

## Lecture 1.

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### Chapter I. Generators & Relations.

1. Free semigroups. Let  $X$  be a nonempty set. We will call  $X$  an alphabet and each of its elements a letter in the alphabet.

A word in  $X$  is either a finite sequence of letters written juxtaposed next to one another or the empty word.

For example, if  $X = \{x_1, x_2, \dots\}$  then  $x_3 x_4 x_3 x_1$  is a word in  $X$ .

We denote the set of all words in the alphabet  $X$  as  $X^*$ .

Then  $X^*$  is a multiplicative semigroup

with multiplication defined by concatenation of words.

Example. Let  $v = x_3 x_4 x_1$ ,  $w = x_3$ .

Then  $v.w = x_3 x_4 x_1 x_3$ ,  $w.v = x_3^2 x_4 x_1$ .

Proposition <sup>I.1.1</sup> Let  $S$  be a semigroup. An arbitrary mapping  $X \xrightarrow{\varphi} S$  uniquely extends to a homomorphism  $X^* \rightarrow S$ .

Proof. Given  $\varphi : X \rightarrow S$  define  $\bar{\varphi}$  sending a word  $v = x_{i_1} x_{i_2} \dots x_{i_k}$  to

$$\bar{\varphi}(v) = \underbrace{\varphi(x_{i_1}) \cdot \varphi(x_{i_2}) \dots \varphi(x_{i_k})}_{\text{here multiplication in } S}$$

Def. Let  $S$  be a semigroup. An equivalence relation  $\sim \subseteq S \times S$  is called a congruence if  $a \sim b, c \sim d \Rightarrow ac \sim bd$ .

In this case we can define the semigroup on  $S/\sim = \{\text{equivalence classes}\}$ . Moreover  $S \rightarrow S/\sim, a \rightarrow a/\sim$  is a homomorphism of semigroups.

Given a homomorphism  $\varphi: S_1 \rightarrow S_2$  of semigroups

$a \sim b$  if and only if  $\varphi(a) = \varphi(b)$

is a congruence.

Congruence = analog of a normal subgroup in GROUPS and an ideal in RINGS.

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Suppose that a semigroup  $S$  is generated by a subset  $\{a_i, i \in I\} \subset S$ .

Consider the alphabet

$$X = \{x_i, i \in I\}$$

and the mapping

$$\varphi : x_i \rightarrow a_i, i \in I.$$

The mapping  $\varphi$  extends to an epimorphism

$$\bar{\varphi} : X^* \rightarrow S$$

Let  $a \sim b$  iff  $\bar{\varphi}(a) = \bar{\varphi}(b)$ ;  $a, b \in X^*$ .

Then  $S \cong X^* / \sim$

So, every semigroup is a homomorphic image of a free semigroup of an appropriate rank.

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Remember:  $\sim \subset X^* \times X^*$

We say that  $R \subset \sim$  generates  $\sim$  if  $\sim$  is the smallest congruence that contains  $R$ , so

$\sim = \cap$  (all congruences on  $X^*$  that contain  $R$ )

Then  $R$  uniquely determines  $\sim$  and, hence, uniquely determines the semigroup  $S$  up to isomorphism.

Let  $R = \{a_j \times b_j\}_{j \in J} \subset X^* \times X^*$ ;  $a_j, b_j$  are words. We write:

$$S = \langle X \mid a_j = b_j, j \in J \rangle.$$

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why is it important to have a nice presentation by generators and relations?

Let  $S$  be a semigroup and let  $s_1, \dots, s_n$  be a set of generators of  $S$ .

Let  $T$  be another semigroup.

Not every mapping  $s_i \rightarrow t_i \in T, 1 \leq i \leq n$ , extends to a homomorphism  $S \rightarrow T$ .  
How can we find out if it extends or not?

Suppose that we know a presentation of the semigroup  $S$  in these generators.

$$S = \langle x_1, \dots, x_n \mid a_1(x) = b_1(x), \dots, a_m(x) = b_m(x) \rangle$$

It means the following:

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let  $\sim$  be the congruence ~~that~~ on  $X^*$  that corresponds to the homomorphism  $x_i \rightarrow \Delta_i, 1 \leq i \leq n$ . The congruence  $\sim$  is generated by  $a_1 \times b_1, \dots, a_m \times b_m$ .

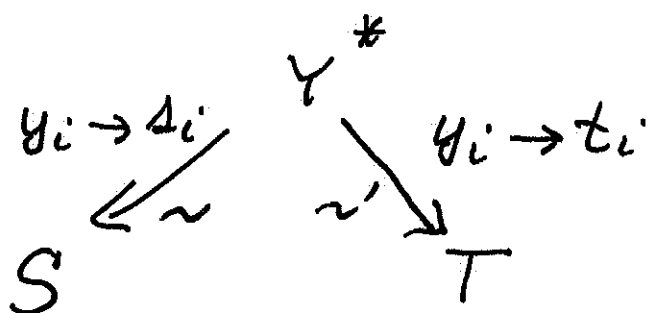
Proposition I.1.2. Let  $\varphi: \Delta_i \rightarrow t_i \in T, 1 \leq i \leq n$  be a mapping. This mapping extends to a homomorphism  $S \rightarrow T$  if and only if  $a_i(t_1, t_2, \dots, t_n) = b_i(t_1, \dots, t_n), 1 \leq i \leq m$ .

Proof. In one direction the assertion is clear: if  $\varphi$  extends to a homomorphism then  $a_i(t_1, t_2, \dots, t_n) = b_i(t_1, \dots, t_n)$ .

Now suppose that these equalities

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hold. Consider the homomorphisms



Let  $\sim'$  be the congruence that corresponds to the homomorphism  $y_i \rightarrow t_i, 1 \leq i \leq n$ .

We have  $a_i(t_1, \dots, t_n) = b_i(t_1, \dots, t_n), 1 \leq i \leq n$ , hence  $a_i(Y) \sim' b_i(Y), 1 \leq i \leq n$ , hence  $\sim'$  contains  $a_i \times b_i, 1 \leq i \leq n$ . Hence

$$\sim \subseteq \sim'$$

This implies that the mapping

$$u(a_1, \dots, a_n) \rightarrow u(t_1, \dots, t_n), u \in Y^*$$

is well defined.



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Indeed, if  $u(a_1, \dots, a_n) = v(a_1, \dots, a_n)$  then  $u \sim v$ . This implies  ~~$u(t_1, \dots, t_n) = v(t_1, \dots, t_n)$~~  that  $u \sim' v$ , i.e.  $u(t_1, \dots, t_n) = v(t_1, \dots, t_n)$ . This mapping is a homomorphism. This completes the proof of the Proposition.

Equalities  $a_j = b_j, j \in J$ , are called defining relations.

Let  $v, w$  be words in the alphabet  $X$ .

We say that  $w$  is obtained from  $v$  by substitution if some word  $a_j$  is a subword of  $v$ ,  $v = v' a_j v''$  and  $w = v' b_j v''$  or some word  $b_j$  is a subword of  $v$ ,  $v = v' b_j v''$  and  $w = v' a_j v''$ .

In this case we write  $v \rightarrow w$ .

This relation is symmetric, if  $v \rightarrow w$ , then  $w \rightarrow v$ .

Proposition <sup>I.1.3.</sup> Let  $S = \langle X \mid a_j = b_j, j \in J \rangle$ .

Words  $v, w$  are equal in  $S$  if and only if there is a finite sequence

$$v = v_0 \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_k = w.$$

Proof. Let  $\sim$  be the congruence in  $X^*$

generated by all elements  $a_j \times b_j, j \in J$ .

We need to prove that  $v \sim w$  if and only if there is a finite sequence

$$v = v_0 \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_k = w.$$

Define another congruence  $\approx$  as follows:

$v \approx w$  if and only if there exists a finite sequence

$$v = v_0 \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_k = w.$$

It is easy to see that  $\approx$  is a congruence.

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All elements  $a_j \times b_j, j \in J$ , belong to  $\approx$ .

Since  $\sim$  is a minimal congruence containing all  $a_j \times b_j, j \in J$ , we conclude that

$$\sim \subseteq \approx$$

On the other hand if

$$v = v_0 \rightarrow v_1 \rightarrow \dots \rightarrow v_k = w$$

then  $v \sim v_1, v_1 \sim v_2, \dots, v_{k-1} \sim w$ , hence

$$v \sim w.$$

In other words,  $\approx \subseteq \sim$ . We proved that

$$\sim = \approx.$$

A semigroup  $S$  is called finitely presented if it has a finite presentation

$$S \cong \langle x_1, \dots, x_n \mid a_1 = b_1, \dots, a_m = b_m \rangle, \\ a_i, b_i \in X^*.$$

It means that  $S$  is generated by a finite subset  ~~$\{x_1, \dots, x_n\}$~~   $\Delta_1, \dots, \Delta_n$  and the congruence on  $X^*$ ,  $X = \{x_1, \dots, x_n\}$ , defined by the homomorphism  $x_i \rightarrow \Delta_i$ ,  $1 \leq i \leq n$ , is finitely generated.

Proposition <sup>I.1.4</sup>. Let  $S$  be a semigroup. Let  $\{\Delta_1, \dots, \Delta_n\}$  and  $\{\Delta'_1, \dots, \Delta'_k\}$  be finite generating subsets of  $S$ . If  $S$  is finitely presented in  $\{\Delta_1, \dots, \Delta_n\}$  then it is finitely presented also in  $\{\Delta'_1, \dots, \Delta'_k\}$ .

Proof. Let  $\sim$  be the congruence on  $\{x_1, \dots, x_n\}^*$  that corresponds to the homomorphism  $x_i \rightarrow \Delta_i$ ,  $1 \leq i \leq n$ .

Let  $\sim'$  be the congruence on  $\{y_1, \dots, y_k\}^*$  that corresponds to the homomorphism  $y_j \rightarrow \Delta_j', 1 \leq j \leq k$ .

Let  $\sim$  be generated by the relations

$$a_\mu(x_1, \dots, x_n) \sim b_\mu(x_1, \dots, x_n),$$

$\mu$  runs over a finite set of numbers.

Let  $\Delta_i = C_i(\Delta_1', \dots, \Delta_k'), 1 \leq i \leq n, C_i \in \{y_1, \dots, y_k\}^*$

~~We claim that the congruence  $\sim'$  is generated by~~

~~$$a_\mu(C_1(y_1, \dots, y_k), \dots, C_n(y_1, \dots, y_k)) \sim'$$~~

~~$$b_\mu(C_1(y_1, \dots, y_k), \dots, C_n(y_1, \dots, y_k)).$$~~

~~First, we notice that~~

Let  $\Delta'_j = d_j(\Delta_1, \dots, \Delta_n), 1 \leq j \leq K, d_j \in \{x_1, \dots, x_n\}^*$ .

Then

$$a_\mu(c_1(y_1, \dots, y_K), \dots, c_n(y_1, \dots, y_K)) \sim' \quad (I)$$

$$b_\mu(c_1(y_1, \dots, y_K), \dots, c_n(y_1, \dots, y_K)),$$

$$y_j \sim' d_j(c_1(y_1, \dots, y_K), \dots, c_n(y_1, \dots, y_K)) \quad (II)$$

We claim that the congruence  $\sim'$  is generated by (I) and (II).

In other words: if  $\sim''$  is a congruence on  $\mathcal{Y}^*$  and

$$a_\mu(c_1(y), \dots, c_n(y)) \sim'' b_\mu(c_1(y), \dots, c_n(y)),$$

$$y_j \sim'' d_j(c_1(y), \dots, c_n(y))$$

then  $\sim' \subseteq \sim''$ . Let  $T = Y^*/\sim''$ .

Consider the mapping  $\Delta_i \rightarrow C_i(Y)/\sim''$ .

By the Proposition this mapping extends to a homomorphism  $\varphi: S \rightarrow T$ .

Since  $\Delta_j' = d_j(\Delta_1, \dots, \Delta_n)$  ~~and~~ we have

$$\varphi(\Delta_j') = d_j(C_1(Y), \dots, C_n(Y))/\sim'' = y_j/\sim''.$$

Now, if  $u(y_1, \dots, y_k) \sim' v(y_1, \dots, y_k)$  i. e.

$$u(\Delta_1', \dots, \Delta_k') = v(\Delta_1', \dots, \Delta_k')$$

then

$u(y_1, \dots, y_k) \sim'' v(y_1, \dots, y_k)$ . We proved that

$$\sim' \subseteq \sim''.$$