

File Management



➤ File Management

Q. What is file?

❖ User view:

- File is a named collection of logically related information/data.
- File is a container which contains logically related information/data.
- File is a collection of characters/records/lines.
- File is a basic storage unit

❖ System view:

- File is a **stream of bits/bytes**.

- **File = data + metadata**

- data = actual file contents
- metadata = information about the file.



Operating Systems Concepts

- Data of the file exists inside the file whereas information about the file gets stored inside one structure referred as **FCB (File Control Block)**.
- In UNIX environment FCB is also called as an **iNode**.
- **FCB / iNode contains information about the file like:**
 1. iNode number: unique identifier of a file
 2. Name of the file
 3. Type of the file
 4. Size of the file
 5. Parent folder location
 6. Access permissions
 7. Time stamps etc...



- Per file one iNode/FCB gets created by the system and
-i.e. hence no. of iNodes = no. of files onto the disk.
- **data + metadata** of all the files are kept onto the disk, as disk may contain thousands of files, so thousands of iNodes and millions of bytes of data gets stored onto the disk, and hence all this **data + metadata** of all files need be kept onto the disk in an organized manner so that it can be accessed efficiently.

○ **File system** : file system is a way to store data onto the disk in an organized manner so that it can be accessed efficiently and conveniently.

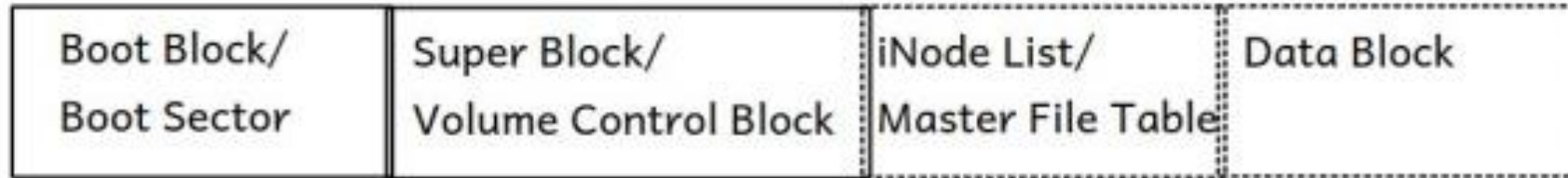
- e.g. Each OS has its own file system like,
- UNIX: UFS(UNIX File system),
- Linux: Extended File system ext2, ext3, ext4,
- Windows: FAT, NTFS etc...,
- MAC OSX: HFS(Hierarchical File system) etc...



Operating Systems Concepts

➤ **Files system Structure:** File system divides disk/partition logically into sectors/blocks, like boot sector/boot block, volume control block/super block, master file table/iNode list block and Data Block.

FILESYSTEM STRUCTURE



1. **Boot Block:** It contains information about booting the system like bootstrap program, bootloader etc...
2. **Super Block:** It contains information about remaining sections, like total no. of data blocks, no. of free data blocks, no. of allocated data blocks etc....
3. **iNode List:** It contains linked list of iNode's of all files exists on a disk.
4. **Data Block:** It contains actual data.

➤ Disk space allocation methods:

- When a file is requesting for free data blocks, then in which manner free data blocks gets allocated for that file and how its information can be kept inside inode of that file is referred as disk space allocation method.

❖ Three disk space allocation methods are there:

1. Contiguous Allocation
2. Linked Allocation
3. Indexed Allocation

Operating Systems Concepts

1. Contiguous Allocation : free data blocks gets allocated for a file in a contiguous manner.

❑ Advantages :

1. Sequential access.
2. Random access .
3. Simple to implement.

❑ Disadvantages :

1. File may not grow(Limitations).
2. **External fragmentation**
-Number of blocks required are available but not contiguous.

❖ Defragmentation :

- Moves files on disk so that maximum contiguous free space is available.

Disk Space Allocation Method:

1. Contiguous Allocation

india.txt

- inode number: 101

- addr of starting data block=0

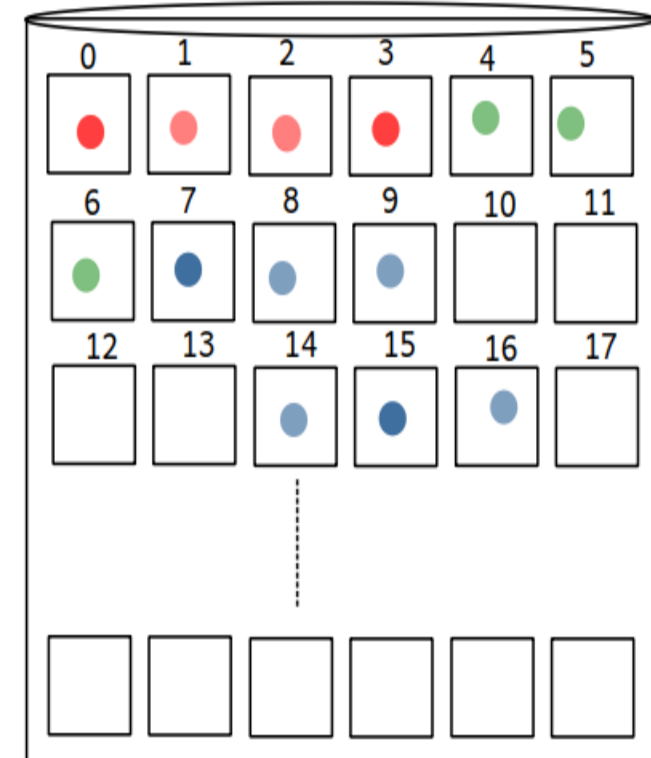
- count=4

pakistan.txt

- inode number: 201

- addr of starting data block=4

- count=3



Operating Systems Concepts

2. **Linked Allocation**: any free data blocks gets allocated for a file in a linked list manner.

❑ **Advantages :**

1. Sequential access.
2. No file grow limit.
3. No external fragmentation.

❑ **Disadvantages :**

1. Slow random access.

❑ **Example : FAT**

SUN

Disk Space Allocation Method:

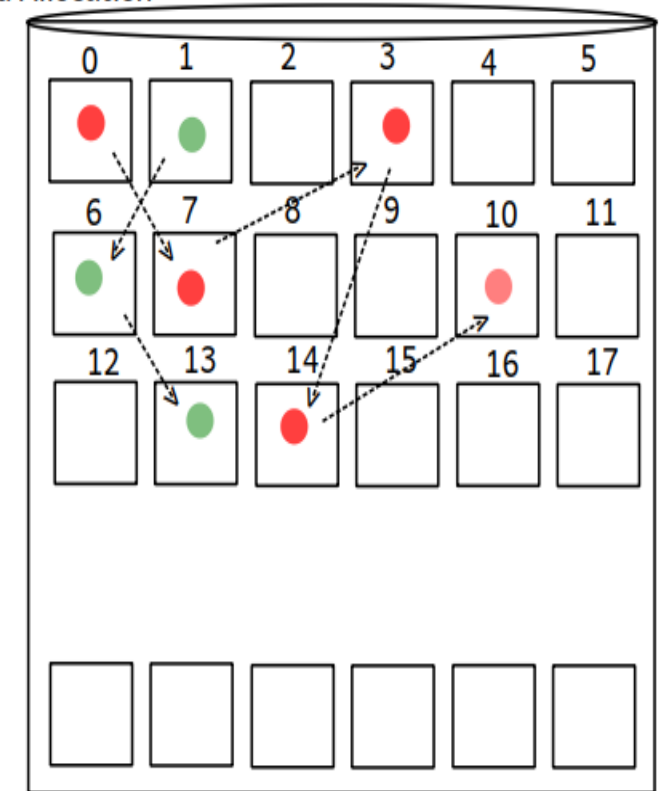
2. Linked Allocation

india.txt

- inode number: 101
- addr of starting data block=0
- addr of end data block=10

pakistan.txt

- inode number: 201
- addr of starting data block=1
- addr of end data block=13



Operating Systems Concepts

3. **Indexed Allocation** : any free data blocks gets allocated for a file, as by maintaining an index data block information about allocated data blocks can be kept inside it.

❑ **Advantages :**

1. Sequential Access.
2. Random Access.
3. No External Fragmentation

❑ **Disadvantages :**

1. File cant grow up to some limit.

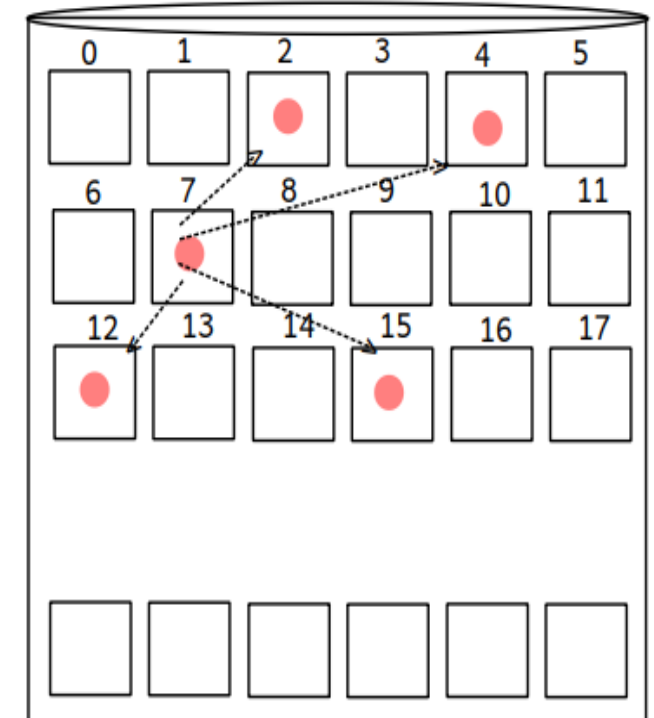
❑ **Example :** UFS , EXT 2 , EXT 3

Disk Space Allocation Method:

3. Indexed Allocation

india.txt

- inode number: 101
- addr of index data block=7



➤ Disk Scheduling Algorithms:

- When system want to access data from a disk, request can sent to disk controller and disk controller accepts one request at a time and complete it.
- There are chances that at a time more than one requests for accessing data from the disk can be made by the processes running in a system, in that case all the requests can **be kept in a waiting queue of the disk maintained by an OS**, and there is need to schedule/select only one request at a time and sent it to the disk controller, to do this there are certain algorithms referred as **disk scheduling algorithms**.

Operating Systems Concepts

1. **FCFS (First Come First Served)**: request which is arrived first gets accepted and completed.
2. **SSTF (Shortest Seek Time First)**: request which is closed to the current position of the head gets accepted and completed.
3. **SCAN**: head keeps scanning the disk from starting cylinder to end cylinder and whichever request came across gets accepted and completed.
4. **C-SCAN (Circular SCAN)**: head scans the disk only in a one direction
5. **LOOK**: this policy can be used either with SCAN/C-SCAN, in this, if there is no request in a waiting queue then movement of the head gets stopped.



Computer Fundamentals and Operating Systems

➤ Memory Technologies

❖ There are four methods by which data can be accessed from the computer memory:

1. Sequential Access: e.g. Magnetic Tape
2. Direct Access: e.g. **Magnetic Disk**
3. Random Access: e.g. RAM Memory
4. Associative Access: e.g. Cache Memory

❖ Magnetic Disk : Hard Disk Drive Structure

- HDD is made up of one or more circular platters arranged like CD rack.
- A Circular platter is made up of non-magnetic substance like aluminum or aluminum alloy, which is coated with a magnetic substance.



➤ Magnetic Disk : Hard Disk Drive Structure

- **HDD** is made up of **one or more circular platters** arranged like CD rack.
- A Circular platter is a made up of **non-magnetic substance like aluminum or aluminum alloy**, which is coated with a magnetic substance.
- Coating of magnetic substance is either from one side to the platter or from both the sides (for increasing its capacity) and hence platter in a magnetic disk may be either single sided platter or double sided platter.
- **Circular platter is divided into the hundred's of concentric rings called as tracks** whereas **each track is divided into thousands of same size of blocks called as sectors**.
- Usually the **size of each sector is 512 bytes**
- There is **one conducting coil referred as head** which is used to access data from the sector i.e. head can read and write data from and into a sector at a time.
- **Head writes and read data sector by sector i.e. block by block**, and magnetic disk is also called as **block device**.



Computer Fundamentals and Operating Systems

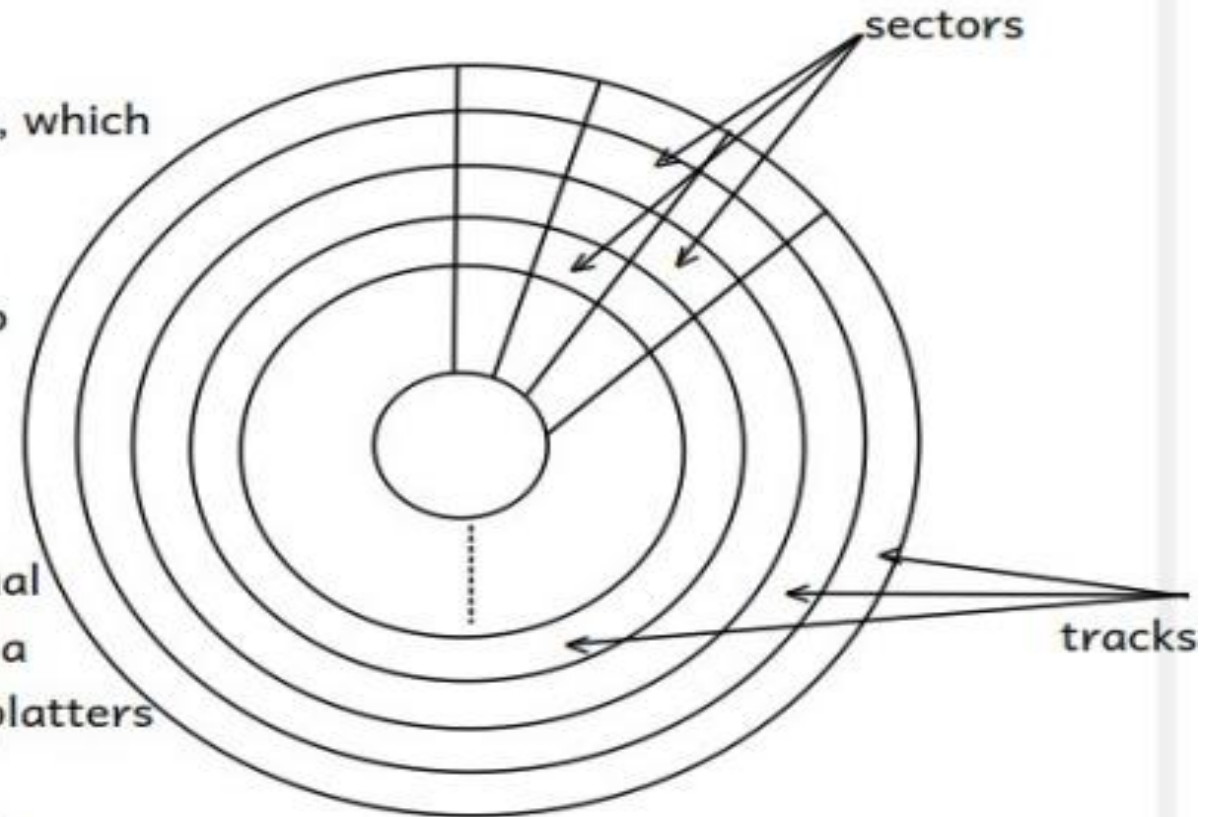
- All the operations like read, write, control etc... in a HDD are controlled by **disk controller**, and hence movement of the head also controlled by it.
- **Seek Time:** time required for the disk controller to move head from its current position to the desired track.
- **Rotational Latency:** after reaching head at desired track, circular platter gets rotated till the head does not comes aligned with the desired sector, and time required for this rotation is referred as rotational latency.
- **Access Time = Seek Time + Rotational Latency.**



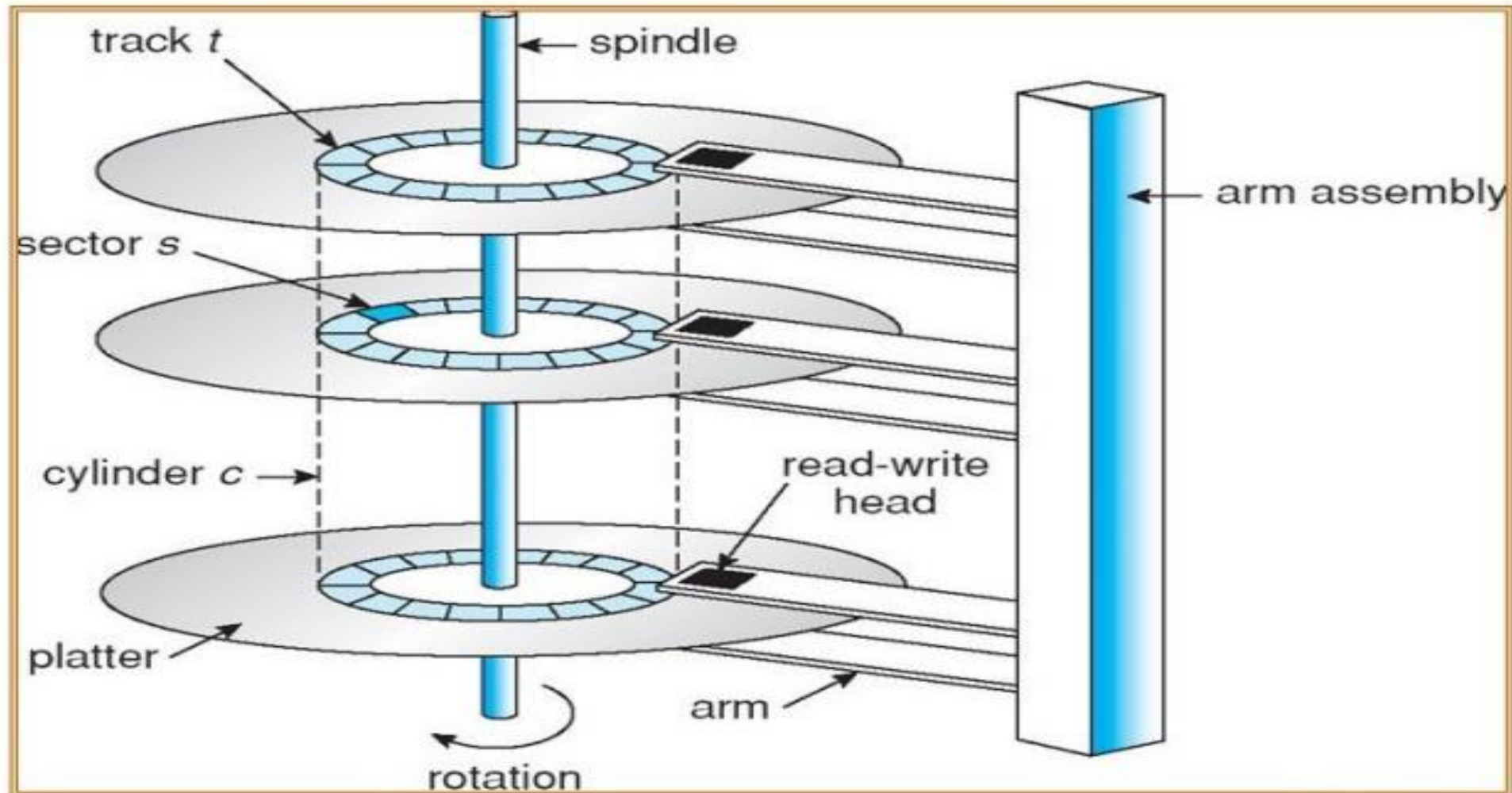
Computer Fundamentals and Operating Systems

Hard Disk Drive:(HDD)

- It contains a "circular platter"(s) made up of non-magnetic material like aluminum or aluminium alloy, which is coated with a magnetic substance.
- Each platter is divided into hundreds of concentric rings called as "tracks" and each track is divided into fixed size of blocks called as "sectors".
- Size of each sector on each track is same usually size of the sector = 512 bytes.
- Cylinder: A cylinder is any set of all of tracks of equal diameter in a hard disk drive. It can be visualized as a single, imaginary, circle that cuts through all of the platters (and both sides of each platter) in the drive.
- Seek Time: it is the time required for a disk controller to move head from its current position to desired cylinder.
- Rotational Latency: Once head moved at desired cylinder, time required to rotate the platter to get aligned with desired sector is called as rotational latency.



Moving-head Disk Mechanism





Thank you!

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