Quizlet

OS Chapter 5: Concurrency: Mutual Exclusion and Synchronization

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 2 solutions to readers/writers problem 3 Control Problems 	aders/writers to starvation) oblem Writers Have Priority Control Problems th Competing allowed access to critical resource at once) 2. Deadlock	7. Basic requirement for support of concurrent processes is ability to enforce	mutual exclusion
with Competing Processes		8. Binary semaphore	semaphore that only takes value 0 or 1
		9 busy waiting/spin	process can do nothing until it gets
3. 3 Levels of Process Inter-awareness	3. Starvation 1. Processes unaware of each other (multiprogramming of multiple independent processes) 2. Processes indirectly aware of each other (share access to some object (like an I/O buffer); exhibit cooperation in sharing the object) 3. Processes directly aware of each other (communicate with each other using process IDs, work jointly; exhibit cooperation)	waiting	permission to enter its critical section but continues to execute an instruction (or set of them) that tests the appropriate variable to gain entrance
		problems occur when there is a single processor?	Yes, an interrupt can stop instruction execution anywhere in a process.
		1). Central themes of OS design (3)	Multiprogramming Multiprocessing Distributed processing
4. 6 Requirements for Mutual Exclusion	 Mutual exclusion must be enforced A process that halts in non critical section must do so without interfering with other processes. It must not be possible for a process requiring access to a critical section to be delayed indefinitely; no deadlock, no starvation allowed. When no process currently in critical section, any process requesting access to critical section must be permitted to enter without delay. No assumptions made about number of processors or relative process speeds. Process remains inside its critical section for finite time only. 	12. Concurrency arises in 3 different contexts	 multiple applications structured applications operating system structure
		competition inevitably involves the because it is the that allocates resources	Operating System
		14. Difficulties with multiprocessing AND multiprogramming	 difficult to locate programming errors because results not deterministic or reproducible difficult for OS to manage the allocation of resources optimally sharing of global resources fraught with peril
5. advantages of enforcing mutual exclusion with special machine instruction	1. Applicable to any number of processes (either single processor or multiple processors sharing main memory) 2. Simple, easy to verify 3. Can be used to support multiple critical sections	15. Dijkstra's treatise	Two or more processes can cooperate by means of simple signals (semaphores), such that a process can be forced to stop at a specified place until it has received a specific signal
		16. disadvantages of enforcing mutual	busy waiting is employed (wasted processor time)
6. Basic problem in multiprogramming systems?	relative speed of execution of processes cannot be predicted	exclusion with special machine instruction	starvation possible (selection of waiting process arbitrary) Deadlock is possible
		17. Distributed processing	management of multiple processes executing on multiple, distributed computer systems

 18. Event flag 19. Fundamental to all OS themes is 20. How are processes 	memory word used as synchronization mechanism; each bit represents event, threads check bits to check status of certain events; thread blocked until all required bits set (AND) or until one is set (OR) concurrency	31. Producer/Consumer problem	one or more producers generating some type of data and placing these in a buffer; a single consumer is taking items out of buffer one at a time; system is to be constrained to prevent overlap of buffer operations. (only one agent may access buffer at any one time); ensure that producer doesn't add data if buffer full, consumer doesn't take out if empty
removed from the queue of processes waiting on semaphore?	"strong semaphore"); guarantee freedom from starvation	32. Race condition when multi read and w final result	when multiple processes or threads read and write data items so that the final result depends on the order of execution of instructions in the
21. In a multiprocessor architecture does not guarantee	preventing interruption		multiple processes; "loser" of the race to write determines the final value
mutual exclusion 22. In a uniprocessor system, concurrent processes cannot have execution; they can only be	overlapped; interleaved	33. Readers/Writers Problem	there is a data area shared among a number of processes. The data area could be a file, block of main memory, or even a bank of processor registers. There are a number of processes that only read the data area (readers) and a number that only write to the data area (writers). The conditions that must be satisfied are: 1. any number of readers may simultaneously read 2. only one writer at a time may write 3. if a writer is writing to the file, no reader may read it
23. Key terms related to concurrency	atomic operation critical section deadlock livelock mutual exclusion race condition		
24. Monitor	programming language construct that encapsulates variables, access procedures and initialization code within an abstract data type	34. Semaphore	an integer value used for signaling among processes; may be initialized to nonnegative int, decremented, or incremented; decrement may result in process blocking, increment may result in unblocking.
25. Multiprocessing	management of multiple processes within a multiprocessor	35. semSignal	operation that increments semaphore value; if resulting value less than or equal to 0, then process blocked by semWait operation, if any, is unblocked
26. Multiprogramming	management of multiple processes within a uniprocessor system		
27. Mutex	mutual exclusion lock; similar to binary semaphore, but process that locks the MUTEX (sets it to 0) must be the one to unlock it (sets it to 1)	36. semSignalB	operation that checks if any processes are blocked on the binary semaphore (if equal 0); if so, process blocked by semWaitB operation is now unblocked;
28. Processes directly aware of each other, relationship?	cooperation by communication		if none blocked, value of semaphore set to 1
29. Processes indirectly aware of each other, relationship?	cooperation by sharing	37. semwait	operation that decrements semaphore value; if value becomes negative, process executing semWait is blocked; otherwise, continue execution
30. Processes unaware of each other,	competition		

relationship?

38. semWait and semSignal primitives are assumed to be	atomic
39. semWaitB	operation that checks the binary semaphore value; if 0, process executing semWaitB is blocked; if 1, value changed to 0 and process continues execution
40. spinlock	mutual exclusion mechanism in which a process executes in an infinite loop waiting for value of a lock variable to indicate availability
41. To guarantee mutual exclusion, it is sufficient to	prevent a process from being interrupted
42. What design and management issues are raised by concurrency?	 OS must track various processes (with process control blocks) OS must allocate and deallocate various resources for each active process. Multiple processes may want same resources simultaneously OS must protect data and physical resources of each process against unintended interference by other processes Function and output of a process must be independent of the speed at which its execution is carried out relative to speed of other concurrent processes
43. Why do concurrency problems occur in multiprocessor systems?	Two processes may be executing simultaneously and both trying to access the same global variable
44. With cooperation among processes by sharing, only operations must be mutually exclusive	writing