

Pseudo code

DFS()

Same as ppt

Difference:

If conflict = **true**

Then return

DFS_VISIT(Vertex &u)

Same as ppt

Difference:

If visit a gray node

Then conflict → **true**

Return

Topological_sort

When a vertex is finished, push_back in **out** vector

main()

save vertex into **arr** by information in first line

every vertex has a **near** vector

push_back the index of classes which have to be taken after the current class

ex, 1 0 → vertex0.push_back(1)

DFS()

If conflict = **true**

Then print "Conflict"

Else

Then print **out** vector from behind

Result and Analysis

The sooner a vertex turns into black, the later the class should be, the last vertex turn into black should be the root, the most essential class that should be taken before any class is taken.

At first, when we build up the tree(before entering the DFS), we spend $\theta(V)$ to push_back every vertex, and $\theta(E)$ to push_back every edge, so the time complexity building will be $\theta(V+E)$.

Every vertex will only be visited once, and after that, all the edges(ex, if 2 in vertex0.near == 0 \rightarrow 2, \rightarrow is an edge) pointing out from it will never be visit again, as an edge will only exist in one vertex's **near** vector, if exist in both, it means that there are two edges between the two vertex, which will cause a conflict, the bonus part that I'll explain in the next paragraph. Furthermore, the vertex will become black to keep others from visiting, so DFS() will at most be **$O(V+E)$** , as all vertices and edges will only be visited once.

Bonus:

Any connected component should get black from the leaf all the way to root, moreover, when we are working on one connected component, the other component should be all black(already finished) or all white(hasn't start working on); therefore, the colors a edge can find on the other side should only be white and black, if the edge find a gray vertex on the other side, that will be the current vertex's ancestor, which will cause a conflict.