

## Problem #1

A)

Divide array to 2 parts and compare max & min of both to get max & min of whole thing

MaxMin(arr, arr-size)

if arr-size = 1:

return elem as both max & min

else if arr-size = 2

compare to determine min & max

return max & min

else

recur for max & min of left half

recur for right half

one finds max comparison

one finds min comparison

return max & min

B)

$$T(n) = T(\text{floor}(n/2)) + T(\text{ceil}(n/2)) + 2$$

C)

$$T(2) = 1 \quad T(1) = 0$$

$$T(n) = 2T(n/2) + 2$$

$$T(n) = 3n/2 - 2$$

\* uses  $n$  steps to get max/min cause it traverses the array linearly

## Problem #2

A)

Mergesort3 (A[0..n-1])

if  $n \leq 1$

return A[0..n-1]

Let  $k = \lceil n/3 \rceil$  &  $m = \lceil 2n/3 \rceil$

return merge3(Mergesort3)

merge3(L0, L1, L3)

return merge(L0, merge(L1, L3))

B)

$$T(n) = 3T(n/3) + O(n)$$

C)

$$T(n) = O(n \log n)$$

\* Solution of recurrence in part B

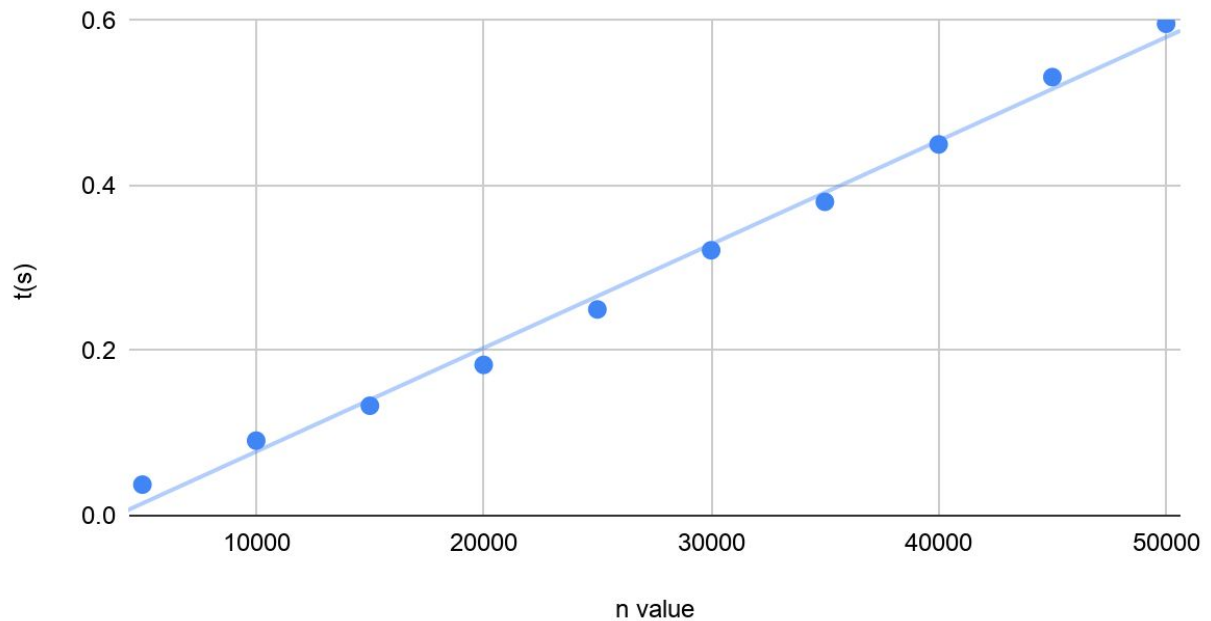
Problem #4:

B.

N value	t(s)
5000	0.03757095337
10000	0.09090304375
15000	0.1331989765
20000	0.1827030182
25000	0.2499029636
30000	0.3215858936
35000	0.3803420067
40000	0.4499289989
45000	0.5314919949
50000	0.5961170197

C.

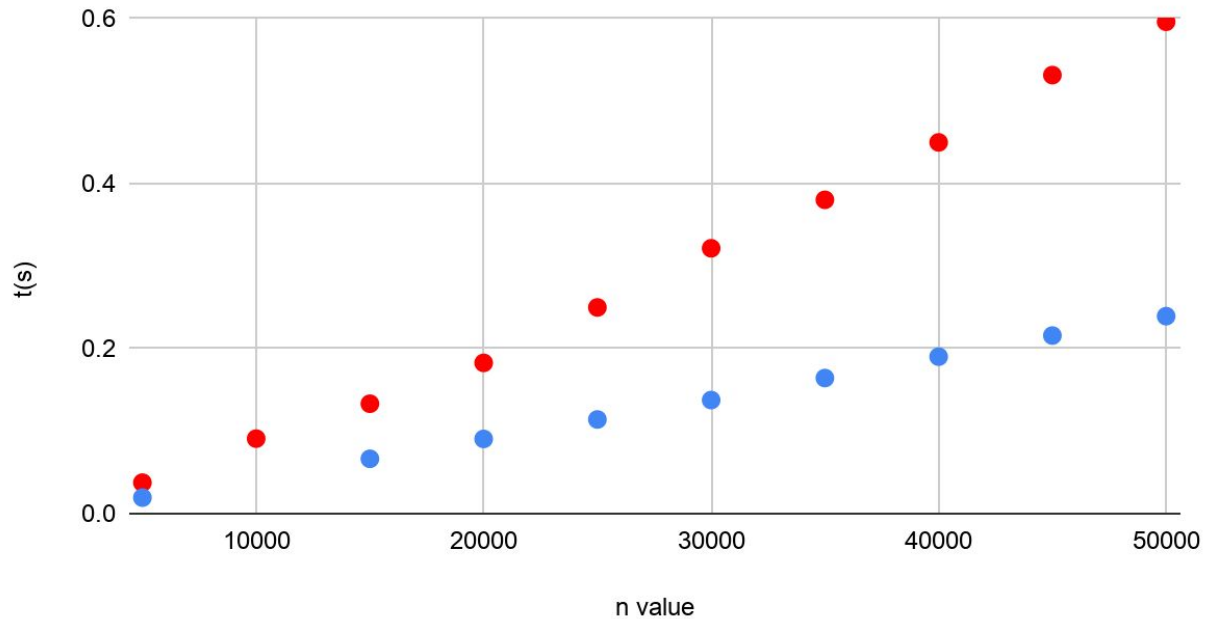
t(s) vs. n value



Judging from my data and the graph, The equation it mostly follows would be a polynomial curve just like the mergesort from assignment 1. Possible  $n \log(n)$ .

E.

t(s) vs. n value



This graph shows red being the mergesort3 and blue as the mergesort from assignment 1. When comparing the two I was really surprised because I figured mergesort3 would be faster. I averaged five separate time collections for each sorting technique and it remained the same. When comparing the two I see that they both seem to follow a similar curve. The only thing that I could see making the difference is the base of the  $n \log(n)$  for each curve. Mergesort3 has a log base of three compared to mergesorts log base 2.