

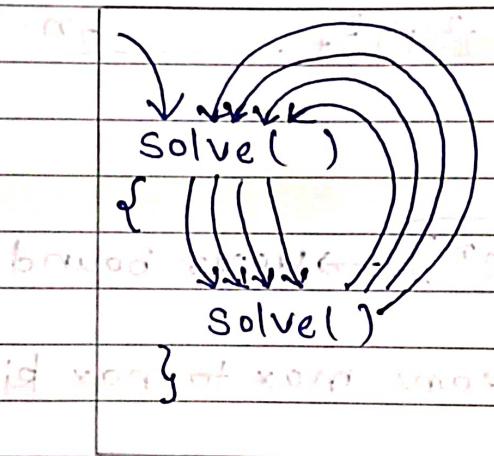
Recursion is the game of trust. So, trust the recursion.

09/10/2023
Monday

RECURSION

- Bookish language → when a function calls itself directly/indirectly.

- In-Depth → Solution of bigger problem depends on Solution of chotti problem of same type.



→ Solve() function is calling it again & again jab tak base case nahi mil jata.

- Recursion be like, 1 case tum solve kar do, baki mai sambhaal lunga.

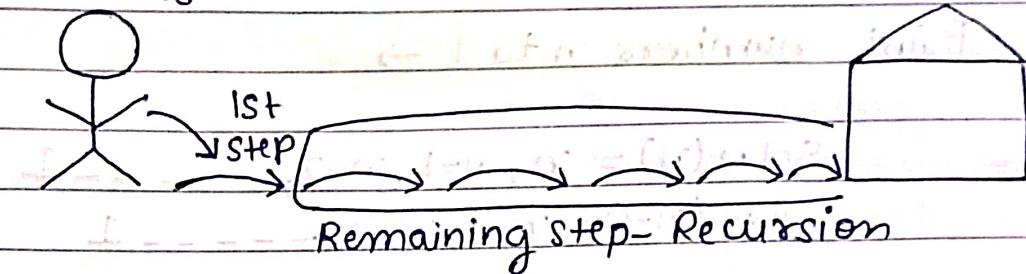
How Recursion works?

- * Assume, ek person ko place A se place B par Jana hai, jiske liye 6 steps travel karna pdta hai.

So, 1 step person chalega and remaining 5 steps recursion chalega.

Source

Destination



$$\text{Step}(6) = 1 + \text{step}(5)$$

(1 step)

Bigger
problem

if top
resolution

chotti problem
handled by
recursion

Single Step
travelled by person

So, here, OR kitna chalna hai \rightarrow function calls

Kab zukna hai \rightarrow Base case

Left/right janahai \rightarrow Processings

Function recursive calls, base case and processings
collectively make a recursive function.

* calculate $2^n \rightarrow$

$$\text{Solve}(n) = 2^n$$

$$\text{Solve}(n-1) = 2^{n-1}$$

$$2^n = 2 * 2^{n-1}$$

$$\text{Solve}(n) = 2 * \text{Solve}(n-1)$$

Bigger
problem

chotti
problem

* Point numbers n to 1 →

$$\text{Solve}(n) = n, n-1, n-2, \dots, 1$$

$$\text{Solve}(n-1) = n-1, n-2, \dots, 1$$

$$\text{Solve}(n) = n, [n-1, n-2, \dots, 1]$$

Solve(n-1)

$$\text{Solve}(n) = n, \text{Solve}(n-1)$$

Bigger
problem

Chotti
problem

* Find factorial →

$$\text{solve}(5) = 5 \times 4 \times 3 \times 2 \times 1$$

$$\text{solve}(4) = 4 \times 3 \times 2 \times 1$$

$$\text{solve}(5) = 5 \times \text{solve}(4)$$

Bigger
problem

chotti
problem

• Recursion code comprises →

(1) Base Case (mandatory)

(2) Recursive call (mandatory)

(3) Processing (optional)

LL

* Calculate factorial →

$$\text{fact}(n) = n * (n-1) * (n-2) * \dots * 3 * 2 * 1$$

$$\text{fact}(n-1) = (n-1) * (n-2) * \dots * 3 * 2 * 1$$

$$\text{fact}(n) = n * \underbrace{(n-1) * (n-2) * \dots * 3 * 2 * 1}_{\text{fact}(n-1)}$$

fact(n-1)

$$\boxed{\text{fact}(n) = n * \text{fact}(n-1)}$$

Code → #include <bits/stdc++.h>

using namespace std;

int fact(int n){

// base case

if (n==0 || n==1) return 1;

// recursive call

int recAns = fact(n-1);

// processing

int finalAns = n * recAns;

return finalAns;

}

int main(){

cout << fact(7) << endl;

O/P → 5040

Behind the scenes,

fact(4) → [n=4]

if (n==0 || n==1) → false
return 1;

int secAns = fact(n-1);

int finalAns = n * secAns;

return finalAns;

fact(3) → [n=3]

if (n==0 || n==1) → false
return 1;

int secAns = fact(n-1);

int finalAns = n * fact(n-1);

return finalAns;

finalAns = 24

(return)

6

2

fact(2) → [n=2]

if (n==0 || n==1) ←
return 1; ↑ True

int secAns = fact(n-1);

int finalAns = n * secAns;

return finalAns;

fact(1) → [n=2]

if (n==0 || n==1) ← → false
return 1;

int secAns = fact(n-1);

int finalAns = n * secAns;

return finalAns;

function call stack →

1. fact(4)
2. fact(3)
3. fact(2)
4. fact(1)

fact(1)

fact(2)

fact(3)

fact(4)

main()

1. fact(4)

2. fact(3)

3. fact(2)

4. fact(1)

5. main()

Entries get removed after returning, some value and stack gets empty.

* Print counting n to 1 → ~~postorder traversal~~

Print(n) = $n, n-1, n-2, \dots, 2, 1$

Print($n-1$) = $n-1, n-2, \dots, 2, 1$

($\text{Print}(n) = n, n-1, n-2, \dots, 2, 1$)

↓
Print($n-1$)

[Print(n) = $n, \text{print}(n-1)$]

Code → `#include <bits/stdc++.h>`

`using namespace std;`

`void print(int n){`

// base case

`if(n==0) return;`

// processings

`cout << n << " ";`

// recursive call

`print(n-1);`

`int main(){`

`print(5);`

`}`

O/P → 5 4 3 2 1

Note - To understand recursion, you need to first understand the recursion.

* Tail and Head Recursion →

If recursive relation comes after processing, then it is tail recursion.

If recursive relation comes before processing, then it is head recursion.

* Calculate $n^m \rightarrow$

$$\text{power}(n, m) = n^m$$

$$\text{power}(n, m-1) = n^{m-1}$$

$$\text{power}(n, m) = n * \boxed{n^{m-1}}$$

$$\text{power}(n, m-1)$$

$$[\text{power}(n, m) = n * \text{power}(n, m-1)]$$

Code → #include <bits/stdc++.h>

```
using namespace std;
```

```
int power(int n, int m){
```

```
    if (m == 0) return 1; // base case
```

```
    int pow = power(n, m-1); // recursive relation
```

```
    return n * pow;
```

```
}
```

```
int main(){
```

```
    cout << power(5, 3) << endl;
```

```
}
```

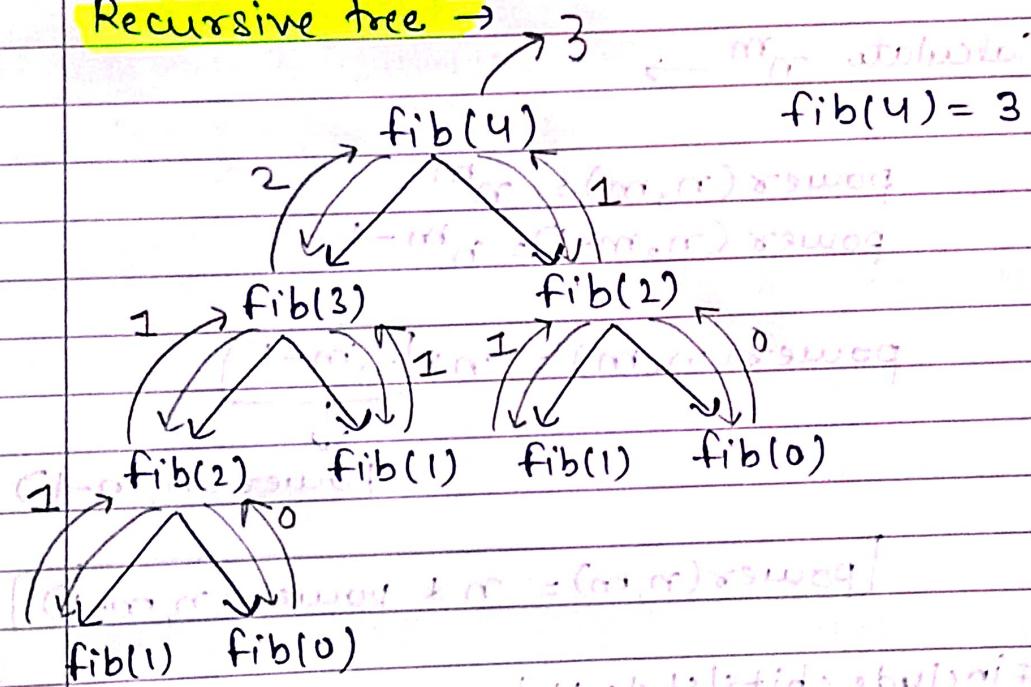
O/P → 125

* Fibonacci Series →

0 1 2 3 5 8 13 21 34 ...

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$$

Recursive tree →



Code → #include <bits/stdc++.h>

using namespace std;

```
int fib(int n){
```

// base case: if n=0 or n=1 return n;

// recursive call

```
int ans = fib(n-1) + fib(n-2);
```

return ans;

}

```
int main(){
```

cout << fib(8) << endl; → O/P = 21

}

* find Sum →

$$\text{sum}(n) = n + (n-1) + (n-2) + \dots + 2 + 1$$

$$\text{sum}(n-1) = (n-1) + (n-2) + \dots + 2 + 1$$

$$\boxed{\text{sum}(n) = n + \text{sum}(n-1)}$$

Code → #include <bits/stdc++.h>

using namespace std;

int sum(int n){

 // base case

 if(n==0) return 0;

 // recursive call

 int ans = n + sum(n-1);

 return ans;

}

int main(){

 cout << sum(5) << endl;

O/P → 15

Recursive tree →

6 → capitalise whenever

15 ↓

sum(5)

3+3=6

↓ sum(4)

1+2=3

↓ sum(3)

↓ sum(2)

↓ sum(1)