MET CS662 - Assignment #2

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1. Let $L = \{ab, aa, baa\}$. Which of the following strings are in L^* . abaabaaabaa, aaaabaaaa, baaaaabaaaab, baaaaabaa? **Ans**:

abaabaaabaa, aaaabaaaa, baaaaabaa

- **2.** Find grammars for $\Sigma = \{a, b\}$ that generate the set of:
 - **a.** All strings with exactly one a.

Ans:

 $\mathsf{S} \to \mathsf{BaB}$

 $B \rightarrow bB \mid \lambda$

In the process of derivation, S is called only once, hence the generated sentence contains exactly one a.

b. All strings with no more than three a's.

Ans:

 $S \rightarrow B \mid BaB \mid BaBaB \mid BaBaBaB$

$$B \rightarrow bB \mid \lambda$$

In the process of derivation, S is called only once, hence the generated sentence contains no more than three a's.

Note that the textbook solution would say "S \rightarrow BaB | BaBaB | BaBaBaB" since λ is not supposed to be in the language in the textbook, but the professor said in the class told us the language can contain λ , so this answer is correct.

In each case, give convincing arguments that the grammar you give indeed generate the indicated language.

3. Give a simple description of the language generated by the productions:

$$S \rightarrow aA$$

$$A \rightarrow bS$$

$$S \rightarrow \lambda$$

Ans:

S -> aA -> abS -> abaA -> ababS -> ababA ->
$$\cdots$$
 -> λ -> ab

Hence,
$$L = \{ (ab)^n \mid n >= 0 \}$$

4. What language does the grammar with these productions?

 $S \rightarrow Aa$

 $A \rightarrow B$

 $B \rightarrow Aa$

Ans:

S -> Aa -> Ba -> Aaa -> Baa -> Aaaa -> ...

This grammer doesn't generate any string.

5. Find a grammar that generates the language

$$L = \{\omega\omega^R | \omega \in \{a, b\}^+\}$$

Ans:

S -> aAa | bAb

A -> aAa | bAb | λ

6. Are the two grammars with respective productions

$$S \rightarrow aSb \mid ab \mid \lambda$$

and

$$S \rightarrow aAb \mid ab$$

$$A \rightarrow aAb \mid \lambda$$

equivalent? Assume that *S* is the start symbol in both cases.

Ans:

Not equivalent. The counter example is as follows.

Grammer 1: S -> λ

Grammer 2: S -> (cannot generate λ).

7. So far we have given examples of only relatively simple grammars; every production had a single variable on the left side. As we will see, such grammars are very important, but the definition of grammar allows more general forms.

Consider the grammar $G = (\{A, B, C, D, E, S\}, \{a\}, S, P)$, with productions

$$S \rightarrow ABaC$$

$$Ba \rightarrow aaB$$

$$BC \rightarrow DC|E$$

$$aD \rightarrow Da$$

$$AD \rightarrow AB$$

$$aE \rightarrow Ea$$

$$AE \rightarrow \lambda$$

Derive three different sentences in L(G). From these, make a conjecture about L(G). **Ans:**

S -> ABaC -> AaaBC -> AaaE -> AaEa -> AEaa ->
$$\lambda$$
aa -> aa S -> ABaC -> AaaDaC -> ADaaC -> ABaaC -> AaaBaC -> AaaaaBC -> AaaaaE -> AaaaEa -> AaaEaa -> AaEaaa -> AEaaaa -> λ aaaaa -> aaaa S =>* aaaaaaaaa We can conjecture that L(G) = { a^{2^n} | n >= 1 }