as argument to the main set of three addresscode
- void backpatch(backPatchList * list, int gotoL)
 It is to add the line number "gotoL" to backPatchList

backPatchList* mergelists(backPatchList * a, backPatchList * b)

It is to merge to back patch list if it is resolved.

 backPatchList* appendToBackPatch(backPatchList * list, threeAddressCode * entry)

It is to append a single line three address code entry to the back patch list.

 tokenList* appendToSymbolTable(char *name,tokenType type,tokenReturnType returnType,long size,long line,char *scope,long parameter)

It is used add token "name" to the symbol table along with all the properties

void writeCode(FILE *icgOut)

This is write the three address code generated till now into a file provided as argument

void writeSymbolTable(FILE *symOut)

This is to write the symbol table entries till now to the file provided as argument

These functions are called from respective productions in syntaxAnalyzer.y whenever needed.All the code returned are stored in a global pointer such that can write it to a file at the end. The source code for the same has been provided below.

SOURCE CODE

lexicalAnalyzer.l

This is the lex program that contains regular expressions and returns whatever tokens are required.

syntaxChecker.y

This is the main parser code, a yacc file which contains the declarations, rules and programs and defines the actions which should be taken for various cases.

```
### Struct
### Struct BackpatchList* nextList;
### Struct BackpatchList* nextList;
### Struct Struct
### Struct
### Struct Struct Struct
### Struct
### Struct Struct
### Struct
###
```

```
backpatch($2.nextList, nextquad());

function

function
```

```
paramCount++;|
addSymbolToQueue($4, FLOAT_type, paramCount);
      VOID
var_type
: VOID
              $$ = Return_VOID;
     | INT
              $$ = Return_INT;
     | FLOAT
              $$ = Return_FLOAT;
;
statement_list
: statement
     | statement_list marker statement
              backpatch($1.nextList,$2.quad);
$$.nextList = $3.nextList;
;
statement
     : matched_statement
              $$.nextList = $1.nextList;
     }
| unmatched_statement
              $$.nextList = $1.nextList;
matched_statement
: IF '(' assignment ')' marker matched_statement jump_marker ELSE marker matched_statement
              backpatch($3.trueList,$5.quad);
```

```
backpatch($3.falseList,$9.quad);
$5.nextList = mergeLists($5.nextList,$10.nextList);
$$.nextList = mergeLists($5.nextList,$6.nextList);

$$.nextList = NULL;

$$.nextList =
```

```
| backpatch($10.nextList, $5.quad);
| IF'(' assignment ')' marker matched_statement jump_marker ELSE marker unmatched_statement
| backpatch($3.talselist, $5.quad);
| call backpatch($3.talselist, $5.quad);
| call backpatch($3.talselist, $5.quad);
| backpatch($3.talselist, $5.quad);
```

```
sprintf(tcgoud, 'if' 'is' co 0) coro', $4.value);
$4. truetist = appendToackstch(NULL, appendCode(tcgOuad));
$5. falseList = appendToackstch(NULL, appendCode(tcgOuad));

$5. falseList = appendToackstch(NULL, appendCode(tcgOuad));

$5. falseList = $6. falseList;
$6. falseList = $6. falseList;
$6. falseList = $6. falseList;
$6. falseList = $6. falseList;
$7. falseLis
```

```
| Signature of Transfale Only|";
| Strain of Transfale Only|
```

```
char* var = NULL;

sitch(type){
    case Ini_type: var = nextIntVar();break;
    case Ini_type: var = nextIntVar();break;

case Ini_type: var = nextIntVar();break;

case Ini_type: var = nextIntVar();break;

case Ini_type: var = nextIntVar();break;

case Ini_type: var;

ss. value = var;

ss. value = var;

ss. type = type;

ss. truet.ts = NULL;

yerror();

int type = s;

tf(si.type is Ini_type && $1.type! = FLOAT_type && $3.type i = Ini_type && $3.type i = FLOAT_type){
    printf('show: Only integer and float values allowed when substracting numbers.\n');

yerror();

int type = s;

tf(si.type = s3.type){
    type = FLOAT_type;
}

case Ini_type: var = nextIntVar();break;

spendcode(tcg)uad;
ss. value = var;
ss. type = type;
ss. truet.ts = NULL;

ss. clope = VAR_type;
ss. truet.ts = NULL;
ss. clope = VAR_type;
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ss. truet.ts = NULL;
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ss. truet.ts = NULL;
ss. clope = VAR_type;
ss. clope = V
```

```
if($1.type != INT_type && $1.type!= FLOAT_type && $3.type!= FLOAT_type){
    printf("ERROR! Only integer and float values allowed when multiplicating numbers.\n");

    yyerror();
}
type = 0;
tr($1.type = $3.type){
    type = $1.type;
}
else{
    type = FLOAT_type;
}
acis FLOAT_type; var = nextIntVar();break;
case FLOAT_type; var = nextIntVar();break;
}
case FLOAT_type; var = nextIntVar();break;
case FLOAT_type; var = nextIntVar();break;
}
case FLOAT_type = Var = nextIntVar();break;
}
case FLOAT_type;
case FLOAT_type;
}
case FLOAT_type;
case FLOAT_type;
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case FLOAT_type;
case FLOAT_type;
case FLOAT_type;
}
case FLOAT_type;
cas
```

icg.c

```
interclude "teg.h"
if therclude "y.tab.h"

tokenList *symPtr = NULL;
stokenList *bufferPtr = NULL;
threeAddresScode *codePtr = NULL;
extern int yyltneno;

int globalOffset = 0, tempIntCounter = 0,tempFloatCounter= 0,currentLine = -1;

void yyerror() {
    printf( ERRGR: Md : Syntactical Error\n",yylineno);
    extern int yyltneno;
extern int yyltneno;

int nadioffset = 0, tempIntCounter = 0,tempFloatCounter= 0,currentLine = -1;

void yyerror() {
    printf( ERRGR: Md : Syntactical Error\n",yylineno);
    extern int yyltneno;
extern int yyltneno
```

```
| clse if(p == mul.) {
| backPatchList* temp = malloc(sizeof(backPatchList));
| temp->entry = newCode;
| temp->next = temp = newCode;
| temp->next = temp;
| p->next = temp;
| p->next = temp;
| p->next = temp;
| p->next = temp;
| temp->preturn temp;
| tokenList* appendToSymbolTable(char *name,tokenType type,tokenReturnType returnType,long size,long line,char *scope,long parameter)
| tokenList* temp = nalloc(sizeof(tokenList));
| temp->name = strdup(name);
| temp->tope = typeurnType;
| temp->elpeurnType = typeurnType;
| temp->elpeurnType = typeurnType;
| temp->elpeurnType = size;
| temp->elpeurnType
```

```
p->stze,
p->scope == NULL ? "None" : p->scope

p->scope == NULL ? "None" : p->scope

p->scope == NULL ? "None" : p->scope

p-p->next;

p-p->next;

tokenList 'tenp = nalloc(stzeof(tokenList));

tenp->name = strdup(name);

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```
else if(|prototype){

if(bufferPtr == NULL && parancount != 0){
    printf("ERROR: Md : Paraneter mismatch of function %s\n",yylineno,name);
    exit();
}

clse{
    int cnt = !;
    int s = 0;
    int size = symbol = bufferPtr;
    while(symboll=NULL){
        int size = symbol->size;
        symbol>size = localOffSet;
        localOffSet+size;
        globaloffSet+size;
        globaloffSet+size;
        globaloffSet+size;
        symbol>soze = name;
    if(parancount != n){
        symbol>sparameter=ont++;
    }

    is = symbol>parameter=ont++;
    }

    is = symbol>parameter=ont++;
    is = symbol>parameter=ont++;
```

```
char* nextFloatVar(){
che but(off);
che
```

icg.h

```
I strictude existion.h>

2 strictude existion.h>
3 strictude existion.h>
5 strictude existion.h>
5 strictude existion.h>
6 strictude existion.h>
7 strictude existion.h>
7 strictude existion.h>
8 soot, type,
8 soot, type,
10 soot, type,
11 int type,
12 FLOAT type,
13 FLOAT type,
14 FLOAT type,
15 FLOAT type,
16 Strictude existion.h

18 caturn_void,
18 caturn_void,
19 Return_tin,
20 Return_tin,
21 Return_tin,
22 Return_tin,
23 tokeneturntype;
24 typedef enum
25 CONST_type,
26 CONST_type,
27 VAR_type,
28 NONE_type
29 NONE_type
20 tokenconstrype;
20 tokenconstrype;
21 tokentype type;
22 tokentype type;
23 tokentype type;
24 tokentype type;
25 tokentype type;
26 tokenturntype returntype;
27 tokentype type;
28 tokentype type;
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44 tokenturntype returntype;
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46 tokenturntype returntype;
47 tokenturntype returntype;
48 tokenturntype returntype;
49 tokenturntype returntype;
40 tokenturntype returntype;
40 tokenturntype returntype;
40 tokenturntype returntype;
41 tokenturntype returntype;
42 typedef struct tokenturntype returntype;
43 typedef struct tokenturntype returntype;
44 typedef struct tokenturntype returntype;
45 typedef struct tokenturntype returntype;
46 typedef struct tokenturntype;
47 typedef struct tokenturntype;
48 typedef struct tokenturntype;
49 typedef struct tokenturntype;
40 typedef struct tokenturntype;
40 typedef struct tokenturntype;
40 typedef struct tokenturntype;
40 typedef struct tokenturntype;
41 typedef struct tokenturntype;
42 typedef struct tokenturntype;
43 typedef struct tokenturntype;
44 typedef struct tokenturntype;
45 typedef struct tokenturntype;
46 typedef struct tokentur
```

```
struct threeAddressCode
{
    char *code;
    int gotoLine;

struct threeAddressCode *next;

struct backPatchList
{
    threeAddressCode *next;

struct backPatchList
{
    threeAddressCode *entry;
    struct backPatchList *next;

}

threeAddressCode *entry;
    struct backPatchList *next;

struct backPatchList *list, int gotoL);

backPatchList *nergeLists(backPatchList *gueue, char *name, long patch gotoLists *gueue, gotoLists gotoLis
```

EXECUTION OF THE CODE

The following is a shell script to automate the compilation and execution of the mini compiler

consist of all phases. There is sample input source code given . Execution of the same is done and output is shown as sample output and respective symbol table is generated.

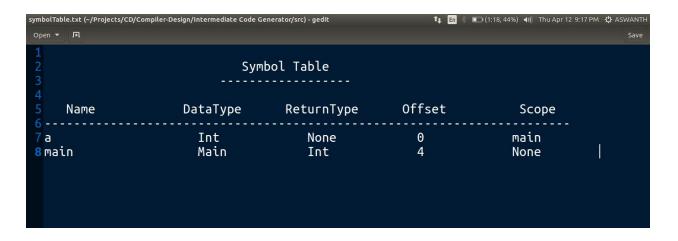
Shell Script

```
1#!/bin/sh
2lex lexicalAnalyzer.l
3 yacc -d syntaxChecker.y
4gcc lex.yy.c y.tab.c icg.c -w -g
5./a.out input.c
6rm y.tab.c y.tab.h lex.yy.c
```

Sample Input

```
1 #include <stdlib.h>
2 int main(){
3         int a;
4         a=40;
5         if(a<50){
6             a=60;
7         }
8         else{
9             a=100;
10         }
11         return 1;
12 }
13</pre>
```

Symbol Table



Sample Three-address Code

TEST CASES

Different test cases are added to check implementation of ICG phase. Source code for the 6 test cases are added as screenshot below along with their three-address code.

Case1.c

```
1//Testcase generate simple expresion statements
2 #include<stdio.h>
3 int main(){
4     int a;
5     float t;
6     a=10;
7     t=t*10;
8     return 1;
9}
```

```
Intermediate Code Generated

GOTO 1

a := 10

Tf_1 := t * 10

RETURN 1
```

Case2.c

```
Intermediate Code Generated

GOTO 1

1 a := 20

1 IF (a < 30) GOTO 4

3 GOTO 6

4 t := 3.0

9 GOTO 7

10 6 t := 4.0

11 RETURN 2
```

Case3.c

Case4.c

```
Intermediate Code Generated

GOTO 3

S := 56

2 a := 20

RETURN 1
```

Case5.c

```
1//testcase to check main function errors
2 #include<stdio.h>
3 int main();
4 int main(){
5     int y;
6     y=20;
7}
```

```
Intermediate Code Generated

GOTO 1

y := 20
```

CONCLUSION

Intermediate Code Generation (ICG) is the fourth phase of a compiler. It is the final phase of the front end of a compiler. It generates intermediate code, that is a form which can be readily executed by machine. Intermediate code is converted to machine language using the last two phases which are platform dependent. We can take the intermediate code from the already existing compiler and build the last two parts. After this, the intermediate code is ready and it is then sent to the **Code Optimizer** phase which transforms the code so that it consumes fewer resources and produces more speed.