# Report 1-B

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Video link: https://youtu.be/HIbxGNPOhXs

# Questions:

1. What is a kernel function? What is a system call?

Kernel function: function that is called in the kernel operation, including access computer resource, memory management, and so on.

System call: syscall, the interface between process and OS, when the user processes need the service from OS, it will use syscall to send request to the OS kernel, call kernel function, each kernel function has its syscall number

2. What is KASLR? What is it for?

kernel address space layout randomization, to randomly place the address of the kernel, having a little offset from the link address, to make sure the outside won't have direct access to the kernel

3. What are GDB's non-stop and all-stop modes?

Non-stop mode: when the thread stops to report a debugging event, only the thread is stopped

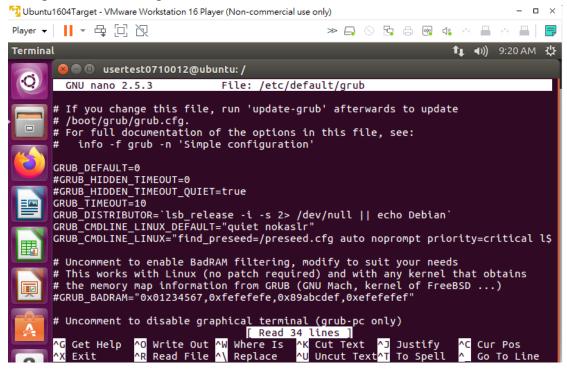
All-stop mode: when the program stops under GDB for any reason, all threads of execution stop

4. Explain what the command echo g > /proc/sysrq-trigger does.

The command gives back the control of the vm back to the machine which is debugging the current machine with gdb

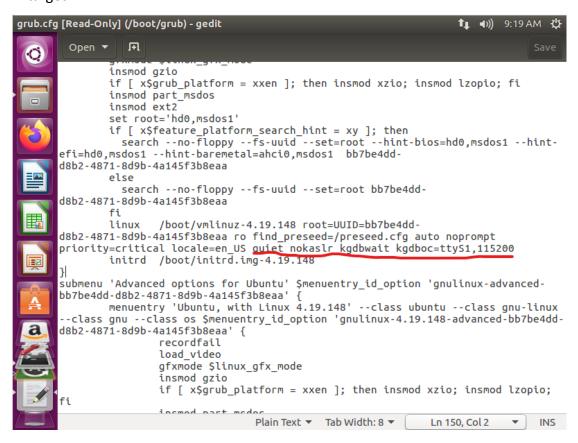
#### Screenshot 1:

Add "nokaslr" in the comment to make sure that the host machine can directly debug the kernel in target machine



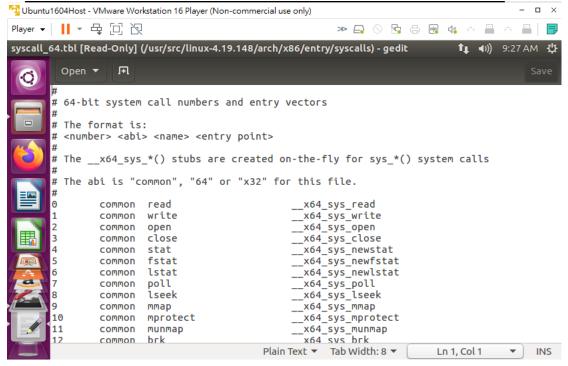
#### Screenshot 2:

Same as screenshot 1, to make sure that the host can get direct access to the kernel in target



## Screenshot 3:

# Directly check the kernel functions in the kernel package



# Screenshot 4:

# Make sure the syscall tables find online are coherent to one another

79	getcwd	man/ cs/	0x4f	char *buf	unsigned long size
80	chdir	man/ cs/	0x50	const char *filename	-
81	fchdir	man/ cs/	0x51	unsigned int fd	-
82	rename	man/ cs/	0x52	const char *oldname	const char *newname
83	mkdir	man/ cs/	0x53	const char *pathname	umode_t mode
84	rmdir	man/ cs/	0x54	const char *pathname	-
85	creat	man/ cs/	0x55	const char *pathname	umode_t mode

# Screenshot 5:

76	sys_truncate	const char *path	long length	
77	sys_ftruncate	unsigned int fd	unsigned long length	
78	sys_getdents	unsigned int fd	struct linux_dirent *dirent	unsigned int count
79	sys_getcwd	char *buf	unsigned long size	
80	sys_chdir	const char *filename		
81	sys_fchdir	unsigned int fd		
82	sys_rename	const char *oldname	const char *newname	
83	sys_mkdir	const char *pathname	int mode	
84	sys_rmdir	const char *pathname		
85	sys_creat	const char *pathname	int mode	
86	sys_link	const char *oldname	const char *newname	
87	sys_unlink	const char *pathname		

# Screenshot 6:

76	truncate	sys_truncate	fs/open.c
77	ftruncate	sys_ftruncate	fs/open.c
78	getdents	sys_getdents	fs/readdir.c
79	getcwd	sys_getcwd	fs/dcache.c
80	chdir	sys_chdir	fs/open.c
81	fchdir	sys_fchdir	fs/open.c
82	rename	sys_rename	fs/namei.c
83	mkdir	sys_mkdir	fs/namei.c
84	rmdir	sys_rmdir	fs/namei.c
85	creat	sys_creat	fs/open.c
86	link	sys_link	fs/namei.c
87	unlink	sys_unlink	fs/namei.c

## Screenshot 7:

As we are making a break point in mkdir, we have to find the actual function that the syscall call, which is do\_mkdirat, so we use command "break do\_mkdirat" to set up the break point

```
Ubuntu1604Host - VMware Workstation 16 Player (Non-commercial use only)
Player ▼ | | | ▼ 뒂 🖸 🕏
                                                                    namei.c [Read-Only] (/usr/src/linux-4.19.148/fs) - gedit
                                                                                           1 ■1) 8:21 AM 😃
                    mode &= ~current_umask();
                                                                      Q mkdir
                                                                                           error = security_path_mkdir(&path, dentry, md
                error = vfs_mkdir(path.dentry->d_inode, dentry, mode);
done_path_create(&path, dentry);
if (retry_estale(error, lookup_flags)) {
        lookup_flags |= LOOKUP_REVAL;
        goto retry;
}
                return error;
       SYSCALL_DEFINE3(mkdirat, int, dfd, const char __user *, pathname, umode_t, mode)
                return do_mkdirat(dfd, pathname, mode);
        SYSCALL_DEFINE2(mkdir, const char __user *, pathname, umode_t, mode)
                return do_mkdirat(AT_FDCWD, pathname, mode);
       int vfs_rmdir(struct inode *dir, struct dentry *dentry)
                int error = may_delete(dir, dentry, 1);
                if (error)
                          return error;
```

```
😑 💷 root@ubuntu: /boot/kgdb-image
              [Inferior 1 (Remote target) exited with code 03]
              (gdb) quit
              root@ubuntu:/boot/kgdb-image# exit
              exit
              userhost0710012@ubuntu:/boot/kgdb-image$ sudo su
             [sudo] password for userhost0710012:
root@ubuntu:/boot/kgdb-image# gdb ./vmlinux
GNU gdb (Ubuntu 7.11.1-0ubuntu1~16.5) 7.11.1
             CODYRIGHT (C) 2016 Free Software Foundation, Inc. License GPLv3+: GNU GPL version 3 or later <a href="http://gnu.org/licenses/gpl.html">http://gnu.org/licenses/gpl.html</a> This is free software: you are free to change and redistribute it. There is NO WARRANTY, to the extent permitted by law. Type "show copying" and "show warranty" for details. This GDB was configured as "x86_64-linux-gnu". Type "show configuration" for configuration details.
rype "snow configuration" for configuration details.

For bug reporting instructions, please see:
<a href="http://www.gnu.org/software/gdb/bugs/">http://www.gnu.org/software/gdb/bugs/</a>.

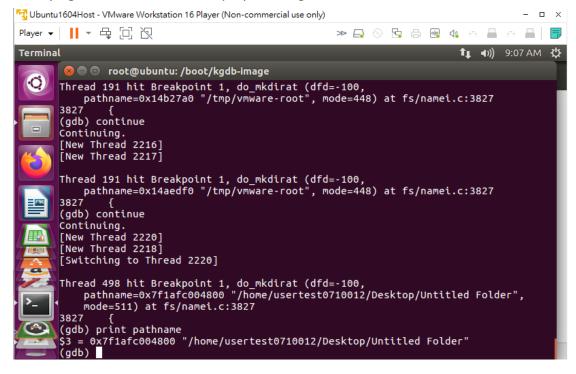
Find the GDB manual and other documentation resources online at:
<a href="http://www.gnu.org/software/gdb/documentation/">http://www.gnu.org/software/gdb/documentation/</a>.

For help, type "help".

Type "apropos word" to search for commands related to "word"...
Į.
             Reading symbols from ./vmlinux...done.
(gdb) target remote /dev/tty
Display all 200 possibilities? (y or n)
             (gdb) target remote /dev/ttyS1
Remote debugging using /dev/ttyS1
kgdb_breakpoint () at kernel/debug/debug_core.c:1086
                                                           wmb(); /* Sync point after breakpoint */
              1086
              (gdb) continue
               .
Continuing.
```

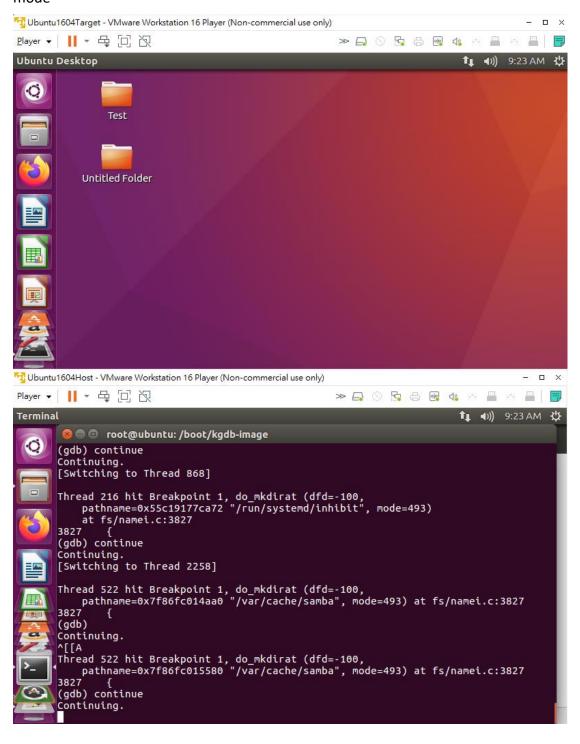
## Screenshot 8:

When the break point is hit, the path will show where it is hit, the example is that we are trying to call mkdir in desktop by creating a new folder



## Screenshot 9:

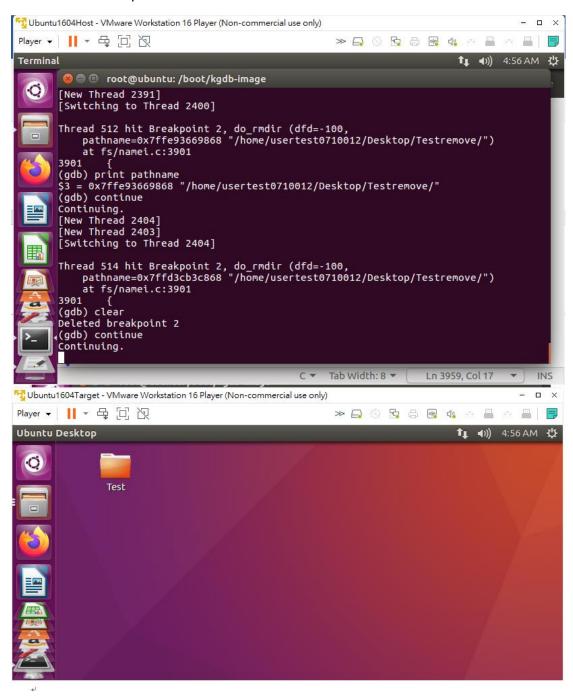
Actually we get through many break point before we reach the status that the target can back to normal mode, as the mkdir is called in many process that we don't know, however, at last, the folder is create successfully and the target resume to work mode



# Do it yourself:

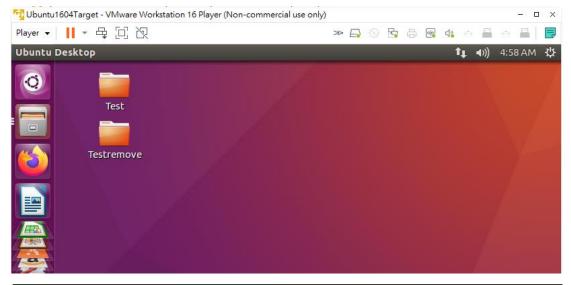
#### Screenshot 10:

The machine is responsive



#### Screenshot 11:

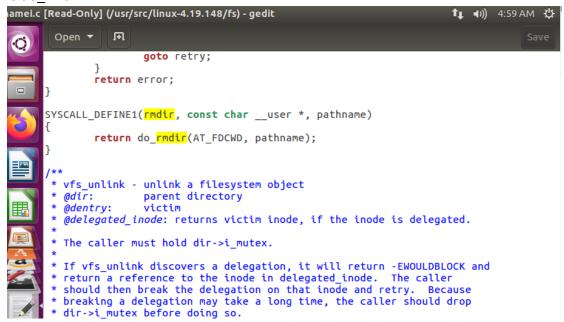
I choose the "rmdir" syscall function as the "do it yourself" example, and I plan to display the example by deleting a folder from the desktop, therefore, there is a folder called "Testremove" to help me explain the idea, and enter the rmdir syscall to hit the breakpoint



```
🔞 🖯 🗊 root@ubuntu: /
root@ubuntu:/home/usertest0710012/bin# cd ../
root@ubuntu:/home/usertest0710012# cd desktop
bash: cd: desktop: No such file or directory
root@ubuntu:/home/usertest0710012# cd Desktop
root@ubuntu:/home/usertest0710012/Desktop# ls
root@ubuntu:/home/usertest0710012/Desktop# mkdir Testremove root@ubuntu:/home/usertest0710012/Desktop# cd ../ root@ubuntu:/home/usertest0710012# cd ../ root@ubuntu:/home# cd ../ root@ubuntu:/home# cd ../ root@ubuntu:/# ls bin dev initrd.img lib64 mnt root snap
                                                                                                                                        vmlinuz
                           initrd.img.old lost+found opt
boot
               etc
                                                                                                                STV
                                                                                                                              UST
                                                                                                                                        vmlinuz.old
cdrom home lib
                                                              media
                                                                                                   sbin
                                                                                                                sys
                                                                                                                              var
root@ubuntu:/# rmdir /home/usertest0710012/Desktop/Testremove/
root@ubuntu:/# rmdir /home/usertest0710012/Desktop/Testremove/
rmdir: failed to remove '/home/usertest0710012/Desktop/Testremove/': No such fil
e or directory
root@ubuntu:/# rmdir /home/usertest0710012/Desktop/Testremove/
```

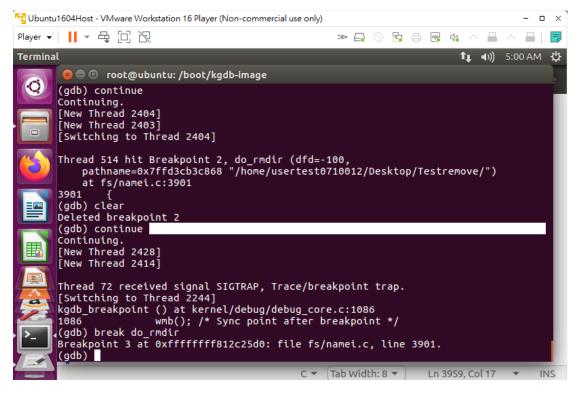
#### Screenshot 12:

Find the rmdir function in the syscall file and find out that the actual function it calls is do rmdir



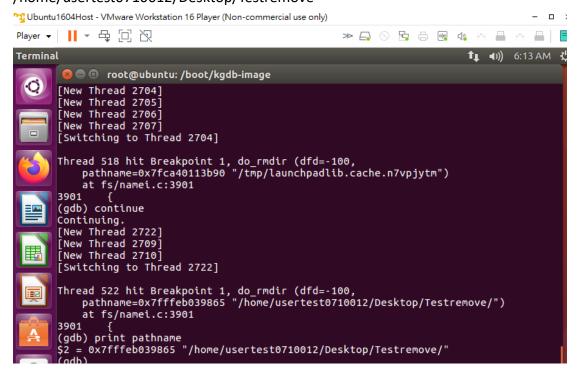
## Screenshot 13:

Create a break point at do rmdir



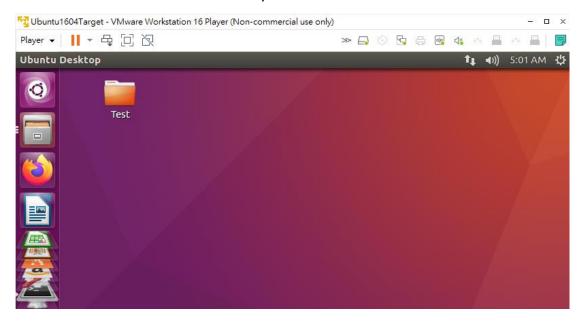
#### Screenshot 14:

The path name showed that the breakpoint in rmdir function is hit in /home/usertest0710012/Desktop/Testremove



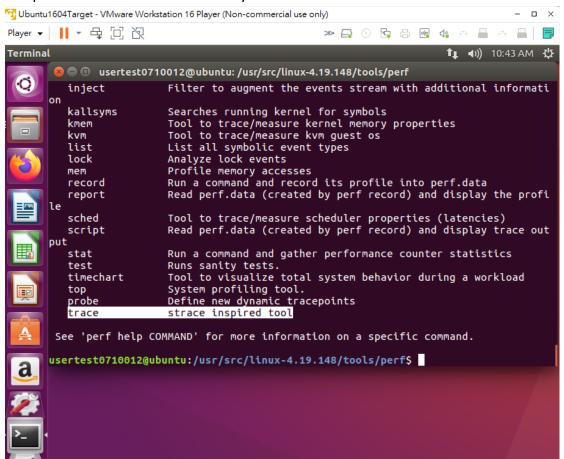
## Screenshot 15:

The "Testremove" folder is successfully removed



#### Screenshot 16:

The perf trace function is successfully installed



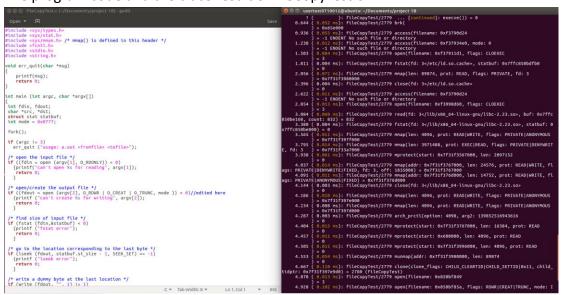
#### Screenshot 17:

The program code and the trace result of emptyTest.c

```
### Comparison of Comparison o
```

# Screenshot 18

The program code and the trace result of fileCopyTest.c



# Screenshot 19:

We can see that the trace result of fileCopyTest.c has a bulk of text more than emptyTest.c, which should be the actual code that is executed in fileCopyTest.c

1. What are these functions: clone, mmap, write and open?

Clone: create new(child) process

Mmap: create a new mapping in the virtual address space of the calling process

Write: write the context in the buffer into the file that is referred to

Open: open the file that is referred to

2. Why is there no fork system call? What is the difference between fork and clone?

As it is unnecessary to create a new process that is totally independent to the origin process, clone can save more resources

The child process created by clone allows them to share parts of its execution context with the calling process

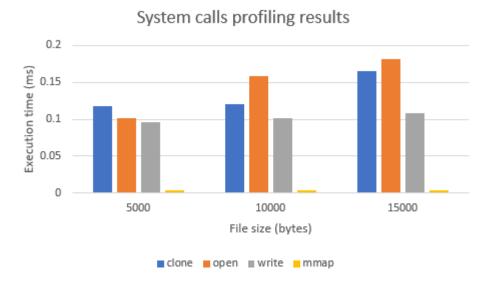
3. Will the functions' execution time be longer if the file is bigger?

Yes, as the data the program take in and print out is different. More data should take more time.

#### Screenshot 20:

-from fastest to slowest-:mmap, write, clone, open

Though open takes less time when executing the 5000 bytes file, the average time is longer than clone



# Screenshot 21:

a.read the perf.data file and display the context

b.The percentage under the overhead column shows the overall samples collected in the corresponding function and its children, therefore, the fileCppyTest function and its chilren, which is dynamically linked and have the shared library named "mptbase" contain 100% of the overall samples, while the fileCopyTest function(the one that come from the kernel) and its children contains 0% of the overall samples.

Actually can'y really understand the chart, as it looks totally different with the one I find online, no matter it is the parents function display or the percentage. Moreover, I think my perf.data may have some problem as the percentage is shown in red font, which is not normal, however, no matter how many time I tried to rerun the program or "record" a new perf.data, the result is always the same.

```
■ usertest0710012@ubuntu: ~/Documents/project 1B
# To display the perf.data header info, please use --header/--header-only opti
#
# Total Lost Samples: 0
# Samples: 1 of event 'cpu-clock'
# Event count (approx.): 250000
#
               Self Command
  Children
                                   Shared Object
                                                              Symbol
                        fileCopyTest [mptbase]
                                                           [k] mpt_put_msg_frame
             ---entry_SYSCALL_64_after_hwframe
                do_syscall_64
0xffffffff8109a6d8
                do_group_exit
do_exit
                 task_work_run
                 ____fput
__fput
                 ext4_release_file
ext4_alloc_da_blocks
                 filemap_flush
                  _filemap_fdatawrite_range
                 do_writepages
                ext4_writepages
blk_finish_plug
                 blk_flush_plug_list
                queue_unplugged
__blk_run_queue
scsi_request_fn
                 scsi_dispatch_cmd
                mptspi_qcmd
mptscsih_qcmd
                 mpt_put_msg_frame
                 0.00% fileCopyTest [kernel.kallsyms] [k] entry_SYSCALL_64_af
              ---entry_SYSCALL_64_after_hwframe
                do_syscall_64
0xffffffff8109a6d8
                 do_group_exit
                 do exit
                 task_work_run
                   __fput
                 ____fput
                 ext4_release_file
                 ext4_alloc_da_blocks
                 filemap_flush
                 __filemap_fdatawrite_range
do_writepages
                 ext4 writepages
usertest0710012@ubuntu:~/Documents/project 1B$
```

# The example I found online:

```
Children
             Self Symbol
            0.00% __libc_start_main
100.00%
          -- __libc_start_main
            0.00% main
100.00%
          --- main
             __libc_start_main
100.00%
           40.00% bar
             main
             __libc_start_main
 60.00%
           60.00% foo
         --- foo
             main
               _libc_start_main
```

## Screenshot 22:

I choose "perf timechart" as my example, the function is simply record(perf.data) the time table (system behavior) of the program and print

(Tool to visualize total system behavior during a workload)

The function will output a file called output.svg(graphic file), which is shown in below, we can see the execution of the function in the graph, showing the relative time and feature of parallel among the functions.

