

数字逻辑设计

Digital Logic Design

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Unit 3 Boolean Algebra

- 逻辑运算
- 布尔表达式和真值表
- 逻辑代数定理及规则
- 代数化简法



George Boole

各种逻辑运算



■ 基本逻辑运算 (Basic Operations)

➤ 与 (AND)

➤ 或 (OR)

➤ 非 (NOT)

■ 复合逻辑运算 (Other Operations)



基本运算（Basic Operations）

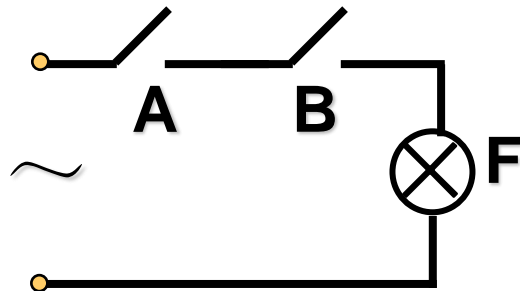
1. AND（逻辑“与”）

$$F=A \cdot B$$

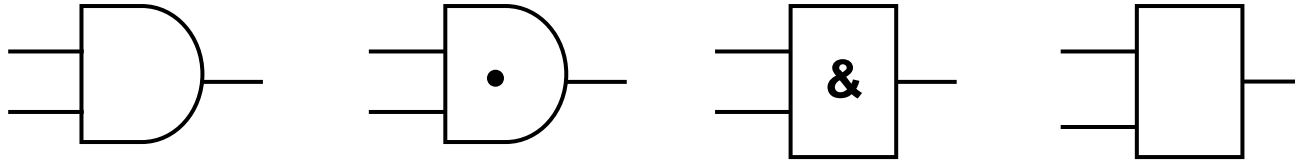
Truth Table

AB	F
0 0	0
0 1	0
1 0	0
1 1	1

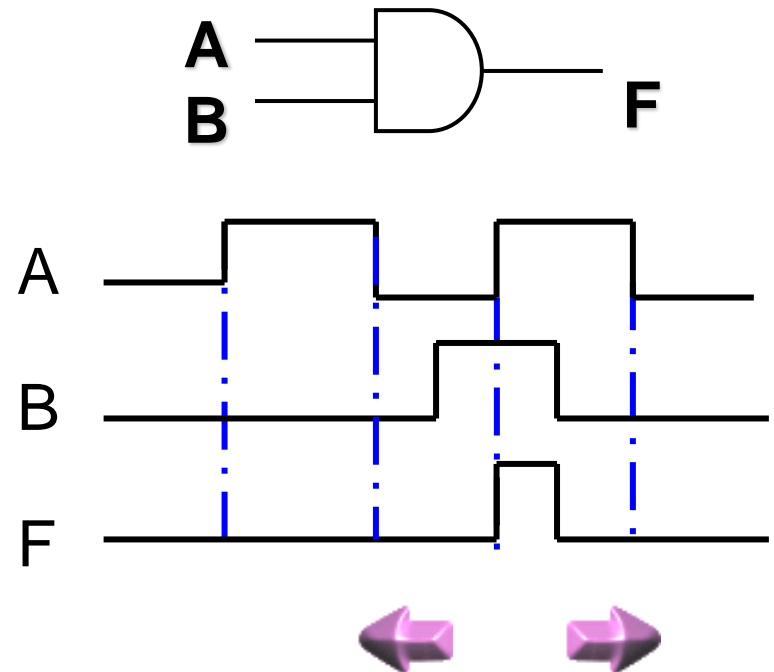
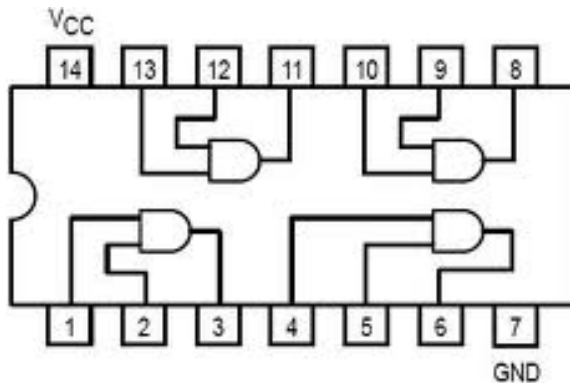
① 也称为：逻辑“乘”



② AND gate (与门) 逻辑符号



③ Typical Chip: 74LS08

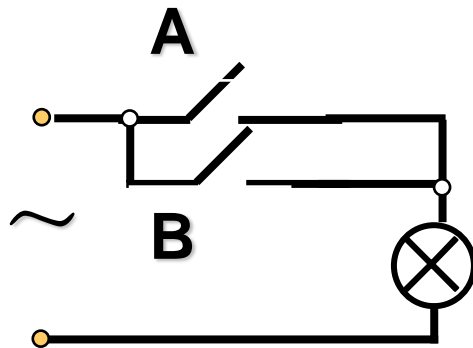


基本运算（Basic Operations）

2. OR（逻辑“或”）

$$F=A+B$$

①也称为：逻辑“加”

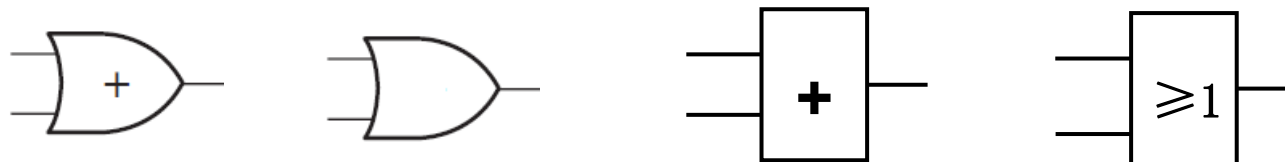


Truth Table

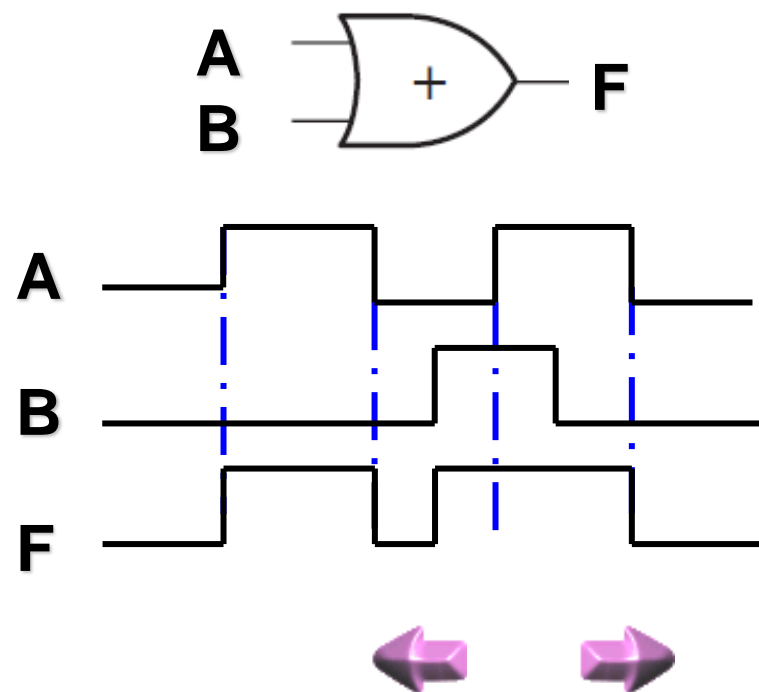
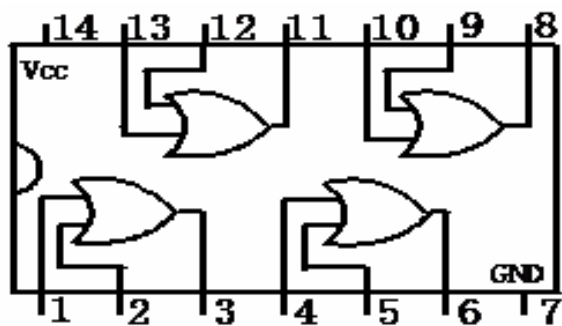
AB	F
0 0	0
0 1	1
1 0	1
1 1	1



② OR gate (或门) 逻辑符号



③ Typical Chip: 74LS32



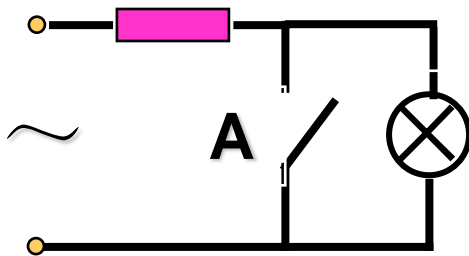
基本运算 (Basic Operations)

3. NOT (逻辑 “非”)

$$F = \overline{A}$$

(or $F = A'$)

①也称为：反相器

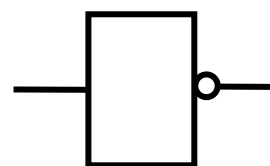
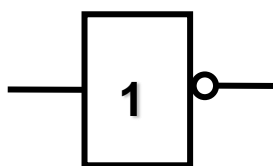
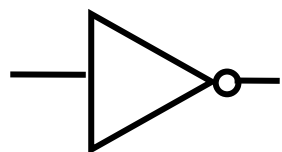


True table

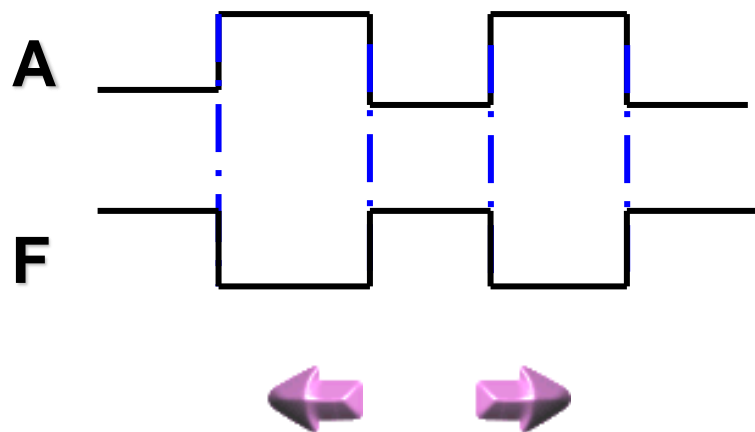
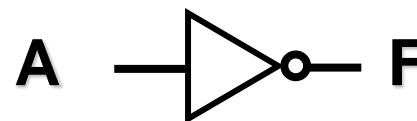
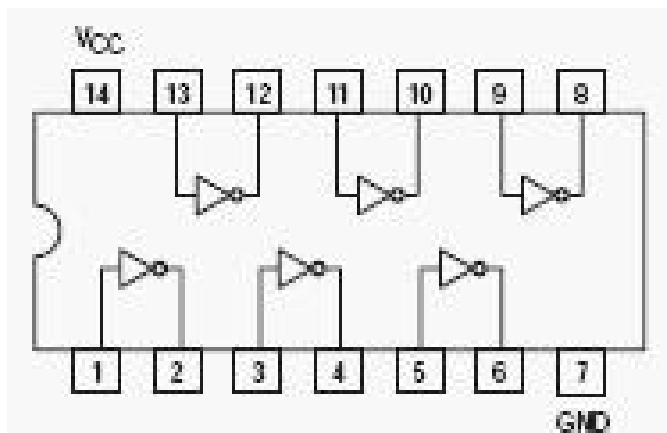
A	F
0	1
1	0



② NOT gate (非门) 逻辑符号



③ Typical Chip: 74LS04



复合运算（Other Operations）

■ 基本逻辑运算（Basic Operations）

➤ 与（AND）

➤ 或（OR）

➤ 非（NOT）



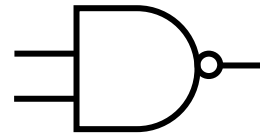
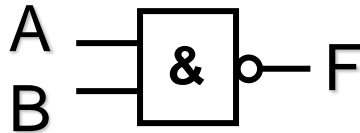
■ 复合逻辑运算（Other Operations）



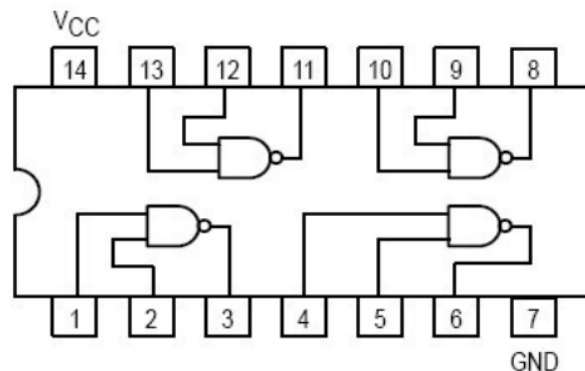
复合运算 (Other Operations)

4. 与非门 (NAND gate)

$$F = \overline{AB}$$



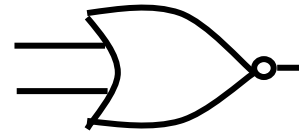
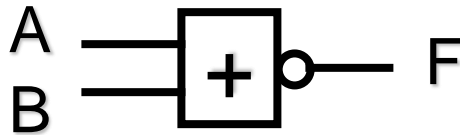
■ Typical Chip: 74LS00



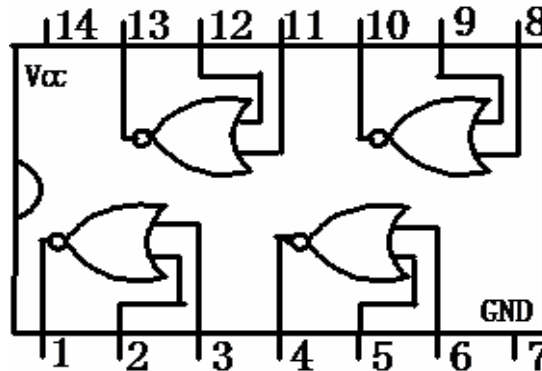
复合运算（Other Operations）

5. 或非门（NOR gate）

$$F = \overline{A+B}$$



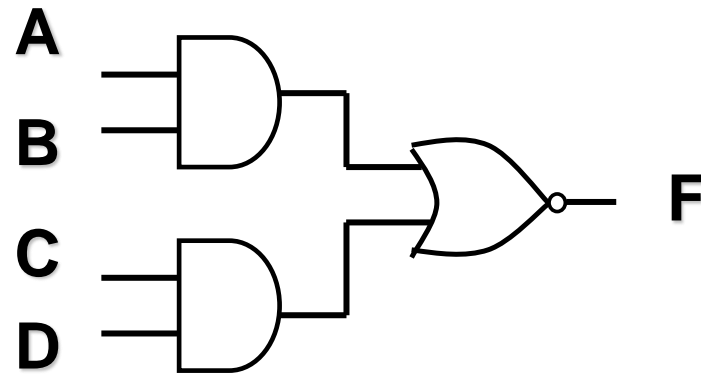
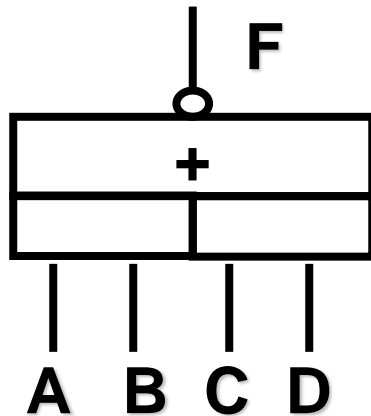
- Typical Chip: 74LS02



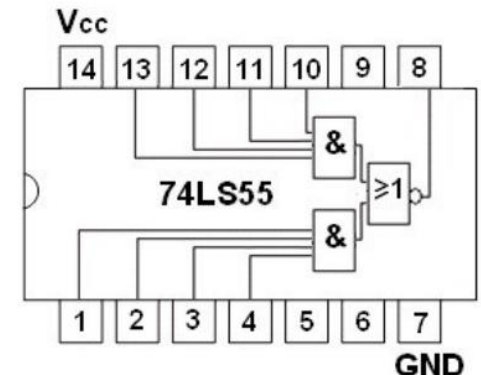
复合运算 (Other Operations)

6. 与或非门 (AND-OR-NOT gate)

$$F = \overline{AB + CD}$$



■ Typical Chip: 74LS51, 74LS55



复合运算 (Other Operations)

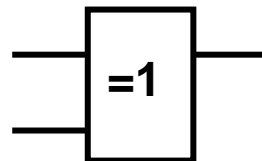
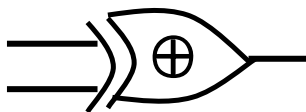
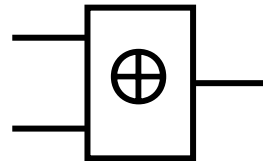
7. 异或门 (\oplus , Exclusive-OR operation)

$$F = A \oplus B = \bar{A}B + A\bar{B}$$

Truth Table

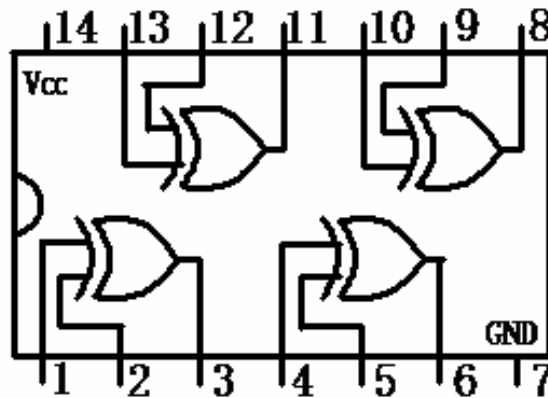
AB	F
0 0	0
0 1	1
1 0	1
1 1	0

① 逻辑符号



复合运算（Other Operations）

② Typical Chip: 74LS86



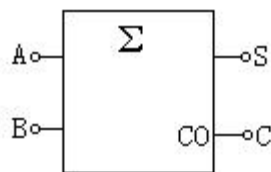
③ Applications

- 全加器 (Full adder)
- 半加器 (Half-adder)



复合运算 (Other Operations)

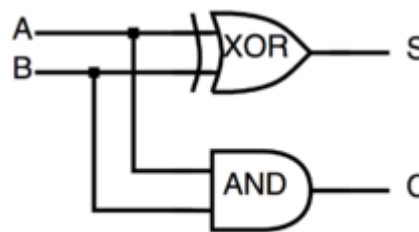
■ 半加器 (Half-adder)



半加器逻辑符号

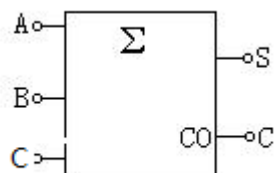
输入		输出	
A	B	C	S
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

逻辑表达式: $S = A \oplus B$; $C = A \cdot B$ 。



半加器的逻辑实现

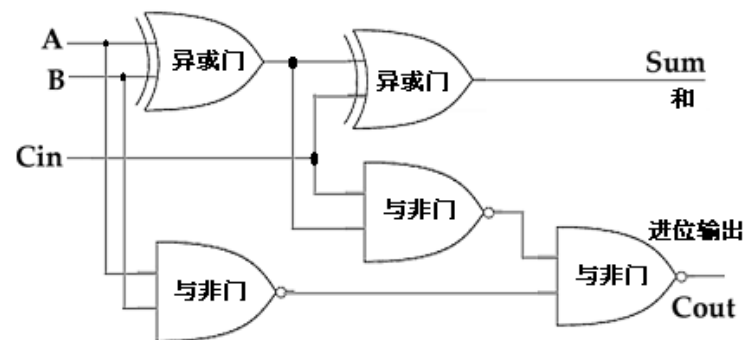
■ 全加器 (Full adder)



全加器逻辑符号

输入			输出	
Ci-1	Ai	Bi	Si	Ci
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

$$S = A \oplus B \oplus C_{in}$$
$$C_{out} = (A \cdot B) + (C_{in} \cdot (A \oplus B))$$



复合运算 (Other Operations)

8. 同或门 (\odot or \equiv , Equivalence operation)

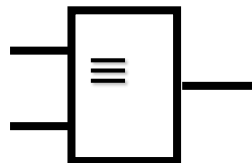
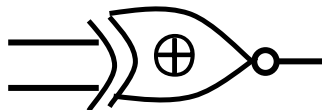
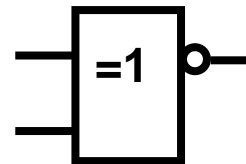
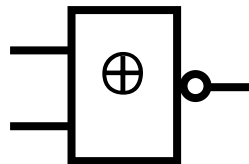
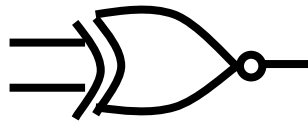
$$F = A \equiv B \quad \text{or}$$

$$F = A \odot B = \bar{A}\bar{B} + AB$$

Truth Table

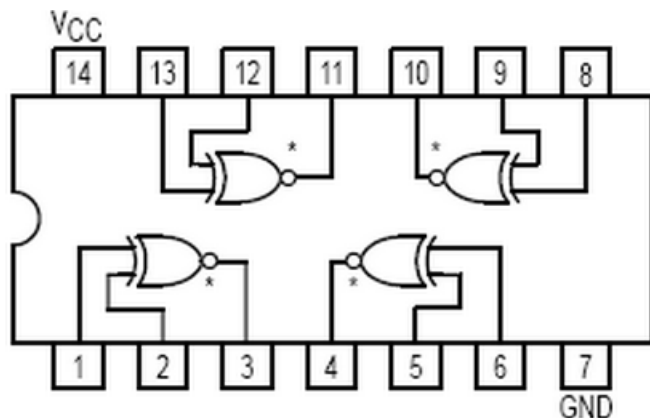
AB	F
0 0	1
0 1	0
1 0	0
1 1	1

① 逻辑符号



复合运算（Other Operations）

② Typical Chip: 74LS266



如何构造1位等值比较器??

—— 利用异或门（同或门）



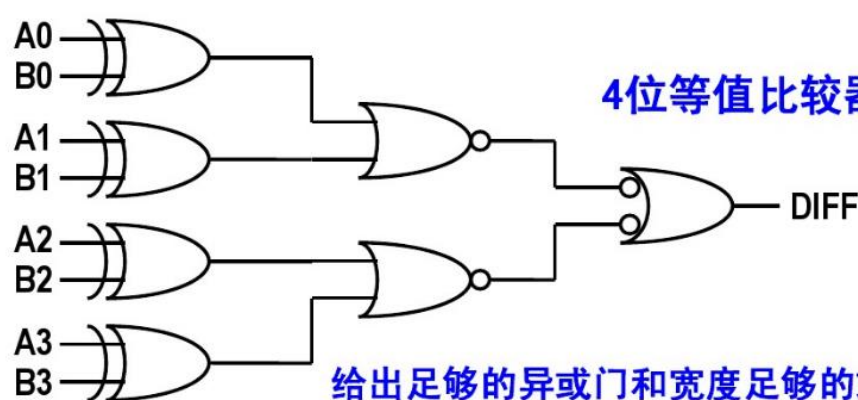
DIFF : different



EQ : equal

③ Applications

■ 等值比较器



给出足够的异或门和宽度足够的或门，
可以搭建任意输入位数的等值比较器。

复合运算 (Other Operations)

④ 性质

$$A \oplus 1 = \bar{A}$$

$$A \odot 1 = A$$

$$A \oplus 0 = A$$

$$A \odot 0 = \bar{A}$$

$$A \oplus A = 0$$

$$A \odot A = 1$$

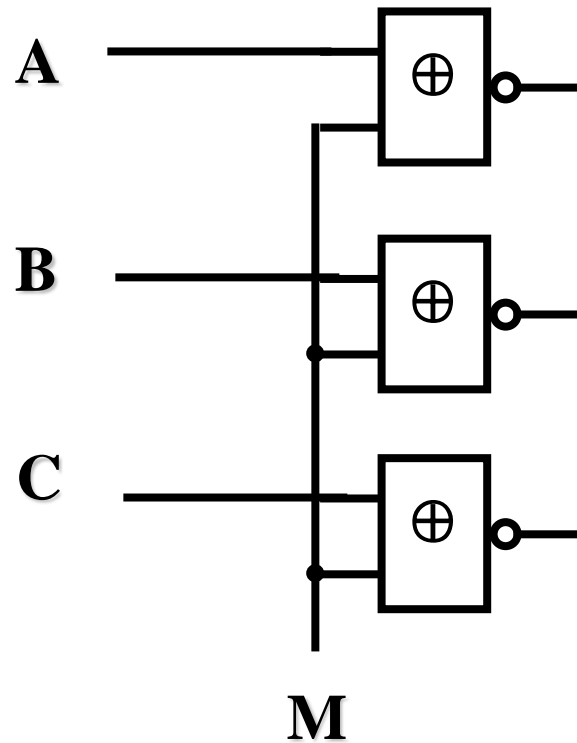
$$A \oplus \bar{A} = 1$$

$$A \odot \bar{A} = 0$$



复合运算 (Other Operations)

Applications



Unit 2 Boolean Algebra

- 各种逻辑运算
- 布尔表达式和真值表
- 逻辑代数定理及规则
- 代数化简法

布尔表达式和真值表

1. 布尔表达式 (Boolean Expressions)

Example

$$F = AB + \bar{A}\bar{B}$$

$$F = [A(C + D)]' + BE$$

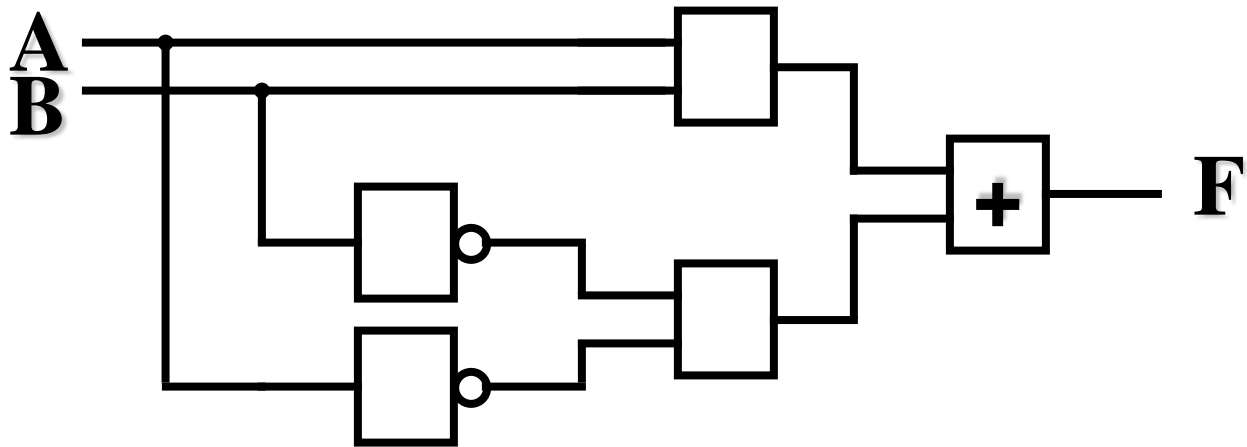
- *Boolean expressions* are formed by application of the basic operations (**and**, **or**, **not**) to one or more variables or constants.



布尔表达式和真值表

逻辑图

Example $F = AB + \bar{A}\bar{B}$



- Each expression corresponds directly to a circuit of logic gates



布尔表达式和真值表

2. 真值表 (Truth Table)

Truth Table

Example $F = AB + \bar{A}\bar{B}$

AB	F
0 0	1
0 1	0
1 0	0
1 1	1

- A *truth table* specifies the values of a Boolean expression for every possible combination of values of the variables in the expression.
- n 个输入变量有 2^n 种取值组合



布尔表达式和真值表

- 如果两个逻辑表达式的真值表相等，则这两个逻辑表达式相等.

$$AB' + C = (A + C)(B' + C)$$

适用情况：逻辑表达式简单，逻辑变量较少

A B C	$AB' + C$	$(A + C)(B' + C)$
0 0 0	0	0
0 0 1	1	1
0 1 0	0	0
0 1 1	1	1
1 0 0	1	1
1 0 1	1	1
1 1 0	0	0
1 1 1	1	1

Unit 2 Boolean Algebra

- 各种逻辑运算
- 布尔表达式和真值表
- 逻辑代数定理及规则
- 代数化简法

Laws and Theorems

1. 公理 (Axiom)

$$(1) \quad 0 \cdot 0 = 0$$

$$(1)' \quad 0+0=0$$

$$(2) \quad 0 \cdot 1 = 1 \cdot 0 = 0$$

$$(2)' \quad 1+0 = 0+1=1$$

$$(3) \quad 1 \cdot 1 = 1$$

$$(3)' \quad 1+1 = 1$$

$$(4) \quad \bar{0} = 1$$

$$(4)' \quad \bar{1} = 0$$

$$(5) \quad \text{If } A \neq 0 \text{ then } A=1$$

$$(5)' \quad \text{If } A \neq 1 \text{ then } A=0$$



Laws and Theorems

2. 基本定理 (Basic Theorems)

■ *single variable is involved*

$$(6) \quad A + 0 = A$$

$$(6)' \quad A \cdot 0 = 0$$

$$(7) \quad A + 1 = 1$$

$$(7)' \quad A \cdot 1 = A$$

0 - 1 律

$$(8) \quad A + \bar{A} = 1$$

$$(8)' \quad A \cdot \bar{A} = 0$$

互补律

$$(9) \quad A + A = A$$

$$(9)' \quad A \cdot A = A$$

重叠率

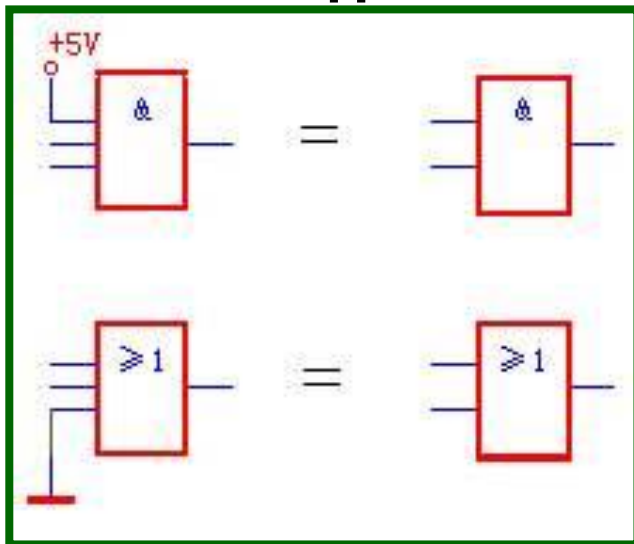


Laws and Theorems

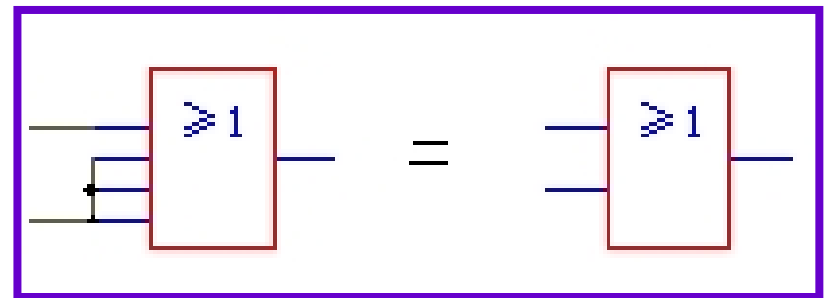
2. 基本定理 (Basic Theorems)

Applications——

0 - 1 律



重叠率



Laws and Theorems

■ 与普通代数相似的定理

交换律

$$(10) \mathbf{A+B=B+A}$$

$$(10)' \mathbf{A \cdot B = B \cdot A}$$

结合律

$$(11) \mathbf{(A+B)+C=A+(B+C)}$$

$$(11)' \mathbf{(A \cdot B) \cdot C = A \cdot (B \cdot C)}$$

分配律

$$(12) \mathbf{A \cdot (B+C) = AB+AC}$$

$$(12)' \mathbf{A+BC=(A+B) \cdot (A+C)}$$



普通代数
不支持

第二分配律



Laws and Theorems

■ 第二分配律证明

$$(A+B) \cdot (A+C)$$

$$= A \cdot (A+C) + B \cdot (A+C) \quad \text{分配律}$$

$$= A + AC + AB + BC$$

$$= A \cdot 1 + AC + AB + BC \quad \text{0-1律}$$

$$= A \cdot (1 + C + B) + BC \quad \text{分配律}$$

$$= A + BC \quad \text{0-1律}$$

Laws and Theorems

■ 第二分配律证明 (第2种证明)

$$(A+B) \cdot (A+C)$$

$$= A \cdot (A+C) + B \cdot (A+C) \quad \text{分配律}$$

$$= A + AC + AB + BC$$

$$= A + A + AC + AB + BC \quad \text{重叠率}$$

$$= (A + AC) + (A + AB) + BC \quad \text{交换律}$$

$$= A(1+C) + A(1+B) + BC \quad \text{分配律}$$

$$= A + A + BC \quad \text{0-1律}$$

$$= A + BC$$

Laws and Theorems

■ 特殊定理

德摩根定理 (DeMorgan's Laws)

$$(13) \quad \overline{A+B} = \bar{A} \cdot \bar{B} \quad (13)' \quad \overline{A \cdot B} = \bar{A} + \bar{B}$$

$$(14) \quad \overline{\overline{A}} = A$$



Laws and Theorems

■ 特殊定理

DeMorgan's Laws 😊

◆ Applications: 表达式化简

$$(1) \quad \overline{X_1 X_2 \dots X_n} = \bar{X}_1 + \bar{X}_2 + \dots + \bar{X}_n$$

$$(2) \quad \overline{X_1 + X_2 + \dots + X_n} = \bar{X}_1 \bar{X}_2 \dots \bar{X}_n$$



Laws and Theorems

■ 特殊定理

对偶规则 😊

◆ Applications: Algebraic Simplification

{ 变量:
运算符:

不变

\cdot	\longrightarrow	$+$
$+$	\longrightarrow	\cdot
\oplus	\longrightarrow	\odot
\odot	\longrightarrow	\oplus

不能改变原
来的优先级



Laws and Theorems

■ 特殊定理

对偶规则

Example

$$F = A \cdot (B + C) \xrightarrow{\text{对偶}} (F)^D = A + B \cdot C$$

$$F = A \cdot \bar{B} + AC \xrightarrow{\text{对偶}} (F)^D = (A + \bar{B}) \cdot (A + C)$$

$$F = \overline{\bar{A} \cdot B \cdot \bar{C}} \xrightarrow{\text{对偶}} (F)^D = \overline{\bar{A} + B + \bar{C}}$$



Laws and Theorems

■ 特殊定理

Inference of Dual Rule

① $F \xleftrightarrow{\text{Dual Rule}} (F)^D$

② 两个逻辑表达式相等，它们的对偶也相等

Example

$$A + BCD = (A + B)(A + C)(A + D)$$



Dual Rule



Dual Rule

$$A \cdot (B + C + D) = AB + AC + AD$$

Laws and Theorems

From (18) :

3. 常用公式

$$(15) \quad \mathbf{AB + A\bar{B} = A} \quad (\text{合并律})$$

$$(16) \quad \mathbf{A + AB = A} \quad (\text{吸收律})$$

$$(17) \quad \mathbf{A + \bar{A}B = A + B} \quad (\text{消除律})$$

$\mathbf{A + \bar{A}B}$ 第二分配律

$$= (\mathbf{A + \bar{A}}) (\mathbf{A + B})$$

$$= \mathbf{A + B}$$

$$= \mathbf{AB + \bar{A}C}$$

$$(18) \quad \mathbf{AB + \bar{A}C + BC = AB + \bar{A}C} \quad (\text{蕴含律})$$

$$(18)' \quad \mathbf{AB + \bar{A}C + BCD = AB + \bar{A}C}$$

$$(18)'' \quad (\mathbf{A + B})(\mathbf{B + C})(\mathbf{A' + C}) = (\mathbf{A + B})(\mathbf{A' + C})$$



Laws and Theorems

$$(19) \quad \overline{A \bar{B} + \bar{A} B} = \bar{A} \bar{B} + AB$$

$$\begin{aligned} & \overline{A \bar{B} + \bar{A} B} \\ &= \overline{A \bar{B}} \cdot \overline{\bar{A} B} \\ &= (\bar{A} + B) \cdot (A + \bar{B}) \\ &= \bar{A} \bar{B} + AB \end{aligned}$$



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Algebraic Simplification

■ 一个逻辑函数有多种不同的表达式

$$F=AB+A\bar{C} \quad \dots\dots \text{与-或}$$

$$\overline{\overline{AB+A\bar{C}}}$$

$$=\overline{\overline{AB}} \cdot \overline{\overline{A\bar{C}}} \quad \dots\dots \text{与非-与非}$$

$$=(\overline{\overline{A+B}}) \cdot (\overline{\overline{A+C}}) \quad \dots\dots \text{或-与非}$$

$$=(\overline{\overline{A+B}}) + (\overline{\overline{A+C}}) \quad \dots\dots \text{或非-或}$$

$$F=(A+B) \cdot (A+\bar{C}) \quad \dots\dots \text{或-与}$$

$$\overline{\overline{(A+B) \cdot (A+\bar{C})}}$$

$$=\overline{\overline{(A+B)}} + \overline{\overline{(A+\bar{C})}} \quad \dots\dots \text{或非-或非}$$

$$=\overline{\overline{A}} \cdot \overline{\overline{B}} + \overline{\overline{A}} \cdot \overline{\overline{C}} \quad \dots\dots \text{与-或非}$$

$$=\overline{\overline{A}} \overline{\overline{B}} \cdot \overline{\overline{A\bar{C}}} \quad \dots\dots \text{与非-与}$$

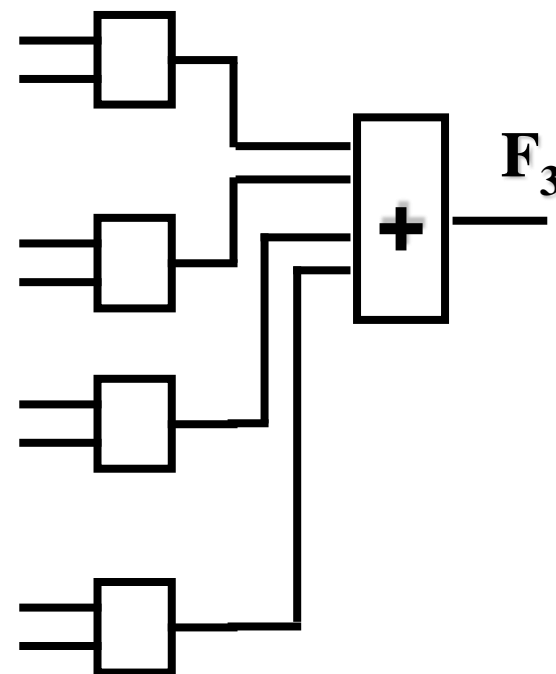
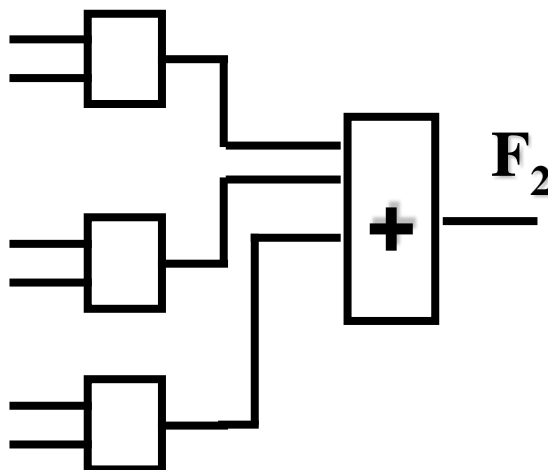
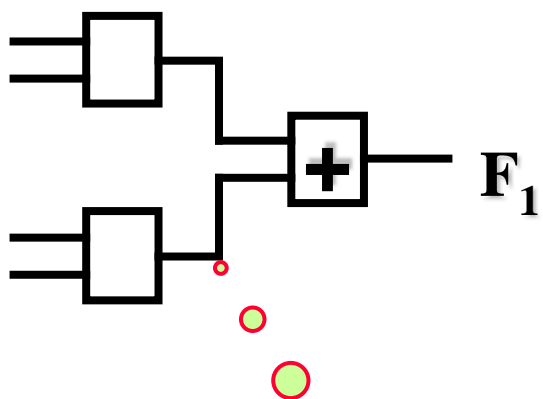


■ 同一类型的表达式也不是唯一的

$$F=AB+\bar{A}C \quad \dots\dots\dots \textcircled{1} F_1$$

$$=AB+\bar{A}C+BC \quad \dots\dots\dots \textcircled{2} F_2$$

$$=ABC+AB\bar{C}+\bar{A}BC+\bar{A}\bar{B}C \quad \dots\dots\dots \textcircled{3} F_3$$



最简，元件少，可靠



Algebraic Simplification



最简 (Minimum Expressions) ?

① 与项 (和项) 的个数最少

② 每个与项 (和项) 中变量的个数最少

minimum cost

① 逻辑门的数量最少

② 逻辑门的输入个数最少

目的:

① 降低成本

② 提高可靠性

Methods

■ 代数法 (Algebraic techniques)

■ 卡诺图法 (K. map method)



Simplification Methods

代数化简法

example

$$F = \underbrace{A + A\bar{B}\bar{C} + \bar{A}CD + \bar{C}E + \bar{D}E}$$

$$= \underbrace{A + \bar{A}CD + \bar{C}E + \bar{D}E}$$

$$= A + CD + \underbrace{\bar{C}E + \bar{D}E}$$

$$= A + CD + E(\bar{C} + \bar{D})$$

$$= \underbrace{A + CD + E\bar{C}\bar{D}}$$

$$= A + CD + E$$



example

$$F = \underline{AB} + \underline{A\bar{C}} + \bar{B}C + \bar{B}\bar{C} + \bar{B}D + \bar{B}\bar{D} + \underline{ADE(F+G)}$$

$$= \underline{A(\bar{B}C)} + \bar{B}C + \bar{B}\bar{C} + \bar{B}D + \bar{B}\bar{D} + \underline{ADE(F+G)}$$

$$= \underline{A} + \bar{B}C + \bar{B}\bar{C} + \bar{B}D + \bar{B}\bar{D} + \underline{ADE(F+G)}$$

$$= A + \underline{\bar{B}C + \bar{B}\bar{C} + \bar{B}D + \bar{B}\bar{D}} + C\bar{D}$$

$$= A + \bar{B}C + \bar{B}\bar{C} + \bar{B}D + \bar{B}\bar{D} + C\bar{D}$$

$$= A + \bar{B}C + \bar{B}\bar{C} + \bar{B}D + C\bar{D}$$

$$= A + \bar{B}\bar{C} + \bar{B}D + C\bar{D}$$



Simplification Methods

example

$$F = (\bar{B} + D)(\bar{B} + D + A + G)(C + E)(\bar{C} + G)(A + E + G)$$

Dual Rule:

$$J = \bar{B}D + \bar{B}DAG + CE + \bar{C}G + AEG$$

$$= \bar{B}D + \underbrace{CE + \bar{C}G}_{\text{blue bracket}} + \cancel{AEG}$$

$$= \bar{B}D + CE + \bar{C}G$$

Dual Rule:

$$F = (\bar{B} + D)(C + E)(\bar{C} + G)$$

example

$$F = A + AB + \bar{A}C + BD + ACEF + \bar{B}E + DEF$$

$$= A + C + BD + \bar{B}E$$



代数化简法

优点——

- 不受变量数目的约束；
- 对公理、定理和规则十分熟练时，化简较方便。

缺点——

- 技巧性强
- 在很多情况下难以判断化简结果是否最简



Unit 2 Boolean Algebra

- 各种逻辑运算
- 布尔表达式和真值表
- 逻辑代数定理及规则
- 代数化简法