

# MIS 102: Computer Programming

## Unit 10 – C Data Structures

Yihuang K. Kang, PhD

*“Simplicity and elegance are unpopular  
because they require hard work and  
discipline to achieve and education to be  
appreciated.”*

– Edsger W. Dijkstra

# Introduction

- We've been discussing that C structure can be used to create complex and aggregate data types. This unit introduces popular uses of the structures—dynamic data structures that can grow or shrink at execution time.
- These techniques will enable us to build these data types in a dramatically different manner designed for producing software that's much easier to maintain and reuse.

# Introduction<sub>(cont.)</sub>

- **Linked lists** are collections of data items “lined up in a row”—insertions and deletions are made anywhere in a linked list.
- **Stacks** are important in compilers and operating systems—insertions and deletions are made only at one end of a stack—its top.
- **Queues** represent waiting lines; insertions are made only at the back/tail of a queue and deletions are made only from the front/head of a queue.
- **Binary trees** facilitate high-speed searching and sorting of data, efficient elimination of duplicate data items, representing file system directories and compiling expressions into machine language.

# Self-Referential Structure Review

- Recall that a **self-referential structure** contains a pointer member that points to a structure of the same structure type. For example, the definition:

```
struct node {  
    int data;  
    struct node *nextPtr;  
}; // end struct node
```

defines a type, *struct node*. This structure of type has two members—integer member *data* and pointer member *nextPtr*.

# Self-Referential Structure Review<sub>(cont.)</sub>

- Member *nextPtr* points to a structure of type struct node—a structure of the same type as the one being declared here, hence the term “self-referential structure.”
- Member *nextPtr* is referred to as a **link**—i.e., it can be used to “tie” a structure of type struct node to another structure of the same type.
- Self-referential structures can be linked together to form useful data structures such as lists, queues, stacks and trees.

# Self-Referential Structure Review<sub>(cont.)</sub>

- Here is an example. Objects linked together to form a list. Note that a slash is placed in this example to represent a NULL. The link of 2<sup>nd</sup> self-referential structure does not point to another structure.



- A NULL pointer normally indicates the end of a data structure just as the null character indicates the end of a string.

# Dynamic Memory Allocation

- Creating and maintaining dynamic data structures requires **dynamic memory allocation**—the ability for a program to obtain more memory space at execution time to hold new nodes, and to release space no longer needed.
- Functions *malloc()* and *free()*, and operator *sizeof*, are essential to dynamic memory allocation.
- One disadvantage of using dynamic memory allocation is that it incurs the overhead of function calls.



# Dynamic Memory Allocation<sub>(cont.)</sub>

- Function *malloc()* takes as an argument the number of bytes to be allocated and returns a pointer of type **void \*** (pointer to void) to the allocated memory.
- As you recall, a void \* pointer may be assigned to a variable of any pointer type.
- Function *malloc()* is normally used with the *sizeof()*.

# Dynamic Memory Allocation<sub>(cont.)</sub>

- For example, the statement

```
newPtr = malloc( sizeof( struct node ) );
```

evaluates *sizeof*(struct node) to determine the size in bytes of a structure of type struct node, allocates a new area in memory of that number of bytes and stores **a pointer to the allocated memory** in variable *newPtr*.

- Note that the allocated memory is not initialized. If no memory is available, *malloc()* returns NULL.

# Dynamic Memory Allocation<sub>(cont.)</sub>

- Function *free()* de-allocates memory—i.e., the memory is returned to the system so that it can be reallocated in the future.
- To free memory dynamically allocated by the preceding *malloc()* call, use the statement

`free( newPtr );`

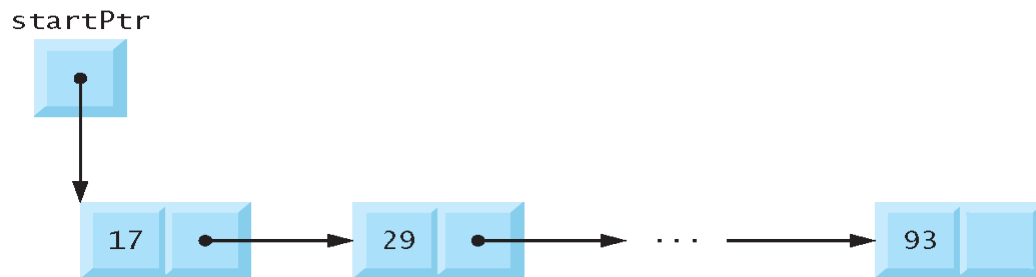
- Remember to free/return allocated memory when it's no longer needed. Otherwise the system may run out of memory prematurely, which is sometimes called “memory leak”.

# Linked Lists

- A **linked list** is a linear collection of self-referential structures, called **nodes**, connected by pointer links—hence, the term “linked” list.
- A linked list is accessed via a pointer to the first node of the list. Subsequent nodes are accessed via the link pointer member stored in each node.
- By convention, the link pointer in the last node of a list is set to NULL to mark the end of the list. Data is stored in a linked list dynamically—each node is created as necessary.

# Linked Lists<sub>(cont.)</sub>

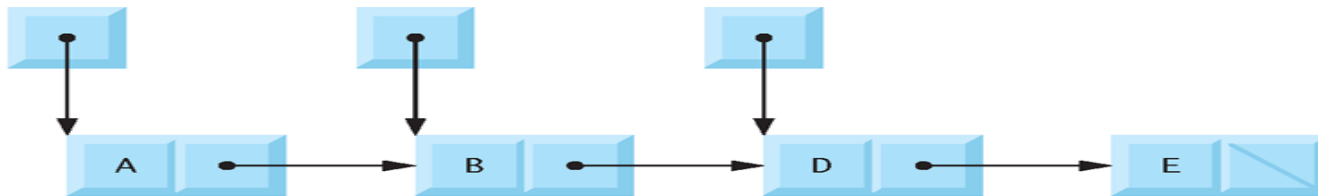
- Lists of data can be stored in arrays, but linked lists provide several advantages. For example, Linked lists are dynamic, so the length of a list can increase or decrease as necessary. A linked list is appropriate when the number of data elements to be represented in the data structure is unpredictable.
- Unlike arrays, linked-list nodes are normally not stored contiguously in memory. However, the nodes appear to be logically contiguous.



# Linked-List Operations

```
1  // Fig. 12.3: fig12_03.c
2  // Inserting and deleting nodes in a list
3  #include <stdio.h>
4  #include <stdlib.h>
5
6  // self-referential structure
7  struct ListNode {
8      char data; // each ListNode contains a character
9      struct ListNode *nextPtr; // pointer to next node
10 }; // end structure ListNode
11
12 typedef struct ListNode ListNode; // synonym for struct ListNode
13 typedef ListNode *ListNodePtr; // synonym for ListNode*
14
15 // prototypes
16 void insert( ListNodePtr *sPtr, char value );
17 char delete( ListNodePtr *sPtr, char value );
18 int isEmpty( ListNodePtr sPtr );
19 void printList( ListNodePtr currentPtr );
20 void instructions( void );
21
```

# Linked-List Operations – Print List



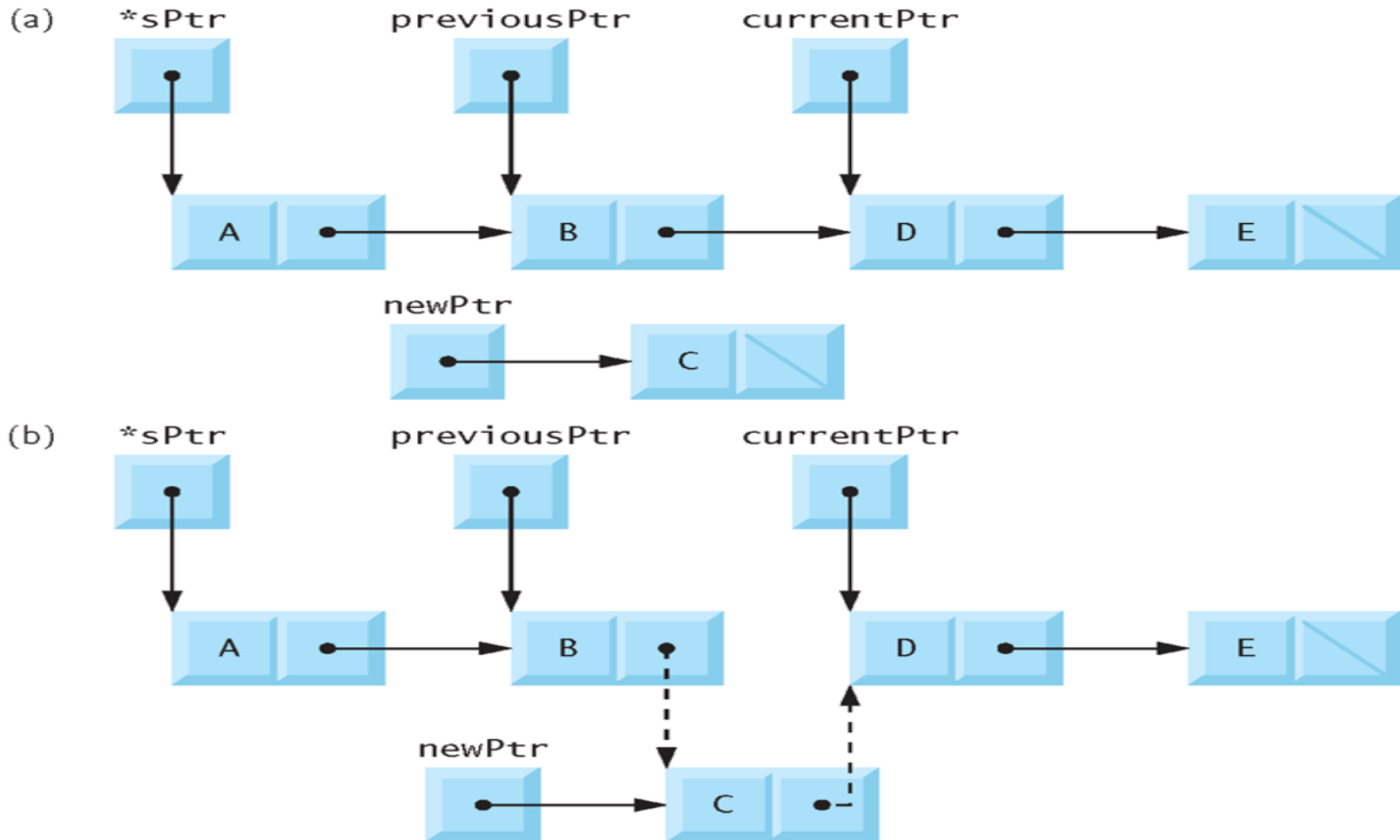
```
163 // print the list
164 void printList( ListNodePtr currentPtr )
165 {
166     // if list is empty
167     if ( isEmpty( currentPtr ) ) {
168         puts( "List is empty.\n" );
169     } // end if
170     else {
171         puts( "The list is:" );
172
173         // while not the end of the list
174         while ( currentPtr != NULL ) {
175             printf( "%c --> ", currentPtr->data );
176             currentPtr = currentPtr->nextPtr;
177         } // end while
178
179         puts( "NULL\n" );
180     } // end else
181 } // end function printList
```

# Linked-List Operations – Insertion

```
84 // insert a new value into the list in sorted order
85 void insert( ListNodePtr *sPtr, char value )
86 {
87     ListNodePtr newPtr; // pointer to new node
88     ListNodePtr previousPtr; // pointer to previous node in list
89     ListNodePtr currentPtr; // pointer to current node in list
90
91     newPtr = malloc( sizeof( ListNode ) ); // create node
92
93     if ( newPtr != NULL ) { // is space available
94         newPtr->data = value; // place value in node
95         newPtr->nextPtr = NULL; // node does not link to another node
96
97         previousPtr = NULL;
98         currentPtr = *sPtr;
99
100        // loop to find the correct location in the list
101        while ( currentPtr != NULL && value > currentPtr->data ) {
102            previousPtr = currentPtr; // walk to ...
103            currentPtr = currentPtr->nextPtr; // ... next node
104        } // end while
105
106        // insert new node at beginning of list
107        if ( previousPtr == NULL ) {
108            newPtr->nextPtr = *sPtr;
109            *sPtr = newPtr;
110        } // end if
111        else { // insert new node between previousPtr and currentPtr
112            previousPtr->nextPtr = newPtr;
113            newPtr->nextPtr = currentPtr;
114        } // end else
115    } // end if
116    else {
117        printf( "%c not inserted. No memory available.\n", value );
118    } // end else
119 } // end function insert
```



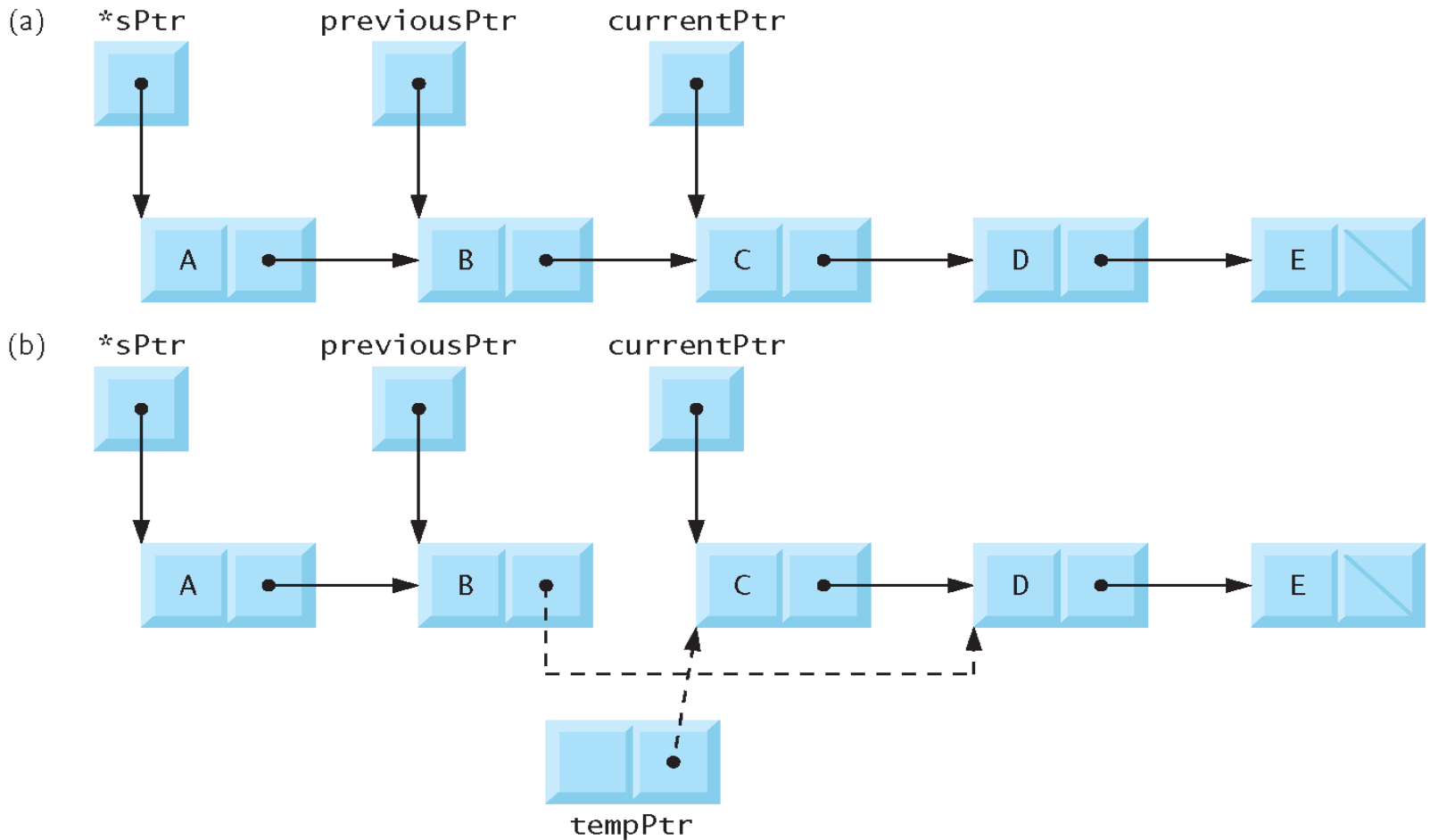
# Linked-List Operations – Insertion



# Linked-List Operations – Deletion

```
121 // delete a list element
122 char delete( ListNodePtr *sPtr, char value )
123 {
124     ListNodePtr previousPtr; // pointer to previous node in list
125     ListNodePtr currentPtr; // pointer to current node in list
126     ListNodePtr tempPtr; // temporary node pointer
127
128     // delete first node
129     if ( value == ( *sPtr )->data ) {
130         tempPtr = *sPtr; // hold onto node being removed
131         *sPtr = ( *sPtr )->nextPtr; // de-thread the node
132         free( tempPtr ); // free the de-threaded node
133         return value;
134     } // end if
135     else {
136         previousPtr = *sPtr;
137         currentPtr = ( *sPtr )->nextPtr;
138
139         // loop to find the correct location in the list
140         while ( currentPtr != NULL && currentPtr->data != value ) {
141             previousPtr = currentPtr; // walk to ...
142             currentPtr = currentPtr->nextPtr; // ... next node
143         } // end while
144
145         // delete node at currentPtr
146         if ( currentPtr != NULL ) {
147             tempPtr = currentPtr;
148             previousPtr->nextPtr = currentPtr->nextPtr;
149             free( tempPtr );
150             return value;
151         } // end if
152     } // end else
153
154     return '\0';
155 } // end function delete
156
157 // return 1 if the list is empty, 0 otherwise
158 int isEmpty( ListNodePtr sPtr )
159 {
160     return sPtr == NULL;
161 } // end function isEmpty
162
```

# Linked-List Operations – Deletion



# Linked-List Operation(cont.)

```
22  int main( void )
23  {
24      ListNodePtr startPtr = NULL; // initially there are no nodes
25      unsigned int choice; // user's choice
26      char item; // char entered by user
27
28      instructions(); // display the menu
29      printf( "%s", "? " );
30      scanf( "%u", &choice );
31
32      // loop while user does not choose 3
33      while ( choice != 3 ) {
34
35          switch ( choice ) {
36              case 1:
37                  printf( "%s", "Enter a character: " );
38                  scanf( "\n%c", &item );
39                  insert( &startPtr, item ); // insert item in list
40                  printList( startPtr );
41                  break;
42              case 2: // delete an element
43                  // if list is not empty
44                  if ( !isEmpty( startPtr ) ) {
45                      printf( "%s", "Enter character to be deleted: " );
46                      scanf( "\n%c", &item );
```

# Linked-List Operation(cont.)

```
47
48         // if character is found, remove it
49         if ( delete( &startPtr, item ) ) { // remove item
50             printf( "%c deleted.\n", item );
51             printList( startPtr );
52         } // end if
53         else {
54             printf( "%c not found.\n\n", item );
55         } // end else
56     } // end if
57     else {
58         puts( "List is empty.\n" );
59     } // end else
60
61     break;
62 default:
63     puts( "Invalid choice.\n" );
64     instructions();
65     break;
66 } // end switch
67
68 printf( "%s", "? " );
69 scanf( "%u", &choice );
70 } // end while
71
```

# Linked-List Operation(cont.)

Enter your choice:

- 1 to insert an element into the list.
- 2 to delete an element from the list.
- 3 to end.

? 1

Enter a character: B

The list is:

B --> NULL

? 1

Enter a character: A

The list is:

A --> B --> NULL

? 1

Enter a character: C

The list is:

A --> B --> C --> NULL

? 2

Enter character to be deleted: D

D not found.

? 2

Enter character to be deleted: B

B deleted.

The list is:

A --> C --> NULL

? 2

Enter character to be deleted: C

C deleted.

The list is:

A --> NULL

? 2

Enter character to be deleted: A

A deleted.

List is empty.

? 4

Invalid choice.

Enter your choice:

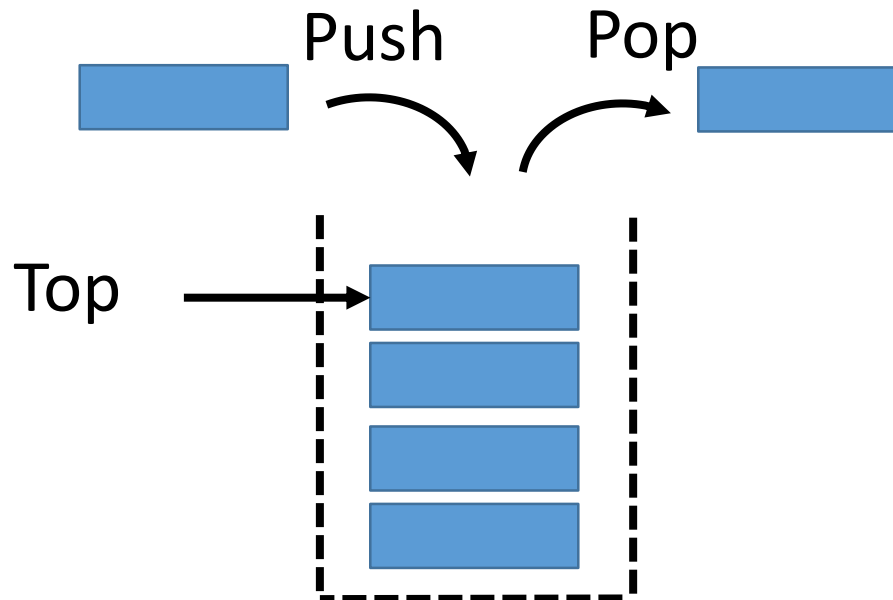
- 1 to insert an element into the list.
- 2 to delete an element from the list.
- 3 to end.

? 3

End of run.

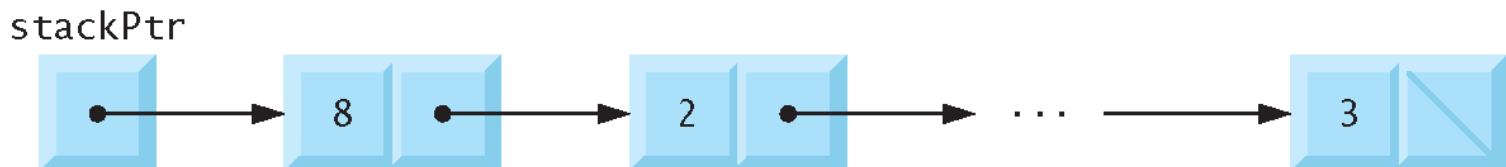
# Stacks

- A **Stack** can be considered a constrained version of the Linked list, as new nodes can only be added/removed from the **top** of a stack. For this reason, a stack is referred to as a **last-in, first-out (LIFO)** data structure.



# Stacks<sub>(cont.)</sub>

- A stack is referenced via a **pointer to the top** element of the stack. The link member in the last node of the stack is set to NULL to indicate the bottom of the stack.
- Stacks and linked lists are represented identically. The difference between stacks and linked lists is that insertions and deletions may occur anywhere in a linked list, but only at the top of a stack.





# Stacks(cont.)

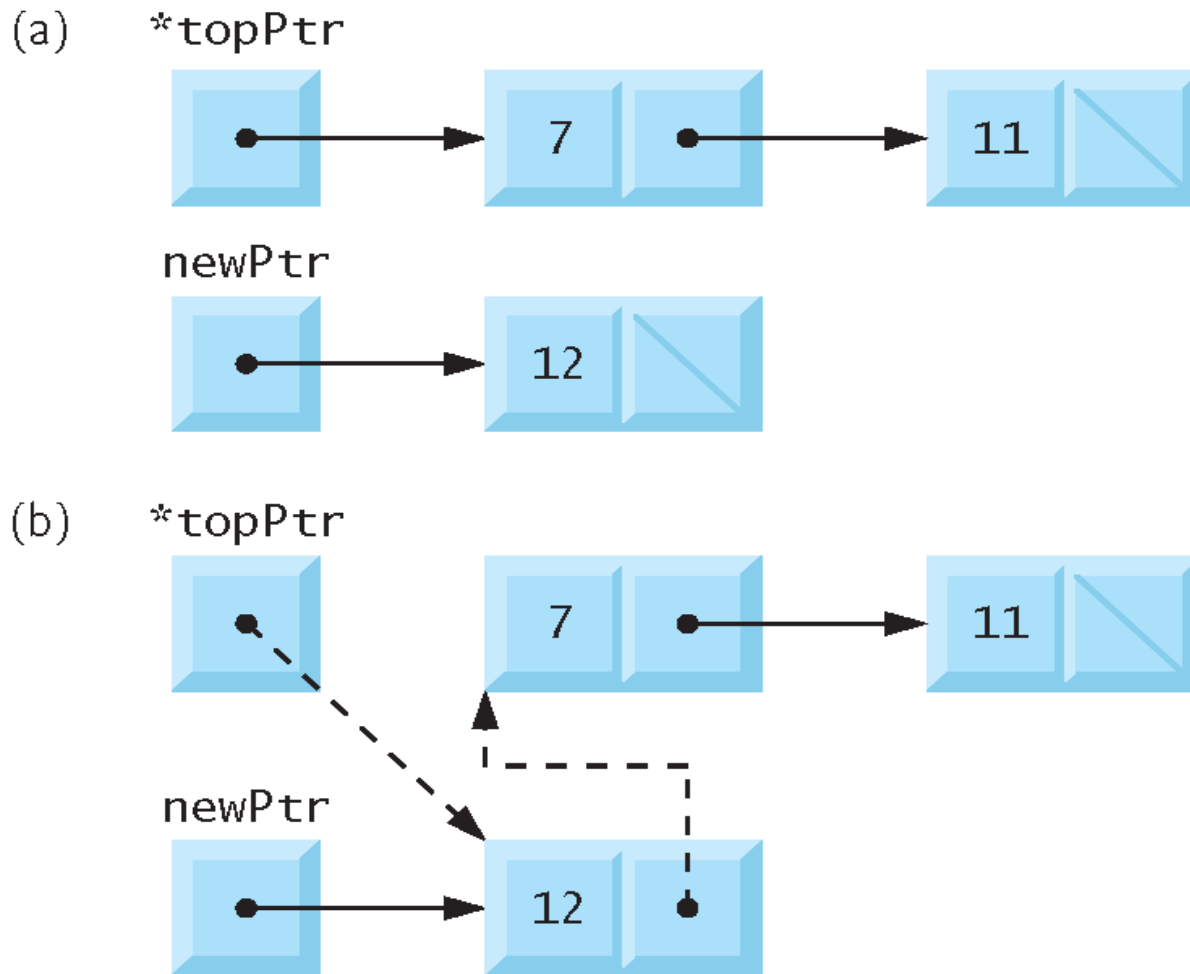
```
1 // Fig. 12.8: fig12_08.c
2 // A simple stack program
3 #include <stdio.h>
4 #include <stdlib.h>
5
6 // self-referential structure
7 struct stackNode {
8     int data; // define data as an int
9     struct stackNode *nextPtr; // stackNode pointer
10 }; // end structure stackNode
11
12 typedef struct stackNode StackNode; // synonym for struct stackNode
13 typedef StackNode *StackNodePtr; // synonym for StackNode*
14
15 // prototypes
16 void push( StackNodePtr *topPtr, int info );
17 int pop( StackNodePtr *topPtr );
18 int isEmpty( StackNodePtr topPtr );
19 void printStack( StackNodePtr currentPtr );
20 void instructions( void );
```

# Stacks(cont.)

- Function ***push()*** creates a new node and places it on top of the stack.

```
75 // insert a node at the stack top
76 void push( StackNodePtr *topPtr, int info )
77 {
78     StackNodePtr newPtr; // pointer to new node
79
80     newPtr = malloc( sizeof( StackNode ) );
81
82     // insert the node at stack top
83     if ( newPtr != NULL ) {
84         newPtr->data = info;
85         newPtr->nextPtr = *topPtr;
86         *topPtr = newPtr;
87     } // end if
88     else { // no space available
89         printf( "%d not inserted. No memory available.\n", info );
90     } // end else
91 } // end function push
```

# Stacks(cont.)



# Stacks(cont.)

- Function ***pop()*** removes a node from the top of the stack, frees the memory that was allocated to the popped node and returns the popped value.

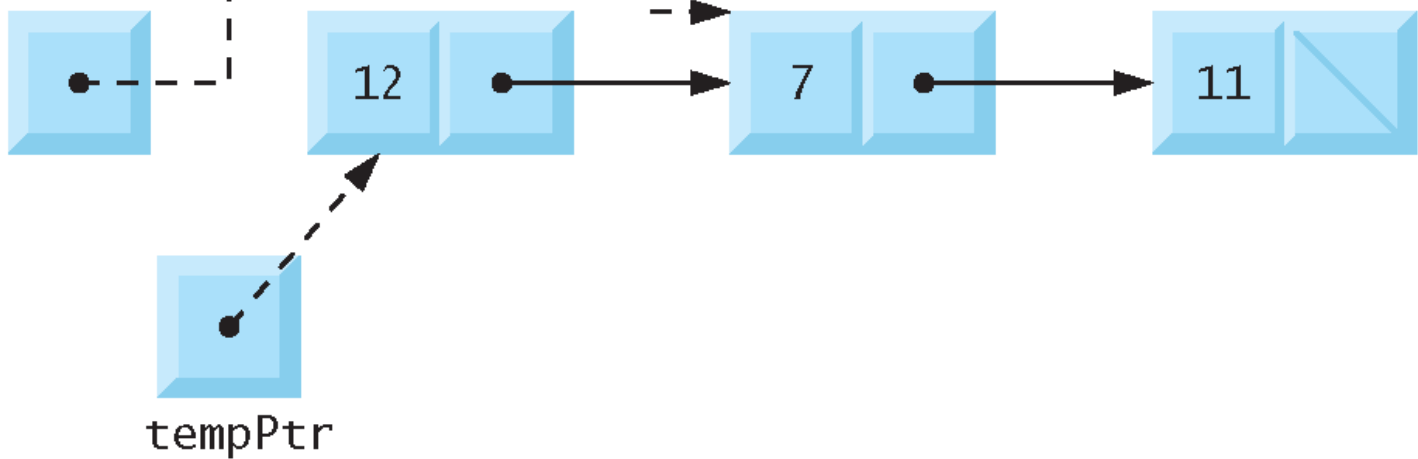
```
93 // remove a node from the stack top
94 int pop( StackNodePtr *topPtr )
95 {
96     StackNodePtr tempPtr; // temporary node pointer
97     int popValue; // node value
98
99     tempPtr = *topPtr;
100    popValue = ( *topPtr )->data;
101    *topPtr = ( *topPtr )->nextPtr;
102    free( tempPtr );
103    return popValue;
104 }
```

# Stacks(cont.)

(a) \*topPtr



(b) \*topPtr



# Stacks(cont.)

```
Enter choice:
1 to push a value on the stack
2 to pop a value off the stack
3 to end program
? 1
Enter an integer: 5
The stack is:
5 --> NULL

? 1
Enter an integer: 6
The stack is:
6 --> 5 --> NULL

? 1
Enter an integer: 4
The stack is:
4 --> 6 --> 5 --> NULL
```

```
? 2
The popped value is 4.
The stack is:
6 --> 5 --> NULL
```

```
? 2
The popped value is 6.
The stack is:
5 --> NULL
```

```
? 2
The popped value is 5.
The stack is empty.
```

```
? 2
The stack is empty.
```

```
? 4
Invalid choice.
```

```
Enter choice:
1 to push a value on the stack
2 to pop a value off the stack
3 to end program
? 3
End of run.
```

# Application of Stacks

- Stacks have many interesting applications. For example, whenever a function call is made, the called function must know how to return to its caller, so the return address is pushed onto a stack.
- If a series of function calls occurs, the successive return values are pushed onto the stack in last-in, first-out order so that each function can return to its caller.
- When the function returns to its caller, the space for that function's automatic variables is popped off the stack, and these variables no longer are known to the program.

# Queues

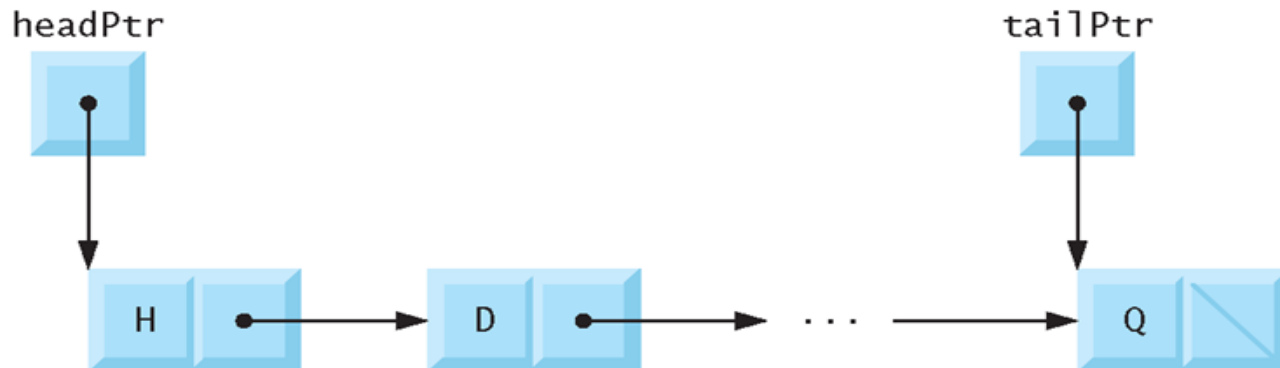
- Another common data structure is the **queue**. A queue is similar to a checkout line in a store—the first person in line is serviced first, and other customers enter the line only at the end and wait to be serviced.
- Queue nodes are removed only from the **head** of the queue and are inserted only at the **tail** of the queue. For this reason, a queue is referred to as a **first-in, first-out (FIFO)** data structure.

The insert and remove operations are known as **enqueue** and **dequeue**, respectively.



# Queues(cont.)

- Queues have many applications in computer systems. For example, queues are used to support print spooling. If a printer is busy, other outputs may still be generated. These are spooled to disk where they wait in a queue until the printer becomes available.
- Here is a graphical representation of Queues.



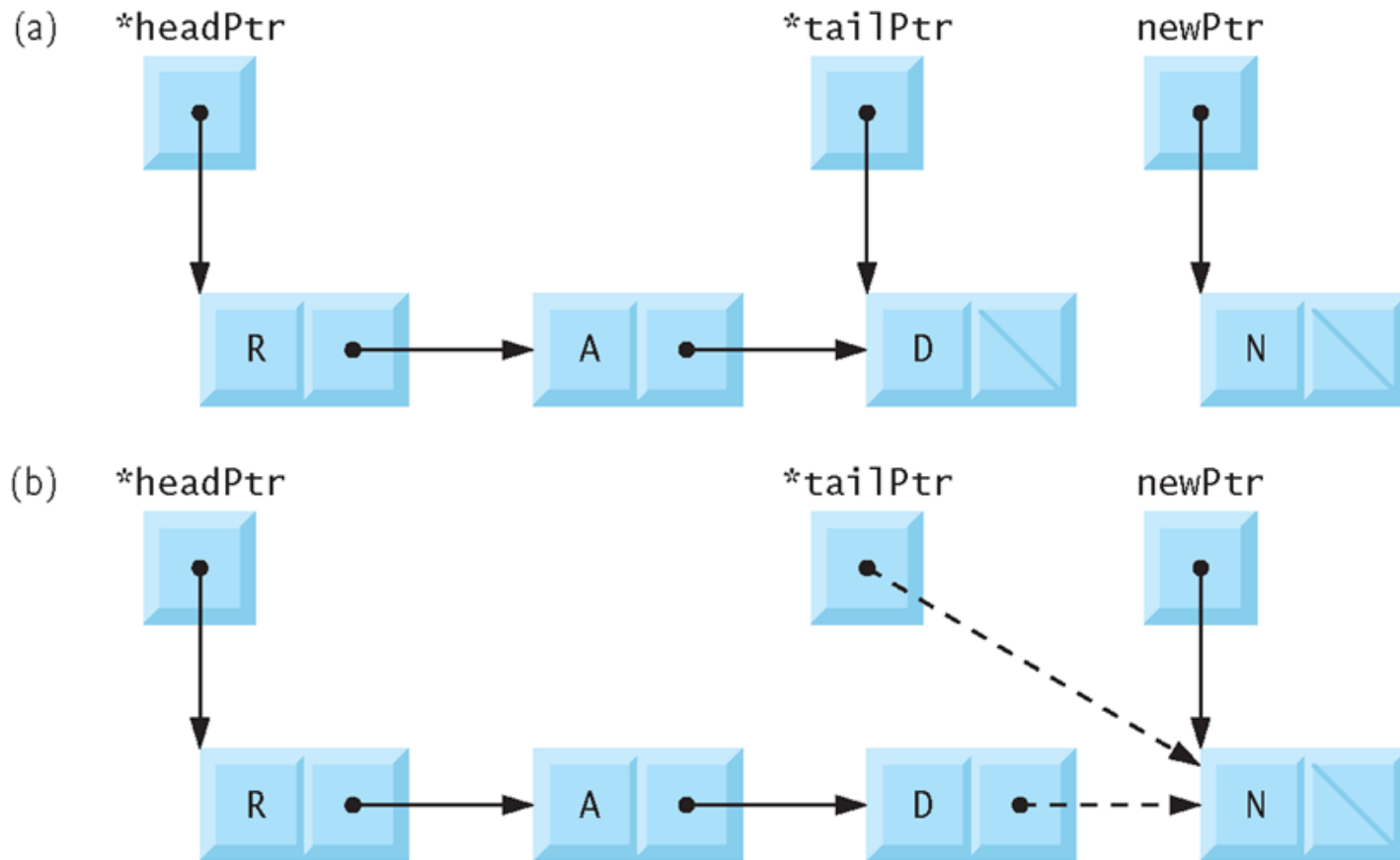
# Queues(cont.)

```
1 // Fig. 12.13: fig12_13.c
2 // Operating and maintaining a queue
3 #include <stdio.h>
4 #include <stdlib.h>
5
6 // self-referential structure
7 struct queueNode {
8     char data; // define data as a char
9     struct queueNode *nextPtr; // queueNode pointer
10 }; // end structure queueNode
11
12 typedef struct queueNode QueueNode;
13 typedef QueueNode *QueueNodePtr;
14
15 // function prototypes
16 void printQueue( QueueNodePtr currentPtr );
17 int isEmpty( QueueNodePtr headPtr );
18 char dequeue( QueueNodePtr *headPtr, QueueNodePtr *tailPtr );
19 void enqueue( QueueNodePtr *headPtr, QueueNodePtr *tailPtr,
20     char value );
21 void instructions( void );
22
```

# Queues(cont.)

```
78 // insert a node in at queue tail
79 void enqueue( QueueNodePtr *headPtr, QueueNodePtr *tailPtr,
80     char value )
81 {
82     QueueNodePtr newPtr; // pointer to new node
83
84     newPtr = malloc( sizeof( QueueNode ) );
85
86     if ( newPtr != NULL ) { // is space available
87         newPtr->data = value;
88         newPtr->nextPtr = NULL;
89
90         // if empty, insert node at head
91         if ( isEmpty( *headPtr ) ) {
92             *headPtr = newPtr;
93         } // end if
94         else {
95             ( *tailPtr )->nextPtr = newPtr;
96         } // end else
97
98         *tailPtr = newPtr;
99     } // end if
100     else {
101         printf( "%c not inserted. No memory available.\n", value );
102     } // end else
103 } // end function enqueue
```

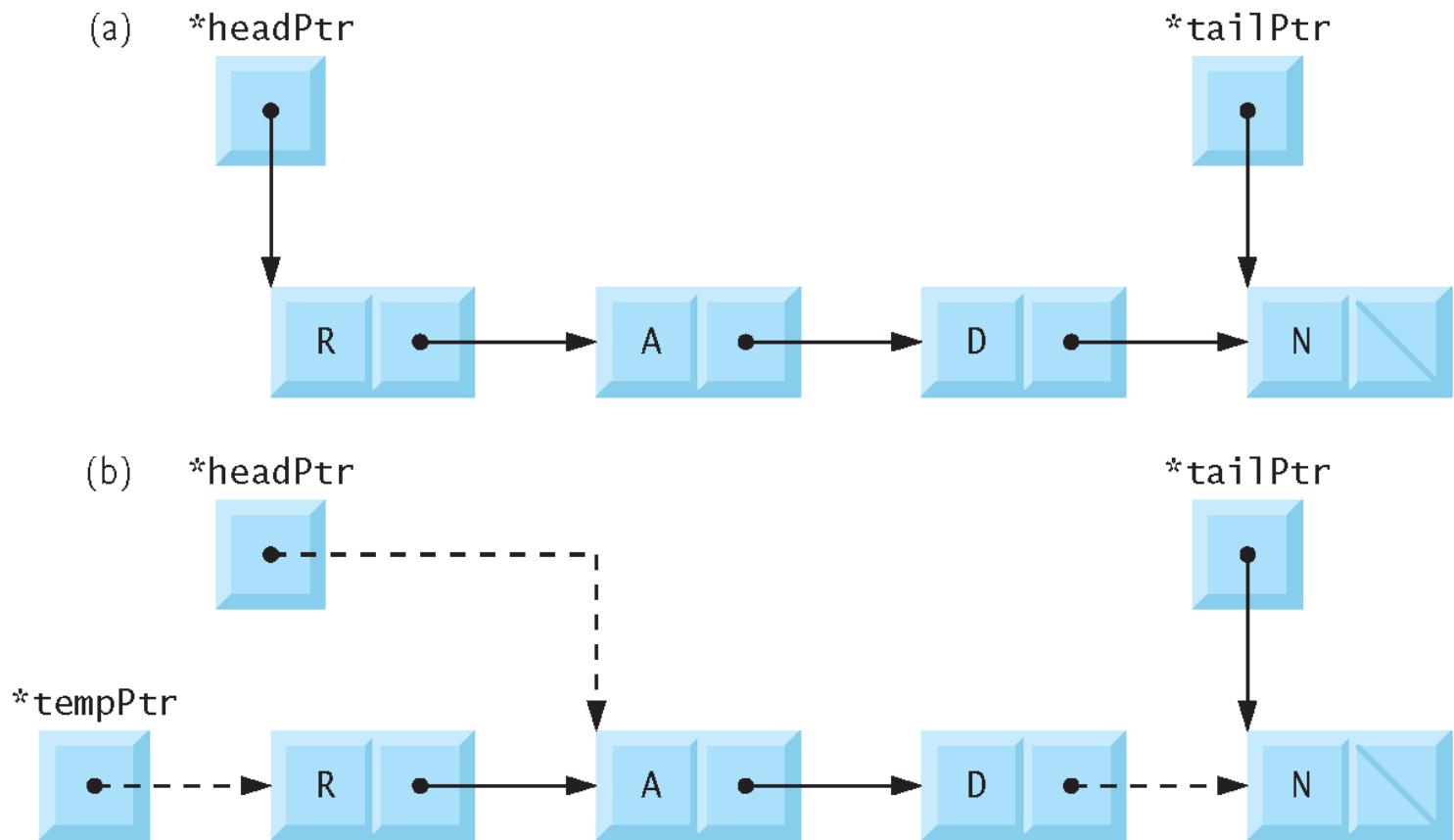
# Queues(cont.)



# Queues(cont.)

```
104
105 // remove node from queue head
106 char dequeue( QueueNodePtr *headPtr, QueueNodePtr *tailPtr )
107 {
108     char value; // node value
109     QueueNodePtr tempPtr; // temporary node pointer
110
111     value = ( *headPtr )->data;
112     tempPtr = *headPtr;
113     *headPtr = ( *headPtr )->nextPtr;
114
115     // if queue is empty
116     if ( *headPtr == NULL ) {
117         *tailPtr = NULL;
118     } // end if
119
120     free( tempPtr );
121     return value;
122 } // end function dequeue
123
124 // return 1 if the queue is empty, 0 otherwise
125 int isEmpty( QueueNodePtr headPtr )
126 {
127     return headPtr == NULL;
128 } // end function isEmpty
```

# Queues(cont.)



# Queues(cont.)

```
Enter your choice:
  1 to add an item to the queue
  2 to remove an item from the queue
  3 to end
? 1
Enter a character: A
The queue is:
A --> NULL

? 1
Enter a character: B
The queue is:
A --> B --> NULL

? 1
Enter a character: C
The queue is:
A --> B --> C --> NULL

? 2
A has been dequeued.
The queue is:
B --> C --> NULL
```

```
? 2
B has been dequeued.
The queue is:
C --> NULL
```

```
? 2
C has been dequeued.
Queue is empty.
```

```
? 2
Queue is empty.
```

```
? 4
Invalid choice.
```

```
Enter your choice:
  1 to add an item to the queue
  2 to remove an item from the queue
  3 to end
? 3
End of run.
```

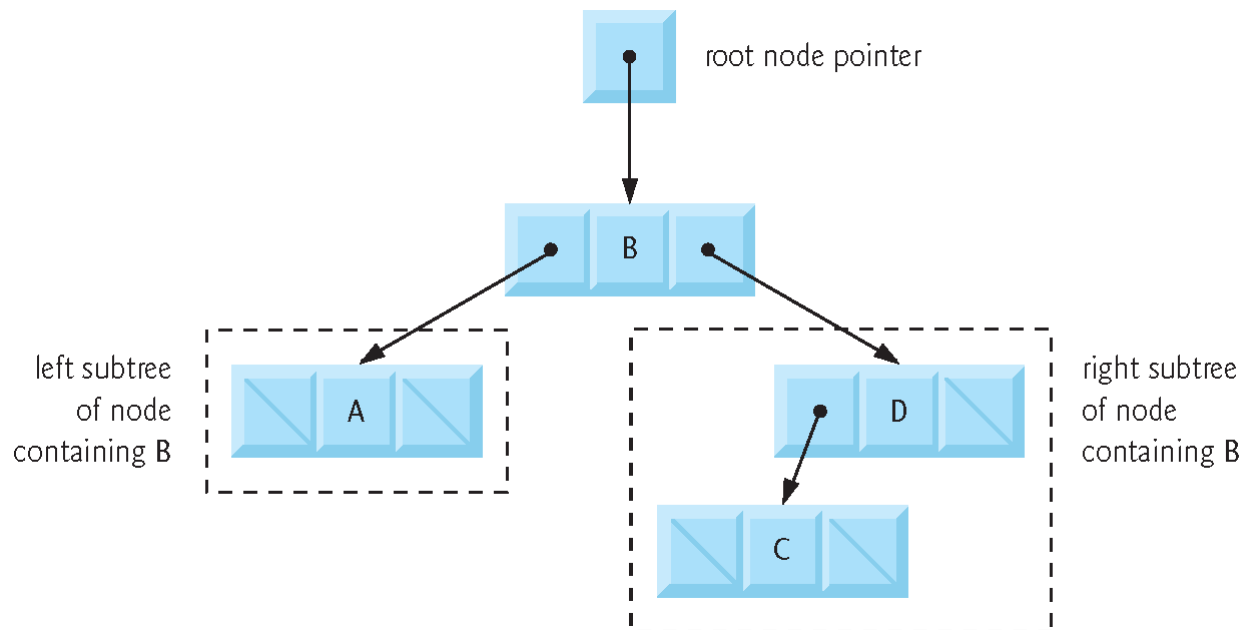
# Trees

- Linked lists, stacks and queues are **linear data structures**. A **tree** is a nonlinear, two-dimensional data structure with special properties.
- A Tree node contains two or more links. Here, we only discuss **binary search trees**—trees whose nodes all contain two links (none, one, or both of which may be NULL).
- Computer scientists normally draw trees from the root node down—exactly the opposite of trees in nature.



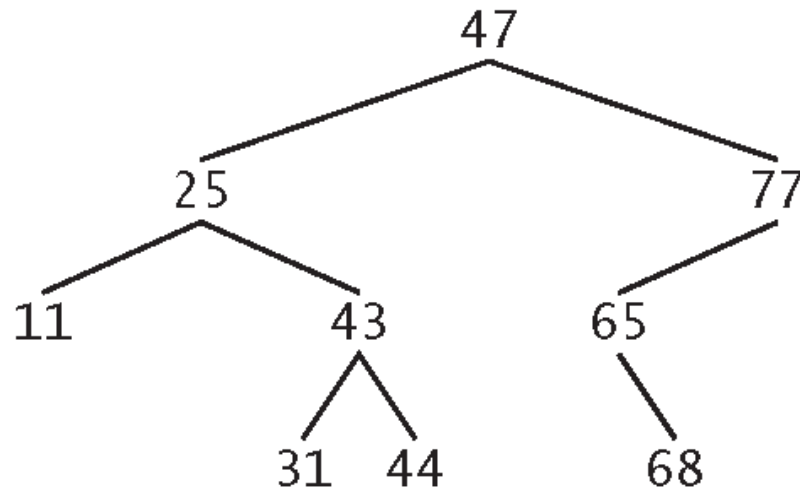
# Trees<sub>(cont.)</sub>

- The **root** node is the first node in a tree. Each link in the root node refers to a child. The left child is the first node in the left subtree, and the right child is the first node in the right subtree. The children of a node are called **siblings**. A node with no children is called a **leaf** node.



# Trees<sub>(cont.)</sub>

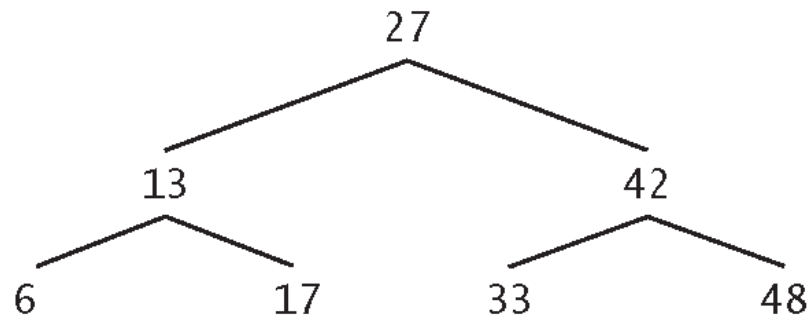
- A **binary search tree (BST)** has the characteristic that the values in *any left subtree are less than the value in its parent node*, and the values in *any right subtree are greater than the value in its parent node*. Below is an example of the tree.



# Trees<sub>(cont.)</sub>

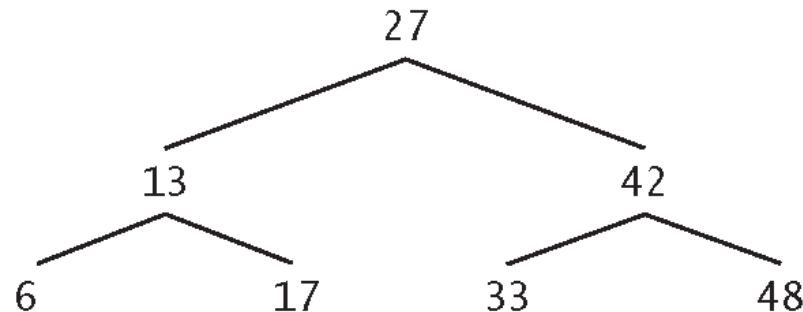
- The following example shows how to create BST and traverse it in three ways—**inOrder**, **preOrder** and **postOrder**.
- The below BST is built from the following sequence of numbers.

**27, 13, 42, 6, 17, 33, 48**



# Trees<sub>(cont.)</sub>

- The steps for an **inOrder** traversal of below BST are: traverse the left subtree inOrder (L), process the data value in the node (D), and then traverse the right subtree inOrder (R).

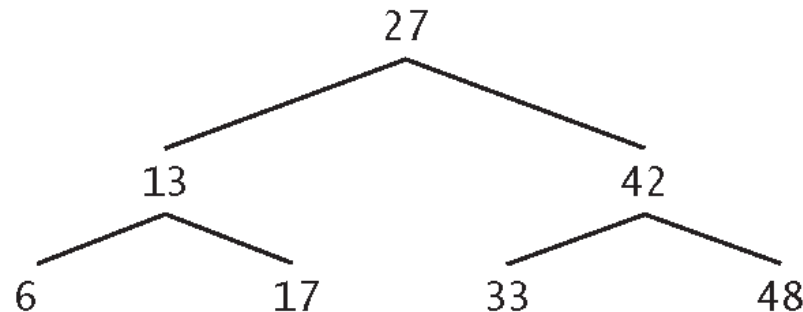


- So the inOrder (LDR) traversal of the tree is:

**6 13 17 27 33 42 48**

# Trees<sub>(cont.)</sub>

- The steps for a **preOrder** traversal of below BST are: process the value in the node (D), traverse the left subtree preorder (L), and then traverse the right subtree preorder (R).

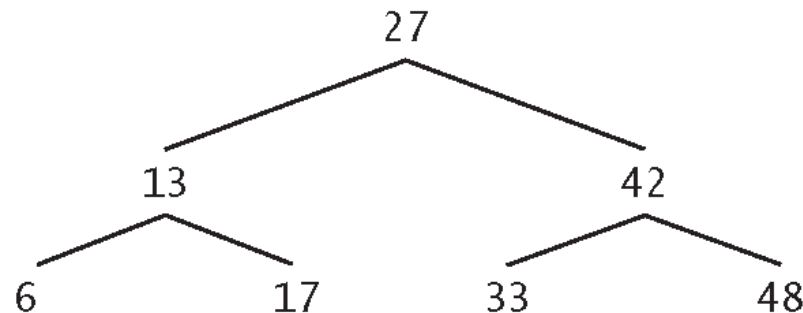


- So the preOrder (DLR) traversal of the tree is:

**27 13 6 17 42 33 48**

# Trees<sub>(cont.)</sub>

- The steps for a **postOrder** traversal of below BST are: traverse the left subtree postOrder (L), traverse the right subtree postOrder (R), and then process the value in the node (D).



- So the postOrder (LRD) traversal of the tree is:

**6 17 13 33 48 42 27**

# Binary Tree Search

- You may notice an advantage of using BST—once the binary search tree has been created, its elements can be retrieved by recursively in-order traversing the tree. It allows fast lookup as it keeps the values in sorted order, so that lookup and other operations can use the principle of **binary search**.
- An BST with  $n$  elements would have a maximum of  $\log_2(n)$  levels, and thus a maximum of  $\log_2(n)$  comparisons would have to be made either to find a match or to determine that no match exists. It means, for example, that when searching a 1,000,000-element binary search tree, no more than 20 comparisons need to be made because  $2^{20} > 1,000,000$ .

# Tree Implementation

```
1 // Fig. 12.19: fig12_19.c
2 // Creating and traversing a binary tree
3 // preorder, inorder, and postorder
4 #include <stdio.h>
5 #include <stdlib.h>
6 #include <time.h>
7
8 // self-referential structure
9 struct treeNode {
10     struct treeNode *leftPtr; // pointer to left subtree
11     int data; // node value
12     struct treeNode *rightPtr; // pointer to right subtree
13 }; // end structure treeNode
14
15 typedef struct treeNode TreeNode; // synonym for struct treeNode
16 typedef TreeNode *TreeNodePtr; // synonym for TreeNode*
17
18 // prototypes
19 void insertNode( TreeNodePtr *treePtr, int value );
20 void inOrder( TreeNodePtr treePtr );
21 void preOrder( TreeNodePtr treePtr );
22 void postOrder( TreeNodePtr treePtr );
```



# Tree Implementation-Insert Node

```
54 // insert node into tree
55 void insertNode( TreeNodePtr *treePtr, int value )
56 {
57     // if tree is empty
58     if ( *treePtr == NULL ) {
59         *treePtr = malloc( sizeof( TreeNode ) );
60
61         // if memory was allocated, then assign data
62         if ( *treePtr != NULL ) {
63             ( *treePtr )->data = value;
64             ( *treePtr )->leftPtr = NULL;
65             ( *treePtr )->rightPtr = NULL;
66         } // end if
67         else {
68             printf( "%d not inserted. No memory available.\n", value );
69         } // end else
70     } // end if
71     else { // tree is not empty
72         // data to insert is less than data in current node
73         if ( value < ( *treePtr )->data ) {
74             insertNode( &( ( *treePtr )->leftPtr ), value );
75         } // end if
76
77         // data to insert is greater than data in current node
78         else if ( value > ( *treePtr )->data ) {
79             insertNode( &( ( *treePtr )->rightPtr ), value );
80         } // end else if
81         else { // duplicate data value ignored
82             printf( "%s", "dup" );
83         } // end else
84     } // end else
85 } // end function insertNode
```

# Tree Implementation-Traversal

```

87 // begin inorder traversal of tree
88 void inOrder( TreeNodePtr treePtr )
89 {
90     // if tree is not empty, then traverse
91     if ( treePtr != NULL ) {
92         inOrder( treePtr->leftPtr );
93         printf( "%3d", treePtr->data );
94         inOrder( treePtr->rightPtr );
95     } // end if
96 } // end function inOrder

98 // begin preorder traversal of tree
99 void preOrder( TreeNodePtr treePtr )
100 {
101     // if tree is not empty, then traverse
102     if ( treePtr != NULL ) {
103         printf( "%3d", treePtr->data );
104         preOrder( treePtr->leftPtr );
105         preOrder( treePtr->rightPtr );
106     } // end if
107 } // end function preOrder

108 // begin postorder traversal of tree
109 void postOrder( TreeNodePtr treePtr )
110 {
111     // if tree is not empty, then traverse
112     if ( treePtr != NULL ) {
113         postOrder( treePtr->leftPtr );
114         postOrder( treePtr->rightPtr );
115         printf( "%3d", treePtr->data );
116     } // end if
117 } // end function postOrder
118
```

# Tree Implementation(cont.)

```
25  int main( void )
26  {
27      unsigned int i; // counter to loop from 1-10
28      int item; // variable to hold random values
29      TreePtr rootPtr = NULL; // tree initially empty
30
31      srand( time( NULL ) );
32      puts( "The numbers being placed in the tree are:" );
33
34      // insert random values between 0 and 14 in the tree
35      for ( i = 1; i <= 10; ++i ) {
36          item = rand() % 15;
37          printf( "%3d", item );
38          insertNode( &rootPtr, item );
39      } // end for
40
41      // traverse the tree preOrder
42      puts( "\n\nThe preOrder traversal is:" );
43      preOrder( rootPtr );
44
45      // traverse the tree inOrder
46      puts( "\n\nThe inOrder traversal is:" );
47      inOrder( rootPtr );
48
49      // traverse the tree postOrder
50      puts( "\n\nThe postOrder traversal" );
51      postOrder( rootPtr );
52  } // end main
```

# Tree Implementation<sub>(cont.)</sub>

---

The numbers being placed in the tree are:

6 7 4 12 7dup 2 2dup 5 7dup 11

The preOrder traversal is:

6 4 2 5 7 ~~12 11~~  
11 12

---

The inOrder traversal is:

2 4 5 6 7 11 12

The postOrder traversal is:

2 5 4 11 12 7 6