

1. 6.1 Give examples of reusable and consumable resources.	<p>Examples of reusable resources are processors, I/O channels, main and secondary memory, devices, and data structures such as files, databases, and semaphores.</p> <p>Examples of consumable resources are interrupts, signals, messages, and information in I/O buffers.</p>	<p>6. 6.6 How can the circular wait condition be prevented?</p> <p>The circular-wait condition can be prevented by defining a linear ordering of resource types.</p> <p>If a process has been allocated resources of type R, then it may subsequently request only those resources of types following R in the ordering.</p>
2. 6.2 What are the three conditions that must be present for deadlock to be possible?	<p>Mutual exclusion. Only one process may use a resource at a time. Hold and wait</p> <p>A process may hold allocated resources while awaiting assignment of others. No preemption.</p> <p>No resource can be forcibly removed from a process holding it</p>	<p>7. 6.7 What is the difference among deadlock avoidance, detection, and prevention?</p> <p>Deadlock prevention constrains resource requests to prevent at least one of the four conditions of deadlock; this is either done indirectly, by preventing one of the three necessary policy conditions (mutual exclusion, hold and wait, no preemption), or directly, by preventing circular wait</p> <p>Deadlock avoidance allows the three necessary conditions, but makes judicious choices to assure that the deadlock point is never reached.</p> <p>Deadlock detection, requested resources are granted to processes whenever possible.; periodically, the operating system performs an algorithm that allows it to detect the circular wait condition.</p>
3. 6.3 What are the four conditions that create deadlock?	<p>The above three conditions, plus: Circular wait A closed chain of processes exists,</p> <p>such that each process holds at least one resource needed by the next process in the chain.s</p>	
4. 6.4 How can the hold-and-wait condition be prevented?	<p>The hold-and-wait condition can be prevented by requiring that a process request all of its required resources at one time,</p> <p>and blocking the process until all requests can be granted simultaneously</p>	
5. 6.5 List two ways in which the no-preemption condition can be prevented.	<p>First, if a process holding certain resources is denied a further request, that process must release its original resources and, if necessary, request them again together with the additional resource.</p> <p>Alternatively, if a process requests a resource that is currently held by another process, the operating system may preempt the second process and require it to release its resources.</p>	