

OS Chapter 5 - Concurrency: Mutual Exclusion & Synchronization

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 Addressing - Direct Addressing - Indirect 	requires destination requires sources also implicit put in queue for receivers to check "port"	13. mutual exclusion - interrupt disabling	+process will run until interrupt +disabling interrupt guarantees mutual exclusion -efficiency could be degraded -does not work in multi-proc architecture
 atomic operation Can relative speed of execution of processes be predicted on a uniprocessor? 	not breakable into smaller parts no	14. mutual exclusion - special machine instruction	+applicable to any # of processors +easy to verify -busy-waiting employed, cpu time wasted -starvation possible -deadlock possible
5. common concurrency mechanisms	semaphore mutex monitor event flags messages spinlocks	15. O.S. Concerns for concurrency	OS must be able to keep track of various processes OS must be able to allocate and deallocate resources for each active process OS must protect the data and physical resources of each process against interference by other processes OS must ensure the processes and outputs are independent of the
6. critical section	section of code requiring shared resources and will fail if run when another process is in corresponding section of code		
7. deadlock	two or more processes are each waiting for the other to do something		processing speed Interleaving
8. Difficulties of concurrency	sharing of global resources mgmt of optimal resource allocation results are not deterministic & reproducible, so difficult to locate	77. producer/consumer problem 18. race condition	overlappting producer places data into buffer consumer takes from buffer only one may access at a time
9. livelock	programming errors two processes continually change states cancelling each other out	19. Resource Competition	- must be enforced - halts must not interfere with other
10. message passing	synchronization & communication 3 types both sender and receiver are blocked until msg delivered - tight synch nonblocking sender, blocking receiver - most useful (e.g. service provides resource to other services) neither blocked		processes - no deadlock or starvation - must not be denied access to critical section when nothing else is using it - no assumptions about processor speed - processor remains inside critical section for finite time
		20. semaphore operations	 initialization decrease (semWait) increase (semSignal)
11. monitors	data structure - control variable like semaphore, but easier to control	21. semaphore status levels	sv>0 = free sv==0 = used, no waiting sv<0 = used, process waiting
12. mutual exclusion	when one process is in crit sect, no other proc can be in a crit section accessing shared resources	22. semSignal & semWait	atomic primitives implemented in firmware also via software algorithm