Interpreting Cubical Type Theory in Appropriate Presheaf Toposes

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Cubical Type Theory

- CTT: Extension of dependent type theory (with Σ , Π -types and a universe) by an **interval**, a **lattice of faces**, **path types** and certain operations for type families (**composition** and **glueing**).
- Devised by Bezem, Cohen, Coquand, Huber and Mörtberg [BCH14, CCHM16] in 2014-2016 as an intensional type theory which validates Voevodsky's Univalence Axiom and has computational meaning.
- Further developments by Orton and Pitts [OP16] as well as Birkedal, Bizjak, Clouston, Gratwohl, Spitters and Vezzosi [BBCGSV16] in 2016.
- Goal of this talk: Present semantics of CTT in Set $^{\mathbb{C}^{\mathrm{op}}}$ and $\mathcal{E}^{\mathcal{C}^{\mathrm{op}}}$, with \mathcal{E} a model of extensional type theory (ETT) and \mathcal{C} a category internal to \mathcal{E} .

Interval and face lattice I

The interval is a pretype $\mathbb I$ with constants $0,1:\mathbb I$ and operations

$$\begin{split} &\sqcap,\sqcup:\mathbb{I}\to\mathbb{I}\to\mathbb{I},\\ &1-\cdot:\mathbb{I}\to\mathbb{I}, \end{split}$$

endowing \mathbb{I} with the structure of a **de Morgan algebra** where 1 is indecomposable.

We obtain the face lattice \mathbb{F} from \mathbb{I} by factoring modulo $x \sqcap (1-x) = 0$.

Interval and face lattice II

Example:
$$\varphi = (i = 0) \sqcup (i = 1) \sqcup (j = 0) : \mathbb{F}$$

$$(i=0)$$

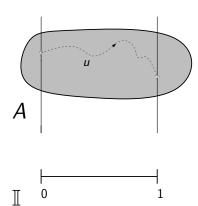
$$(i=1)$$

$$(j=0)$$

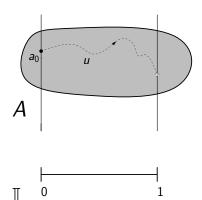
$$(j=0)$$

For $\Gamma \vdash \varphi : \mathbb{F}$, $\Delta \vdash \psi : \mathbb{F}$, ... we may form the **restricted contexts** (Γ, φ) , (Δ, ψ) , ...

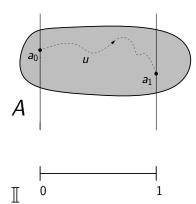
$$\frac{\Gamma \vdash \varphi \quad \Gamma, i : \mathbb{I} \vdash A \quad \Gamma, \varphi, i : \mathbb{I} \vdash u : A \quad \Gamma \vdash a_0 : A(i0)[\varphi \mapsto u(i0)]}{\Gamma \vdash a_1 = \mathsf{comp}^i(A, \varphi, u, a_0) : A(i1)[\varphi \mapsto u(i1)]}$$



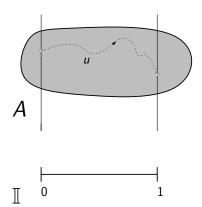
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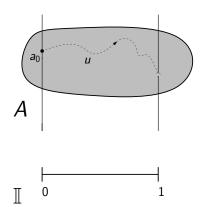
$$\frac{\Gamma \vdash \varphi \qquad \Gamma, i : \mathbb{I} \vdash A \qquad \Gamma, \varphi, i : \mathbb{I} \vdash u : A \qquad \Gamma \vdash a_0 : A(i0)[\varphi \mapsto u(i0)]}{\Gamma \vdash a_1 = \mathsf{comp}^i(A, \varphi, u, a_0) : A(i1)[\varphi \mapsto u(i1)]}$$



Composition is actually equivalent to filling:

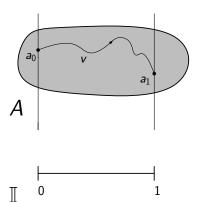


Composition is actually equivalent to filling:



Composition is actually equivalent to **filling**:

$$\Gamma, i : \mathbb{I} \vdash v = \text{fill}^i(A, \varphi, u, a_0)$$



Glueing operation I

$$\frac{\Gamma \vdash A \qquad \Gamma, \varphi \vdash T \qquad \Gamma, \varphi \vdash w : T \to A}{\Gamma \vdash \mathsf{Glue}_{\Gamma}(\varphi, T, A, w)}$$

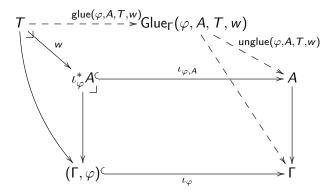
$$\frac{\Gamma \vdash b : \mathsf{Glue}_{\Gamma}(\varphi, T, A, w)}{\Gamma \vdash \mathsf{unglue}(b) : A[\varphi \mapsto w(b)]}$$

$$\frac{\Gamma, \varphi \vdash w : T \to A \qquad \Gamma, \varphi \vdash t : T \qquad \Gamma \vdash a : A[\varphi \mapsto w(t)]}{\Gamma \vdash \langle \mathsf{glue}(\varphi, t), a \rangle : \mathsf{Glue}_{\Gamma}(\varphi, T, A, w)}$$

s.t. judgmentally:

$$\begin{aligned} \mathsf{Glue}_{\Gamma}(1,T,A,w) &= T & \mathsf{unglue}(\mathsf{glue}(\varphi,t),a) &= t \\ \mathsf{glue}(1,T,A,w)(t) &= t & \langle \mathsf{glue}(\varphi,t), \mathsf{unglue}(b) \rangle &= b \end{aligned}$$

Glueing operation II



s.t. judgmentally:

$$\mathsf{Glue}_{\Gamma}(1, T, A, w) = T$$
$$\mathsf{glue}(1, T, A, w)(t) = t$$

$$\mathsf{Glue}_{\Gamma}(1, T, A, w) = T$$
 $\mathsf{unglue}(\mathsf{glue}(\varphi, t), a) = t$ $\mathsf{glue}(1, T, A, w)(t) = t$ $\mathsf{glue}(\varphi, t), \mathsf{unglue}(b) = b$

Cubical Type Theory

Let $\mathbb C$ be a category such that:

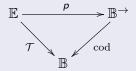
- lacktriangle $\Bbb C$ has finite products
- ② $\widehat{\mathbb{C}}$ has an interval object \mathbb{I} which is representable and connected in the sense that the canonical morphism $\mathbf{2} \to \mathbf{2}^{\mathbb{I}}$ is an isomorphism.
- $\textbf{3} \ \ \, \text{The weakening morphism } \mathbb{F} \to \mathbb{F}^{\mathbb{I}} \text{ has a right adjoint } \\ \forall \colon \mathbb{F}^{\mathbb{I}} \to \mathbb{F}.$

We call $\mathbb C$ "the" category of cubes and $\widehat{\mathbb C}=\mathbf{Set}^{\mathbb C^\mathrm{op}}$ "the" category of cubical sets.

Modelling CTT in presheaves I

Definition (Type category)

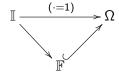
A type category consists of categories $\mathbb B$ and $\mathbb E$ with a discrete Grothendieck fibration $\mathcal T\colon \mathbb E\to \mathbb B$ and a cartesian functor $p\colon \mathbb E\to \mathbb B^{\to}$ from $\mathcal T$ to the codomain fibration:



Set $\mathbb{B}=\widehat{\mathbb{C}}$ and $\mathbb{E}_{\Gamma}=\widehat{\int_{\mathbb{C}}\Gamma}\simeq\widehat{\mathbb{C}}/\Gamma$ for $\Gamma\in\widehat{\mathbb{C}}.$

Modelling CTT in presheaves II

• Consider the classifying map $(\cdot = 1) : \mathbb{I} \to \Omega$ of the global element 1 of \mathbb{I} . Interpret the face lattice as $\mathbb{F} := \operatorname{im}(\cdot = 1)$:



• Consider $\varphi \colon \Gamma \to \mathbb{F}$. Restricted contexts arise as pullbacks

$$\begin{array}{ccc}
\Gamma, \varphi & \longrightarrow & \mathbb{F} \\
\operatorname{Id}_{\mathbb{F}}(\varphi, 1) & & & \operatorname{Id}_{\mathbb{F}} \\
\Gamma & & & & & \operatorname{F} \times \mathbb{F}
\end{array}$$

and we write $[\varphi] := \mathrm{Id}_{\mathbb{F}}(\varphi, 1)$.

Universe of pretypes

Construct generic family $E \to U$ in $\widehat{\mathbb{C}}$ à la [HS97, Str05]:

ullet U Grothendieck universe hosting the category $\mathbb C$

Modelling CTT in Presheaves

• Define $U \in \widehat{\mathbb{C}}$ by

$$U(I) := \mathcal{U}^{(\mathbb{C}/I)^{\mathrm{op}}} \quad \text{for} \quad I \in \mathbb{C},$$
 $U(u: J \to I)(A) := A \circ (\Sigma_u)^{\mathrm{op}} \quad \text{for} \quad u: J \to I.$

• Define E over U as the presheaf

$$E(\langle I, A \rangle) := A(\mathrm{id}_I),$$

 $E(u : \langle J, u^*A \rangle \to \langle I, A \rangle)(a) := A(u : u \to \mathrm{id}_I)(a).$

• N.B.: We get Ω when choosing $\mathcal{U} = \{0, 1\}$.

Cubical Type Theory

Definition (Composition structure (cf. [CCHM16], [OP16]))

Define the family Comp : $\mathcal{T}(U^{\mathbb{I}})$ of composition structures as

$$\mathsf{Comp}(A:U^{\mathbb{I}}) := (\Pi\varphi : \mathbb{F})(\Pip : [\varphi] \to (\Pi i : \mathbb{I})A(i))$$

$$\{a \in A(0) \mid \forall u : [\varphi] . p(u)(0) = a\}$$

$$\to \{a \in A(1) \mid \forall u : [\varphi] . p(u)(1) = a\}.$$

Externally composition structures "put lids on open boxes".

Definition (Fibration structure (cf. [CCHM16], [OP16]))

Define the family Fib : $\mathcal{T}(U^{\Gamma})$ of **fibration structures** as

$$\mathsf{Fib}_{\Gamma}(A) := (\mathsf{\Pi} p : \Gamma^{\mathbb{I}}) \mathsf{Comp}(A \circ p).$$

Universe of types

• Define $U_f \in \widehat{\mathbb{C}}$ by

$$U_f(I) := \{A = \langle |A|, \operatorname{fib}(A) \rangle \mid |A| \colon I \to U, \operatorname{fib}(A) \colon \operatorname{Fib}_I(|A|) \}.$$

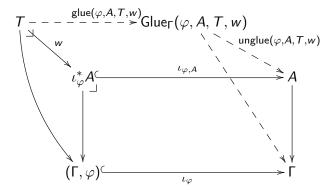
• The universe $E_f \to U_f$ of fibrant types is obtained from $E \to U$ by pulling back along the forgetful map $U_f \to U$:

$$E_f \longrightarrow E$$

$$\downarrow^{J} \qquad \downarrow$$

$$U_f \longrightarrow U$$

Interpreting glueing I



Strictness: Glue_{Γ}(1, A, T, w) = T and glue $(1, A, T, w) = id_T$



Interpreting glueing II

Strictness issue: difficult in general toposes, cf. [OP16]. But can be achieved in $\widehat{\mathbb{C}}$, cf. [CCHM16], as follows. Write $G := \mathsf{Glue}_{\Gamma}(\varphi, A, T, w)$ and for $I \in \mathbb{C}, \gamma \in \Gamma(I)$ let

$$G(I,\gamma) := egin{cases} T(I,\langle\gamma,*
angle) & ext{if } arphi_I(\gamma) = 1 \ ig\{\langle t,a
angle \mid a \in A(I,\gamma), t \colon [arphi]_I(\gamma)
ightarrow T, \ a_{|arphi_I(\gamma)} = \iota_{arphi,A} \circ w \circ tig\} & ext{otherwise} \end{cases}$$

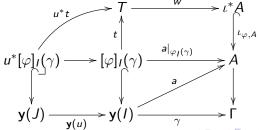
where $a_{|\varphi_I(\gamma)}$ is the restriction of a along the inclusion of $\varphi_I(\gamma)$ into y(I).

Interpreting glueing III

Reindexing: Let $u: J \to I$ in \mathbb{C} . Reindexing along $u: \langle J, u^*\gamma \rangle \to \langle I, \gamma \rangle$ is by case distinction:

$$arphi_I(\gamma) = 1$$
: $T(I, \langle \gamma, * \rangle) \ni t \longmapsto u^*t \in T(J, \langle u^*\gamma, * \rangle)$
 $\varphi_I(\gamma) \neq 1$: $\langle t, a \rangle \longmapsto \begin{cases} u^*a & \text{if } \varphi_J(u^*\gamma) = 1 \\ \langle u^*t, u^*a \rangle & \text{otherwise} \end{cases}$

Here $u^*a = a \circ y(u)$, and u^*t arises as in:



Interpreting glueing IV

• The map glue(φ , A, T, w): $T \to G$ is defined by:

$$\mathsf{glue}(arphi, A, T, w)_{I,\gamma}(t) := egin{cases} t & arphi_I(\gamma) = 1 \ \langle t, w_{I,\gamma}(t)
angle & \mathsf{otherwise} \end{cases}$$

The map

$$unglue(\varphi, A, T, w): G \rightarrow A$$

over Γ is given by:

$$\operatorname{unglue}(\varphi,A,T,w)_{I,\gamma}(b) := \begin{cases} w_{I,\gamma}(b) & \varphi_I(\gamma) = 1\\ \operatorname{pr}_2(b) & \text{otherwise} \end{cases}$$

Composition for glueing I

For a type $p_A: A \to \Gamma$ let P_A be the **type of paths** in A interpreted as follows:

$$P_{A} = A^{\mathbb{I}} \xrightarrow{\langle A^{0}, A^{1} \rangle} A \times_{\Gamma} A$$

$$(p_{A})^{\Gamma^{*}\mathbb{I}} \qquad \Gamma$$

Define weak equivalences á la Voevodsky:

Modelling CTT in Presheaves

- Contractability: isContr(A : U) := $(\Sigma x : A)(\Pi y : A)P_A(x, y)$
- Homotopy fiber: $\mathsf{hfib}(A, B : U)(f : A \to B)(y : B) := (\Sigma x : A)P_B(f(x), y)$
- Weak equivalences: isWeq $(A, B : U)(f : A \rightarrow B) := (\Pi_V : A \rightarrow B) := (\Pi_V$ B)isContr(hfib(A, B, f, γ)), $Weq(A, B : U) := (\Sigma f : A \rightarrow B)isWeq(A, B, f)$

Composition for glueing II

Theorem ([CCHM16], Sec. 6.2)

Let $\varphi \colon \Gamma \to \mathbb{F}$, $A \in \mathcal{T}(\Gamma \times \mathbb{I})$ and $T \in \mathcal{T}((\Gamma, \varphi) \times \mathbb{I})$ have a composition structure, and $w : Weq(T, \iota_{\varphi}^*A)$. Then $G := \mathsf{Glue}_{\Gamma \times \mathbb{I}}(\varphi, A, T, w) \in \mathcal{T}(\Gamma \times \mathbb{I})$ also has a composition structure.

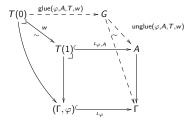
The proof makes crucial use of $\sqcap, \sqcup \colon \mathbb{I}^2 \to \mathbb{I}$ and the map $\forall \colon \mathbb{F}^{\mathbb{I}} \to \mathbb{F}$, which is right adjoint to weakening.

Theorem ([CCHM16], Sec. 7.1)

Modelling CTT in Presheaves

The universe U_f is fibrant.

Proof idea: Paths in the universe can be transformed into weak equivalences between their endpoints. Now given a partial path $T: (\Gamma, \varphi) \to U^{\mathbb{I}}$ and a total extension A of T(1), by glueing we get a total type G which is a total extension of T(0):



Concrete instances

Our (re)construction of the cubical set model has not required too much about the site $\mathbb C$ and the interval object $\mathbb I$. Nevertheless, the only known examples are the algebraic theories of distributive lattices and de Morgan algebras, resp.

Coquand has pointed out that in these particular cases the construction of the cubical set model can be performed in a fairly weak meta-theory, e.g. ETT with a sufficiently well-behaved universe.

This allows one to construct a whole bunch of new models replacing **Set** by models of sufficiently well-behaved models of ETT.

Modelling CTT in internal presheaves

Let \mathcal{E} be a model of ETT with a universe \mathcal{U} containing a natural numbers object (nno) and exact quotients of ¬¬-closed equivalence relations.

Let \mathcal{C} be the category internal to \mathcal{E} which is the opposite of the category of finitely presented free de Morgan algebras and homomorphisms.

Then our interpretation of CTT carries over to the category of internal presheaves $\mathcal{E}^{\mathcal{C}^{\mathrm{op}}}$ since we have never made any substantial use of the subobject classifier in $\widehat{\mathbb{C}}$.

The universe U_f is impredicative in $\mathcal{E}^{\mathcal{C}^{\mathrm{op}}}$ whenever \mathcal{U} is impredicative.



Realizability models for CTT

Cubical Type Theory

Thus, in particular, we can perform the construction of the cubical set model within $\mathcal{E} = \mathbf{Asm}(\mathcal{A})$ for any $\mathbf{pca}\ \mathcal{A}$ instantiating \mathcal{U} with the universe $\mathbf{Mod}(\mathcal{A})$.

Since $\mathbf{Mod}(\mathcal{A})$ is impredicative in \mathcal{E} the ensuing universe \mathcal{U}_f in $\mathcal{E}^{\mathcal{C}^{\mathrm{op}}}$ is impredicative as well (cf. also recent work by Awodey, Frey and Hofstra [Awo17, Fre17]).

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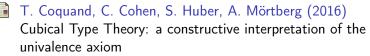
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