

## Peak Index in Mountain Array - LeetCode 852

The problem finds the index of the "peak" element in a mountain array - an array that increases to a single peak, then decreases. Guaranteed to have exactly one peak, array length  $\geq 3$ .

Example: [0,3,8,9,5,2] → Peak is 9 at index 3

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## Mountain Array Properties

Peak element characteristics:

- Left side (before peak): Strictly increasing →  $\text{arr}[i] > \text{arr}[i-1]$
- Right side (after peak): Strictly decreasing →  $\text{arr}[i] > \text{arr}[i+1]$
- Peak itself:  $\text{arr}[\text{peak}] > \text{arr}[\text{peak}-1]$  AND  $\text{arr}[\text{peak}] > \text{arr}[\text{peak}+1]$

Important constraints:

- Peak cannot be at index 0 or n-1
  - Array length  $\geq 3$ , so safe to access mid-1 and mid+1
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## Why Binary Search?

Mountain array has sorted structure (increasing → decreasing), perfect for binary search

Time complexity:  $O(\log n)$  instead of  $O(n)$  linear scan

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## Step-by-Step Algorithm

### 1. Initialize Search Space

```
start = 1, end = n-2 // Skip edges (never peak)
```

Why? Eliminates boundary checks for mid-1 / mid+1

## 2. Binary Search Loop ( `while start <= end` )

```
mid = start + (end - start) / 2 // Avoid overflow
```

## 3. Check if mid is Peak

```
if (arr[mid] > arr[mid-1] && arr[mid] > arr[mid+1])  
    return mid; // Found peak!
```

## 4. Decide Search Direction

### Case A: Increasing slope ( `arr[mid-1] < arr[mid]` )

```
start = mid + 1; // Peak must be RIGHT
```

### Case B: Decreasing slope ( `arr[mid-1] >= arr[mid]` )

```
end = mid - 1; // Peak must be LEFT
```

## Dry Run: [0,3,8,9,5,2]

```
n=6, start=1, end=4 (n-2)
```

Iteration 1:

$mid = 1 + (4-1)/2 = 2$  (`arr[2]=8`)

Check:  $8 > 3?$  ✓  $8 > 9?$  ✗ → Not peak

$arr[1]=3 < arr[2]=8?$  ✓ → Increasing → start=3

Iteration 2:

$start=3, end=4$

$mid = 3 + (4-3)/2 = 3$  (`arr[3]=9`)

Check:  $9 > 8?$  ✓  $9 > 5?$  ✓ → \*\*PEAK FOUND! Return 3\*\*

## Complete Code

```
int peakIndexInMountainArray(vector<int>& arr) {  
    int start = 1, end = arr.size() - 2; // Optimized bounds  
  
    while (start <= end) {  
        int mid = start + (end - start) / 2; // Safe mid  
  
        // Peak check  
        if (arr[mid] > arr[mid-1] && arr[mid] > arr[mid+1])  
            return mid;  
  
        // Increasing slope → search right  
        else if (arr[mid-1] < arr[mid])  
            start = mid + 1;  
  
        // Decreasing slope → search left  
        else  
            end = mid - 1;  
    }  
    return -1; // Never reaches (guaranteed peak exists)  
}
```