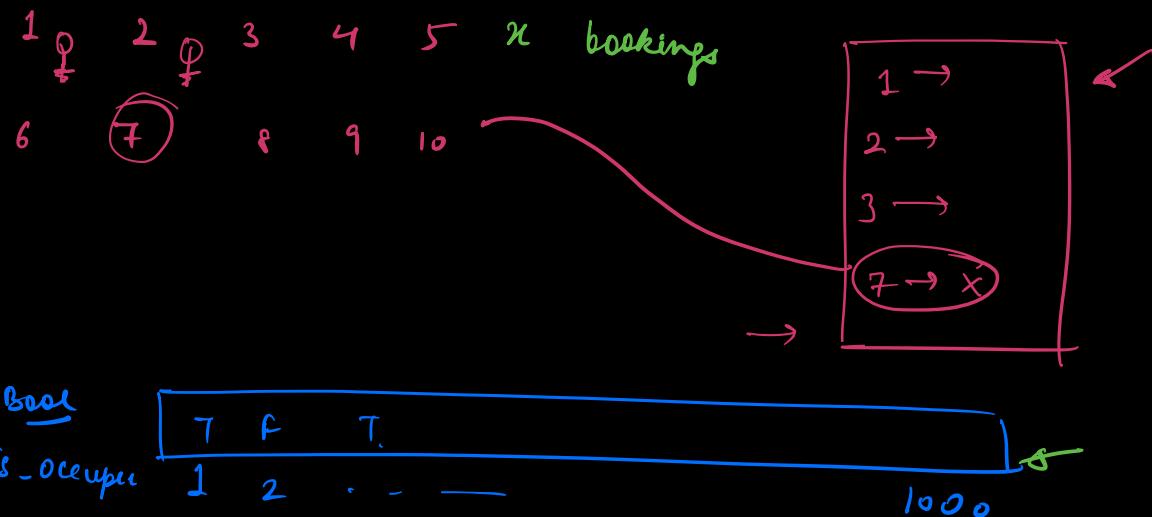
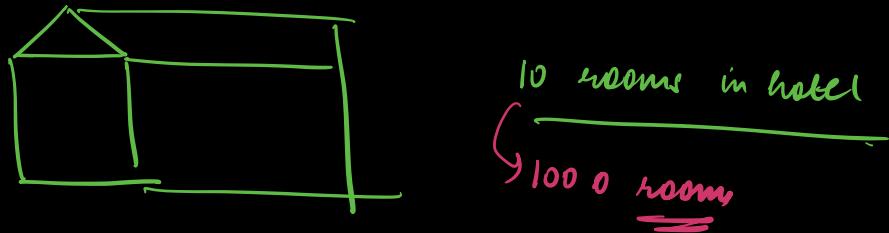
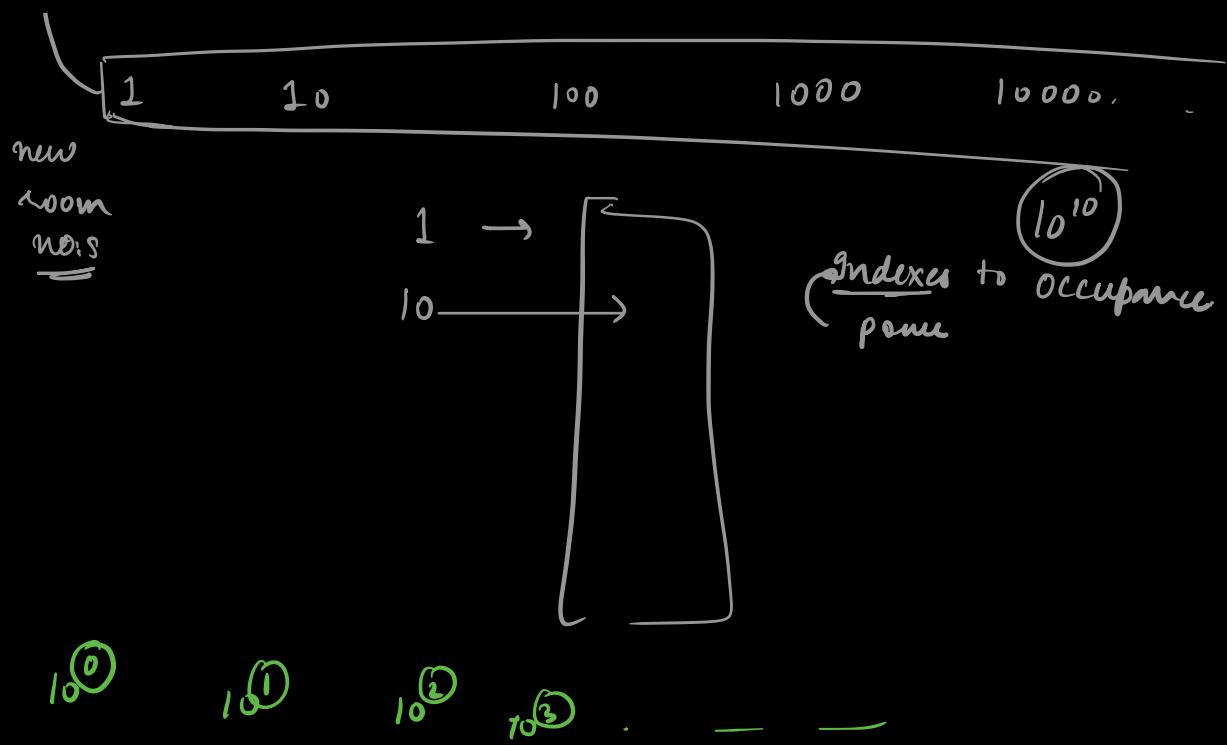


Agenda

- ① Why one more DS?
- ② Implementation
- ③ Python specific
- ④ Problems ↗



[COVID] → Baba



Room No.	Index	Occupier
10^0	0	T
10^1	1	A
10^2	2	T
10^3	3	T

Baba - 2 \rightarrow Name Room after Stars,
Planets

Mercury
 Venus
 Mars
Jupiter

	Occupancy
1 ↗ Mercury	T
2 ↗ Venus	F
3 ↗ Mars	T

*Mapping room names
to numbers*

Just numbers

Index ↗ that could be mapped to values

$\ell[0]$

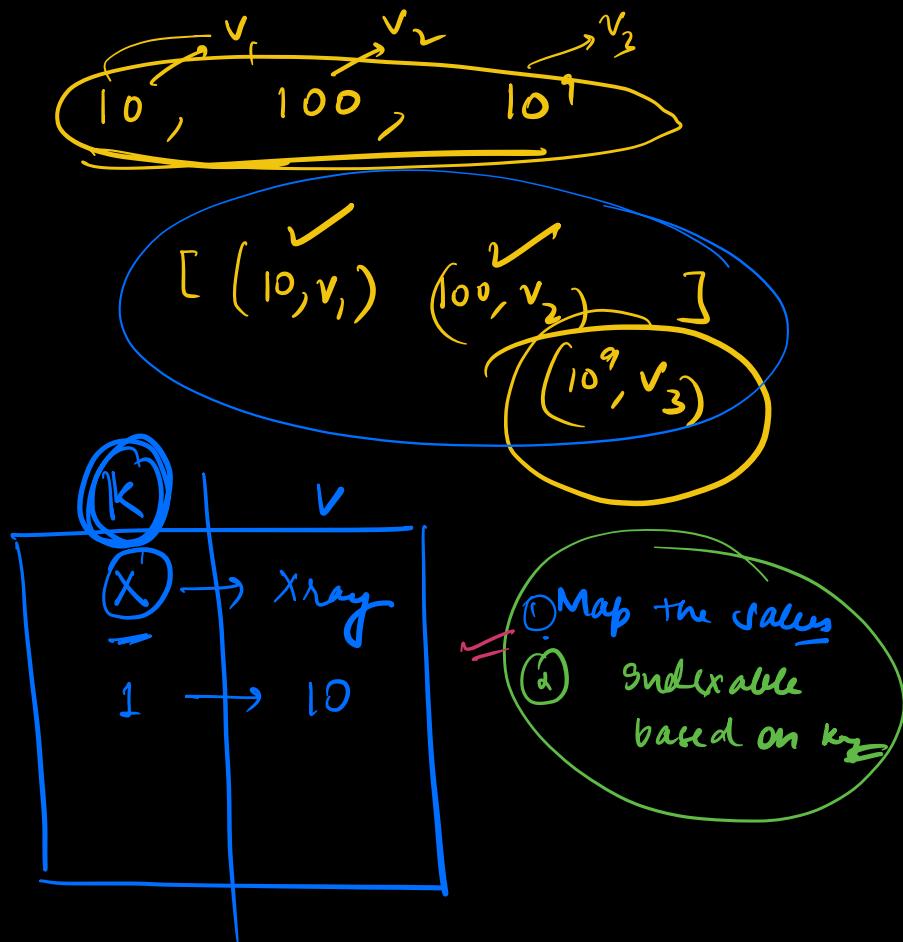
$\underline{\ell[1]} = \underline{v}$

Map keys to some values

Maps [' '] → value

Hash Maps

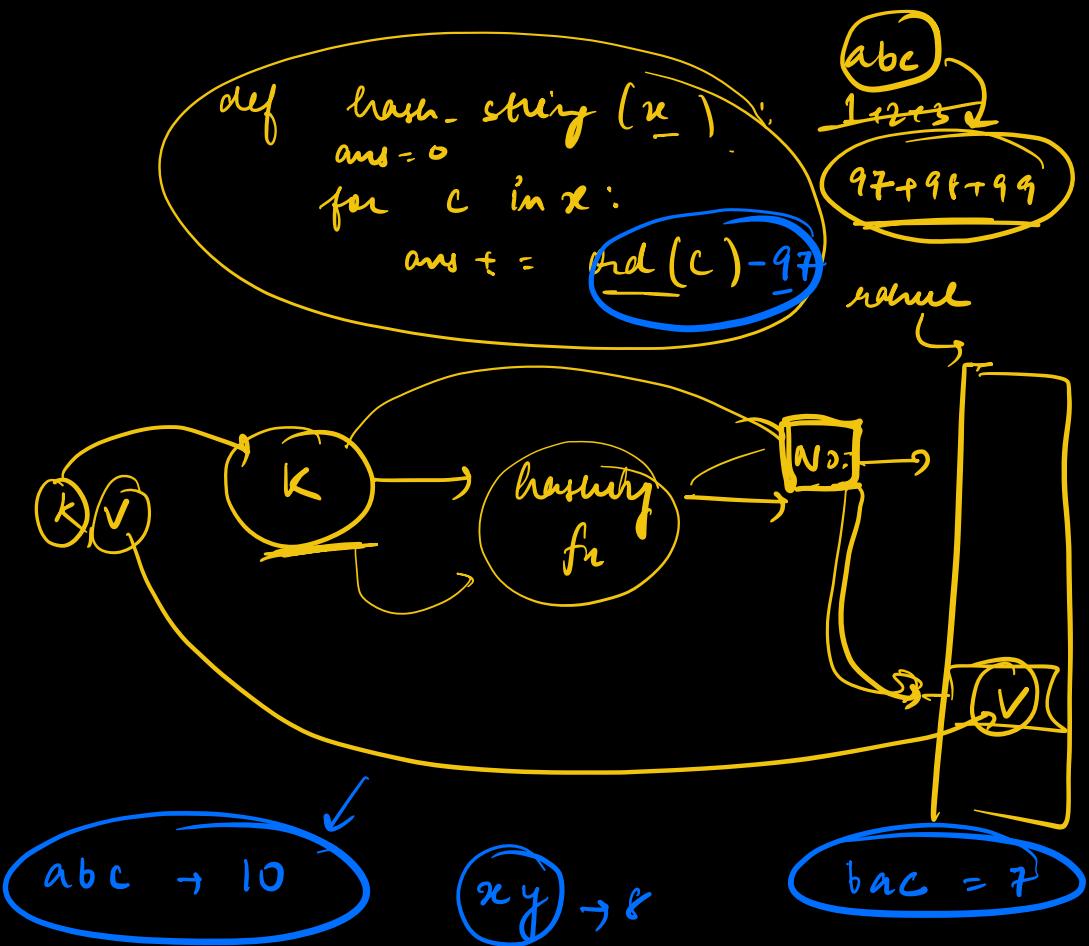
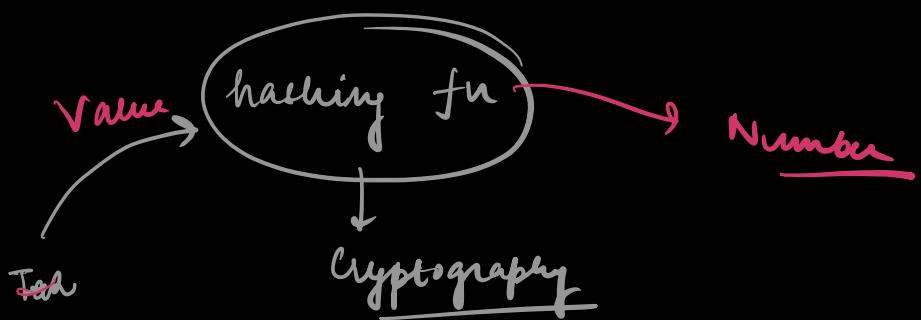
- Space efficient solution to store K-V pair
- Fast access of data ✓
- Check if a key is present or not
- Insert K, V pairs
- Update . . .



Dictionary

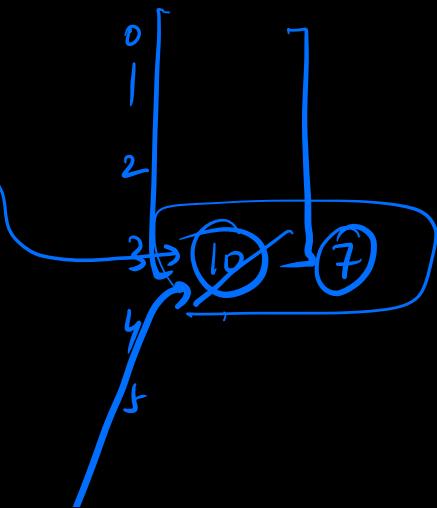
① Mapping

Student Name	Marks
A	10
B	9
C	7
D	8
E	10



$\text{abc} \rightarrow \text{hash}(\text{abc}) \rightarrow 3$
 $0+1+2$

$\text{xy} \rightarrow \text{hash}(\text{xy}) \rightarrow 23+24$
 $= 47$



$$bac \rightarrow \text{hash}(bac) = 1 + 0 + 2 = 3$$

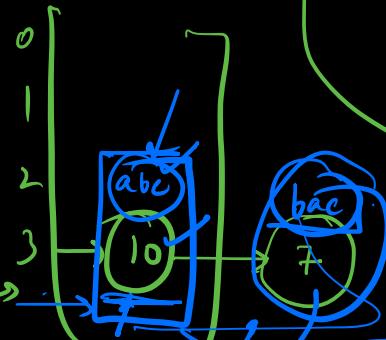
Collisions

For 2 diff keys you get same value

$$abc \rightarrow 3$$

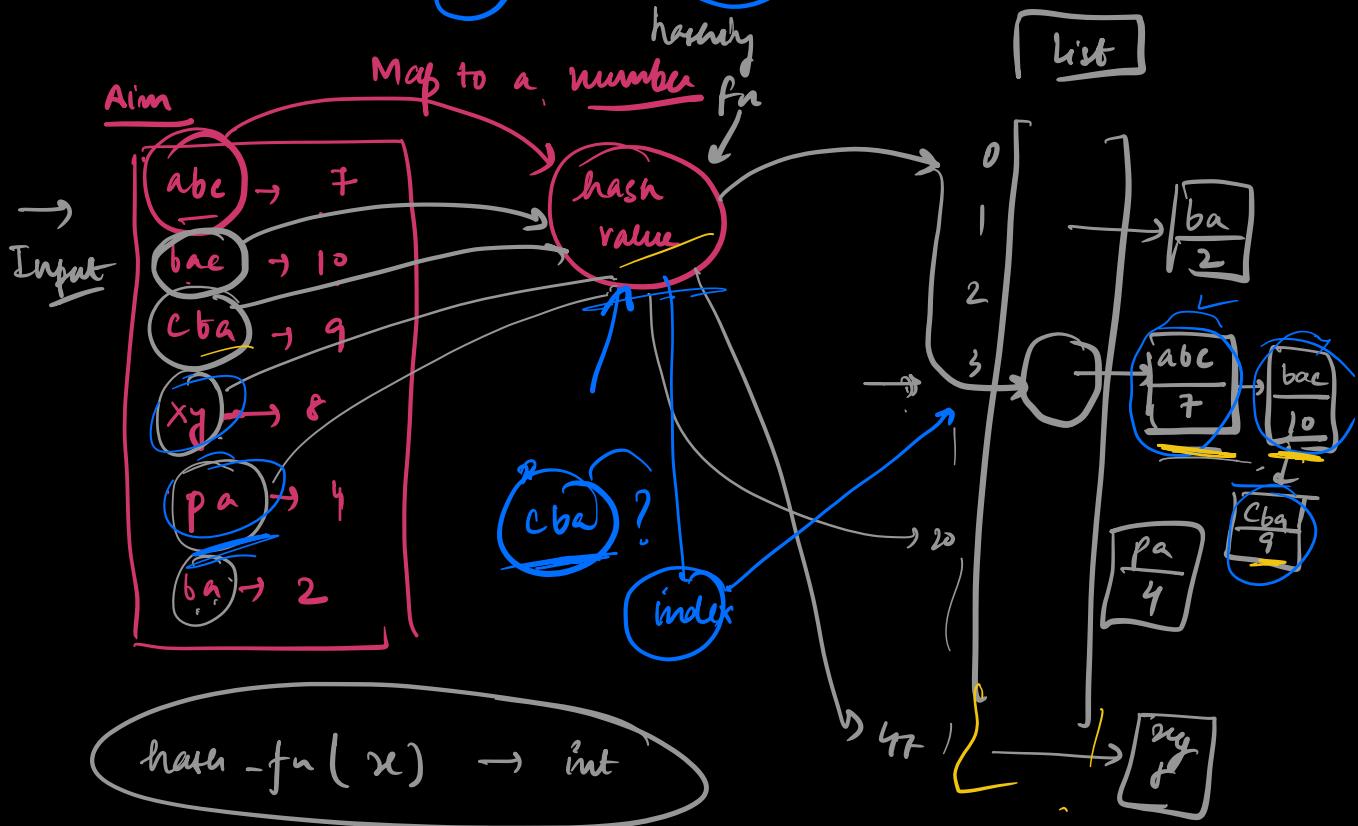
$$bac \rightarrow 3$$

$$\begin{array}{l} \text{get } abc \\ \text{get } bac \end{array}$$



Chaining

$$\begin{array}{l} \text{get } bac \\ k, v \end{array}$$



Need \rightarrow space efficient
 Yet in $O(1)$ time] \rightarrow improving
 $\underline{\text{H.F}}$ \rightarrow unique

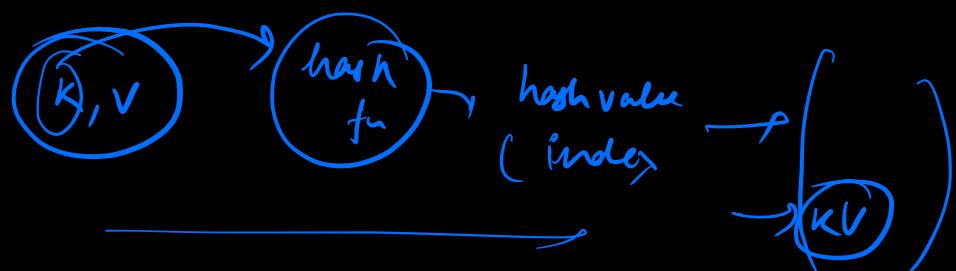
Size of list on right depends on the range of hash values!

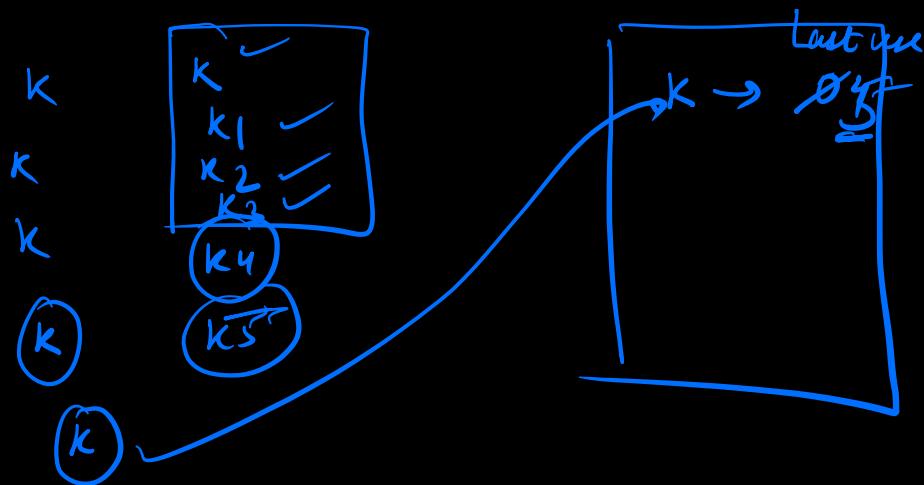
W.l.o.g \rightarrow $O(n)$ \rightarrow All keys are mapped to same \rightarrow hash

Access / fetch time \rightarrow $O(1)$

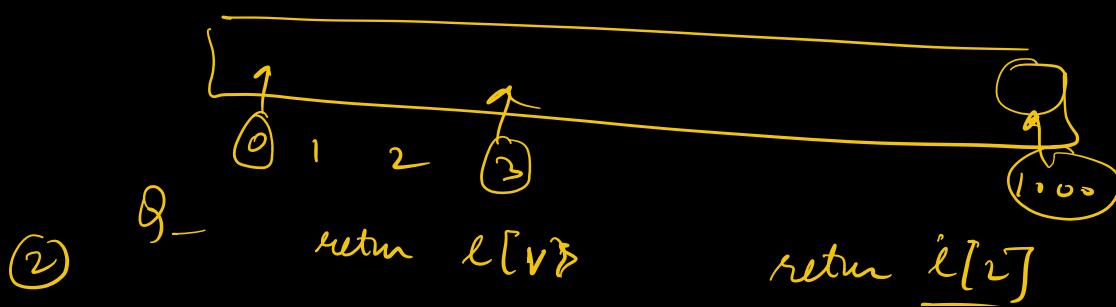
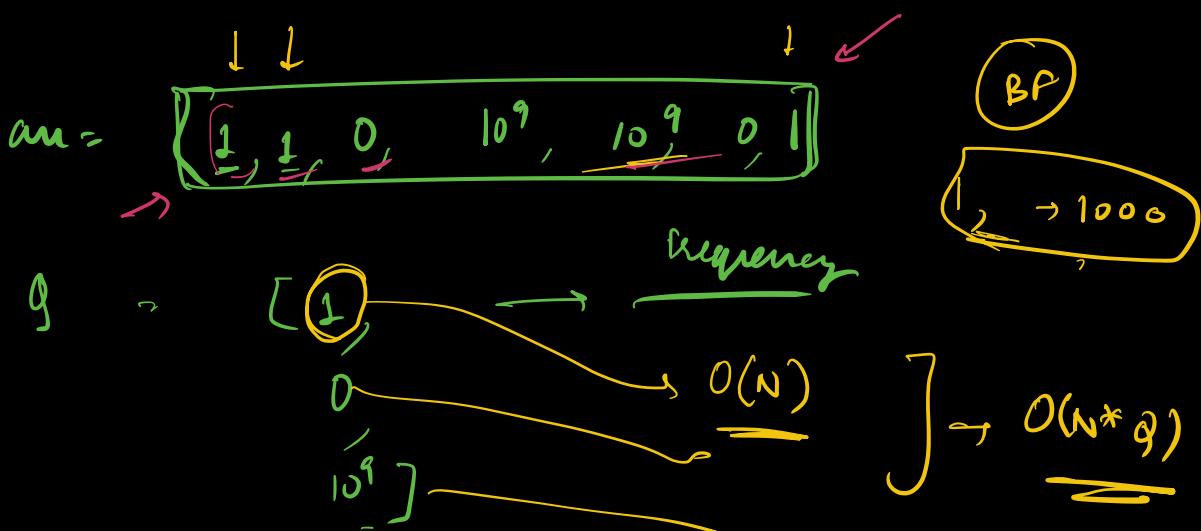
Worst case
 $\underline{\mathcal{O}(n)}$

Avg Case





Q Given an array. You have Q queries
Return the count of elements in array
for each query



val	freq
1	2
0	2
10^9	2

return dict[0]

def solve(arr, queries)

```

freq = {}
for val in arr:
    if val in freq:
        freq[val] += 1
    else:
        freq[val] = 1

```

$O(N)$

$O(Q)$

$O(N+Q)$

```

for query in queries:
    if query in freq:
        return freq[query]
    else:
        return 0

```

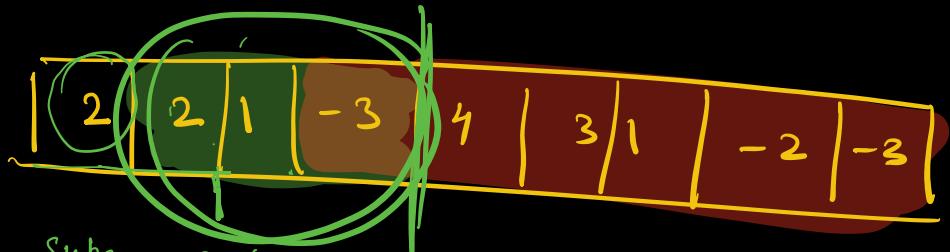
$ans = [1]$

$Q = [1, 1, \dots, 1]$
 $\underbrace{1000}$

$ans = [1, 2, 3, 1, \dots, 1]$
 $\underbrace{1000}$
 $Q \rightarrow 2$

Break → 5 min

Q Given an array . check if there exists a subarray with $\text{sum} = 0$



Subarray = 2, 4

$$\text{sum} = (4) \rightarrow 5$$

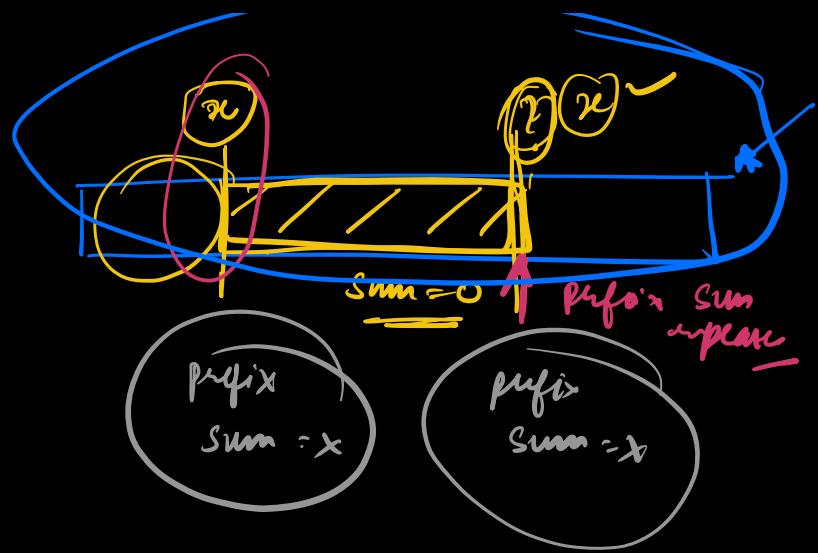
$O(N^2)$
 $O(1)$

$i \rightarrow \text{ending index}$
for i in range(N):
 startly ← for j in range($0, i+1$) ←
 if $\text{pref}[i] - \text{pref}[j] == 0$

At index i is there any subarray
ending at index i with $\text{sum}=0$

|

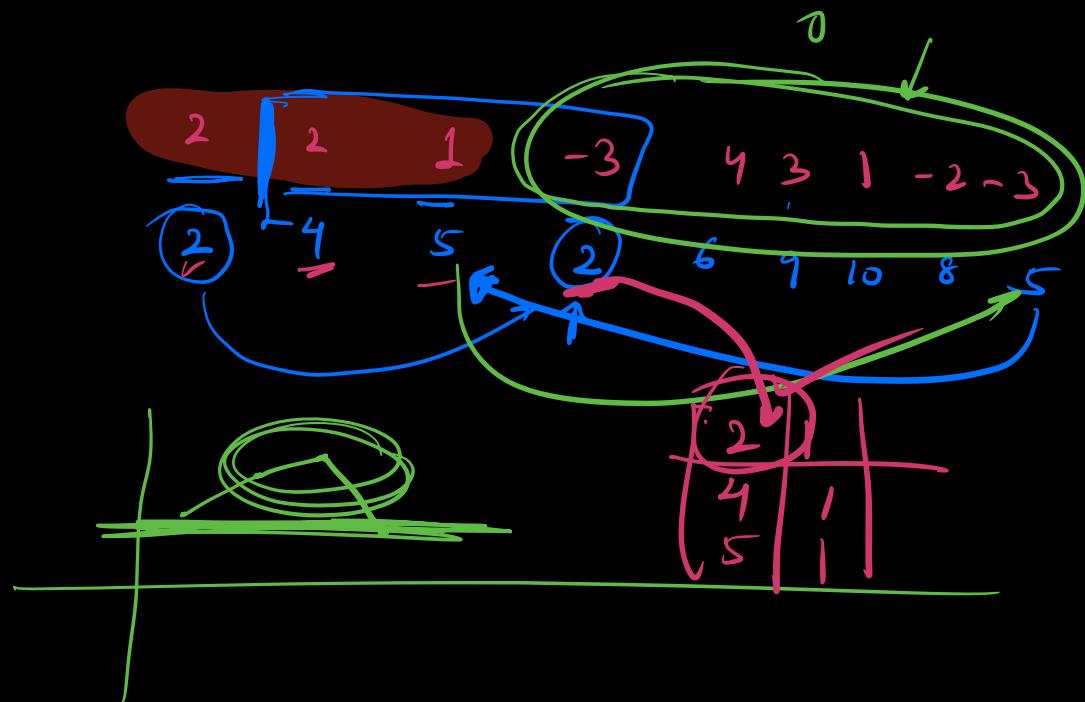
→



Repeats itself \Rightarrow there

exist
a subarr.
with sum 0

a



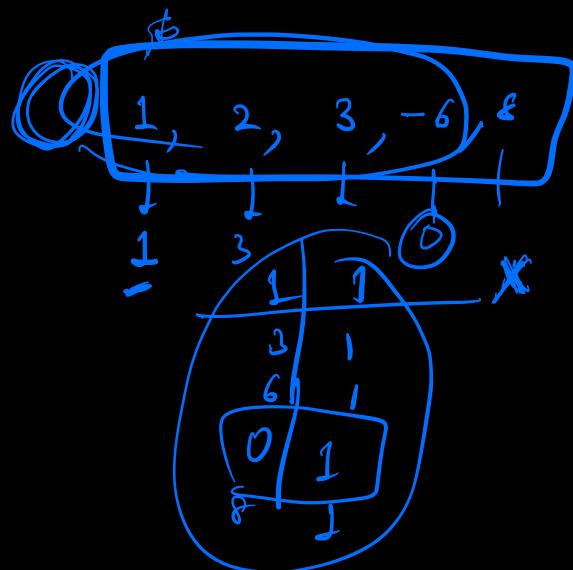
if a prefix sum value is repeated
or not

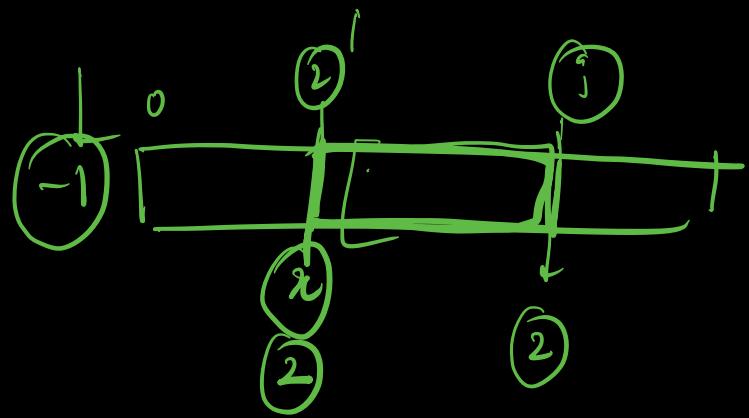
```

dict[0] = 1
for i in range(n):
    sum += arr[i]
    if sum in dict:
        return True
    dict[sum] = i
return False

```

TC $\rightarrow O(n)$
 SC $\rightarrow O(n)$





$\text{dict}[2] = \circledast 1$

$\text{dic}[\text{sum}] = [\circledast 1 \rightarrow j, 0 \rightarrow j]$ → subarray

$\boxed{\text{dic}[0] = \circledast -1}$