

STL

## Standard Template Library

#include <bits/stdc++.h> → this contains all the library.

4 parts

- ① Algorithms ② Containers ③ Functions ④ Iterators

### Pairs

It is a part of utility library.  
pair <int, int> p;

→ data type of 1<sup>st</sup> element

→ data type of 2<sup>nd</sup> element.

to access 1<sup>st</sup> element in a pair  
we use pair.first (p. first here)

& for second we use pair.second  
(p. second here)

as our pair  
name is p.

Initializing pair;

pair<int, int> p = {1, 3};

### Nesting of pairs

pair<int, pair<int, int>> p1  
= {1, {3, 4}};

p1.first

p1.second.first

Here suppose we have to  
store 1, 3, 4 into a pair

{-, {-, -}}

{-, {-, -}}

p1.second.second

~~pair<int, int> arr[ ] = { {1, 2}, {2, 5}, {5, 13};~~

the address of arr → 0 1 2  
arr[0] = {1, 2} ; arr[1] = {2, 5} ; arr[2] = {5, 13} ;

Pairs can be treated as a data type

¶

arr[1].second.

Vectors → It is a container which is dynamic in nature. When we do not know about the size of the data ~~type~~ then we use vector.

Declaration of vector

vector<int> v;

v.push\_back(1); → Inserts 1 in the vector.

v.emplace\_back(2);

Generally emplace-back is faster than push-back.

Declaration of vector of pair

vector<pair<int, int>> vec;

vec.push\_back({1, 2});

vec.emplace\_back(1, 2);

While inserting a pair we have use curly braces if we are using push-back.

And curly braces are not necessary while using emplace-back because it automatically assumes them a pair.

`vector<int> v(5, 100);`  $\xrightarrow{\text{size}}$  Initialize with  $\{ \underline{100}, \underline{100}, \underline{100}, \underline{100}, \underline{100} \}$   
of vector containing 5 instances of 100

`vector<int> v(5);`  $\rightarrow \{ \underline{-}, \underline{-}, \underline{-}, \underline{-}, \underline{-} \}$   
some garbage value

`vector<int> v1(5, 20);`  $\rightarrow \{ \underline{20}, \underline{20}, \underline{20}, \underline{20}, \underline{20} \}$

Copying a vector into another vector

`vector<int> v2(v1);`

$v_2 = \{ \underline{20}, \underline{20}, \underline{20}, \underline{20}, \underline{20} \}$

Accessing elements in a vector

Using iterator

Syntax:

`vector<int>::iterator it = v.begin();`

data structure

↓  
the data type

$\star(v.begin())$

It gives value  
at the memory/  
address at  
`v.begin()`.

↓  
name of iterator

→ It points  
directly to  
the memory  
means pointing  
`v.begin` will  
give the address  
of first element  
of a vector.

↓  
name of iterator

`vector<int> v; iterator it = v.begin();`

$it++$

$\{10, 15, 6, 7\}$

$it++;$

`cout << *it << " ";` O.P.  $\rightarrow 10.$   $v.begin$  is the address of the first element of the vector

$it = it + 2$

`cout << *it << " ";` O.P.  $\rightarrow 6.$

$v.begin$  is the address of the first element of the vector

### Other iterators

`end, rend (reverse end) & rbegin (reverse begin)`

`vector<int> v; iterator it = v.end();`

`vector<int> v; iterator it = v.rend();`

`vector<int> v; iterator it = v.rbegin();`

$v = \{10, 20, 30, 40\}$

$v.end$  will not point to 40.

$v.end$  points to a memory element which is eight after the last element.

$v.end()$

(Never Used)

$\{10, 20, 30, 40\}$

$v.rbegin$

(Never Used)

$\{10, 20, 30, 40\}$

$it++$  it comes to 30.

```
cout << v[0] << " " << v.at(0);    {10, 20, 30}  
cout << v.back() << " ";
```

Points to the last element in a vector.

### Printing the entire vector

```
for(vector<int>::iterator it = v.begin(); it != v.end();  
     it++)
```

```
{ - cout cout << *(it) << " ";
```

}

→ Here auto automatically assigns it to a vector iterator  
for (auto it = v.begin(); it != v.end(); it++)

```
{ cout << *(it) << " ";
```

}

→ it on v.

for (auto it : v) ex- auto b = <sup>66</sup>array<sub>11</sub>; ex- auto a = 5 int.

```
{ cout << it << " ";
```

→ Here the compiler knows 5 is an integer so it automatically assigns the data type.

It iterates on the data type not an iterator here. We don't

use value at (\*)

operator here.

## Erasing elements from a vector

- `erase (iterator)` Used for deleting a single element  
↓  
location of the address which we want to delete.

Ex → `v.erase(v.begin() + 1);` {10, 20, 30, 40}

It deletes 20. New v = {10, 30, 40}.

## Deleting multiple elements at the same time

- `erase (start element's address , end address)`

↓  
It is after the element. Which

Ex - `v = {10, 20, 30, 40, 50}`. we want to delete.  
Suppose we want to delete 20, 30 & 40

So,

`v.erase (v.begin() + 1 , v.begin() + 4);`

at the end we have to give address of the next element.

it is like [Begin, end).

## Inserting into a vector

vector <int> v {2, 100}; // {100, 100}

v.insert(v.begin(), 300); // {300, 100, 100};  
→ value which  
↳ address where we want  
we want to insert  
to insert

ex - {10, 20, 30, 40}  
↑  
⑤

If we want to insert 5 after 10  
we will use

v.insert(v.begin() + 1, 5);

## Inserting multiple elements

v.insert(v.begin(), 2, 10);  
→ No. of elements we want to  
insert  
↳ address where → The number we  
want to insert.  
we want to  
insert.

So doing this on  $v = \{40, 20, 30\}$   
will make  $v = \{10, 40, 20, 30\}$ .

For inserting a vector into another vector.

$$v = \{30, 10, 10, 100, 100\}$$

vector<int> copy(2, 50); // {50, 50}

v.insert(v.begin(), copy.begin(), copy.end());

$$\text{New } v = \{50, 50, 30, 10, 10, 100, 100\}.$$

If we want to insert a portion of the vector,

This inserts the entire vector copy into the vector v.

v.insert(v.begin(), copy.begin(), copy.end());

↓  
index where  
we want to  
add copy  
vector.

The elements  
from which  
we want to

insert.

The element #  
after the element  
which we wanted  
to insert  
because  
~~(at [start, end])~~

v.size() → gives the number of elements  
in a vector.

v.pop\_back() → removes the last element  
from the vector.

v. swap(v2);

It swaps the entire vectors

ex → initially

$$v_1 = \{10, 20\}$$

$$v_2 = \{30, 40\}$$

v1.swap(v2);

$$v_1 = \{30, 40\} \quad v_2 = \{10, 20\}$$

v. erase → clears the entire vector

cout << v.empty(); It is a boolean func<sup>n</sup>

Returns 1 if v is empty  
else returns 0.

### Lists

It is similar to vector.

The only difference is it gives front operations as well. It is a container & dynamic in nature.

list <int> ls; → declaration

ls.push\_back(2); // {2}

ls.emplace\_back(4); // {2, 4}

ls.push\_front(5); // {5, 2, 4}

ls.emplace\_front(6); // {6, 5, 2, 4}

// rest are same as vector.

// begin, end, rbegin, rebegin, clear

// insert, swap and size.

an insert in a vector is more costlier (in terms of Time complexity) as compared to push front in a list.

## Degue

exactly similar to list & vector

deque < int > dq;

dq.push-back(1); // {1}

dq.emplace-back(2); // {1, 2}

dq.push-front(4); // {4, 1, 2}

dq.emplace-front(3); // {3, 4, 1, 2} shows the last element

dq.pop-back(); // {3, 4, 1} dq.back();

dq.pop-front(); // {4, 1} dq.front(); shows the first element.

// just same as vector

// begin, end, rbegin, rend, clear, insert, sized swap.

## Stack

stack has LIFO

↳ last in first out

stack < int > st;

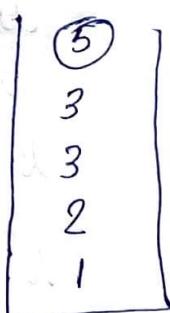
st.push(1); // {1}

st.push(2); // {2, 1}

st.push(3); // {3, 2, 1}

st.push(3); // {3, 3, 2, 1}

st.emplace(5); // {5, 3, 3, 2, 1}



`cout << st.top();` ~~sets~~ points the top element of the stack. it does nothing to the stack. 5 in our case. top element.

Index based accessing is not allowed in stack.

`st.pop(); {3,3,2,1}`

`cout << st.top(); 113`

`cout << st.size(); 114`

`cout << st.empty(); false`

`stack<int> st1, st2;`

`st1.swap(st2);`

push } all are  
pop }  $O(1)$  in  
top time complexity

## Queue

similar to stack

FIFO

↳ First In First Out.

(Imagine queue as a line of people for buying ticket)

`queue<int> q;`

`q.push(1); 11{1}`

`q.push(2); 11{1,2}`

`q.emplace(4); 11{1,2,4}`

`q.back() += 5 → adds 5 to the last element`

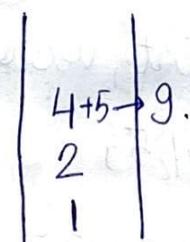
`cout << q.back(); //points 9.`

`//q is {1,2,9}`

`cout << q.front(); //points 1.`

`q.pop(); // {2,9}`

`cout << q.front(); //points 2.`



All the operations are happening in  $O(1)$  Time complexity.

// size swap empty

same as stack

## Priority Queue

The elements which has the largest value stays at the top.

// Max heap

priority-queue<int> pq;

pq.push(5); // {5}

pq.push(2); // {5,2}

pq.push(8); // {8,5,2}

pq.emplace(10); // {10,8,5,2}

cout << pq.top(); // prints 10.

pq.pop(); // removes the greatest element.

cout << pq.top(); // prints 8.

10
8
5
2

A tree is maintained inside.  
Data is not stored in linear fashion

push } main function  
top }  
pop }

push  $\rightarrow O(\log n)$

pop  $\rightarrow O(\log n)$

top  $\rightarrow O(1)$

// Minimum heap  $\rightarrow$  The priority queue which stores minimum element at the top.

priority-queue<int, vector<int>, greater<int>> pq;

pq.push(5); // {5}

pq.push(2); // {2,5}

pq.push(8); // {2,5,8}

pq.emplace(10); // {2,5,8,10}

2
5
8
10

cout << pq.top(); // prints 2.

## Set

stores everything in a sorted order and stores unique values only.

A tree is maintained

```
set <int> st;
st.insert(1); // {1}
st.emplace(2); // {1, 2}
st.insert(2); // {1, 2} → It won't add 2 as it stores only unique elements.
st.insert(4); // {1, 2, 4}
st.insert(3); // {1, 2, 3, 4}
```



```
// {1, 2, 3, 4, 5}
auto it = st.find(3);
```

It returns an iterator which points to the 3. (points to the address)

```
// {1, 2, 3, 4, 5}
```

```
auto it = st.find(6)
```

The 6 is not in the set.

So when the element is not in the set it always returns ~~set~~.st.end().

it points to right after end. {1, 2, 4, 5, 6}

```
st.erase(5); // erases 5 → O(log n)
```

{1, 2, 4, 6} & maintains the sorted order.

int count = st.count(1);  
It can ~~have~~ give only 1 or 0.

0 when it doesn't exist

1 when it is present and only 1 because set contains only unique elements.

we can also erase by giving address of the iterator

```
auto it = st.find(3);
```

```
st.erase(it); → it takes O(1) time.
```

Time Complexity -

// {1, 2, 3, 4, 5}

auto it1 = st.find(2);

auto it2 = st.find(4);

st.erase(it1, it2);

// after erase {1, 4, 5} [first, last).

In set everything  
happens in O(log n)  
time complexity.

// lower\_bound() & upper\_bound() functions  
work in the same way as they work  
in vectors.

### Syntax

auto it = st.lower\_bound(2);

auto it = st.upper\_bound(3);

size, empty, swap everything is similar  
to vector, begin, end etc.

### Multiset

It can store multiple occurrence unlike  
set.

It stores in sorted manner.

multiset < int > ms;

ms.insert(1); // {1}

ms.insert(1); // {1, 1}

ms.insert(1); // {1, 1, 1}

ms.erase(1); // all 1's are erased

If we want to erase 2 1's  
ms.erase(ms.find(1), ms.find(1)+2);

It finds 1 and deletes 2 of  
them.

int count = ms.count(1); → Returns no of  
1's.

ms.erase(ms.find(1));

Here we erase the iterator to delete just one occurrence  
of 1 instead of saying erase element as it deletes all the 1's.

But if we say  
erase address  
it only erases  
that portion

## Unordered Set

Everything same as set. It does not store in sorted order. It only stores unique elements.

All the operations have Time complexity  $O(1)$ .  
All the functions are same  
only the lower bound and upper bound functions do not work.

For worst case - Time complexity  $O(N)$   
Happens in once in a millennium

Map → It is a container  
map stores everything in respect of {key, value}

Keys & Values

Keys must be unique

Key can be any data type

map<int, int> mpp;

data type of key

Values can be same.

Value can also

be of any data type.

Map works in  $O(\log N)$  time

complexity.

map<pair<int, int>, int> mpp;

Here key is a pair of 2 integers

Value is 1 integer

map<int, pair<int, int>> mpp;  
key is 1 integer  
value is a pair of 2 integers

`map<int, int> mpp;`

`mpp[1] = {2, i}`  
    ↓ key   ↑ value

`mpp.emplace({3, 1});`  
`mpp.insert({2, 4});`

$\begin{Bmatrix} \{1, 2\} \\ \{2, 4\} \\ \{3, 1\} \end{Bmatrix}$

These three lines  
will be stored  
like this as  
it should be in  
sorted order.

`mpp[{2, 3}] = 10;` → It is for `map<pair<int, int>, int> mpp;`

{  
    {1, 2}  
    {2, 4}  
    {3, 1}  
}  
}      Output (In the sorted order)  
                of key  
For each loop

`for(auto it : mpp)`

{ `cout << it.first << " " << it.second << endl;`  
}  
}

`cout << mpp[1];` // it says 2.

`cout << mpp[5];` // 0 or null (it doesn't exist).

It gives iterator to 3, 1.

`auto it = mpp.find(3);`

`cout << *it.second;`

`auto it = mpp.lower_bound(2);`

`auto it = mpp.upper_bound(3);`

`auto it = mpp.find(5);`

∴ 5 is not in the map it will point to  
`mpp.end();`

`clear`, `empty`, `size`, `swap` are same as above.

~~auto~~ Multimap similar to map.

It can store duplicate keys.  
everything in sorted order.

## Unordered map

~~It can store not duplicate keys.~~ It will not store in sorted order.

Almost in every case unordered map works in  $O(1)$  time complexity

In worst case  $\rightarrow O(N)$  Time complexity  
(Happens very rarely)

## Algorithms

Algorithms

position of starting iterator [Start, end) .

sort ( $a, a+n$ ) ; //For arrays

sort ( $v.begin(), v.end()$ );

position of last iterator

$\{ 1, 3, 5, 2 \}$

$a+2$        $a+4$

↑      ↑

$\text{sort}(a+2, a+4)$ ;

If I want

{1, 3, 2, 5}.

to sort

For sorting in descending order this portion

sort ( $a$ ,  $a+n$ , greater<int>);

$\text{pair } \langle \text{int}, \text{int} \rangle a[] = \{\{1, 2\}, \{2, 1\}, \{4, 1\}\};$

// Sorting it in order of increasing 2<sup>nd</sup> element.  
and if second element is same sort in  
according to first element but in descending.

sort (a, a+n, comp);

↳ Here comp is a  
self-written boolean comparator.

$\{\{4, 1\}, \{2, 1\}, \{1, 2\}\};$

bool comp ( $\&$  pair  $\langle \text{int}, \text{int} \rangle p_1$ , pair  $\langle \text{int}, \text{int} \rangle p_2$ )

{ If ( $p_1.\text{second} < p_2.\text{second}$ )  
return true;

If ( $p_1.\text{second} > p_2.\text{second}$ )  
return false;

If ( $p_1.\text{second} = p_2.\text{second}$ )

{ If ( $p_1.\text{first} > p_2.\text{first}$ )  
return true;

else

return false;

}

whenever writing  
comp just think  
in terms of  
2 pairs only.

int num = 7; → It returns the no of set bits (1).

int count = \_\_builtin\_popcount();

long long num = 16576578657869;

int count = \_\_builtin\_popcountll();

If its long long .

string s = "123";

sort(s.begin(), s.end());

do {

cout << s << endl;

} while (next\_permutation(s.begin(), s.end()));

→ It returns all the possible permutation .

s = 123

s = 132

s = 213

s = 231

s = 312

s = 321

s → Null.

when there are no more permutations left the while loop cond<sup>n</sup> becomes false .

If s was "123" "231"

s = 231

s = 312

s = 321

So if we want to have all the permutations we should start from the sorted string .

int maxi = \*max\_element(a, a+n); null.

↓  
It gives max element of an array / vector

int mini = \*min\_element(a, a+n);