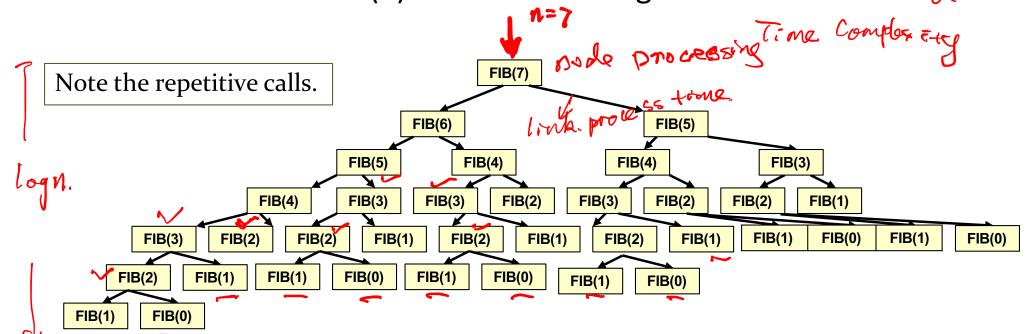
#### Fibonacci Series

The Fibonacci series can be given as:

```
0 1 1 2 3 5 8 13 21 34 55 ...... Fibonacci(n) = Fibonacci(n – 1) + Fibonacci(n – 2)
```

- Each term in the series is the sum of the two previous terms except for the first and second terms of 6 and 1.
- The recursive call FIB(7) has the following recursive call tree. p[nlog<sup>n</sup>]



#### Pros and Cons of Recursion

#### Pros

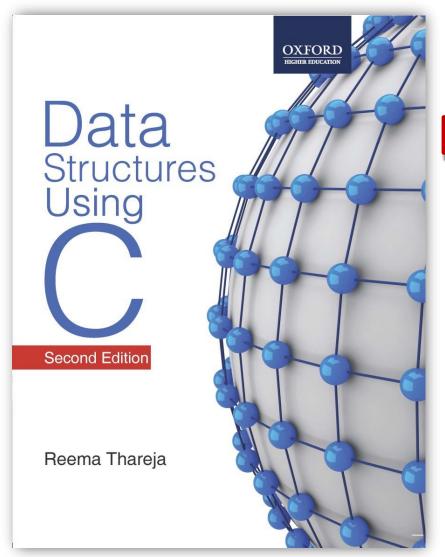
- Recursive solutions often tend to be shorter and simpler than non-recursive ones.
- Code is clearer and easier to use.
- Recursion follows a divide and conquer technique to solve problems.
- In some (limited) instances, recursion may be more efficient.

#### Cons

- Recursion is implemented using system stack. If the stack space on the system is <u>limited</u>, recursion to a deeper level will be difficult to implement.
- Aborting a recursive process in midstream is slow.
- Using a recursive function takes more memory and time to execute as compared to its non-recursive counterpart.
- It is difficult to find bugs, particularly when using global variables.







# Data Structures Using C, 2e

#### Reema Thareja

## Chapter 8

Queues

#### Introduction

- A queue is a FIFO (First-In, First-Out) data structure in which the element that is inserted first is the first one to be taken out.
- The elements in a queue are added at one end called the rear and removed from the other end called the front.
- Queues can be implemented using arrays or linked lists.

Basic operations: Enqueue and Dequeue

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#### Applications of Queues

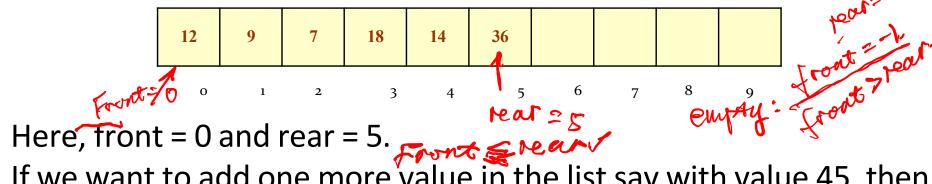
- Queues are widely used as waiting lists for a single shared resource like printer, disk, CPU.
- Queues are used to transfer data asynchronously, e.g. pipes, file I/O, sockets.
- Queues are used in Playlist for youtube to add songs to the end, play from the front of the list.

### Array Representation of Queues

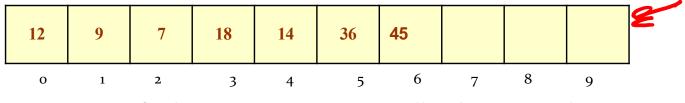
Queues can be easily represented using linear arrays.

Every queue has front and rear variables that point to the position from where deletions and insertions can be done, respectively.

Consider the queue shown in the figure below.



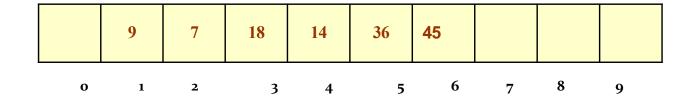
If we want to add one more value in the list say with value 45, then rear would be incremented by 1 and the value would be stored at the position pointed by rear.



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## Array Representation of Queues

- Now, front = 0 and rear = 6.
- Now, if we want to delete an element from the queue, then the value of front will be incremented.



Now, front = 1 and rear = 6.

#### Array Representation of Queues

- Before inserting an element in the queue, we must check for overflow conditions.
- An overflow occurs when we try to insert an element into a queue that is already full, i.e. when rear = MAX 1, where MAX specifies the maximum number of elements that the queue can hold.
- Similarly, before deleting an element from the queue, we must check for underflow condition.
- An underflow occurs when we try to delete an element from a
  queue that is already empty. If front = -1 or front > rear, this
  means there is no element in the queue.

#### Algorithm for Insertion Operation

```
Algorithm to insert an element in a queue
Step 1: IF REAR=MAX-1, THEN
             Write OVERFLOW
             Goto Step 4
        [END OF IF]
Step 2: IF FRONT == -1 and REAR = -1, THEN
        \sim SET FRONT = REAR = 0
        ELSE
            SET REAR = REAR + 1
        [END OF IF]
Step 3: SET QUEUE[REAR] = NUM
Step 4: EXIT
```

### Algorithm for Deletion Operation

```
Algorithm to delete an element from a queue
Step 1: IF FRONT = -1 OR FRONT > REAR, THEN
Write UNDERFLOW

ELSE

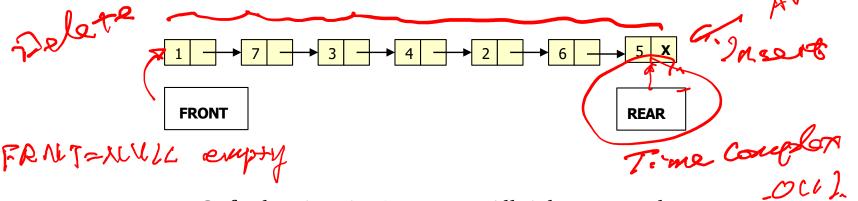
SET VAL = QUEUE[FRONT]
SET FRONT = FRONT + 1

[END OF IF]
Step 2: EXIT
```

### Linked Representation of Queues

- In a linked queue, every element has two parts: one that stores data and the other that stores the address of the next element.
- The START pointer of the linked list is used as FRONT.
- We will also use another pointer called REAR which will store the address of the last element in the queue.
- All insertions will be done at the rear end and all the deletions will be done at the front end.

If FRONT = NULL, then it indicates that the queue is empty.



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