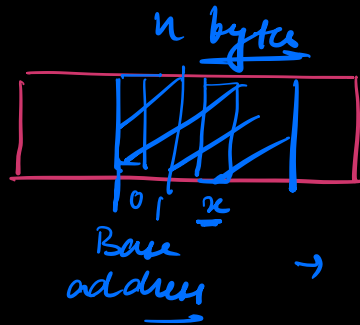


Linked Lists Basics :

Arrays

- Contiguous memory
- Sequential
- Linear

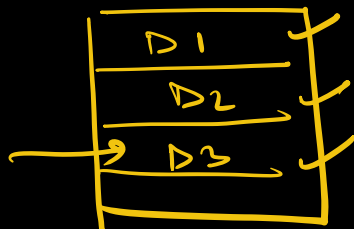
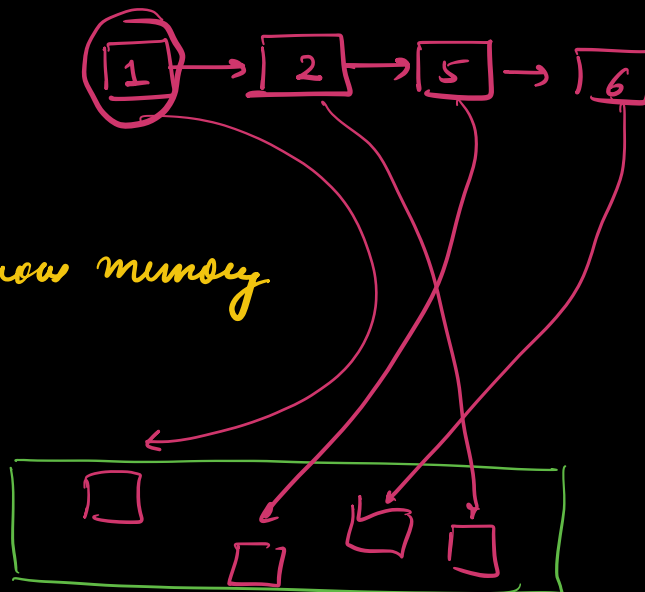
n bytes

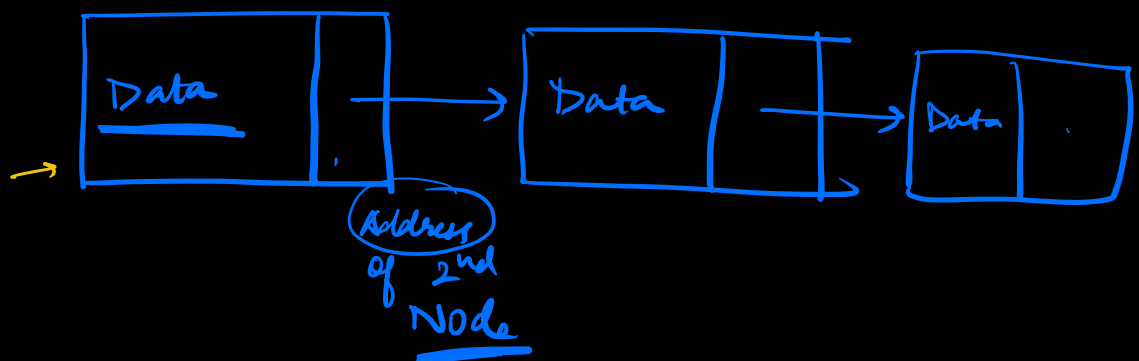
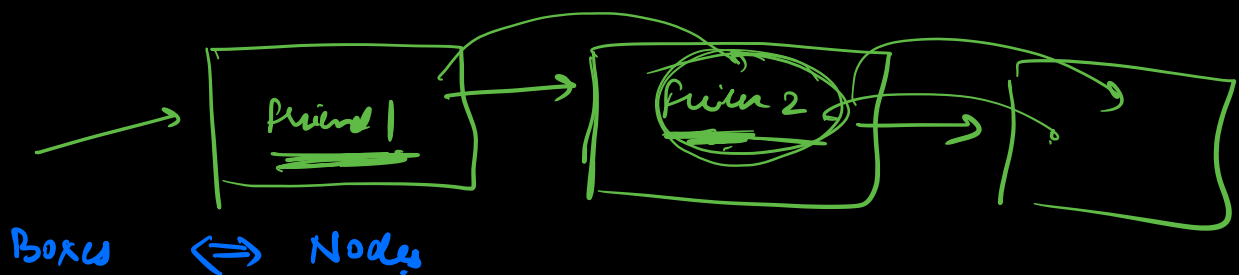
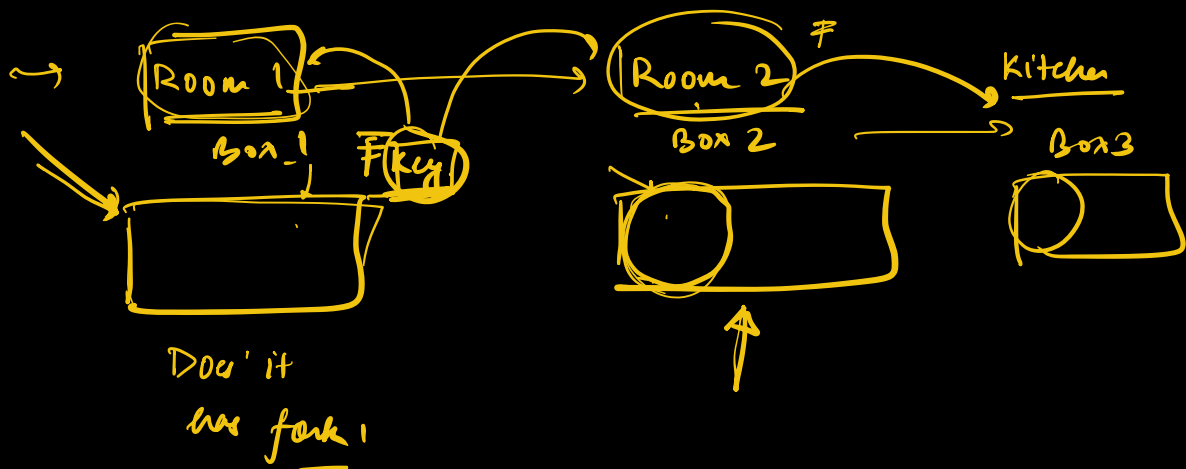


→ constant
time access
 $O(1)$

Linked List

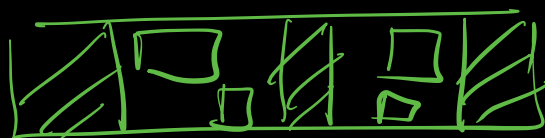
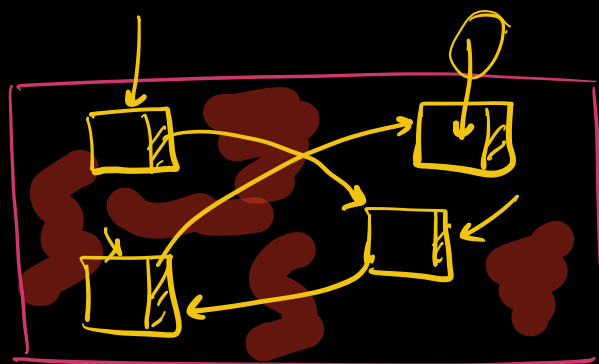
- Linear
- Sequential
- not contiguous memory allocation.





Singly
Linked
List
➤

Elastic Search



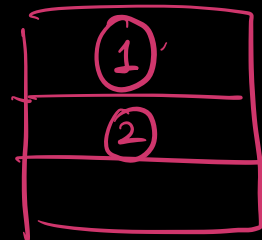
- ① LL are the building block of many data structure \Rightarrow Stacks / Queues / Trees / Trails

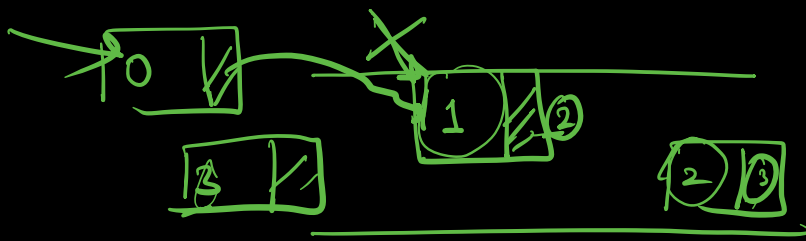
collisions are handled by LL in hash

map

LL + Hashmap

Array

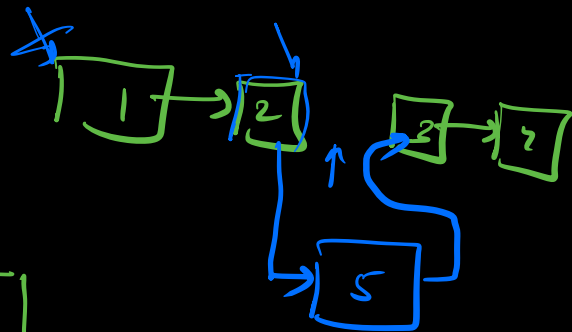
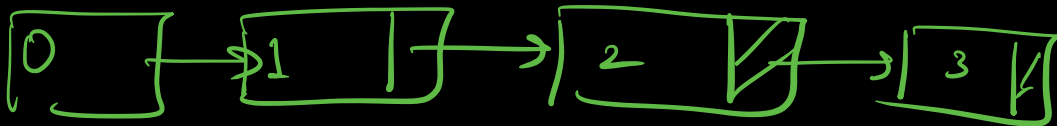




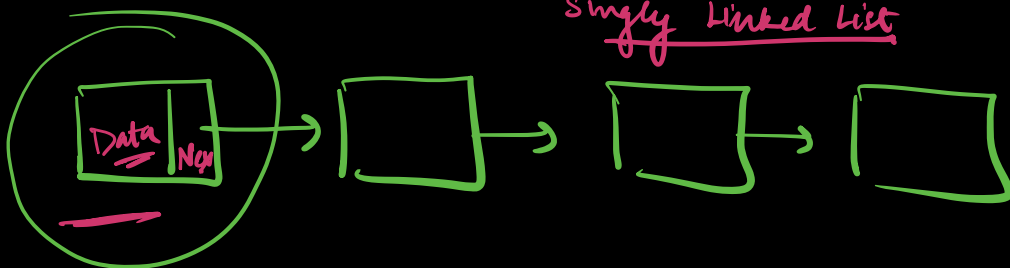
$O(N)$

Time

Insert an element in front.



Singly Linked List



Structure of
Node
of a LL

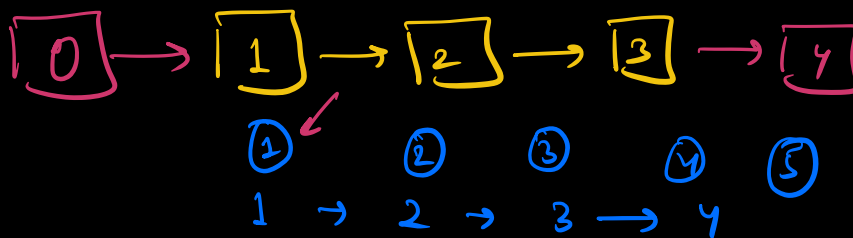
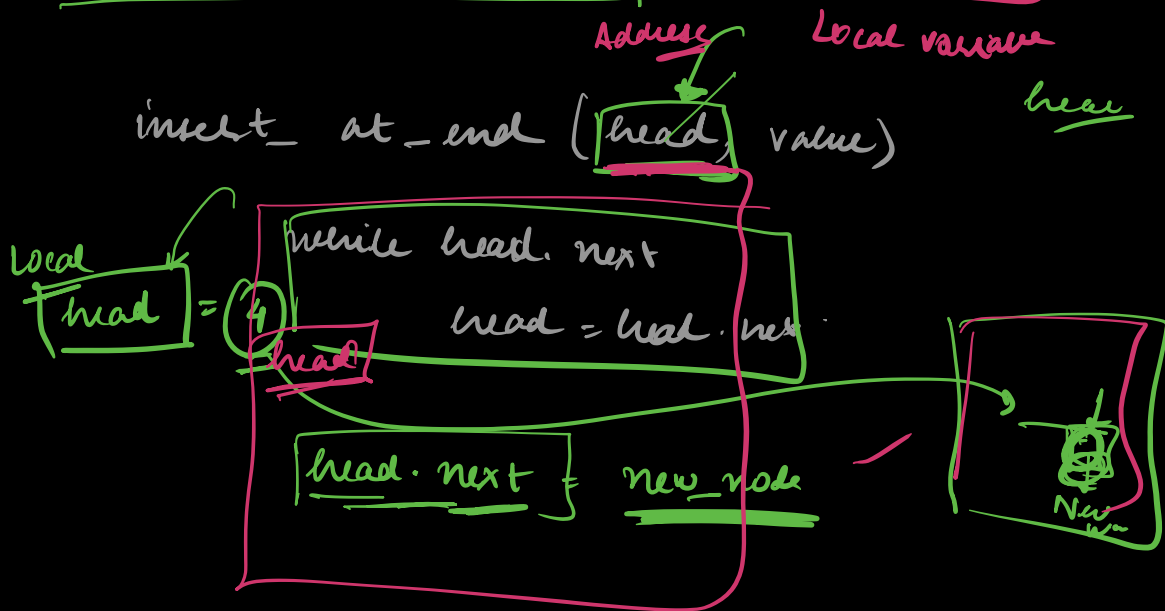
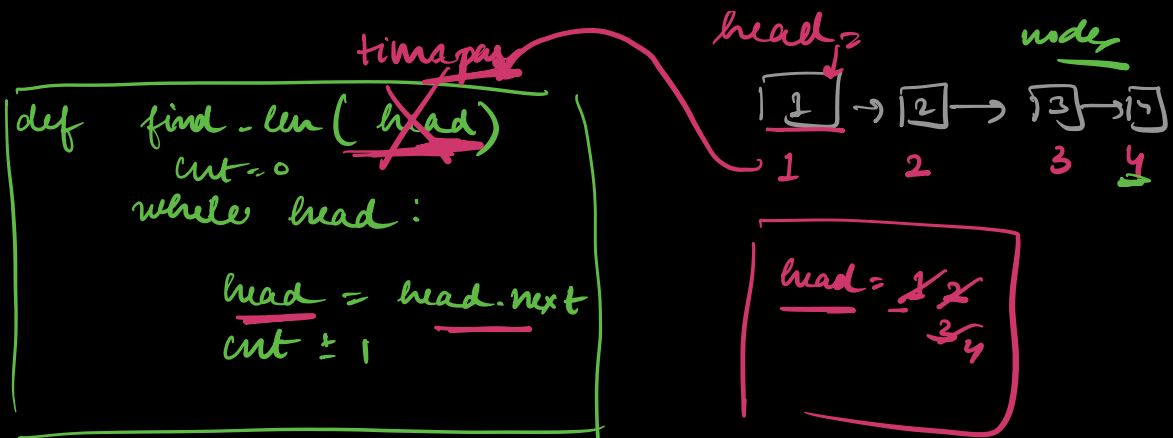
class Node:

def __init__(self, data):

✓ self.data = data :

✓ self.next = None

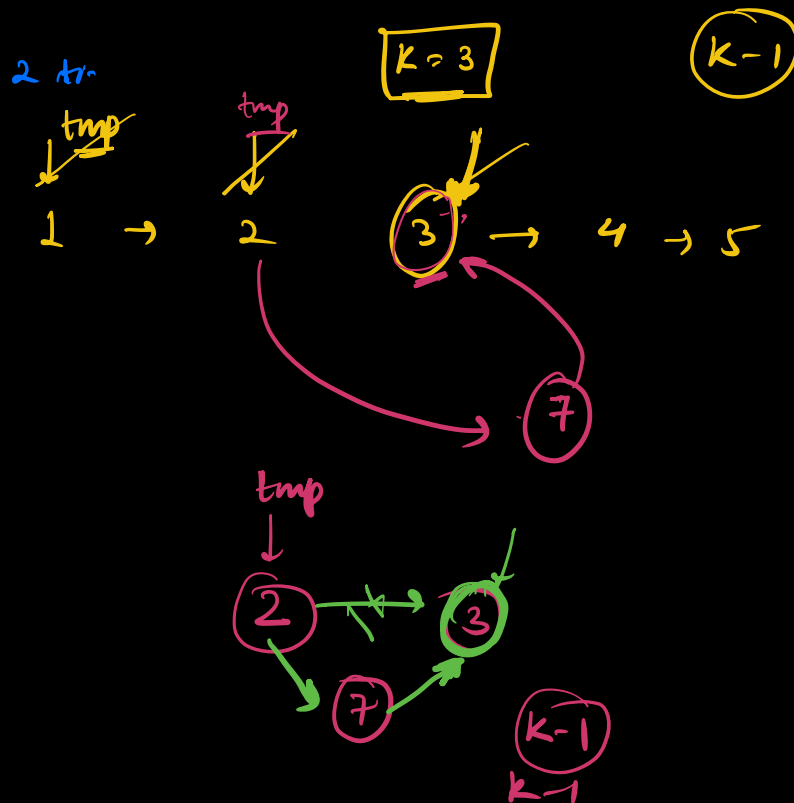
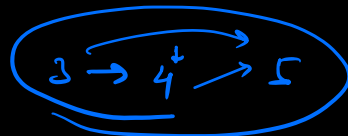
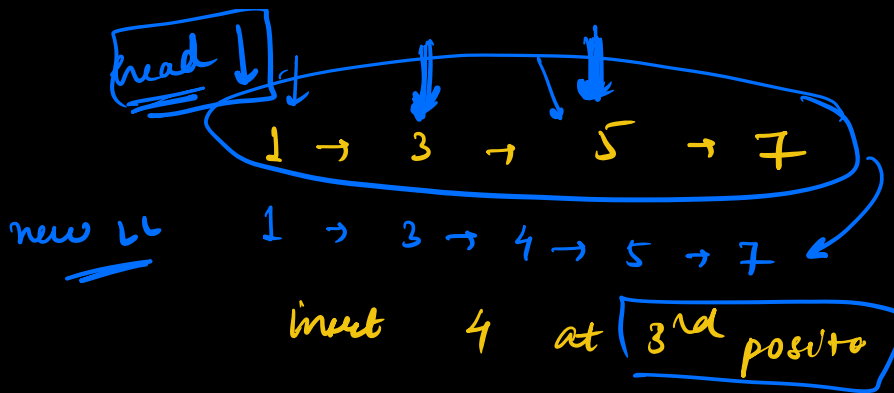
→ Address
of the next



Insert a node at K^{th} position

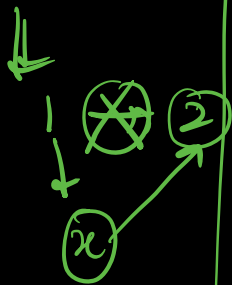
head, value, K

return the new head



① Move head ptr $\frac{k-1}{k-2}$ times.

② $\text{new_node.next} = \text{tmp.next}$
 $\text{tmp.next} = \text{new_node}$



def insert at k (value, head, k):

tmp = head

for i in range(k-2):

tmp = tmp.next

new_node = Node(value)

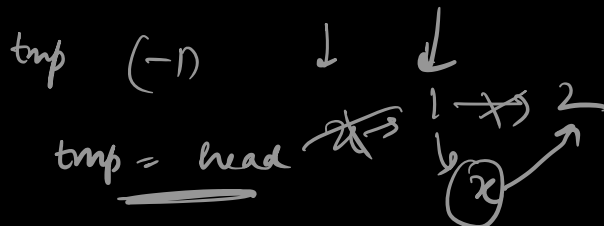
new_node.next = tmp.next

tmp.next = new_node

k=2

k=1

- ① kth position
- ② Last position (in the end.)
- ③ Insert at beginning → k=1



Edge case

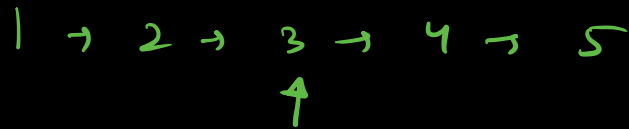
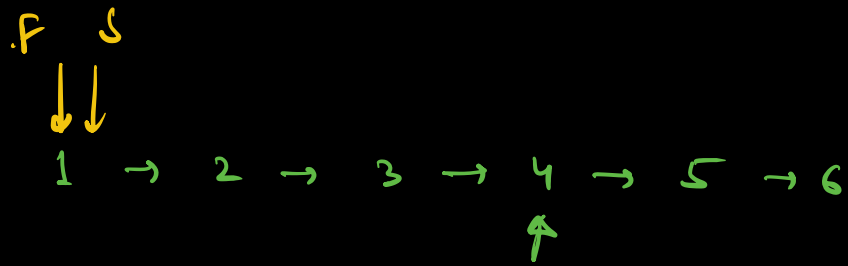
if head == None
new_node
return new_node

if k == 1:

insert at begin

new_node.next = head

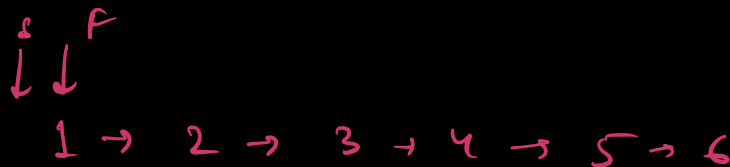
return new_node



① Get len

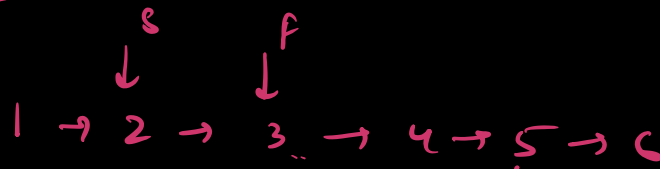
② Go to middle
Print it.

Fast and slow ptr.
Hare and Tortoise

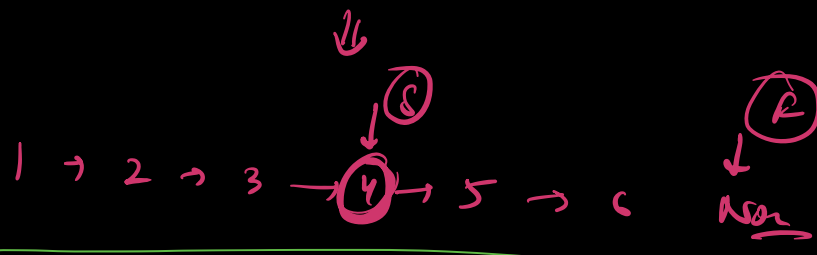
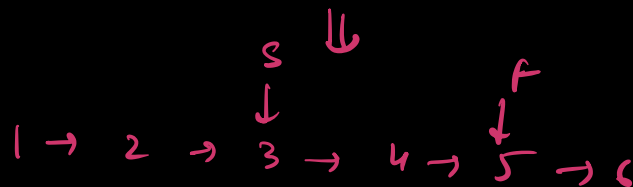


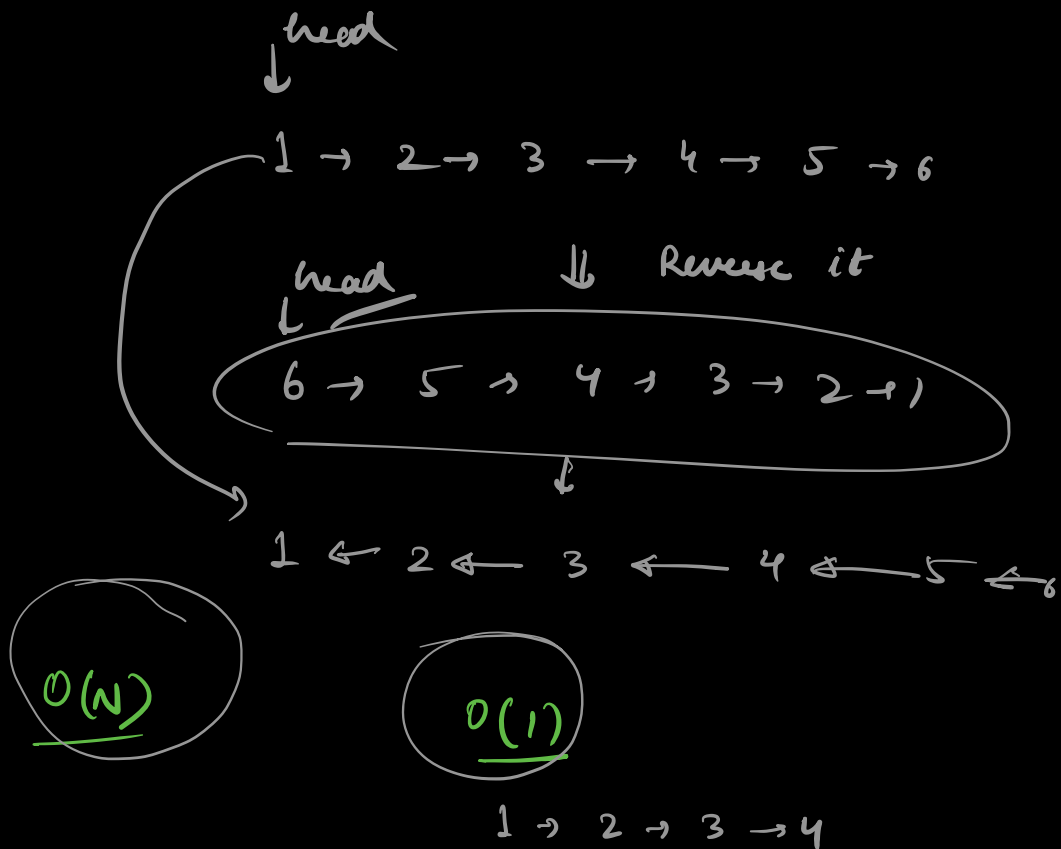
F = F.next.next

S = S.next



middle





Finding length \rightarrow Recursion

```
def size(head):
    if head == None:
        return 0
    return 1 + size(head.next)
```

