



Models and Tools for Embedded Systems

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Models and Tools for Embedded Systems

Organization

- 1. Model-based Development of Embedded Sys.
- 2. Review of models of concurrency in programming languages
- 3. Introduction to Lustre/Scade, (Esterel,) Scilab
- 4. Simple case studies

Development Challenges (Complexity)

- Correct functioning is crucial
- Reactive
- Concurrent
- Realtime
- Stringent resource constraints

Development Challenges

Embedded Systems are complex

1. Correct functioning is crucial

- safety-critical applications
- damage to life, economy can result

2. They are Reactive Systems

- Once started run forever.
- Termination is a bad behavior.
- Compare conventional computing (transformational systems)

Development Challenges

3. Concurrent systems

- System and environment run concurrently
- Multi-functional

4. Real-time systems

- Not only realtime outputs but in realtime
- Imagine delay of minutes in pacemaker system

Development Challenges

5. Stringent resource constraints

- Compact systems
 - simple processors
 - limited memory
- Quick response
- Good throughput
- Low power
- Time-to-market

System Development

- Process of arriving at final product from reqs
- Requirements
 - Vague ideas, algorithms, of-the shelf components, additional functionality etc.
 - Natural Language statements
 - Informal
- Final Products
 - System Components
 - Precise and Formal

System Components

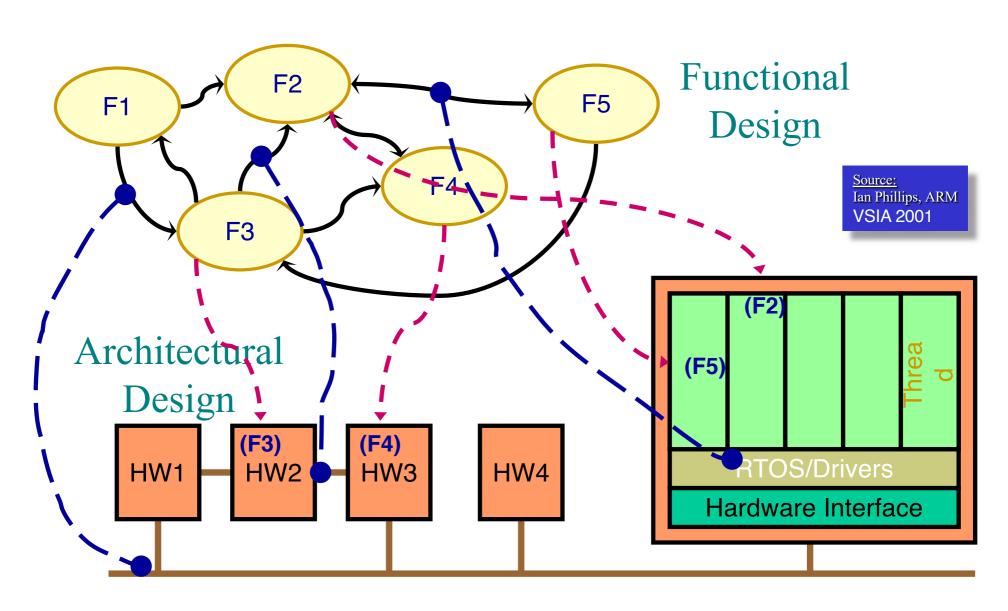
Embedded System Components

- Programmable processors (controllers & DSP)
- Standard and custom hardware
- Concurrent Software
- OS Components:
 - Schedulers, Timers, Watchdogs,
 - IPC primitives
- Interface components
 - External, HW and SW interface

System Development

- Decomposition of functionality
- Architecture Selection:
 Choice of processors, standard hardware
- Mapping of functionality to HW and SW
- Development of Custom HW and software
- Communication protocol between HW and SW
- Prototyping, verification and validation

Functional Design & Mapping



Design Choices

- Choices in Components
 - Processors, DSP chips, Std. Components
- Many different choices in mapping
 - Fully HW solution
 - More speed, cost, TTM (Time to market), less robust
 - Standard HW development
 - Fully SW solution
 - Slow, less TTM, cost, more flexible
 - Standard microcontroller development

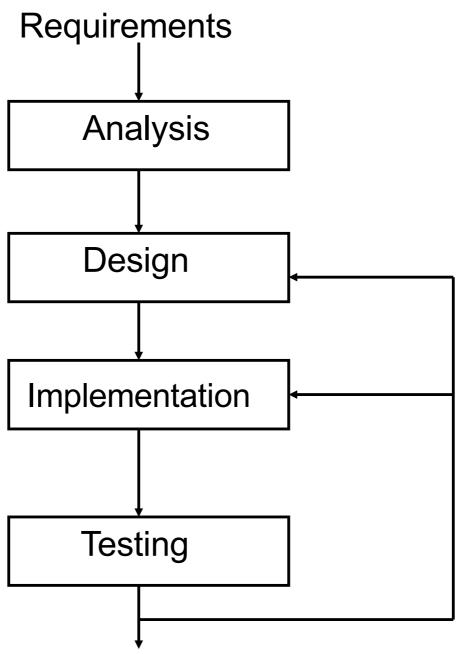
Mixed Solution

Desired Solution is often mixed

- Optimal performance, cost and TTM
- Design is more involved and takes more time
- Involves Co-design of HW and SW
- System Partitioning difficult step
- For optimal designs, design exploration and evaluation essential
- Design practices supporting exploration and evaluation essential
- Should support correctness analysis as it is crucial to ensure high quality

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Classical design methodology



Development Methodology

- Simplified Picture of SW development
 - Requirements Analysis
 - Design
 - Implementation (coding)
 - Verification and Validation
 - Bugs lead redesign or re-implementation

Development Methodology

- All steps (except implementation) are informal
 - Processes/ objects not well defined and ambiguous
 - Design and requirement artifacts not precisely defined
 - Inconsistencies and incompleteness
 - No clear relationship between different stages
 - Subjective, no universal validity
 - Independent analysis difficult
 - Reuse not possible

Classical Methodology

- Totally inadequate for complex systems
 - Thorough reviews required for early bug removal
 - Bugs often revealed late while testing
 - Traceability to Design steps not possible
 - Debugging difficult
 - Heavy redesign cost
- Not recommended for high integrity systems
 - i.e. embedded systems

Formal Methodology

- A methodology using precisely defined artifacts at all stages
 - Precise statement of requirements
 - Formal design artifacts (Models)
 - Formal: Precisely defined syntax and semantics
 - Translation of Design models to implementation

Model-based Development

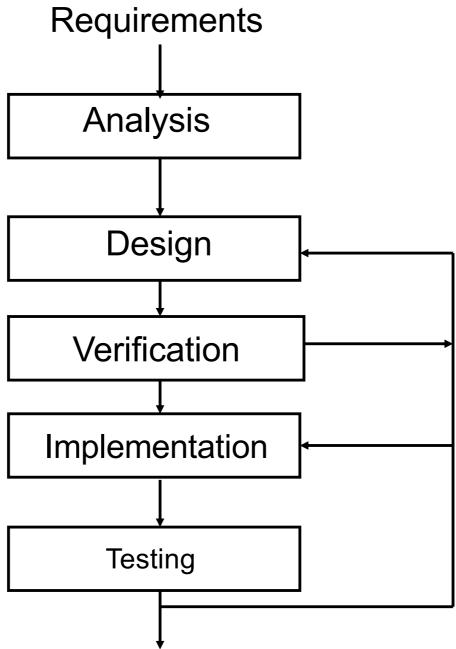
- Models: abstract & high level desc. of design objects
- Focus on one aspect at a time
- Less development and redesign time
- Implementation constraints can be placed on models
- Design exploration, evaluation & quick prototyping possible using models

New Paradigm

- Executable models essential
 - Simulation
- Can be rigorously validated
 - Formal Verification
- Models can be debugged and revised
- Automatic generation of final code
 - Traceability
- The paradigm

Model – Verify – Debug – CodeGenerate

Model-based Methodology



Tools

- Various tools supporting such methodologies
 - commercial and academic
- POLIS (Berkeley), Cierto VCC (Cadence)
- SpecCharts (Irvine)
- STATEMATE, Rhapsody (ilogix)
- Rose RT (Rational)
- SCADE, Esterel Studio (Esterel Technologies)
- Stateflow and Simulink (Mathworks)

Modeling Languages

- Models need to be formal
- Languages for describing models various exist
- High level programming languages (C, C++)
- Finite State Machines, Statecharts, SpecCharts, Esterel, Stateflow
- Data Flow Diagrams, Lustre, Signal, Simulink
- Hardware generation languages (Handel-C)
- Hardware description languages (VHDL, Verilog)
- Unified Modeling Language(UML)

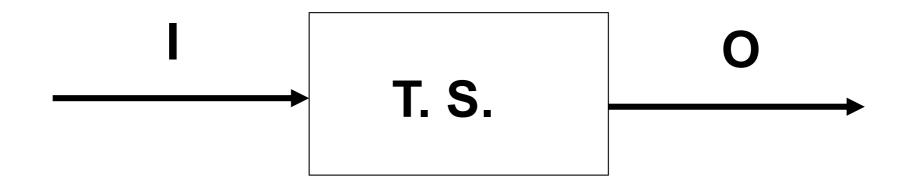
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Modeling Languages

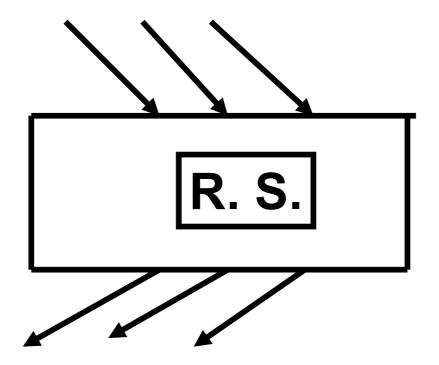
- Choice of languages depends on nature of computations modeled
- Seq. programming models for standard data processing computations
- Data flow diagrams for iterative data transformation
- State Machines for controllers
- HDLs for hardware components

Reactive Systems

- Standard Software is a transformational system
- Embedded software is reactive



Reactive Systems



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RS features

- Non-termination
- Ongoing continuous relationship with environment
- Concurrency (at least system and environment)
- Event driven
- Events at unpredictable times
- Environment is the master
- Timely response (hard and soft real time)
- Safety Critical
- Conventional models inadequate

Finite State Machines

- One of the well-known models
- Intuitive and easy to understand
- Pictorial appeal
- Can be made rigorous
- Standard models for Protocols, Controllers, HW

A Simple Example

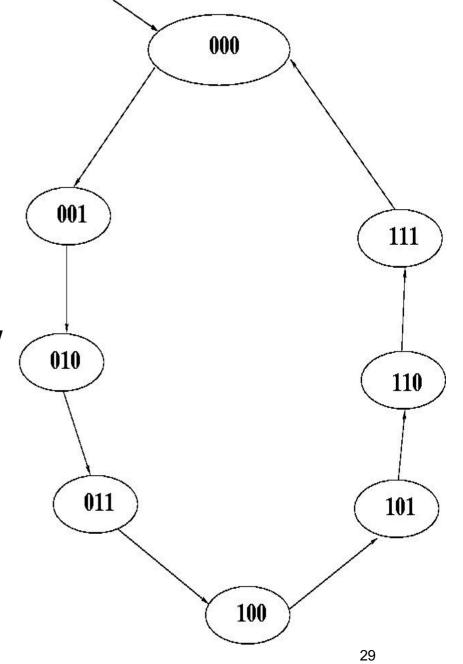
3 bit counter

C – count signal for increments

Resets to 0 when counter reaches maximum value

 Counter can be described by a program with a counter variable (Software Model)

 Or in detail using flip flops, gates and wires (Hardware model)



State Machine Model

- Counter behaviour naturally described by state machine
- States determine the current value of the counter
- Transitions model state changes to the event C.
- Initial state determines initial value of counter
- No final state (why?)

Precise Definition

< Q, q0, S, T>

- Q A finite no. of state names
- q0 Initial state
- S Edge alphabet
 Abstract labels to concrete event, condition and action
- T edge function or relation

Semantics

- Given syntax, a precise semantics can be defined
- Set of all possible sequences of states & edges
- Each sequence starts with the initial state
- Every state-edge-state triples are adjacent states connected by an edge
- Given FSM, unique set of sequences can be associated
- Language accepted by a FSM

Abstract Models

- Finite State machine model is abstract
- Abstracts out various details
 - How to read inputs?
 - How often to look for inputs?
 - How to represent states and transitions?
 - Focus on specific aspects
- Easy for analysis, debugging
- Redesign cost is reduced
- Different possible implementations
 - Hardware or Software
 - Useful for codesign of systems

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Intuitive Models

- FSM models are intuitive
- Visual
 - A picture is worth a thousand words
- Fewer primitives easy to learn, less scope for mistakes and confusion
- Neutral and hence universal applicability
 - For software, hardware and control engineers

Rigorous Models

- FSM models are precise and unambiguous
- Have rigorous semantics
- Can be executed (or simulated)
- Execution mechanism is simple: An iterative scheme

```
state = initial_state
loop
    case state:
        state 1: Action 1
        state 2: Action 2
        end case
end
```

Code Generation

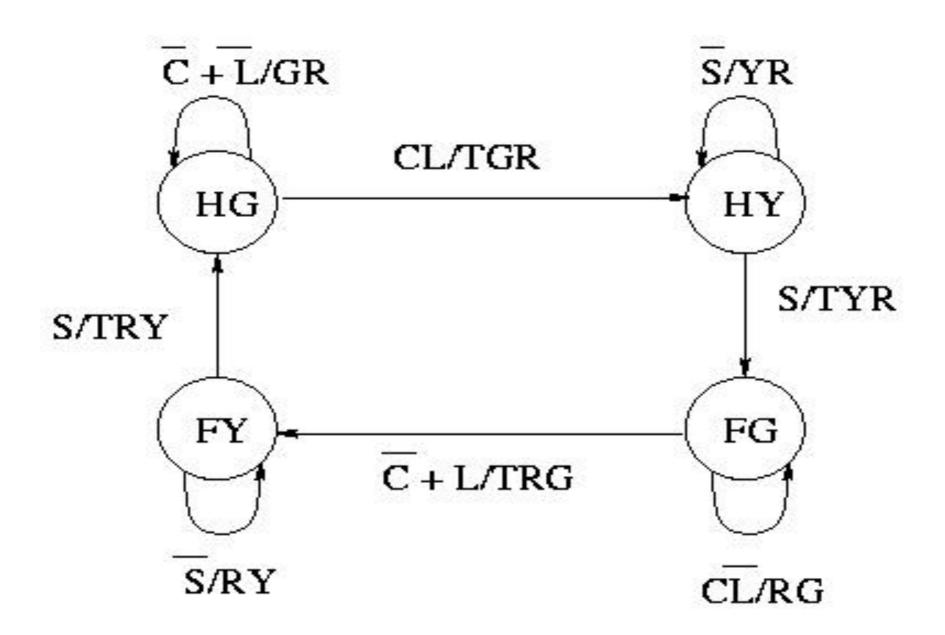
- FSM models can be refined to different implementation
 - Both HW and SW implementation
 - Exploring alternate implementations
 - For performance and other considerations
- Automatic code generation
- Preferable over hand generated code
- Quality is high and uniform

Another Example

A Traffic Light Controller

- Traffic light at intersection of Highway & Farm road
- Farm road sensors (signal C)
- G, R setting signals green and red
- S,L short and long timer signal
- TGR reset timer, set hway green & farm road red

State Machine



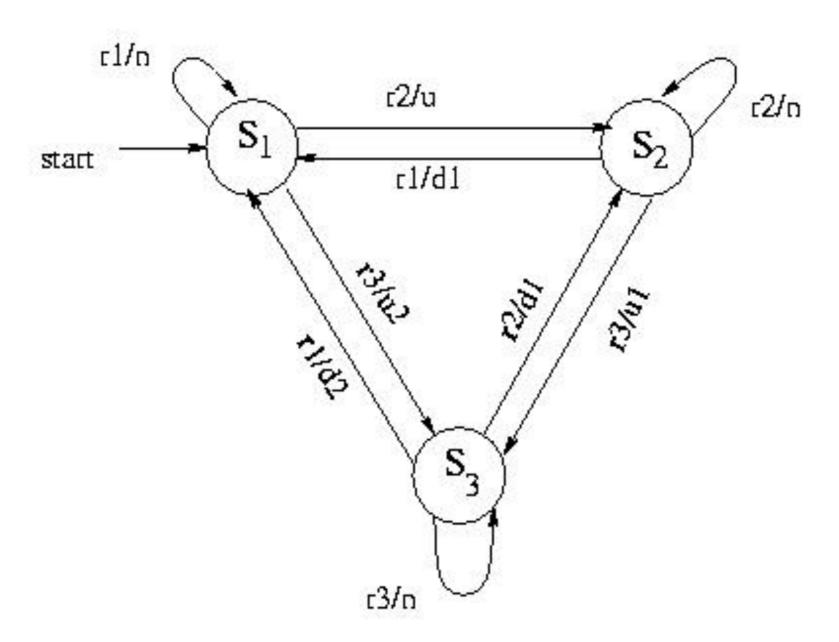
Another Example

A Simple Lift Controller

3-floor lift

- Lift can be in any floor
 - S_i in floor I
- Request can come from any floor
 - ri request from floor I
- Lift can be asked to move up or down
 - uj,dj up/down to jth floor

FSM model



Nondeterminism

- Suppose lift is in floor 2 (State S₂)
- What is the next state when requests r1 and r3 arrive?
 - Go to S₁
 - Or go to S_3
- The model non committal allows both
- More than one next state for a state and an input
- This is called nondeterminism
- Nondeterminism arises out of abstraction
- Algorithm to decide the floor is not modeled
- Models can be nondeterministic but not real lifts!

Nondeterminism

- Models focus attention on a particular aspect
- The lift model focused on safety aspects
- And so ignored the decision algorithm
 - Modeling languages should be expressive
 - Std. Programming languages are not
- Use another model for capturing decision algorithm
- Multiple models, separation of concerns
 - Independent analysis and debugging
 - Management of complexity
- Of course, there should be a way of combining different models

C-model

```
enum floors {f1, f2, f3};
enum State {first, second, third};
enum bool {ff, tt};
enum floors req, dest;
enum bool up, down = ff;
enum State cur floor = first;
req = read_req();
while (1)
{ switch (cur_floor)
 { case first: if (req == f2)
           \{up = tt; dest = f2;\}
         else if (req == f3)
           \{up = tt; dest = f3;\}
         else { up == ff; down = ff;};
         break;
```

C- model

```
case second: if (req == f3)
                  {up = tt; dest = f3;}
               else if (req == f1)
                      { up = ff; down = tt; dest = f1;}
                    else { up == ff; down = ff;};
               break;
case third: if (req == f2)
                up = ff; down = tt; dest = f2;
             else if (req == f1)
                    { up = ff; down = tt; dest = f1;}
                  else { up == ff; down = ff;};
             break; }; /* end of switch */
             req = read req(); } /* end of while */
```

Suitability of C

- C not natural for such applications
- Various problems
 - Events and states all modeled as variables
 - Not natural for event oriented embedded applications
 - States are implicit (control points decide states)
 - No abstract description possible
 - Commitment to details at an early stage
 - Too much work when design likely to be discarded

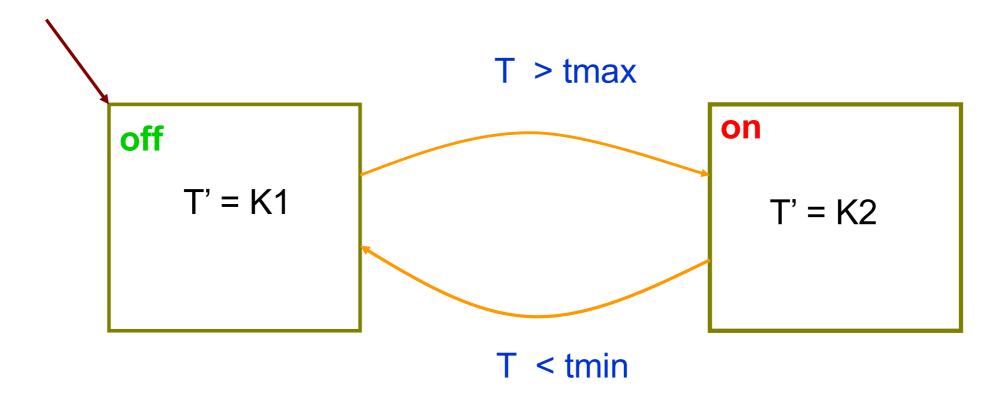
Exercise

Is the C model non-deterministic?

 What happens when two requests to go in different directions arrive at a state?

Yet Another example

A Simple Thermostat controller



Summary

- Finite number of states
- Initial state
- No final state (reactive system)
- Non determinism (result of abstraction)
- Edges labeled with events
- Behavior defined by sequences of transitions
- Rigorous semantics
- Easy to simulate and debug
- Automatic Code generation

Problems with FSMs

- All is not well with FSMs
- FSMs fine for small systems (10s of states)
- Imagine FSM with 100s and 1000s of states which is a reality
- Such large descriptions difficult to understand
- FSMs are flat and no structure
- Inflexible to add additional functionalities
- Need for structuring and combining different state machines

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