

# Homework 7

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## Remarks

Made a mistake while trying to convert markdown to tex. So there are some wrong feedlines for formulas.  
For example:

$$A \vdash A \vee B, B \vee A$$

(T1)

Means:

$$A \vdash A \vee B, B \vee A \quad (T1)$$

## 1 Question

Prove that:

$$(T1) \quad A \vdash A \vee B, B \vee A$$

$$(T2) \quad A \vee B \vdash \neg \neg B \vee A$$

$$(T3) \quad (A \vee B) \vee C \vdash \neg \neg A \vee (B \vee C)$$

$$(T4) \quad (A \vee B) \vdash \neg \neg \neg A \rightarrow B$$

$$(T5) \quad A \rightarrow B \vdash \neg \neg \neg A \vee B$$

$$(T6) \quad \neg(A \vee B) \vdash \neg \neg \neg A \wedge \neg B$$

$$(T7) \quad \neg(A \wedge B) \vdash \neg \neg \neg A \vee \neg B$$

$$(T8) \quad \emptyset \vdash A \vee \neg A$$

## 2 Lemma

First we prove the following lemmas to simplify the proof.

### 2.1 2.1. belong to ( $\in$ ) (in)

If:

$$A \in \Sigma$$

Then:

$$\Sigma \vdash A$$

Proof:

Let  $\Sigma'$  be  $\Sigma - A$ , then:

$$(Ref) \quad A \vdash A$$

$$(+) \quad \Sigma', A \vdash A$$

$$(1,2) \quad \Sigma \vdash A$$

## 2.2 2.2. neg imply

$$\neg A \rightarrow B \vdash \neg B \rightarrow A$$

proof:

$$\neg A \rightarrow B, \neg A, \neg B \vdash \neg A \rightarrow B$$

(in)

$$\neg A \rightarrow B, \neg A, \neg B \vdash \neg A$$

(in)

From 1, 2,  $\rightarrow -$ :

$$\neg A \rightarrow B, \neg A, \vdash B$$

$$\neg A \rightarrow B, \neg A, \neg B \vdash B$$

(3)

$$\neg A \rightarrow B, \neg A, \neg B \vdash \neg B$$

(in)

From  $\neg -$ , 4, 5 :

$$\neg A \rightarrow B, \neg B \vdash A$$

From 6, *rightarrow+*:

$$\neg A \rightarrow B \vdash \neg B \rightarrow A$$

## 2.3 2.3. transitivity1 (trans)

$$A \rightarrow B, B \rightarrow C \vdash A \rightarrow C$$

(trans)

Proof:

$$A \rightarrow B, B \rightarrow C, A \vdash A$$

$$A \rightarrow B, B \rightarrow C, A \vdash A \rightarrow B$$

$$A \rightarrow B, B \rightarrow C, A \vdash B$$

$$A \rightarrow B, B \rightarrow C, A \vdash B \rightarrow C$$

$$A \rightarrow B, B \rightarrow C, A \vdash C$$

$$A \rightarrow B, B \rightarrow C \vdash A \rightarrow C$$

## 2.4 2.4. transitivity2 (Tr)

IF:

$$\Sigma \vdash \Sigma'$$

$$\Sigma' \vdash A$$

Then:

$$\Sigma \vdash A$$

(Tr)

## 2.5 2.5. double neg

	$\neg\neg A \vdash A$
(2neg)	$A \vdash \neg\neg A$
(2neg)	
Proof1:	
	$\neg\neg A, \neg A \vdash A$
	$\neg\neg A, \neg A \vdash \neg A$
From 1, 2, $\neg$ —:	
	$\neg\neg A \vdash A$
Proof2:	
	$\neg\neg\neg A \vdash \neg A$
(2neg)	$A, \neg\neg\neg A \vdash \neg A$
(+)	$A, \neg\neg\neg A \vdash A$
(in)	
From $\neg$ —:	
	$A \vdash \neg\neg A$

## 2.6 2.6. reductio

If:	
	$\Sigma, A \vdash B$
	$\Sigma, A \vdash \neg B$
Then:	
	$\Sigma \vdash \neg A$
(redct)	
Proof:	
	$\Sigma, A \vdash B$
	$\neg\neg A \vdash A$
From 1, 2, Tr:	
	$\Sigma \neg\neg A \vdash B$
Similarly:	
	$\neg\neg A \vdash A$
	$\Sigma, A \vdash \neg B$
	$\Sigma \neg\neg A \vdash \neg B$
Finally, from 3, 6, $\neg$ —:	
	$\Sigma \vdash \neg A$

## 2.7 2.7. confliction

$A, \neg A \vdash B$   
 (conf)  
 $A \vdash \neg A \rightarrow B$   
 (conf)  
 $\neg A \vdash A \rightarrow B$   
 (conf)  
 Proof:  
 $A, \neg A, \neg B \vdash A$   
 (in)  
 $A, \neg A, \neg B \vdash \neg A$   
 (in)  
 $A, \neg A \vdash B$   
 (1,2)  
 From 3,  $\rightarrow +$ :  
 $A \vdash \neg A \rightarrow B$   
 $\neg A \vdash A \rightarrow B$

## 2.8 2.8. and

$A \wedge B \vdash \neg A, B$   
 Proof:  
 $A \wedge B \vdash A \wedge B$   
 (in)  
 From 1,  $\wedge -$ :  
 $A \wedge B \vdash A$   
 $A \wedge B \vdash B$   
 $A \wedge B \vdash A, B$   
 And:  
 $A, B \vdash A$   
 $A, B \vdash B$   
 Then use  $\wedge +$  rule;  
 $A, B \vdash A \wedge B$

## 2.9 2.9. switch

If:

$$A \vdash \neg A'$$

$$\vdash A \rightarrow B$$

$$\vdash C \rightarrow A$$

$$\vdash A \vee D$$

$$\vdash A \wedge E$$

$$\vdash \neg A$$

Then it is trivial that:

$$\vdash A' \rightarrow B$$

$$\vdash C \rightarrow A'$$

$$\vdash A' \vee D$$

$$\vdash A' \wedge E$$

$$\vdash \neg A'$$

Therefore, for A, A' in deduction formula, we can switch A and A'.

## 3 Proof

### 3.1 3.1. T1

$$A \vdash A$$

(Ref)

$$A \vdash A \vee B$$

(1,V+)

$$A \vdash B \vee A$$

(1,V+)

$$A \vdash A \vee B, B \vee A$$

(3,4)

### 3.2 3.2. T2

To prove:

$$A \vee B \vdash \neg B \vee A$$

(T2)

We only need to prove one direction, for this formula is symmetric.

$$A \vdash B \vee A$$

(T1)

$$B \vdash B \vee A$$

(T1)

From 1,2,V-:

$$A \vee B \vdash B \vee A$$

Symmetrically:

$$B \vee A \vdash A \vee B$$

As a result:

$$A \vee B \vdash \neg B \vee A$$

(4,5)

### 3.3 3.3. T3

$$C \vdash (B \vee C)$$

(T1)

$$(B \vee C) \vdash A \vee (B \vee C)$$

(T1)

$$C \vdash A \vee (B \vee C)$$

(1,2)

Similarly, we have:

$$A \vdash A \vee (B \vee C)$$

$$B \vdash A \vee (B \vee C)$$

Then from 4,5,V-:

$$A \vee B \vdash A \vee (B \vee C)$$

Finally:

$$(A \vee B) \vee C \vdash A \vee (B \vee C)$$

(3,6)

### 3.4 3.4. T4

Left to Right:

$$A, \neg A \vdash B$$

(conf)

$$A \vdash \neg A \rightarrow B$$

(conf)

$$B, \neg A \vdash B$$

(in)

From 3,  $\rightarrow +$ :

$$B \vdash \neg A \rightarrow B$$

From 3, 4,  $\vee -$ :

$$(A \vee B) \vdash \neg A \rightarrow B$$

Right to Left:

$$A \vdash A \vee B$$

( $\vee +$ )

From 1, lemma neg imply:

$$\neg(A \vee B) \vdash \neg A$$

$$\neg A \rightarrow B, \neg(A \vee B) \vdash \neg A$$

( $+$ )

$$\neg A \rightarrow B, \neg(A \vee B) \vdash \neg A \rightarrow B$$

(in)

From  $\rightarrow -$ :

$$\neg A \rightarrow B, \neg(A \vee B) \vdash B$$

Using  $\vee +$ :

$$\neg A \rightarrow B, \neg(A \vee B) \vdash (A \vee B)$$

Using in:

$$\neg A \rightarrow B, \neg(A \vee B) \vdash \neg(A \vee B)$$

(in)

From  $\neg -$ :

$$\neg A \rightarrow B \vdash (A \vee B)$$

To sum up:

$$(A \vee B) \vdash \neg A \rightarrow B$$



### 3.5 3.5. T5

$$(\neg A \vee B) \vdash \neg(\neg A) \rightarrow B$$

(T4)

$$\neg\neg A \vdash A$$

From 1, 2, switch:

$$\neg A \vee B \vdash A \rightarrow B$$

### 3.6 3.6. T6

Consider that:

$$\neg(\neg A \rightarrow B), A \vdash \neg(\neg A \rightarrow B)$$

(in)

From :

$$A, \neg A \vdash B$$

Using  $\rightarrow +$ , we gain:

$$\neg(\neg A \rightarrow B), A \vdash (\neg A \rightarrow B)$$

Then use  $\neg -$ :

$$\neg(\neg A \rightarrow B) \vdash \neg A$$

$$\neg(\neg A \rightarrow B) \vdash \neg B$$

Since:

$$\neg(\neg A \rightarrow B) \vdash \neg A$$

$$\neg(\neg A \rightarrow B) \vdash \neg B$$

Then:

$$\neg(\neg A \rightarrow B) \vdash \neg A \wedge \neg B$$

And:

$$\neg(\neg A \rightarrow B) \vdash \neg(\neg A \rightarrow B)$$

(in)

From T4, switch:

$$\neg(A \vee B) \vdash \neg(\neg A \rightarrow B)$$

$$\neg(A \vee B) \vdash \neg A \wedge \neg B$$

(Tr)

Another direction:

$$\neg A \wedge \neg B \vdash \neg A, \neg B$$

(and)

$$\neg A, \neg B, (A \vee B) \vdash \neg B$$

(in)

$$\neg A, \neg B, (A \vee B) \vdash \neg A$$

(in)

$$\neg A, \neg B, (A \vee B) \vdash \neg A \rightarrow B$$

(T4)

From  $\rightarrow \neg$ :

$$\neg A, \neg B, (A \vee B) \vdash B$$

Then from  $\neg \neg$ :

$$\neg A, \neg B \vdash \neg(A \vee B)$$

Using Tr:

$$\neg A \wedge \neg B \vdash \neg(A \vee B)$$

As a result:

$$\neg(A \vee B) \vdash \neg \neg A \wedge \neg B$$

### 3.7 3.7. T7

$$\neg(A \wedge B) \vdash A \rightarrow \neg B$$

And:

$$\neg A \vee \neg B \vdash \neg \neg A \rightarrow \neg B$$

(T4)

From 2neg, switch:

$$\neg A \vee \neg B \rightarrow \neg \neg A \rightarrow \neg B$$

As a result (Tr):

$$\neg(A \wedge B) \vdash \neg A \vee \neg B$$

Another direction:

We have proved that:

$$\neg A \vee \neg B \rightarrow \neg \neg A \rightarrow \neg B$$

For (we have proved it many times):

$$A \rightarrow \neg B, A \vdash \neg B$$

Using Tr, and:

$$\neg A \vee \neg B, (A \wedge B) \vdash \neg B$$

From in, and:

$$\neg A \vee \neg B, (A \wedge B) \vdash B$$

Using  $\neg \neg$ :

$$\neg A \vee \neg B \vdash \neg(A \wedge B)$$

Both 2 directions have been proved:

$$\neg(A \wedge B) \vdash \neg A \vee \neg B$$

(T7)

### 3.8 3.8. T8

$$\neg A \vdash \neg A$$

(in)

From  $\rightarrow +$ :

$$\emptyset \vdash \neg A \rightarrow \neg A$$

From T4:

$$\neg A \rightarrow \neg A \vdash A \vee \neg A$$

As a result:

$$\emptyset \vdash A \vee \neg A$$

(1,3,Tr)