

CHAPTER

9

# I/O SYSTEM



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# What will you learn?

- ☐ I/O devices
  - ☐ I/O System Characteristics
  - ☐ I/O Modules
  - ☐ I/O Register Mapping
  - ☐ I/O Data transfer
- ☐ I/O Command
  - ☐ Life cycle of an I/O request
  - ☐ I/O Bus
  - ☐ Typical x86 PC I/O System

# I/O devices

- Can be typify by:
  - ▣ Behavior: input, output, storage
  - ▣ Partner: human / machine
  - ▣ Data rate: bytes/sec, transfer/sec

- Character & Block devices

The device interface gives the illusion that devices support the same API – character stream and block access

|                   |                                   |   |
|-------------------|-----------------------------------|---|
| application/user: | <i>read character from device</i> | naming, protection, read,write                    |
| operating system: | <i>character &amp; block API</i>  | hardware specific PIO, interrupt handling, or DMA |
| hardware:         | <i>keyboard, mouse, etc.</i>      |   |

# I/O Modules

- ☐ Interface to the processor and memory via the system bus or control switch
- ☐ Interface to one or more peripheral devices

# I/O Register Mapping

## □ Memory-mapped I/O

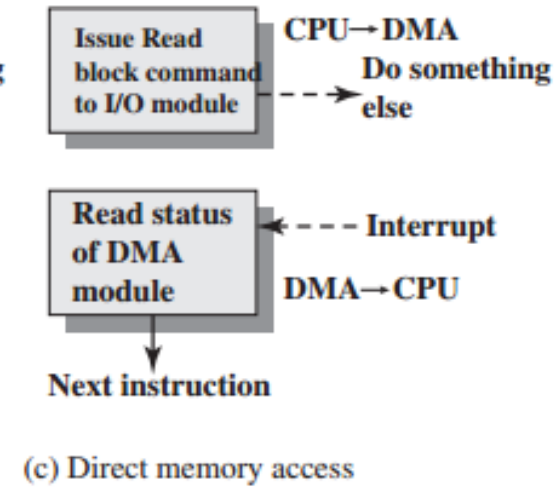
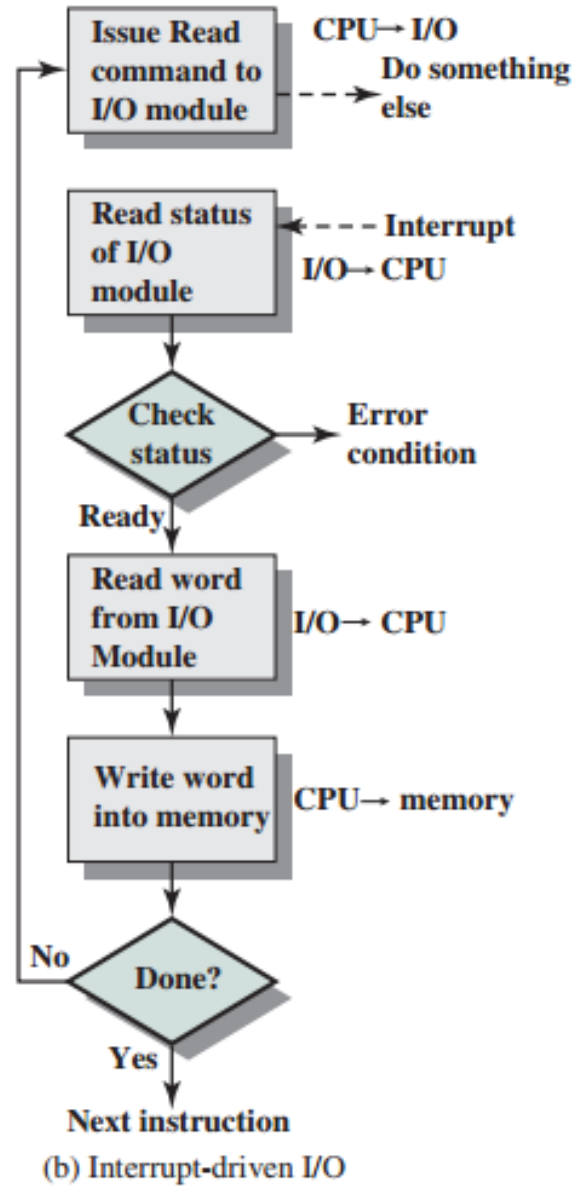
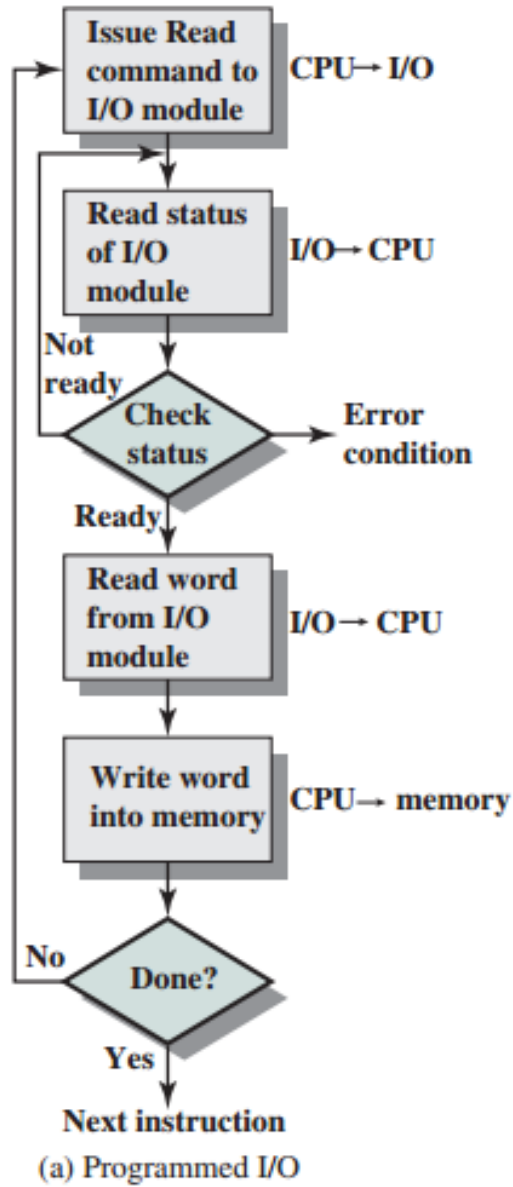
- Registers are addressed in same space as memory
- Address decoder distinguishes between them
- OS uses address translation mechanism to make them only accessible to kernel

## □ Isolated I/O

- Separate instructions to access I/O registers
- Can only be executed in kernel mode

# I/O Data Transfer

- Data transfer between CPU and I/O devices can be handled in generally three types of modes:
  - Programed I/O
  - Interrupt Driven I/O
  - Direct Memory Access



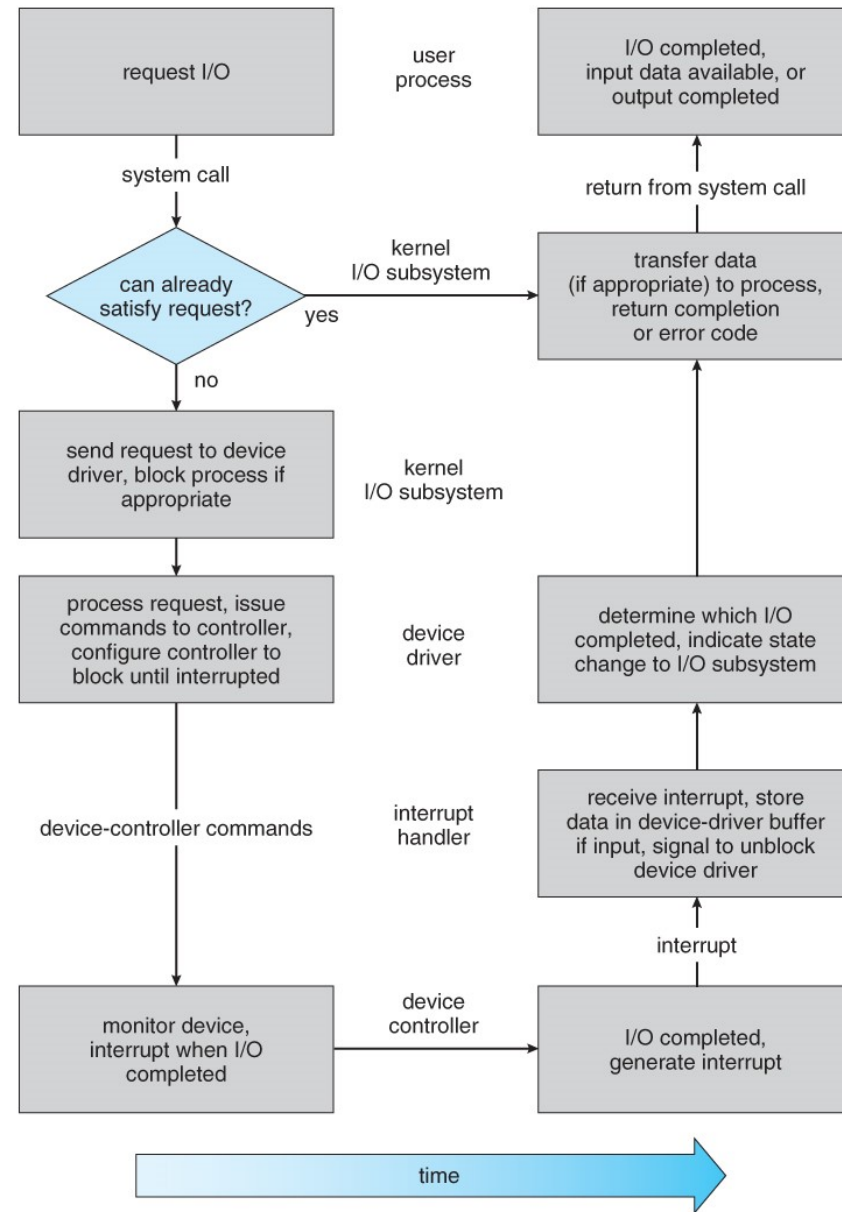
# I/O Commands

- I/O devices are managed by I/O Controller hardware
- The processor issues an address, specifying I/O module and device, and an I/O command. The commands are:
  - Control
  - Test
  - Read
  - Write



## Life cycle of an I/O request

- ❑ Users request data using file names, which must ultimately be mapped to specific blocks of data from a specific device managed by a specific device driver.
- ❑ UNIX uses a *mount table* to map filename prefixes ( e.g. /usr ) to specific mounted devices
- ❑ UNIX uses special *device files*, usually located in /dev, to represent and access physical devices directly



# I/O Bus Types

- ☐ Processor-Memory buses
  - ☐ Short, high speed
  - ☐ Design is matched to memory organization
- ☐ I/O buses
  - ☐ Longer, allowing multiple connections
  - ☐ Connect to processor-memory bus through a bridge

# I/O Bus Example

|                            | Firewire          | USB 2.0                     | PCI Express                              | Serial ATA | Serial Attached SCSI |
|----------------------------|-------------------|-----------------------------|--|------------|----------------------|
| <b>Intended use</b>        | External          | External                    | Internal                                 | Internal   | External             |
| <b>Devices per channel</b> | 63                | 127                         | 1  | 1          | 4                    |
| <b>Data width</b>          | 4                 | 2                           | 2/lane                                   | 4          | 4                    |
| <b>Peak bandwidth</b>      | 50MB/s or 100MB/s | 0.2MB/s, 1.5MB/s, or 60MB/s | 250MB/s/lane<br>1×, 2×, 4×, 8×, 16×, 32× | 300MB/s    | 300MB/s              |
| <b>Hot pluggable</b>       | Yes               | Yes                         | Depends                                  | Yes        | Yes                  |
| <b>Max length</b>          | 4.5m              | 5m                          | 0.5m                                     | 1m         | 8m                   |
| <b>Standard</b>            | IEEE 1394         | USB Implementers Forum      | PCI-SIG                                  | SATA-IO    | INCITS TC T10        |

# Typical X86 PC I/O System

