

# Complexities

## Array

Append  
Sorted Append  
Delete (+shift)  
Binary Search

## Linked List

Insert  
Insert (Enhanced)  $\mathcal{O}(1)$   
Delete (+shift)  $\mathcal{O}(n)$   
Traverse(Search)  $\mathcal{O}(n)$

Enhanced  
Linked Lists  
give pointers  
to both head  
and tail

## Tree/Heap

Insert  
Delete

## Heapify/Heap sort

Top Down  $\mathcal{O}(n \cdot \log_2 n)$   
Bottom up  $\mathcal{O}(n)$

For the  $T(n)$  for the number of sinks run by Heapify, the sequence is:

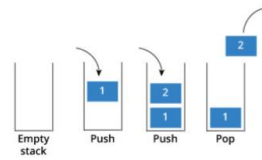
$$T(n) = \frac{n}{2} \cdot 0 + \frac{n}{4} \cdot 1 + \frac{n}{8} \cdot 2 + \dots + 1 \cdot h$$

## Stacks

Push  $\mathcal{O}(1)$   
Pop  $\mathcal{O}(1)$

## Queue

Enqueue  $\mathcal{O}(1)$   
Dequeue  $\mathcal{O}(1)$



## Big O no

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \neq 0, \infty \Rightarrow f \in \Theta(g)$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \neq \infty \Rightarrow f \in \mathcal{O}(g)$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \neq 0 \Rightarrow f \in \Omega(g)$$

worst case  $f(n) \in \mathcal{O}(g(n)) : f(n) \leq Cg(n)$

best case  $f(n) \in \Omega(g(n)) : f(n) \geq Cg(n) \forall n > n_0$

average case  $f(n) \in \Theta(g(n)) : C_1g(n) \leq f(n) \leq C_2g(n)$

## Heaps

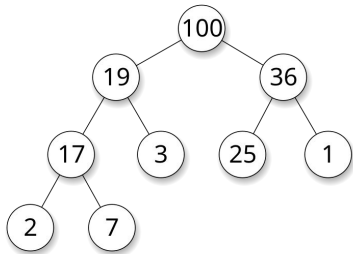
These are the *only* conditions:

### Max Heap

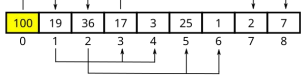
### Min Heap

parent  $\geq$  child    parent  $\leq$  child

Tree representation



Array representation



Array, node at index i:

$$\text{left}(i) = 2i + 1$$

$$\text{right}(i) = 2i + 2$$

$$\text{parent}(i) = \left\lfloor \frac{i-1}{2} \right\rfloor$$

### Height and Size

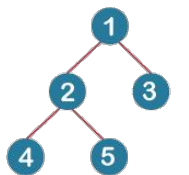
$$h(T) = \begin{cases} -1 & \text{if } T \text{ is null} \\ 1 + \max(h(T_{\text{left}}), h(T_{\text{right}})) & \text{otherwise} \end{cases}$$

$$s(T) = \begin{cases} 0 & \text{if } T \text{ is null} \\ 1 + s(T_{\text{left}}) + s(T_{\text{right}}) & \text{otherwise} \end{cases}$$

### Number of nodes

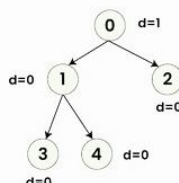
$$N = 2^{h(T)+1} - 1$$

#### Strict (or Full)

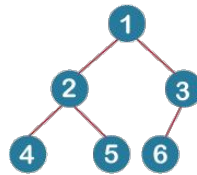


Each node has 0 or 2 children

#### Balanced Binary Tree



#### Complete



Each before the last level is filled

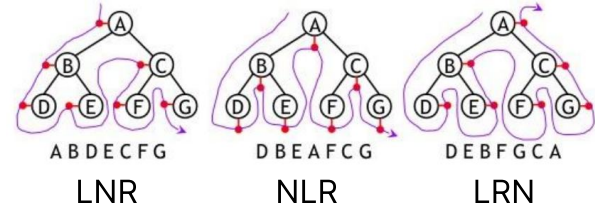
d = Depth of left subtree - Depth of right subtree = 0 or 1

## Trees

Imagine a flag attached to each node:

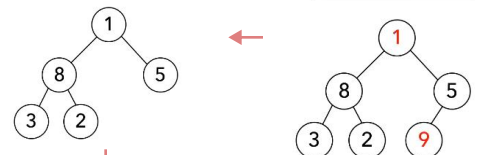
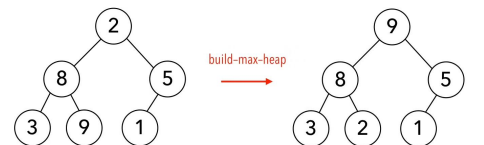


To traverse the tree, collect the flags:



## Heap Sort

### Top Down (siftDown) Approach:



heapify and repeat until sorted

### Bottom up (siftUp)

Pick last unsorted node, heapify up. Repeat until sorted.

