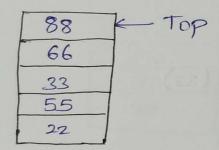
Mame! - M. Hêteshwar Roddy

CSA 0389 192325074 Perform the following operation using stack. Assume the size of stack is s and having value of 22, 58,33,66, 88 in the stack from o position to size-1. Now perform the following operations.

Draw the elements in stack 2, POP[3,3) POP(J,4) Push
Push [90], 5) Push [36], 6) Push [1], 7) Push [88], 8) POP(J, Pop(J)
Draw the diagram of stack & illustrate the above
operations & identify where the top is?

Aus Size of stack: 5

Elements in stack (from bottom to top): 22,55,33,66,88
Top of stack: 88

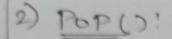


Operations >-

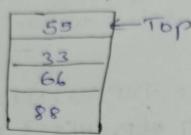
) Invest the elements in the stacking

- . The operation will reverse order of elements in the stack.
- · After Inversion, the stack will look like!

22	Top
55	
33	
66	
88	

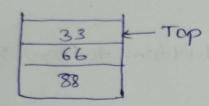


. Remove the top element (22).



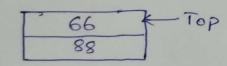
3) POP ()!

· Remove the top element (55).



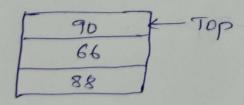
4) POPCI:

· Remove the top element (33). Stack after Pop!



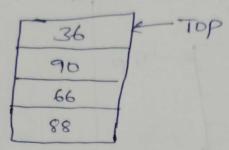
5) Push (90):-

· Push the element · 90 onto the stack. Stack after Push



6) Push (36) !-

. Rush the element 36 onto the stack. stack 'after push!



7) Push (11) !-

- Push the element 11 onto the stack Stack after Push!

11	*	-Tos
36		
90		
66		
88		

8) Push (88):

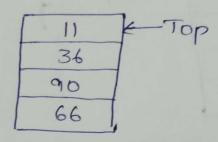
· Push the element 88 onto the stack. Stack after Push:

88	TOP
11	
36	
90	
66	

9) POP():-

· Remove the top element (88):

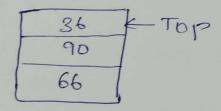
Stack after Pop!



10) POP():

· Remove the top element (11).

Stack after POP



Final stack state

Size of stack: 5

Elements in stack (from bottom to top):

36,90,66

Top of stack: 66

66	* Top
90	1
36	
	4

Develop an algorithm to detect duplicate elements In an unforted acrey using linear search. Determine the time complexity & discuss how you would optimize this process.

Algorithm'

- 1) Initialization:
 - create an Empty set or list to keep track of elements that have already been seen,
 - 2) linear Search'.
 - Iterate through each element of the areay:
 - · For each elements, check if it is already in set
 - of seen elements.
 - . It it is, a duplicate has been found.
 - . It it is found, add it to set of seen elements.
- 3) output:

Return the list of duplicate, or simply indicate that duplicates exit.

C code !-

include estdio.h>

Include LStdbool. h>

int main ()

Port are [] = {4,5,6,7,8,5,4,9,0};

int size = size of (arr) / size of (arr[0]);

bool seen [1000] = {false}

for (int i=0; izhze ji++)

If (seen fars [i])

Printf ('Dupticate tound: y.d(n', au[i]);

else

seen [aur[i]] = true;

return 0;

3

Time complexity

The linear search complexity:

the time complexity for this algorithm is o(n), where in is no of elements in array. This is because each element is checked only once, & operations (checking for membership & adding to a Set) are o(1) on the average.

Space Complexity

the space complexity is o(n) due to additional space used by the 'seen' a (duplicate) sets, which may store up to 'n' elements in the worst case.

Optimo zation

Hashing !-

The use of set for checking duplicates is already efficient because sets provide overage o(i) time complexity for membership tests & Passestions.

Sorting! -

It we are allowed to modify the array, another approach is to fort the array first & then Perform a linear scan to find duplicates.

Sorting would talae $O(n \log n)$ time, and the sub--sequent scan would telae O(n) time. This app--roach uses less space (o(i) additional space if Sorting in - places.