

Harry Langford 4/2/2

Question 4

a)

Proof by <sup>strong</sup> induction

RTP  $x^k \equiv y^k \pmod{m}$

$$x \equiv y \pmod{m} \Rightarrow x^k \equiv y^k \pmod{m}$$

for  $k = 1$

$$x \equiv y \pmod{m}$$

this is true by assumption.

Assume that  $x^n \equiv y^n \pmod{m}$

$$x^n \equiv y^n \pmod{m} \Rightarrow$$

$$\textcircled{1} \exists k, j \in \mathbb{Z} : x^n = y^n + km \Rightarrow$$

since  $x \equiv y \pmod{m}$  by assumption

$$\textcircled{2} \exists i \in \mathbb{Z} : x = y + im$$

Combining  $\textcircled{1}$  and  $\textcircled{2}$  gives

$$\exists i, j \in \mathbb{Z} : x \times x^n = (y + im) \times (y^n + km) \Rightarrow$$

$$\exists i, j \in \mathbb{Z} : x^{n+1} = y^{n+1} + m(jy) + m(iy^n)$$

$$\exists i, j \in \mathbb{Z} : x^{n+1} \equiv y^{n+1} \pmod{m}$$

Since the statement is true for  $k=1$ ,  
and the truth of the statement for  $k=1 \wedge k \leq n \Rightarrow$   
the truth for  $k = n+1$ , by induction by induction

$$x \equiv y \pmod{m} \Rightarrow$$

$$\forall k \in \mathbb{Z} : x^k \equiv y^k \pmod{m}$$

Include the case  $k=0$

$$\forall x \in \mathbb{Z}: x^0 = 1$$

$$\forall y \in \mathbb{Z}: y^1 = y$$

$$\forall x, y \in \mathbb{Z}: x^0 \circ y^0 = (x \circ y)^0$$

$$\text{So } \forall x, y \in \mathbb{Z}: x^0 \circ y^0 = (x \circ y)^0 \Rightarrow$$

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b) i) Alice exponentiates by  $g$  by  $a$  and mods by  $p$ .  
She then sends and records  $a$  and works out the inverse of  $a \pmod{p-1}$  as  $a^{-1}$ .  
 $K_1, a = 1 \pmod{p-1}$

She sends  $g^a$  to bob. ~~the~~  
who exponentiates by  $b$ .

He works out  $K_2: K_2 b \equiv 1 \pmod{p-1}$ .

Sends it back to Alice.

She exponentiates by  $K_1$ , Sends  $g^b$  to bob.

Bob exponentiates by  $K_2$ . This gives him  $g$ .

ii) Alice starts with  $g$ .

sends  $g^a \pmod{p}$

records  $K_1: K_1 a \equiv 1 \pmod{p-1}$

Bob exponentiates  $g^a \pmod{p}$  by  $b$ .

Sends  $g^{ab}$ .

records  $K_2: K_2 b \equiv 1 \pmod{p-1}$

Alice exponentiates by  $K_1$ .

Fig 2:  $g^{abK_1} \equiv g^{b(1+ip-1)} \equiv g^b \times g^{ib(p-1)} \equiv g^b \pmod{p}$   
She then sends this to bob.

Bob receives  $g^b \pmod{p}$   
Exponentiates by  $K_2$

Fig 3:  $g^{K_2 b} \equiv g^{1+j(p-1)} \equiv g \pmod{p}$

Bob now has the original message.

iii) Using Euclid's Extended algorithm

$$K_2 = 40$$

$$39^2 \equiv 21 \pmod{79}$$

$$21^2 \equiv 46 \pmod{79}$$

$$\text{So } 46 \times 39^{20} \equiv 1 \pmod{79}$$

$$46 \div 1, 0 \quad 39 \div 0, 1$$

$$7 \div 1, -1 \quad 39 \div 0, 1$$

$$7 \div 1, -1 \quad 40 \div -5, 6$$

$$3 \div 6, -7 \quad 4 \div -5, 6$$

$$1 \div -11, 13$$

$$\text{So } 39^{20} \equiv -11 \pmod{79}$$

$$39^{20} \equiv 68 \pmod{79}$$

So Alice sends 68.

iv)