



養天地正氣 法古今完人

# 图的遍历



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计算机科学与技术学院,  
苏州大学





## 图的遍历

- ❖ *Depth-first traversal*(深度优先遍历) of a graph is roughly **analogous to preorder traversal of an ordered tree**. Suppose that the traversal has just visited a vertex  $v$ , and let  $w_1; w_2; \dots; w_k$  be the vertices adjacent to  $v$ . Then we shall next visit  $w_1$  and keep  $w_2; \dots; w_k$  waiting. After visiting  $w_1$ , we traverse all the vertices to which it is adjacent before returning to traverse  $w_2; \dots; w_k$ .
- ❖ *Breadth-first traversal*(广度或宽度优先遍历) of a graph is roughly analogous to level-by-level traversal(层序遍历) of an ordered tree. If the traversal has just visited a vertex  $v$ , then it next visits *all* the vertices adjacent to  $v$ , putting the vertices adjacent to these in a waiting list to be traversed after all vertices adjacent to  $v$  have been visited.

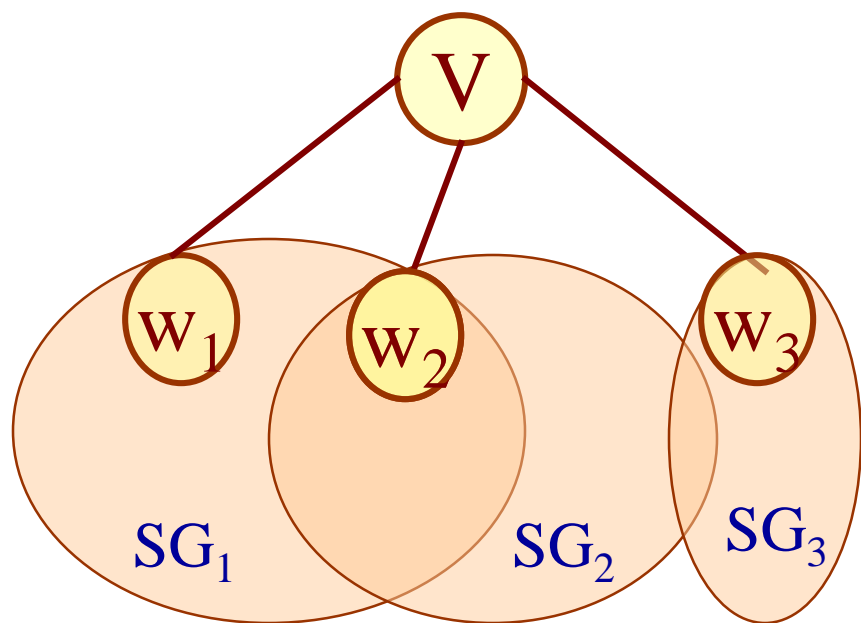




# 连通图的深度优先搜索遍历

从图中某个顶点 $V_0$ 出发，访问此顶点，然后依次从 $V_0$ 的各个未被访问的邻接点出发深度优先搜索遍历图，直至图中所有和 $V_0$ 有路径相通的顶点都被访问到。





$W_1$ 、 $W_2$ 和 $W_3$  均为  $V$  的邻接点,  $SG_1$ 、 $SG_2$ 和 $SG_3$  分别为含顶点 $W_1$ 、 $W_2$ 和 $W_3$  的子图。

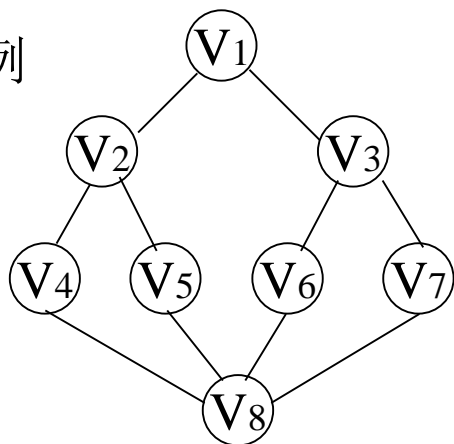
访问顶点  $V$  :

for ( $W_1$ 、 $W_2$ 、 $W_3$ )

若该邻接点 $W$ 未被访问,

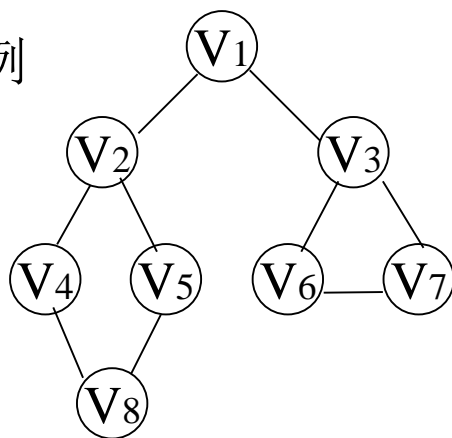
则从它出发进行深度优先搜索遍历。

例



深度遍历:  $V1 \Rightarrow V2 \Rightarrow V4 \Rightarrow V8 \Rightarrow V5 \Rightarrow V6 \Rightarrow V3 \Rightarrow V7$

例



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# 连通图的深度优先搜索遍历

1. 从深度优先搜索遍历连通图的过程类似于树的先根遍历;

2. 如何判别 $V$ 的邻接点是否被访问?

解决的办法是: 为每个顶点设立一个 “访问标志  
 $visited[0..n-1]$ ”。







# 连通图的深度优先搜索遍历



## Depth-First算法

```
template <int max_size>
```

```
void Digraph<max_size> :: traverse(Vertex &v, bool  
    visited[], void (*visit)(Vertex &)) const
```

```
/* Pre: v is a vertex of the Digraph .
```

```
Post: The depth-first traversal, using function*visit ,  
    has been completed for v and for all vertices that  
    can be reached from v .
```

```
Uses: traverse recursively. */
```





# 连通图的深度优先搜索遍历

## ❖ Depth-First算法（续）

{

Vertex w;

visited[v] = **true**;

(\*visit)(v);

**for** (all w adjacent to v)

**if** (!visited[w])

        traverse(w, visited, visit);

}



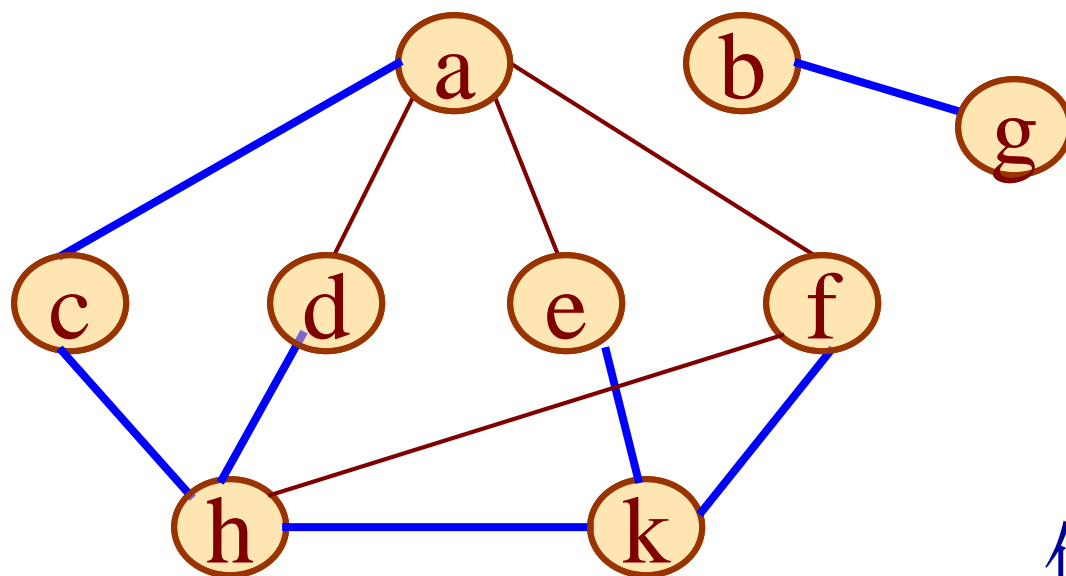




# 非连通图的深度优先搜索遍历

首先将图中每个顶点的访问标志设为 FALSE，之后搜索图中每个顶点，如果未被访问，则以该顶点为起始点，进行深度优先搜索遍历，否则继续检查下一顶点。





例如:

0a 1b 2c 3d 4e 5f 6g 7h 8k

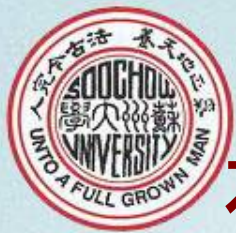
访问标志:

T	T	T	T	T	T	T	T	T
---	---	---	---	---	---	---	---	---

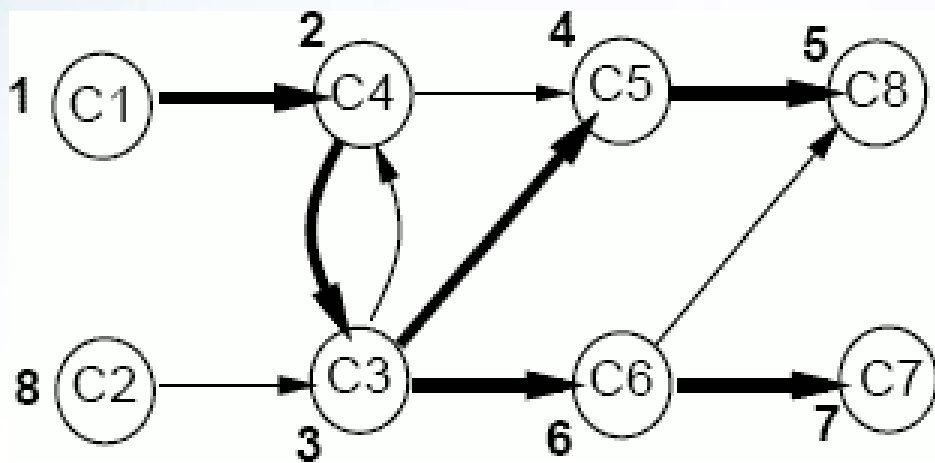
访问次序:

a	c	h	d	k	f	e	b	g
---	---	---	---	---	---	---	---	---





# 有向图的深度优先搜索遍历





# 非连通图的深度优先搜索遍历

❖ Depth-First 算法:

```
template <int max_size>
```

```
void Digraph<max_size> :: depth_first(void (*visit)(Vertex &))  
    const
```

```
/* Post: The function *visit has been performed at each vertex of  
the Digraph in depth-first order.
```

```
Uses: Method traverse to produce the recursive depth-first order. */
```

```
{
```

```
bool visited[max_size];
```

```
Vertex v;
```

```
for (all v in G) visited[v] = false;
```

```
for (all v in G)
```

```
    if (!visited[v])
```

```
        traverse(v, visited, visit);
```

```
}
```

