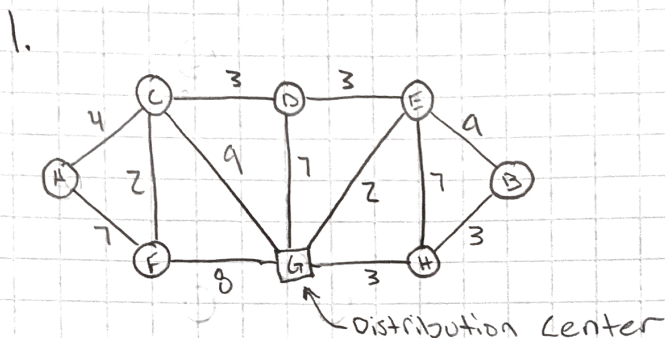


HW # 5

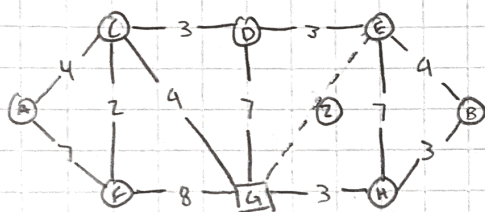
Problem 1.

A. * Best algorithm to find fastest route would be Dijkstra's algorithm.

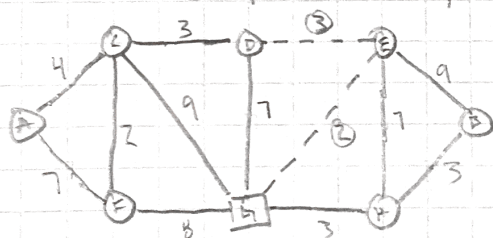
Graph Explanation:



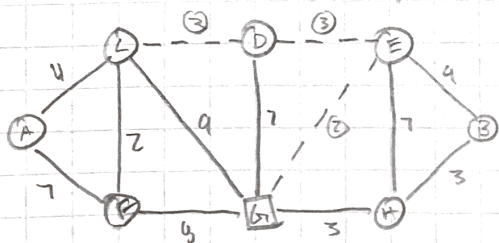
2. * Shortest path is 2, $G \Rightarrow E$



3. * Shortest path is 5, $G \Rightarrow E \Rightarrow D$



4. * Shortest Path is 8, $G \Rightarrow E \Rightarrow D \Rightarrow C$

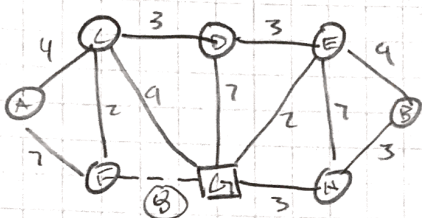


A. continued

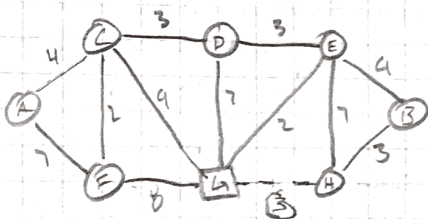
5. * Shortest path is 12, $G \Rightarrow E \Rightarrow D \Rightarrow C \Rightarrow A$



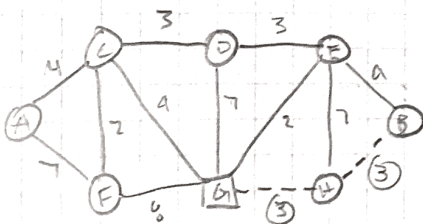
6. * Shortest path is 8, $G \Rightarrow F$



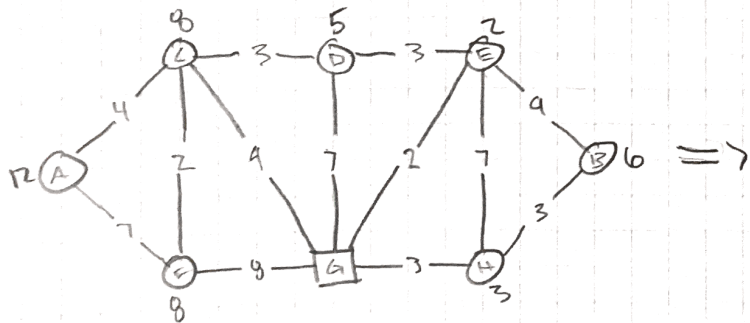
7. * Shortest path is 3, $G \Rightarrow H$



8. * Shortest path is 6, $G \Rightarrow H \Rightarrow B$



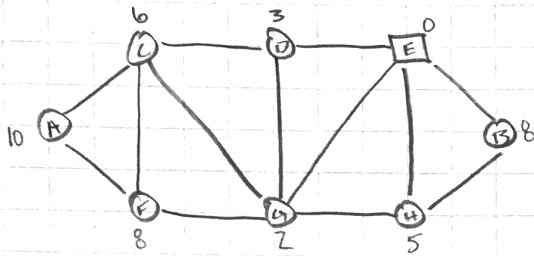
9. * Shortest path to all from G



Vertex	Distance	Parent vertex	Shortest path
$\rightarrow G$	0		G
A	12	C	$G \Rightarrow E \Rightarrow D \Rightarrow C \Rightarrow A$
B	6	H	$G \Rightarrow H \Rightarrow B$
C	9	D	$G \Rightarrow E \Rightarrow D \Rightarrow C$
D	5	E	$G \Rightarrow E \Rightarrow D$
E	2	G	$G \Rightarrow E$
F	8	G	$G \Rightarrow F$
H	3	G	$G \Rightarrow H$

B.

1. * If center is E, graph with shortest paths



* Algorithm for finding optimal solution:

Optimal-Location(City, t):

OptimalLoc $\leftarrow -1$

Sum $\leftarrow \infty$

// find optimal solution

For $i \leftarrow 0$ to $t-1$:

LocalSum $\leftarrow 0$

Distance[] \leftarrow Dijkstra(City[i], t, i)

Roads = Length(Distance)

// find sum of all roads

For $j \leftarrow 0$ to Roads-1:

LocalSum \leftarrow LocalSum + Distance[j]

// End for loop

If Sum > LocalSum then:

Sum \leftarrow LocalSum

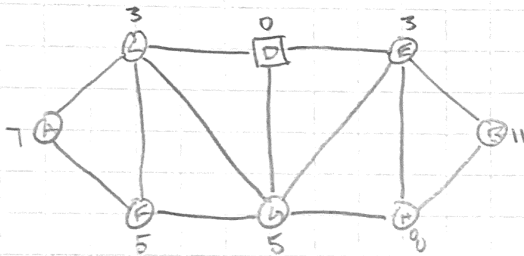
OptimalLoc $\leftarrow i$

// End if

// End for loop

return OptimalLoc

2. * If center is D



* time complexity:

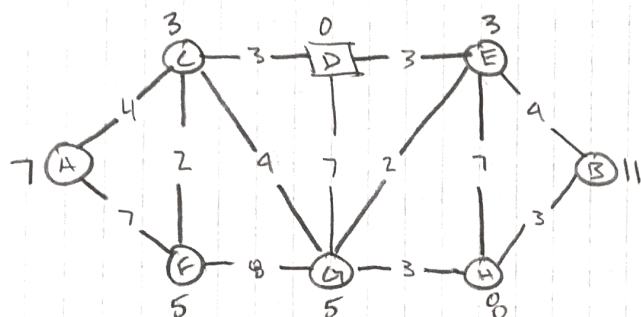
- Dijkstra's algo. runs t times

- time comp. = $tO(r + t(\log t))$

- Converts to, $O(tr + t^2(\log t))$

C.

* Optimal town for distribution center is D



* Cost of spanning tree of city graph is 20

vertex	Distance	Parent vertex	Shortest Path
D	0		D
A	7	C	D → C → A
B	11	H	D → E → G → H → B
C	3	D	D → C
E	3	D	D → E
F	5	C	D → C → F
G	5	E	D → E → G
H	8	G	D → E → G → H

D. * My suggestion would be C & H

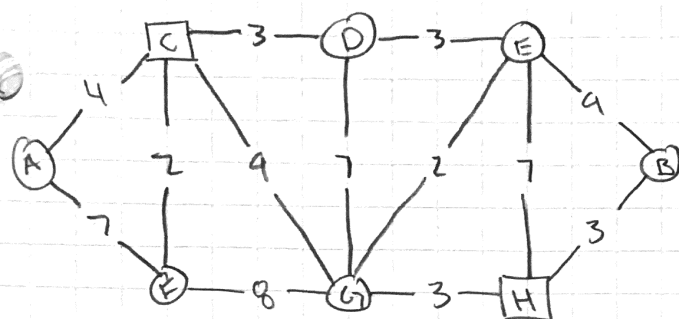
→ For each pair (i, j) run dijkstra's algorithm and track the shortest distances D_i & D_j

$D_i = \text{Dijkstra}(i)$

$D_j = \text{Dijkstra}(j)$

- Compare the distances to the two vertices and track the smallest in the shortest path
- Record maximum path among the shortest paths
- Return a pair (x, y) with the smallest maximum Distance to a town, as required optimal locations to place distribution centers.

Ex.



* Running time:

- Running time of dijkstra = $O(t^2)$ & there are $O(t^2)$ possible pairs so,

= Running time = $O(t^2) \times O(t^2) = O(t^4)$

- C & H have the smallest maximum distance to a location.

Problem 2:

* Algorithm:

Func(graph, vertices)

Create list visited to check vertex is visited or not
Assign distance of every node

Assign 1st node distance to 0

Repeat below until all nodes are visited

Pick vertex that's unvisited & distance min. & add to list

For all adjacent vertices to the current visited list

Update the distance

* time complexity = $O(n^2)$, $n = \#$ of vertices