

# Cryptographic hash functions

## Lecture 5

# Hash functions

- A hash function is a function that is easy to compute, but hard to invert.
- Hash: {arbitrary size documents}  $\rightarrow \{0, 1\}^k$  that satisfies:
  - One-way function: it is easy to compute  $h = \text{Hash}(D)$ , but hard to invert  $h$  from  $H(D)$ .
  - Collision resistance: it is very difficult to find to distinct input  $D$  and  $D'$  whose output  $\text{Hash}(D)$  and  $\text{Hash}(D')$  are the same.

# Hash function implementation

- Common hash: MD5, SHA
- Using an encryption function:  
symmetric/asymmetric encryptions.
- If  $H_1$  and  $H_2$  are two different hashes, then  $H_1 \circ H_2$  and  $H_2 \circ H_1$  are hashes.

# Modular Arithmetic Secure Hash

- MASH-2: generate a  $n$ -bit hash value from a document of size  $b$  bits,  $1 \leq b \leq 2^{n/2}$ .

(1) Generate two  $m$ -bit primes  $p$  and  $q$ , and set  $M = pq$

(2) Get  $n = 16k$ :  $16k \leq m \leq 16(k+1)$

(3)  $H = 0$ ,  $A = 11110000 \dots 0000$

(4) Split document  $D$  to  $t$   $(n/2)$ -bit blocks  $x_1, \dots, x_t$ . Let  $x_{t+1} = \text{binary}(b)$

(5) For  $i=1$  to  $t$ : split  $x_i$  to 4-bit blocks  $x_{i1}, \dots, x_{ij}$ ,

and let  $y_i = 1111x_{i1} \dots 1111x_{ij}$ .

With  $(t+1)$ th block  $x_{t+1}$  to  $y_{t+1} = 1010x_{(t+1)1} \dots 1010x_{(t+1)j}$ .

Let  $y = y_1 \dots y_{t+1}$ .

(6) For  $i=1$  to  $t+1$ :  $F = (H \text{ XOR } y_i) \text{ OR } A)^{2^{57}} \pmod{M}$ .

Let  $G$  be the lowest  $n$  bits of  $F$ , and  $H = G \text{ XOR } H$ ;

(7) Return  $H$

# Application of hash

1. Hash is used to authenticate messages:  
 $MAC = \text{Hashed Code}$ .
2. Hash is used in digital signature schemes:  
 $\text{Signature}(d, D) = E(d, \text{Hash}(D))$  and verification is whether  $\text{Hash}(D) == E(e, \text{Signature})$ , where  $e$  is the public key and  $d$ , the private key of signer.
3. Hash is used to protect password which stored on servers.