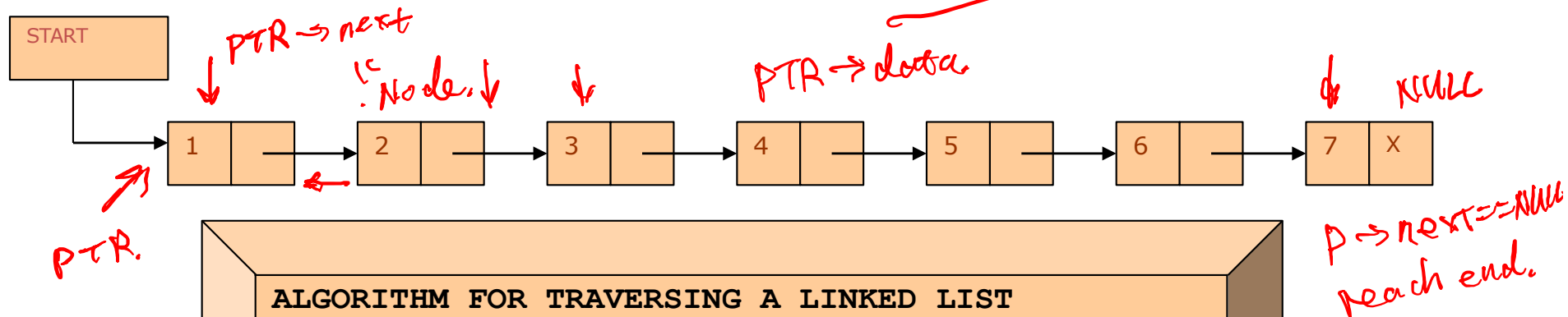


Singly Linked Lists

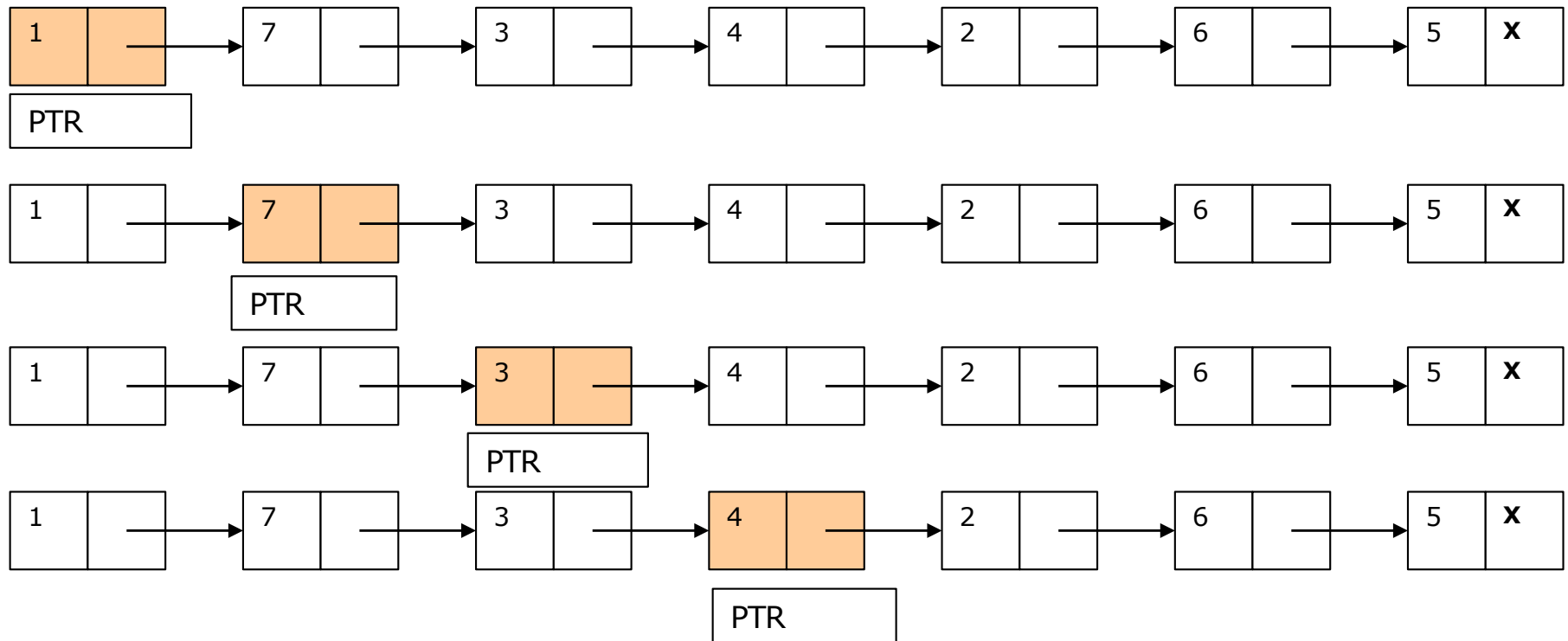
- A singly linked list is the simplest type of linked list in which every node contains some data and a pointer to the next node of the same data type.
- Traversing a linked list means accessing the nodes of the list in order to perform some processing on them.



ALGORITHM FOR TRAVERSING A LINKED LIST

```
Step 1: [INITIALIZE] SET PTR = START
Step 2: Repeat Steps 3 and 4 while PTR != NULL
Step 3:         Apply Process to PTR->DATA
Step 4:         SET PTR = PTR->NEXT
               [END OF LOOP]
Step 5: EXIT
```

Searching for Val 4 in Linked List



Here PTR -> DATA = 4. Since PTR -> DATA = 4, POS = PTR. POS now stores the address of the node that contains VAL.

Search.

```
struct Node
{
    int data;
    struct Node *next;
}
```



$P == NULL$

Time Complexity

$O(n)$, linear

↑ length of linked list.

1) traverse the linked list

while ($P \neq NULL$) $O(n)$

```
{
    if (  $P \rightarrow data == 4$  ) {
        printf("Found value\n");
        return;
    }
    P = P -> next;
}
```

$O(n) - O(1) = O(n)$

→ if $P == NULL$

printf("Not Found\n") → $O(1)$
returns

Searching a Linked List

Searching a linked list means to find a particular element in the linked list.

ALGORITHM TO SEARCH A LINKED LIST

```
Step 1: [INITIALIZE] SET PTR = START
Step 2: Repeat Step 3 while PTR != NULL
Step 3:     IF VAL = PTR->DATA
            SET POS = PTR
            Go To Step 5
        ELSE
            SET PTR = PTR->NEXT
        [END OF IF]
    [END OF LOOP]
Step 4: SET POS = NULL
Step 5: EXIT
```

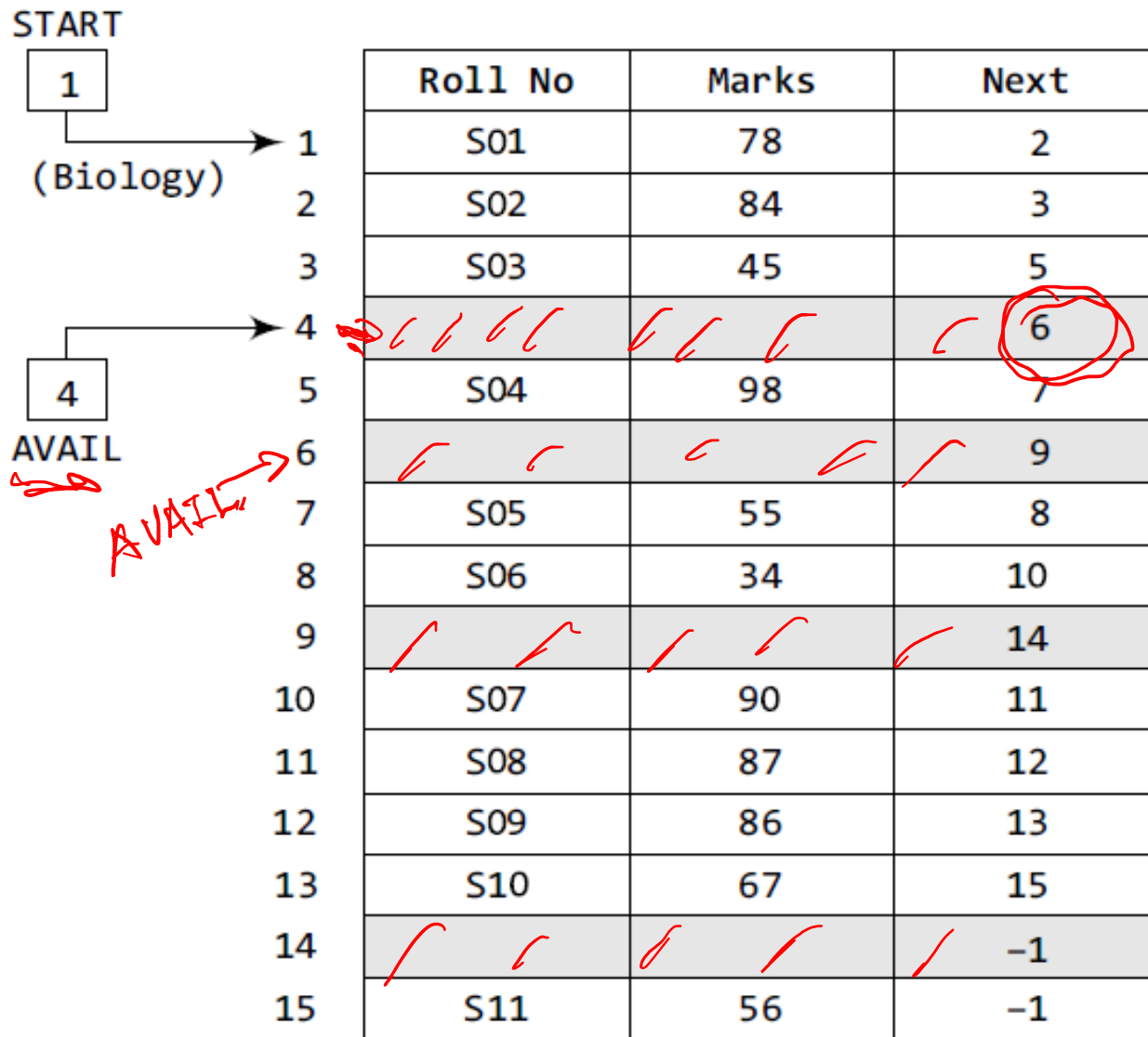
Inserting a New Node in a Linked List

- Four cases
 - Case 1: The new node is inserted at the beginning.
 - Case 2: The new node is inserted at the end.
 - Case 3: The new node is inserted after a given node.
 - Case 4: The new node is inserted before a given node.
- Overflow is a condition that occurs when $AVAIL = NULL$ or no free memory cell is present in the system.
 - When this condition occurs, the program must give an appropriate message.
 - $AVAIL$ is a pointer which stores the address of the first free space.

Memory Allocation and De-Allocation

- If we want to add a node to an already existing linked list in the memory, we first find free space in the memory and then use it to store the information.
- Now, the question is which part of the memory is available and which part is occupied? When we delete a node from a linked list, then who changes the status of the memory occupied by it from occupied to available? The answer is the operating system.
- The computer maintains a list of all free memory cells. This list of available space is called the *free pool*.

Free Pool or Heap



- For the free pool (which is a linked list of all free memory cells), we have a pointer variable AVAIL which stores the address of the first free space.
- Now, when a new student's record is added, the memory address pointed by AVAIL will be taken and used to store the desired information.
- After the insertion, the next available free space's address will be stored in AVAIL.
- For example, when the first free memory space, 4, is utilized for inserting the new node, AVAIL will be set to contain address 6.

Overflow C, Language.

```
struct Node * newnode = (struct Node*) malloc (sizeof (struct Node));
```

```
if (newNode == NULL),
```

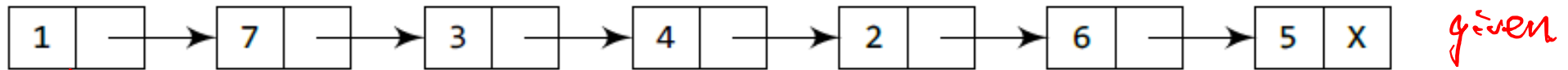
```
{ printf ("Memory overflow");
```

```
return;
```

```
}
```


Inserting a Node at the Beginning

Suppose we want to add a new node with data 9 and add it as the first node of the list.



START

Allocate memory for the new node and initialize its DATA part to 9.



*newNode → data = 9;
newNode → next = NULL;*

Add the new node as the first node of the list by making the NEXT part of the new node contain the address of START.



START

Now make START to point to the first node of the list.



START

Inserting a Node at the Beginning

ALGORITHM TO INSERT A NEW NODE IN THE BEGINNING OF THE LINKED LIST

```
Step 1: IF AVAIL = NULL, then
        Write OVERFLOW
        Go to Step 7
    [END OF IF]
Step 2: SET New_Node = AVAIL
Step 3: SET AVAIL = AVAIL->NEXT
Step 4: SET New_Node->DATA = VAL
Step 5: SET New_Node->Next = START
Step 6: SET START = New_Node
Step 7: EXIT
```

1/ given. newNode.

newNode \rightarrow data = 9;

newNode \rightarrow next = NULL;

2/ step 2 =

newNode \rightarrow next = START;

3/ step 3: Update header

START = newNode;

~~first = new~~

① START == NULL, empty linked list.
✓

Time Complexity

$O(1)$