

# Aufgabe 1

## Aquivalenzumformungen

$\Sigma = 15/16$

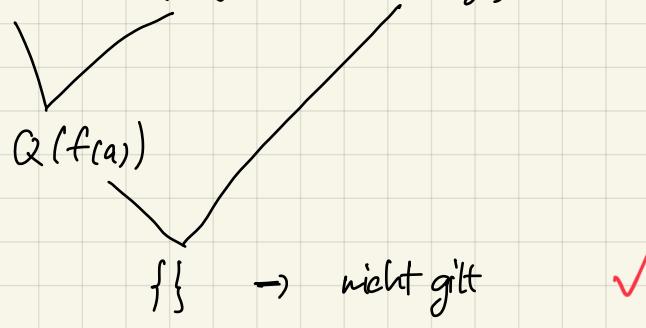
$$\begin{aligned}
 \neg F &= \neg (\forall x. (P \rightarrow Q(x)) \rightarrow P \rightarrow \forall y. Q(y)) \\
 &= \neg (\neg (\exists x. \neg (P \rightarrow Q(x)) \vee (\neg P \vee \forall y. Q(y))) \\
 &= \forall x. (\neg P \vee Q(x)) \wedge \neg (\neg P \vee \forall y. Q(y)) \\
 &= \forall x. (\neg P \vee Q(x)) \wedge (P \wedge \exists y. \neg Q(y)) \\
 &= \forall x. (\neg P \vee Q(x)) \wedge \exists y'. (P_{y' \neq y} \wedge \neg Q(y')) \\
 &= \forall x. \exists y' ((\neg P \vee Q(x)) \wedge (P_{y' \neq y} \wedge \neg Q(y'))) \\
 &= \forall x. (\neg P \vee Q(x)) \wedge (P_{y' \neq y} \wedge \neg Q(f(x)))
 \end{aligned}$$

Warum umbenennen?

$$\{ \{\neg P, Q(x)\}, \{P\}, \{\neg Q(f(x))\} \}$$

$$\{ \{\neg P, Q(x)\}, \{P\}, \{\neg Q(f(x_1))\} \}$$

$$\{ \{\neg P, Q(f(a)), \{P\}, \{\neg Q(f(a))\} \} \}$$



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## Aufgabe 2: Pränexform, Umwandlung in kM

Algorithmus von Martelli - Montanari

$$1. \{P(\underline{g(a)}, \underline{x}, f(h(y))), P(\underline{y}, f(z), f(z))\}$$

$$\begin{array}{lcl} g(a) & = & y \\ x & = & f(z) \\ f(h(y)) & = & f(z) \end{array}$$

Ersetzung durch erste Gleichung

$$\begin{array}{lcl} y & = & g(a) \\ x & = & f(z) \\ h(y) & = & z \end{array}$$

$$\begin{array}{lcl} y & = & g(a) \\ x & = & f(z) \\ z & = & h(g(a)) \end{array}$$

$$\begin{array}{lcl} y & = & g(a) \\ x & = & f(h(g(a))) \\ z & = & h(g(a)) \end{array}$$

Ergebnis ist mgu mit

$$\mu = \{ \begin{array}{l} x \mapsto f(h(g(a))) \\ y \mapsto g(a) \\ z \mapsto h(g(a)) \end{array} \}$$

$$2. \{Q(\underline{x}, f(g(a)), f(x)), Q(\underline{f(a)}, \underline{y}, y)\}$$

$$\begin{array}{lcl} x & = & f(a) \\ f(g(a)) & = & y \\ f(x) & = & y \end{array}$$

$$\begin{array}{lcl} x & = & f(a) \\ y & = & f(g(a)) \\ y & = & f(f(a)) \end{array}$$

$$\begin{array}{lcl} x & = & f(a) \\ y & = & f(g(a)) \\ f(g(a)) & = & f(f(a)) \end{array}$$

$$\begin{array}{lcl} x & = & f(a) \\ y & = & f(g(a)) \\ g(a) & = & f(a) \end{array} \rightarrow \text{nicht unifizierbar}$$

$$3. \{R(\underline{x}, g(f(a)), f(x)), R(\underline{f(y)}, \underline{z}, y)\}$$

$$\begin{array}{lcl} x & = & f(y) \\ g(f(a)) & = & z \\ f(x) & = & y \end{array}$$

$$\begin{array}{lcl} x & = & f(f(x)) \\ z & = & g(f(a)) \\ y & = & f(x) \end{array} \rightarrow \text{occure check}$$

nicht unifizierbar

$$4. \{S(\underline{a}, \underline{x}, f(g(y))), S(\underline{z}, h(z, u), f(u))\}$$

$$\begin{array}{lcl} a & = & z \\ x & = & h(z, u) \\ f(g(y)) & = & f(u) \end{array}$$

$$\begin{array}{lcl} z & = & a \\ x & = & h(a, u) \\ g(y) & = & u \end{array}$$

$$\begin{array}{lcl} z & = & a \\ x & = & h(a, g(y)) \\ u & = & g(y) \end{array}$$

Ergebnis ist mgu mit

$$\mu = \{ \begin{array}{l} z \mapsto a, \\ x \mapsto h(a, g(y)) \\ u \mapsto g(y) \end{array} \}$$

### Aufgabe 3

$\{\{l(e, o)\}, \{l(c(X, Y), s(N)) \cancel{\text{, }}, \neg l(Y, N)\}, \{\neg l(c(e, e), U)\}\}.$

$\{l(e, o)\}$

$\{l(c(x, y), s(N)), \neg l(y, N)\}$

$\{\neg l(c(e, e), U)\}$

$$\mu = \{ \begin{array}{l} x \mapsto e, y \mapsto e \\ u \mapsto s(N) \end{array}$$

$$\begin{aligned} & \neg l(Y, N) \mu \\ &= \neg l(e, N) \end{aligned}$$

$\{l(c(x_1, y_1), s(N_1)), \neg l(y_1, N_1)\}$

$\{l(e, o)\}$

$$\mu' = \{ N \mapsto o \}$$

$$\begin{aligned} & \{ \} = \neg l(e, N) \mu' \\ &= \neg l(e, o) \end{aligned}$$

$$\Rightarrow U = S(o) \quad \checkmark$$

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## Aufgabe 4:

$$1. \forall x.(S(x) \rightarrow (\neg P(x) \wedge \neg Q(x) \rightarrow R(x))) \vdash \forall x.(S(x) \wedge \neg R(x) \rightarrow P(x) \vee Q(x))$$

$$\frac{\text{axiom}}{\begin{array}{c} S(a), P(a) \vdash P(a), Q(a) \\ \neg R(a) \end{array}} \vdash$$

$$\frac{\text{axiom}}{\begin{array}{c} S(a), Q(a) \vdash P(a), Q(a) \\ \neg R(a) \end{array}} \vdash$$

$$\frac{\text{axiom}}{\begin{array}{c} S(a) \vdash \neg P(a) \\ \neg R(a) \end{array}} \vdash$$

$$\frac{\text{axiom}}{\begin{array}{c} S(a) \vdash \neg Q(a) \\ \neg R(a) \end{array}} \vdash$$

$$\frac{\text{axiom}}{S(a), R(a) \vdash R(a), P(a), Q(a)}$$

$$\frac{\begin{array}{c} S(a), \neg R(a) \vdash \neg P(a) \wedge \neg Q(a) \\ P(a), Q(a) \end{array}}{P(a), Q(a)} \quad \frac{\begin{array}{c} R(a) \\ S(a), \neg R(a) \end{array}}{P(a), Q(a)}$$

$$\frac{\text{axiom}}{S(a), \neg R(a) \vdash S(a), P(a), Q(a)} \quad \frac{\begin{array}{c} \neg P(a) \wedge \neg Q(a) \rightarrow R(a) \\ S(a), \neg R(a) \end{array}}{P(a), Q(a)}$$

$$\frac{\begin{array}{c} S(a) \rightarrow (\neg P(a) \wedge \neg Q(a) \rightarrow R(a)) \\ S(a), \neg R(a) \end{array}}{P(a), Q(a)}$$

ersetze  $x$  durch  
..  $a$

$$\frac{\begin{array}{c} \forall x. (S(x) \rightarrow (\neg P(x) \wedge \neg Q(x) \rightarrow R(x))) \\ S(a) \wedge \neg R(a) \end{array}}{P(a), Q(a)}$$

$\vdash v$

$$\frac{\begin{array}{c} \forall x. (S(x) \rightarrow (\neg P(x) \wedge \neg Q(x) \rightarrow R(x))) \\ S(a) \wedge \neg R(a) \end{array}}{P(a) \vee Q(a)}$$

$\vdash \rightarrow$

$$\frac{\begin{array}{c} \forall x. (S(x) \rightarrow (\neg P(x) \wedge \neg Q(x) \rightarrow R(x))) \\ S(a) \wedge \neg R(a) \end{array}}{S(a) \wedge \neg R(a) \rightarrow P(a) \vee Q(a)}$$

$a$  frisch  
 $\vdash \neq$

$$\frac{\forall x. (S(x) \rightarrow (\neg P(x) \wedge \neg Q(x) \rightarrow R(x)))}{\vdash \forall x. (S(x) \wedge \neg R(x) \rightarrow P(x) \vee Q(x))}$$

Für jeden Schritt muss man die benutzte Regel angeben!

$$\frac{\Delta, P(t) \vdash \Gamma}{\Delta, \forall x. P(x) \vdash \Gamma} \quad \text{A} \vdash$$

$$\frac{\Delta \vdash P(a), \Gamma}{\Delta \vdash \forall x. P(x), \Gamma} \quad \vdash A$$

$$\frac{\Delta, P(a) \vdash \Gamma}{\Delta, \exists x. P(x) \vdash \Gamma} \quad \exists \vdash$$

$$\frac{\Delta \vdash P(t), \Gamma}{\Delta \vdash \exists x. P(x), \Gamma} \quad \vdash \exists$$

2.  $\exists x. \forall y. P(x, y) \vdash \forall x. \exists y. P(y, x).$

$$\frac{\underline{P(a,b)} \vdash P(a,b)}{\underline{P(a,b)} \vdash \exists y. P(y,b)} \quad \text{wähle } y=a \quad \exists\text{-R}$$

$\forall y. P(a,y) \vdash \exists y. P(y,b)$       setze  $b$  für ein  $y$   $\forall\text{-L}$       Regel ?

$a$  frisch       $b$  frisch       $\exists\text{-L}$   
 $\exists x. \forall y. P(x,y) \vdash \forall x. \exists y. P(y,x)$        $\forall\text{-R}$

3. Falls  $x \notin \text{Free}(P)$  gilt:  $\forall x. (P \rightarrow Q(x)) \vdash (P \rightarrow \forall y. Q(y)).$

$$\frac{\text{axiom}}{\underline{Q(a), P \vdash Q(a)}} \quad \text{axiom}$$
$$\frac{\underline{\forall x. Q(x)} \vdash Q(a)}{\underline{P \vdash P, Q(a)}} \quad \text{Regel ?}$$

$P \rightarrow \forall x. Q(x), P \vdash Q(a)$        $a$  frisch

$$\frac{\text{(} P \rightarrow \forall x. Q(x) \text{), } P \vdash \forall y. Q(y) }{P \rightarrow \forall x. Q(x) \vdash P \rightarrow \forall y. Q(y)}$$

Für welchen Schritt ist  $x \notin \text{Free}(P)$  nötig?

-0,5 P

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