

1. <b>Atomic</b>	Indivisible	16. <b>Priority Inversion</b>	lower-priority process holds or obtains a lock needed by higher-priority process
2. <b>Atomic Execution</b>	A sequence of linear, indivisible or uninterruptible instructions	17. <b>Process resource utilization system model (3)</b>	1. Request 2. Use 3. Release
3. <b>Atomic Instruction</b>	Instructions that cannot be interrupted in their execution	18. <b>Progress</b>	some process p enters its Critical section next
4. <b>Binary Semaphore</b>	integer value can range only between 0 and 1	19. <b>Properties of Solutions to Critical Section Problem (3)</b>	1. Mutual Exclusion 2. Progress 3. Bounded Waiting
5. <b>Bounded Waiting</b>	process p eventually enters its Critical section	20. <b>Race condition</b>	a shared resource is used or accessed in a non-deterministic way that may lead to multiple possible outcomes  Requires that instructions that ran on it were not atomic
6. <b>Counting Semaphore</b>	integer value can range over an unrestricted domain	21. <b>Resource request can be granted to a process only if</b>	request does not result in the formation of a cycle in the resource allocation graph
7. <b>Critical Section</b>	the portion of code that a process must execute in a exclusive fashion	22. <b>Starvation</b>	lack of progress for at least one process caused by indefinite blocking or lack of scheduling
8. <b>Critical Section Problem (CSP) Components (4)</b>	1. entry section, 2. critical section 3. exit section, 4. remainder section		
9. <b>Deadlock</b>	two or more processes are waiting indefinitely for an event that can be caused by only one of the waiting processes		
10. <b>Deadlock Conditions (4)</b>	1. Mutual Exclusion 2. Hold and Wait 3. No Preemption 4. Circular Wait  (all 4 required for deadlock)		
11. <b>Dining Philosopher Problem</b>	chopstick[5] semaphore initialized to 1  algorithm can have deadlock and starvation because philosophers pick up chopsticks one at a time, however these chopstick semaphores can be taken by a neighbor, resulting in possible cascading deadlock		
12. <b>Mutex Lock Con</b>	Requires busy waiting (spinlock)		
13. <b>Mutex Lock Pros (4)</b>	- Simple solution - Process has no context switching if implemented in user space - Can work for short wait times - Enables for mutual exclusion and progress		
14. <b>Mutual Exclusion Property</b>	if a process is in its critical section for a resource, then no other process can execute in the critical section for that resource		
15. <b>Mutual Exclusion Scheduling Problems (3)</b>	1. Deadlock 2. Starvation 3. Priority Inversion		